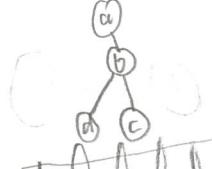


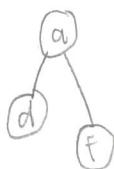
1.

Solution: None of subgraphs have same substructure

 : present in $\frac{3}{4}$ subgraphs, not present in all

E: labeled w/ +1 (not -1) b/c present in majority of decision stumps

An ensemble classifier would look at each subgraph's decision stump and if the subgraphs have the substructure for this instance, it receives a +1. Taken together, the $4 \times$ subgraph doesn't have the substructure.

new graph:  , since  isn't in there its -1

+1 for other 3 subgraphs

Problem 2 Output

Summary motif results

=====

mfinder Version 1.20

MOTIF FINDER RESULTS:

Network name: problem2_input.txt

Network type: Directed

Num of Nodes: 280 Num of Edges: 2170

Num of Nodes with edges: 280

Maximal out degree (out-hub) : 39

Maximal in degree (in-hub) : 45

Roots num: 13 Leaves num: 21

Single Edges num: 1776 Mutual Edges num: 197

Motif size searched 4

Total number of 4-node subgraphs : 875316

Number of random networks generated : 10

Random networks generation method: Switches

Num of Switches range: 100.0-200.0, Success switches Ratio:0.762+-0.00

The following motifs were found:

Criteria taken : Nreal Zscore > 2.00

Pval ignored (due to small number of random networks)

Mfactor > 1.10

Uniqueness >= 4

Top Motifs List

| MOTIF | NREAL | NRAND | | NREAL | NREAL | UNIQ | CREAL |
|-------|-------|-------|--|--------|-------|------|--------|
| ID | | STATS | | ZSCORE | PVAL | VAL | [MILI] |

| | | | | | | | |
|-----|------|---------------|------|-------|----|------|--|
| 204 | 2950 | 2192.8+-101.5 | 7.46 | 0.000 | 32 | 3.37 | |
|-----|------|---------------|------|-------|----|------|--|

0 0 1 1

0 0 1 1

0 0 0 0

0 0 0 0

| | | | | | | | |
|-----|------|------------|--|--------|-------|----|------|
| 222 | 1606 | 73.1+-12.6 | | 121.74 | 0.000 | 17 | 1.83 |
|-----|------|------------|--|--------|-------|----|------|

0 1 1 1

1 0 1 1

0 0 0 0

0 0 0 0

| | | | | | | | |
|-----|------|--------------|-------|-------|----|------|--|
| 904 | 3027 | 1972.4+-61.2 | 17.24 | 0.000 | 27 | 3.46 | |
|-----|------|--------------|-------|-------|----|------|--|

0 0 0 1

0 0 0 1

1 1 0 0

0 0 0 0

| | | | | | | | |
|-----|------|-------------|-------|-------|----|------|--|
| 906 | 1905 | 715.7+-79.0 | 15.05 | 0.000 | 22 | 2.18 | |
|-----|------|-------------|-------|-------|----|------|--|

0 1 0 1

0 0 0 1

1 1 0 0

0 0 0 0

922 2143 98.3+-34.2 59.79 0.000 18 2.45

0 1 0 1

1 0 0 1

1 1 0 0

0 0 0 0

2252 1447 584.4+-62.7 13.75 0.000 28 1.65

0 0 1 1

0 0 1 1

0 0 0 1

0 0 0 0

6552 803 60.9+-20.9 35.48 0.000 11 0.92

0 0 0 1

1 0 0 1

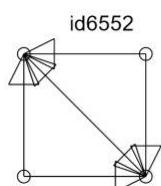
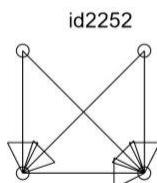
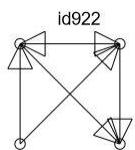
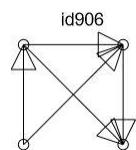
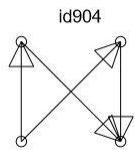
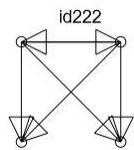
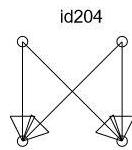
1 0 0 1

1 0 0 0

(Total No. of non-dangling motifs : 75)

Full list includes 102 motifs

Top Motifs Graphs:



3.

Solution:similarities $\leq \alpha = 0.5$ (solid lines)similarities $\geq \beta = 0.6$ (dashed lines)

construct meta-graph $\Gamma(M)$ to measure sim. between graph patterns in M . Each node is a maximal subgraph. Edge exists between 2 nodes if their sim. is $\leq \alpha = 0.5$

Find maximal cliques:

$$R_1 = \{M_1, M_2, M_3, M_4\}$$

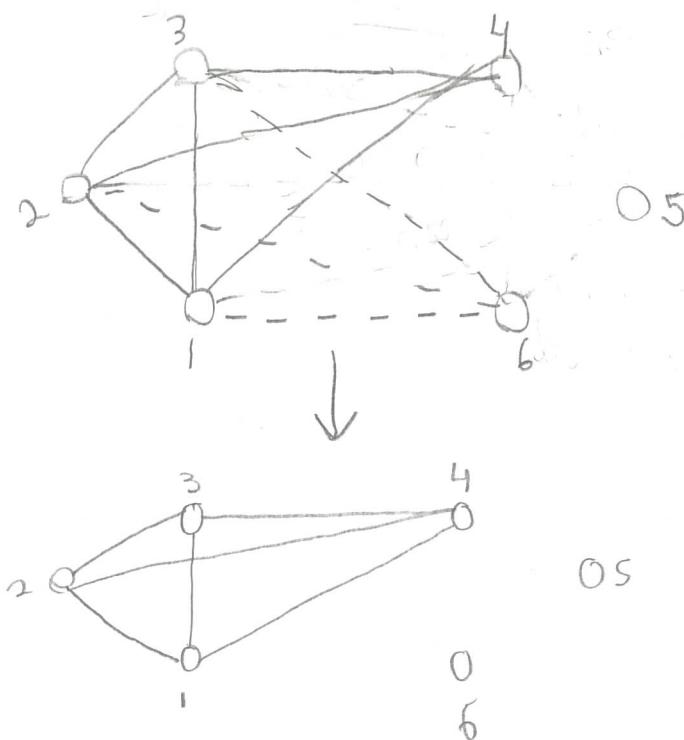
~~$R_2 = \{M_1, M_2, M_3\}$~~

~~$R_3 = \{M_1, M_2, M_4\}$~~

~~$R_4 = \{M_1, M_3, M_4\}$~~

R_2, R_3, R_4 not needed b/c

R_1 captures them



$$\begin{aligned} \delta(R_1, D) &= \{G \mid G \in D \text{ and there exists } G_a \in \{M_1, M_2, M_3, M_4\}, \text{sim}(G_a, G) \geq 0.6\} \\ &= \{M_6\} \end{aligned}$$

Then $\Delta(R_1, D) = D - (R_1 \cup \delta(R_1, D)) = \{M_5\}$ ← unrepresented by R_1

$$\begin{aligned} \delta(R_2, D) &= \{G \mid G \in D \text{ and there exists } G_a \in \{M_1, M_2, M_3\}, \text{sim}(G_a, G) \geq 0.6\} \\ &= \{M_6\} \end{aligned}$$

Then $\Delta(R_2, D) = D - (R_2 \cup \delta(R_2, D)) = \{M_4, M_5\}$

$$\delta(R_3, D) = \{G \mid G \in D \text{ and there exists } G_a \in \{M_1, M_2, M_4\}, \text{sim}(G_a, G) \geq 0.6\}$$

~~$\Delta(R_3, D) = D - (R_3 \cup \delta(R_3, D)) = \{M_1, M_2, M_3\}$~~

3.]

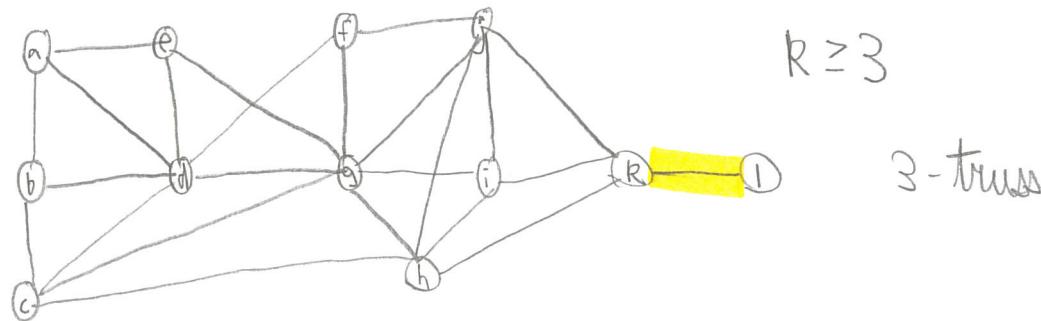
$$\begin{aligned} \delta(R_4, D) &= \{G \mid G \in D \text{ and there exists } G_a \in \{\textcircled{M}_1, \textcircled{M}_3, M_4\}, \text{sim}(G_a, G) \geq 0.6\} \\ &= \{M_6\} \end{aligned}$$

Then $\Delta(R_4, D) = D - (R_4 \cup \delta(R_4, D)) = \{M_2, M_5\}$

\therefore The optimal α -orthogonal, β -representative set is R .

4.

Solution:

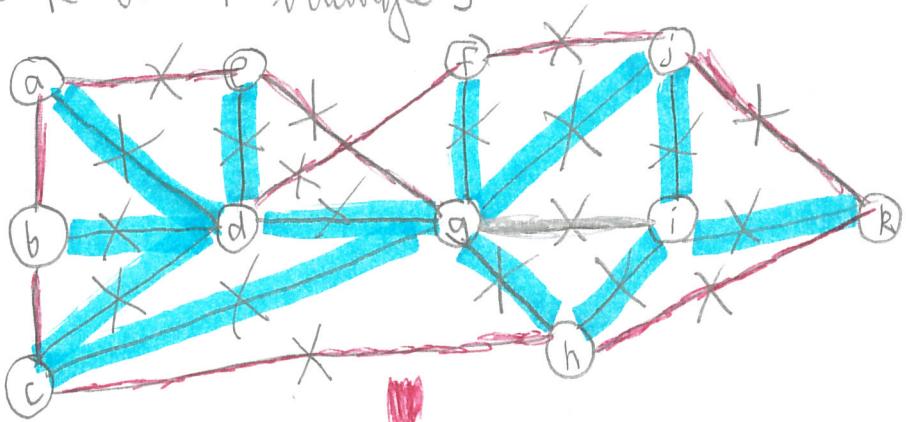


$$R = 3$$

$$(R-2)$$

Look for edges w/ support ≤ 1 : (i, k) i.e., not involved in any triangles
 Removing that edge doesn't cause support of any other edge to change to ≤ 1

Resulting graph is the 3-truss (every remaining edge is contained in $\geq R-2 = 1$ triangle)



$$R = 4$$

no 4-truss

all remaining edges would be removed
 graph is empty, no more trusses

Edges w/ support ≤ 2 : (a,b), (b,c), (c,h), (a,e), (e,g), (d,f), (f,j), (h,k), (j,k) i.e., involved in only 1 triangle

Removing those edges causes support of some other edges to become ≤ 2 : (a,d), (b,d), (c,d), (e,d), (c,a), (d,g), (g,f), (g,h), (g,j), (h,i), (j,i), (i,k)

Removal of those edges causes support of other edges to change to ≤ 2 : (g,i)

5.1 a)

| | degree | neighbor, degree sequence |
|----------|--------|---------------------------|
| g_1 | 2 | 4,4 |
| g_2 | 6 | 4,4,3,5,4,2 |
| g_3 | 4 | 2,6,4,2 |
| g_4 | 4 | 2,6,4,3 |
| g_5 | 3 | 6,4,4 |
| g_6 | 2 | 5,4 |
| g_7 | 5 | 6,3,4,2,2 |
| g_8 | 4 | 6,3,2,5 |
| g_9 | 2 | 4,5 |
| g_{10} | 2 | 6,5 |

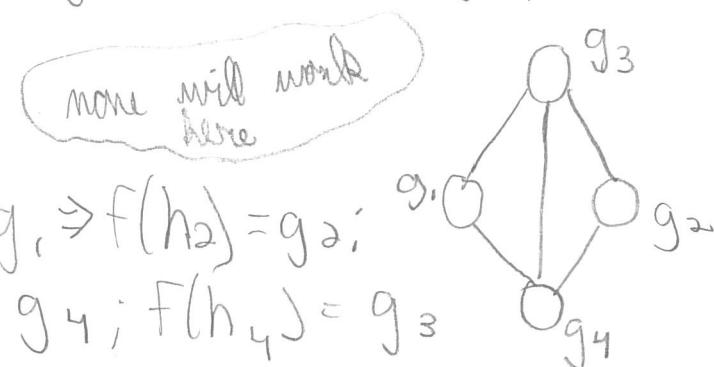
$$g = g_1, h = h_1, f(h_1) = g_1 \Rightarrow f(h_2) = g_2; f(h_3) = g_3$$

$h_1 < h_2, h_3 < h_4;$
 $\checkmark g_1 < g_2, g_3 = g_4 \times$ none will work here

$f(h_1) = g_1 \Rightarrow f(h_2) = g_2;$
 $f(h_3) = g_4; f(h_4) = g_3$

$\checkmark g_1 < g_2, g_4 = g_3 \times$

$g = g_2, h = h_1, f(h_1) = g_2 \Rightarrow f(h_2) = g_6; f(h_3) = g_7$
 $h_1 < h_2, h_3 < h_4;$
 $g_2 < g_6, g_7 < g_8 \times$



5.1 A

$$g = g_5 \quad h = h_1 \quad f(h_1) = g_5 \Rightarrow f(h_2) = g_2; f(h_3) = g_8 \\ f(h_4) = g_7$$

$$h_1 < h_2, h_3 < h_4$$

$$g_5 < g_2, g_8 < g_7$$

only one that will work b/c
 $g_2 > g_5$ and $g_7 > g_8$

$$g = g_4 \quad h = h_1 \quad f(h_1) = g_4 \Rightarrow f(h_2) = g_8; f(h_3) = g_2 \\ f(h_4) = g_5$$

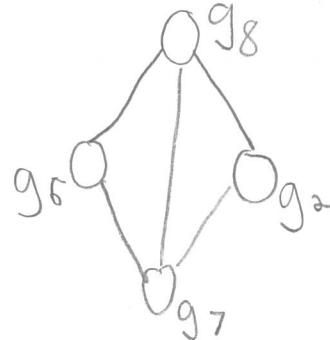
$$h_1 < h_2, h_3 < h_4$$

$$g_4 = g_8, g_2 > g_5$$

X X

∴ none in this substructure will work because g_4 and g_8 have equal degrees

instance found:



frequency of this pattern in the input graph is 1.

b) Z score (S) = $\{ \text{frequency}(S, G) - \text{meanFrequency}(S, G) \} / \sigma$

where σ = standard deviation

$$\begin{aligned} \text{Z score } (S) &= \{ 1 - 2 \} / 0.35 \\ &= -2.85714286 \end{aligned}$$

Prvalno. (S) = (# random graphs in which S occurred more often than in

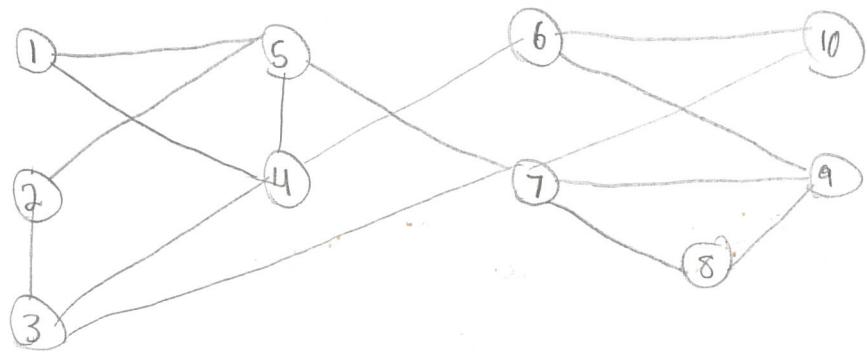
S.1) $\Pr_{\text{H}_0}(\text{Z} \geq 2.8571) = 0.0045$

$$\text{Z-score} < 2 \rightarrow |G_r| = 750 < 1000$$

$$-2.8571 < 2$$

$\therefore \text{H}_0$ wouldn't be considered a motif of G. The calculated Z-score is significantly less than 2.

6.]



d_i = # of neighbors of node v_i :

R_i = # edges among these neighbors

C_i = clustering coefficient of v_i :

$$C_i = \begin{cases} \frac{R_i}{d_i \times (d_i - 1)/2} & d_i > 1 \\ 0 & d_i = 0 \text{ or } 1 \end{cases}$$

The higher the value of C_i , the more likely v_i and its neighbors resemble a clique.

Let v_1 be the node labeled 1; its neighbors are 4 and 5.

$$d_1 = 2$$

$$R_1 = 1$$

$$C_1 = R_1 / (d_1 \cdot (d_1 - 1)/2) = 1 / (2 \cdot (2 - 1)/2) = 1 / (2 \cdot \frac{1}{2}) = 1$$

v_2 's neighbors are 3 and 5

$$d_2 = 2$$

$$R_2 = 0$$

$$C_2 = 0 / (2 \cdot (2 - 1)/2) = 0$$

v_3 's neighbors are 2, 4, 7

$$d_3 = 3$$

$$R_3 = 0$$

$$C_3 = 0 / (3 \cdot (3 - 1)/2) = 0$$

6.1

V_4 's neighbors are 3, 5, 6, and 8

$$d_4 = 4$$

$$R_4 = 0$$

$$C_4 = 0 / (4 \cdot (4-1)/2) = 0$$

V_5 's neighbors are 1, 2, 4, 7

$$d_5 = 4$$

$$R_5 = 1$$

$$C_5 = 1 / (4 \cdot (4-1)/2) = 1 / (4 \cdot 3/2) = 1/6 = 0.\overline{16}$$

V_6 's neighbors are 4, 9, 10

$$d_6 = 3$$

$$R_6 = 0$$

$$C_6 = 0$$

V_7 's neighbors are 5, 8, 9, 10

$$d_7 = 4$$

$$R_7 = 1$$

$$C_7 = 1 / (4 \cdot (4-1)/2) = 1 / (4 \cdot 3/2) = 1/6 = 0.\overline{16}$$

V_8 's neighbors are 4, 7, 9

$$d_8 = 3$$

$$R_8 = 1$$

$$C_8 = 1 / (3 \cdot (3-1)/2) = 1 / (3 \cdot 1) = \frac{1}{3} = 0.\overline{33}$$

V_9 's neighbors are 6, 7, 8

$$d_9 = 3$$

$$R_9 = 1$$

$$C_9 = 1 / (3 \cdot (3-1)/2) = 0.\overline{33}$$

V_{10} 's neighbors are 6, 7

$$d_{10} = 2$$

$$C = \sum_{i=1}^n C_i/n$$

$$= \frac{1+0+0+0+0.\overline{16}+0+0.\overline{16}+0.\overline{33}+0.\overline{33}+0}{10} \\ = 0.20$$

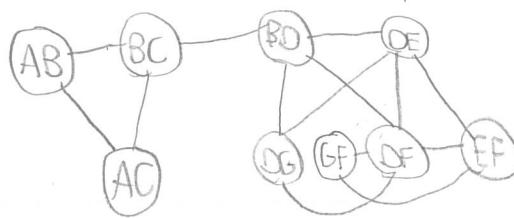
Since the community structure value is low, the nodes are loosely clustered into groups in the network

7.]

a) The interpretation of the edges of the resulting graph is that if Friend X (node X) and Friend Y (node Y) are related in some way through their connecting edge, an edge connecting Friend X and Friend Z also means there will be a relationship in some way for Friend Y and Friend Z through an edge by both of their relationships to Friend X.

b) G' - representing this
edges in G

A-B, A-C, B-C, B-D, D-E, D-F, D-G, E-F, G-F
These will be node pairings in G'

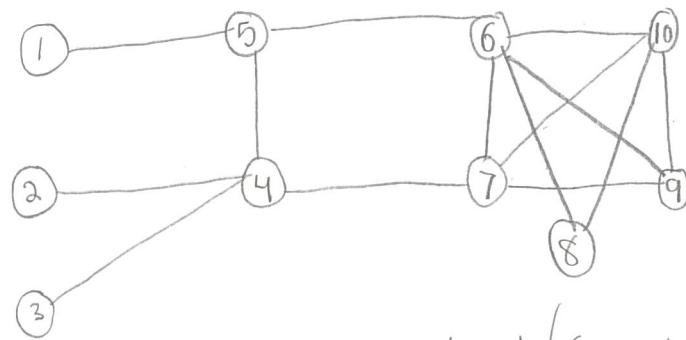


The relationships containing a common letter will all be linked w/ an edge

c) The degree of a node XY in G' compared to the degrees of X and Y in G : up to a certain point, the connections will all relate but as the number of people w/ relationships in the community grows, the network starts to break up into isolated, disconnected components and/or fewer people in each connected component.

The degrees of a node XY in G' should be the same as the degrees X and Y in G . The nodes AB, BC, and AC are a good example to the nodes A, B, and C.

8.



$$V_S = \{6, 7, 8, 9, 10\}$$

$$V_S = \{6, 7, 9, 10\}$$

$$|E_S| / (|V_S| \cdot (|V_S| - 1)/2) = 8/(5 \cdot 4/2)$$

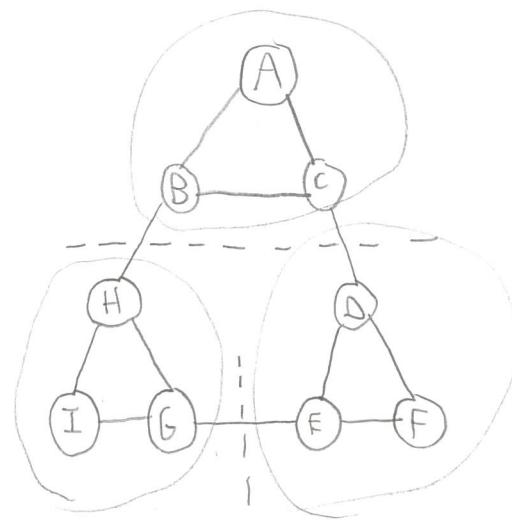
$$= 0.8 > 0.65$$

$$|E_S| / (|V_S| \cdot (|V_S| - 1)/2) = 6/(6 \cdot 5/2)$$

$$= 0.4 < 0.65$$

All other cliques would generate smaller values for their α -density % b/c they will have less edges relative to the number of nodes that are applied to the denominator of the α -density calculation.

$V_S = \{6, 7, 8, 9, 10\}$ is a clique



2 cuts
3 edges removed
3 communities created of equal size

~~Let S be the set of nodes in subgraph S~~

~~Vol(S) = # edges w/ at least one end in S~~

~~Cut(S, T, U) = # edges that connect a node in S to a node in $T \cup U$~~

~~Normalized Cut(S, T, U) = $\frac{\text{Cut}(S, T)}{\text{Vol}(S)} + \frac{\text{Cut}(S, U)}{\text{Vol}(S)}$~~
~~+ $\frac{\text{Cut}(S, T, U)}{\text{Vol}(U)}$~~

~~minimize this~~

~~WRONG~~

~~METHOD~~

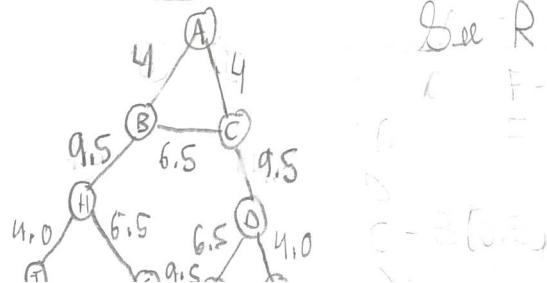
~~Vol(S) = 5, Cut(S, T, U) = 3~~

~~Vol(T) = 5, Cut(S, T, U) = 3~~

~~Vol(U) = 5, Cut(S, T, U) = 3~~

~~Normalized Cuts(S, T, U) = $(3/5) + (3/5) + (3/5) = 1.8$~~

edge betweenness algorithm - Right Method



See R Code

[http://www.johnwittenauer.com/2011/07/01/finding-communities-in-graphs-with-r-and-louvain-algorithm/](#)

[http://www.johnwittenauer.com/2011/07/01/finding-communities-in-graphs-with-r-and-louvain-algorithm/](#)

[http://www.johnwittenauer.com/2011/07/01/finding-communities-in-graphs-with-r-and-louvain-algorithm/](#)

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[http://www.johnwittenauer.com/2011/07/01/finding-communities-in-graphs-with-r-and-louvain-algorithm/](#)

| | A | B | C | D | E | F | G | H | I | |
|---|---|---|---|---|---|---|---|---|---|--|
| A | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | |
| B | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | |
| C | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | |
| D | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | |
| E | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | |
| F | 0 | 0 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | |
| G | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 | 1 | |
| H | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 1 | |
| I | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | |

```

> edges=c(c(0,1,1,0,0,0,0,0,0), c(1,0,1,0,0,0,0,1,0), c(1,1,0,1,0,0,0,0,0), c(0,0,1,0,1,1,0,0,0),
c(0,0,0,1,0,1,1,0,0), c(0,0,0,1,1,0,0,0,0), c(0,0,0,0,1,0,0,1,1), c(0,1,0,0,0,1,0,1), c(0,0,0,0,0,0,1,1,0))

> nodeNames=c("A","B","C","D","E","F","G","H","I")

> adjMatrix=matrix(edges,nrow=9,ncol=9,byrow=TRUE,dimnames=list(nodeNames,nodeNames))

>
>
>
>
>
>
>
>

> adjMatrix

 A B C D E F G H I
A 0 1 0 0 0 0 0 0 0
B 1 0 1 0 0 0 0 1 0
C 1 1 0 1 0 0 0 0 0
D 0 0 1 0 1 1 0 0 0
E 0 0 0 1 0 1 1 0 0
F 0 0 0 1 1 0 0 0 0
G 0 0 0 0 1 0 0 1 1
H 0 1 0 0 0 0 1 0 1
I 0 0 0 0 0 1 1 0 0

> library("igraph")

```

Attaching package: ‘igraph’

The following objects are masked from ‘package:stats’:

decompose, spectrum

The following object is masked from ‘package:base’:

```
union
```

Warning message:

```
package ‘igraph’ was built under R version 3.6.3
```

```
> g=graph_from_adjacency_matrix(adjMatrix, mode="undirected")
```

```
> g
```

```
IGRAPH 1918485 UN-- 9 12 --
```

```
+ attr: name (v/c)
```

```
+ edges from 1918485 (vertex names):
```

```
[1] A--B A--C B--C B--H C--D D--E D--F E--F E--G G--H G--I H--I
```

```
> edge_betweenness(g)
```

```
[1] 4.0 4.0 6.5 9.5 9.5 6.5 4.0 4.0 9.5 6.5 4.0 4.0
```

```
> cluster_edge_betweenness(g,weights=NULL,directed=FALSE)
```

```
IGRAPH clustering edge betweenness, groups: 3, mod: 0.42
```

```
+ groups:
```

```
$`1`
```

```
[1] "A" "B" "C"
```

```
$`2`
```

```
[1] "D" "E" "F"
```

```
$`3`
```

```
[1] "G" "H" "I"
```