CKY Parsing

Ling571
Deep Processing Approaches to NLP
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Roadmap

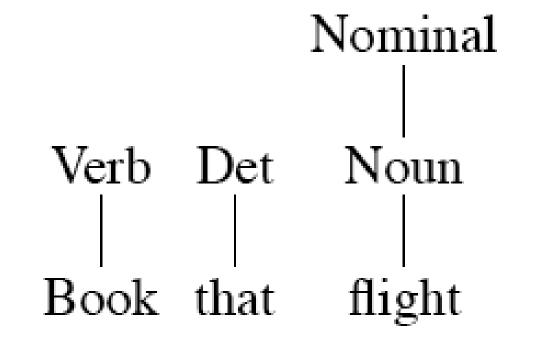
- Motivation:
 - Inefficiencies of parsing-as-search
- Strategy: Dynamic Programming
- Chomsky Normal Form
 - → Weak and strong equivalence
- CKY parsing algorithm

Bottom-Up Parsing

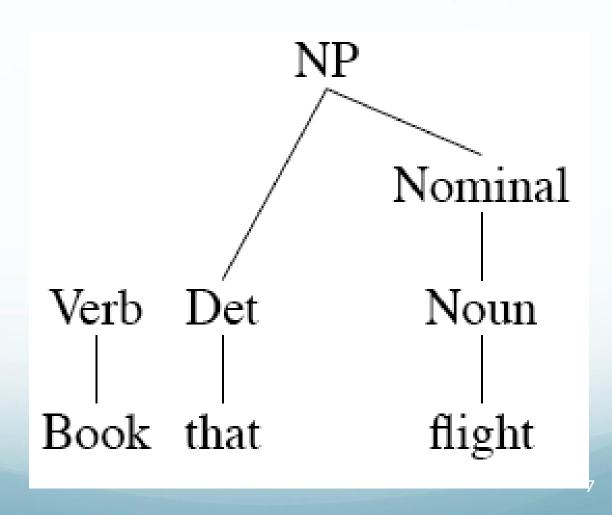
- Try to find all trees that span the input
 - → Start with input string
 - → Book that flight.
 - Use all productions with current subtree(s) on RHS
 - \dashv E.g., N $\rightarrow \rightarrow$ Book; V $\rightarrow \rightarrow$ Book
 - Stop when spanned by S (or no more rules apply)

Book that flight

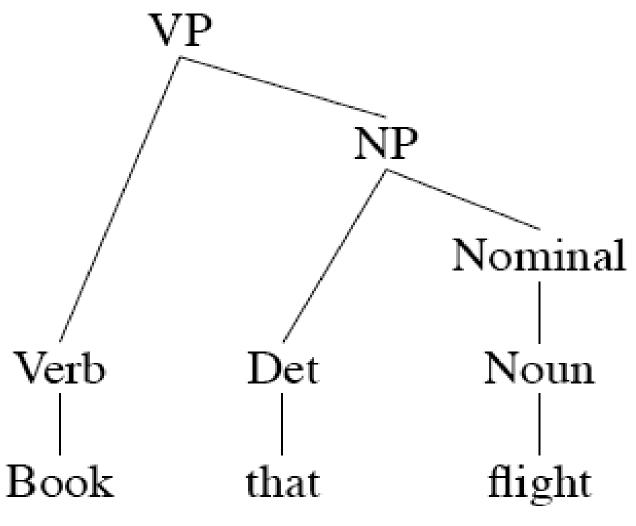
Speech and Language Processing



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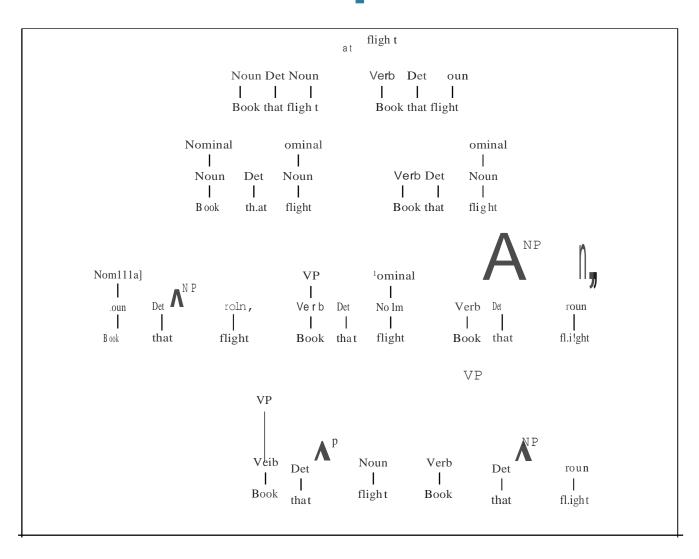


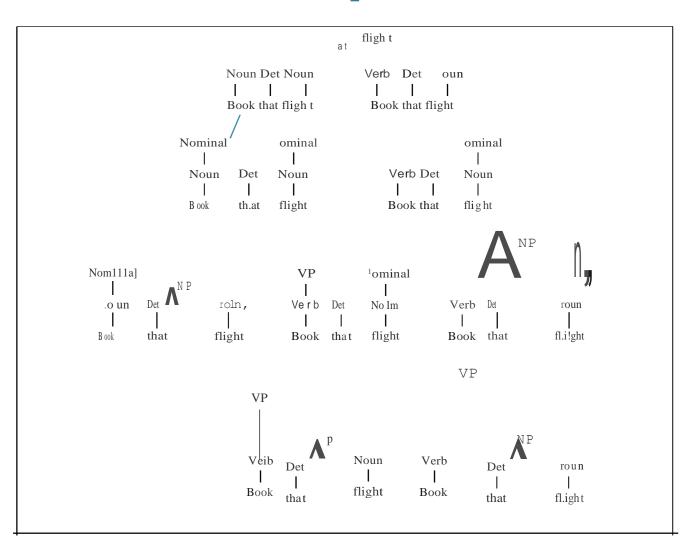
Jurafsky and Martin

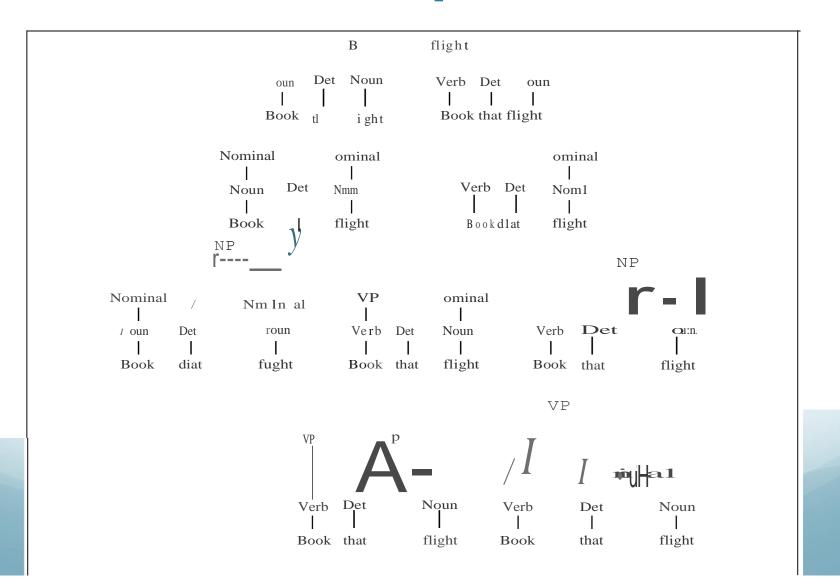


Jurafsky and Martin

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Pros and Cons of Bottom-Up Search

Pros:

- → Will not explore trees that don't match input
- Recursive rules less problematic
- Useful for incremental/ fragment parsing

Cons:

Explore subtrees that will not fit full sentences

Parsing Challenges

Ambiguity

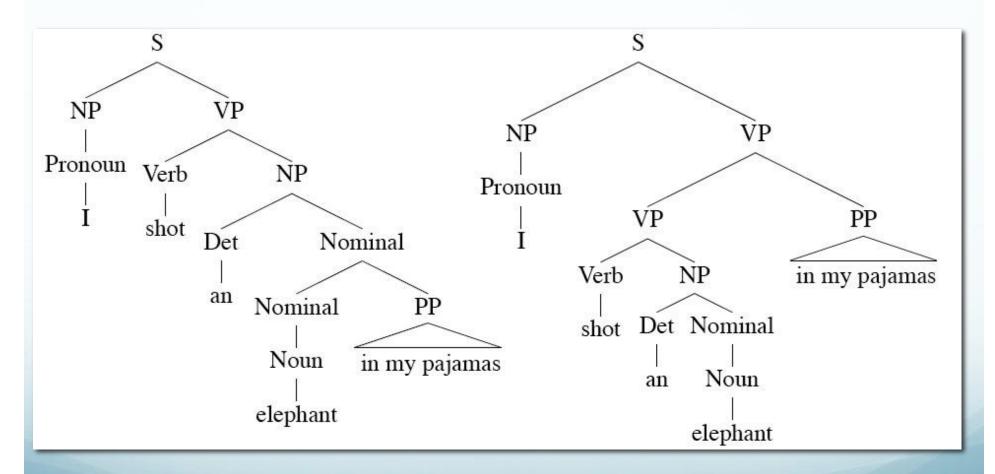
Repeated substructure

Recursion

Parsing Ambiguity

- Many sources of parse ambiguity
 - → Lexical ambiguity
 - → Book/N; Book/V
 - Structural ambiguity: Main types:
 - Attachment ambiguity
 - Constituent can attach in multiple places
 - \dashv I shot an elephant in my pyjamas.
 - Coordination ambiguity
 - Different constituents can be conjoined
 - Old men and women

Ambiguity

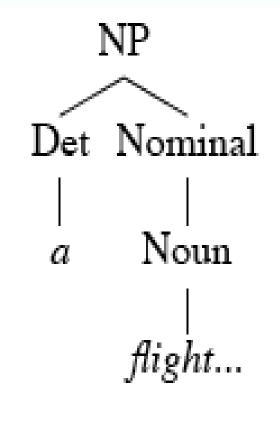


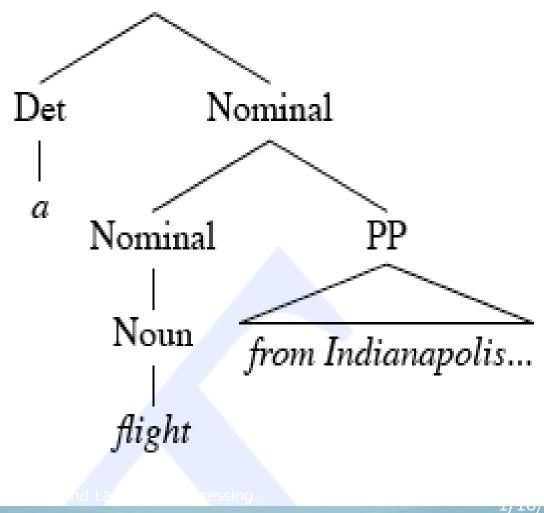
Disambiguation

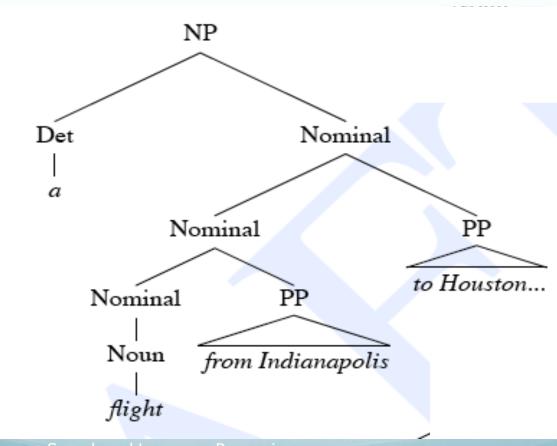
- Global ambiguity:
 - → Multiple complete alternative parses
 - → Need strategy to select correct one
 - Approaches exploit other information
 - Statistical
 - Some prepositional structs more likely to attach high/low
 - ── Some phrases more likely, e.g., (old (men and women))
 - Semantic
 - Pragmatic
 - → E.g., elephants and pyjamas
 - → Alternatively, keep all
- Local ambiguity:
 - Ambiguity in subtree, resolved globally

Repeated Work

- Top-down and bottom-up parsing both lead to repeated substructures
 - Globally bad parses can construct good subtrees
 - But overall parse will fail
 - Require reconstruction on other branch
 - → No static backtracking strategy can avoid
- Efficient parsing techniques require storage of shared substructure
 - Typically with dynamic programming
- Example: a flight from Indianapolis to Houston on TWA

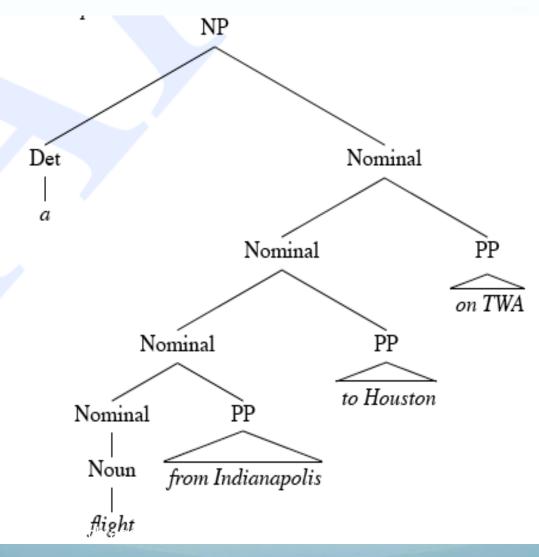






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Recursion

- Many grammars have recursive rules
 - \dashv E.g., S $\rightarrow \rightarrow$ S Conj S
- In search approaches, recursion is problematic
 - Can yield infinite searches
 - Esp., top-down

Dynamic Programming

- → Challenge: Repeated substructure → Repeated work
- → Insight:
 - Global parse composed of parse substructures
 - Can record parses of substructures
- Dynamic programming avoids repeated work by tabulating solutions to subproblems
 - Here, stores subtrees

Parsing w/Dynamic Programming

- Avoids repeated work
- Allows implementation of (relatively) efficient parsing algorithms
 - → Polynomial time in input length
 - \dashv Typically cubic (n^3) or less
- Several different implementations
 - Cocke-Kasami-Younger (CKY) algorithm
 - Earley algorithm
 - Chart parsing

Chomsky Normal Form (CNF)

- CKY parsing requires grammars in CNF
- Chomsky Normal Form
 - → All productions of the form:
 - \dashv A $\rightarrow \rightarrow$ B C, or
 - A $\rightarrow \rightarrow$ a
- However, most of our grammars are not of this form
 - -1 E.g., S $\rightarrow \rightarrow$ Wh-NP Aux NP VP
- Need a general conversion procedure
 - Any arbitrary grammar can be converted to CNF

Grammar Equivalence and Form

Grammar equivalence

- Weak: Accept the same language, May produce different analyses
- Strong: Accept same language, Produce same structure

CNF Conversion

- Three main conditions:
 - Hybrid rules:
 - -INF-VP $\rightarrow \rightarrow$ to VP
 - Unit productions:
 - A $\rightarrow \rightarrow$ B
 - Long productions:
 - A $\rightarrow \rightarrow$ B C D

CNF Conversion

- Hybrid rule conversion:
 - Replace all terminals with dummy non-terminals
 - \dashv E.g., INF-VP $\rightarrow \rightarrow$ to VP
 - \dashv INF-VP $\rightarrow \rightarrow$ TO VP; TO $\rightarrow \rightarrow$ to
- Unit productions:
 - Rewrite RHS with RHS of all derivable non-unit productions
 - \dashv If $A \Rightarrow B$ and $B \rightarrow \Rightarrow$ w, then add $A \rightarrow \Rightarrow$ w

CNF Conversion

- Long productions:
 - Introduce new non-terminals and spread over rules
 - S $\rightarrow \rightarrow$ Aux NP VP
 - \neg S $\rightarrow \rightarrow$ X1 VP; X1 $\rightarrow \rightarrow$ Aux NP
- For all non-conforming rules,
 - Convert terminals to dummy non-terminals
 - Convert unit productions
 - Binarize all resulting rules

2 $_1$ Grammar	S1i in CNF
S+ IP VP	$S \longrightarrow l \backslash TP \ VP$
S+ Aux NP VP	S+ <i>XI VP</i>
	X1+ AuxTP
S+ <i>VP</i>	S+ book include preje r
	S+ VerbNP
	S+ <i>X2PP</i>
	S+VerbPP
	S+ VPPP
NP+ Pronoun	NP - + I + she + rne
l \TP \rightarrow $Proper-Noun$	$i \backslash TP$ + TWA $\mid Houston$
$l \land TP \rightarrow D \ e \ t \ Nominal$	$i \backslash TP$ + Det $l \backslash Tominal$
Norninal+ Noun	Nominal+ book flight meal money
Norninal+ Nominal Noun	Nominal+ Nominal Noun
$l \setminus Torninal$ + $l \setminus Tominal$ PP	$i \backslash Tominal ext{+ } Nominal PP$
VP+ Verb	VP+ book include prefer
VP+ Verb NP	$VP ext{+ } Verb NP$
VP+ Verb NP PP	VP+ X2 PP
	X2+ Verb NP
$VP \rightarrow Verb PP$	VP+ Verb PP
VP+ $VPPP$	VP+ $VPPP$
$PP ext{+} Preposition i \ TP$	$PP ext{+ } Preposition \ l \backslash TP$

CKY Parsing

- Cocke-Kasami-Younger parsing algorithm:
 - (Relatively) efficient bottom-up parsing algorithm based on tabulating substring parses to avoid repeated work
 - Approach:
 - Use a CNF grammar
 - \dashv Build an (n+1) x (n+1) matrix to store subtrees
 - Upper triangular portion
 - Incrementally build parse spanning whole input string

Dynamic Programming in CKY

→ Key idea:

- → For a parse spanning substring [i,j], there exists some k such there are parses spanning [i,k] and [k,j].
 - We can construct parses for whole sentence by building up from these stored partial parses

─ So,

- \dashv To have a rule A $\rightarrow \rightarrow$ B C in [i,j],
 - ── We must have B in [i,k] and C in [k,j], for some i<k<j</p>
 - → CNF grammar forces this for all j>i+1