(a) A possible derivation of the word "abaabaaba" is:

$$1\;(aSa) \rightarrow 2\;(abSba) \rightarrow 1\;(abaSaba) \rightarrow 1\;(abaaSaaba) \rightarrow 4\;(abaabaaba)$$

- (b)  $\mathcal{L}(G)$  is the definition of a Palindrome. The word can always be read either forward or backward and you get the same result.
- (c) The Grammar that represents a representation of binary trees is:

$$G = \left\langle \left\{\,S\,\right\},\; \left\{\,\left[,\circ,\square,\right]\,\right\},\; R,\; S\right\rangle$$

with R:

- 1)  $S \rightarrow \square$
- $2) \ S \to [S \circ S]$

(a) To check if  $G_1 = \langle \{S, X, Y\}, \{a, b\}, R_1, S \rangle$  is a regular language we rewrite the rules set. In a first step we eliminate the start variable from the right-hand side of  $R_1$  and in a second step we eliminate forbidden occurrences of  $\epsilon$  and obtain:

$$R_1' = \{\, S \rightarrow aX, \; S \rightarrow aS', \; S \rightarrow \epsilon, \; S' \rightarrow a, \; S' \rightarrow aX, \; S' \rightarrow aS', \; X \rightarrow ba, \; X \rightarrow bX \,\}$$

So we see that this is a regular grammar and therefore of Type-3, and because of Type-3 also of Type-i for i < 3.

(b) To-Do

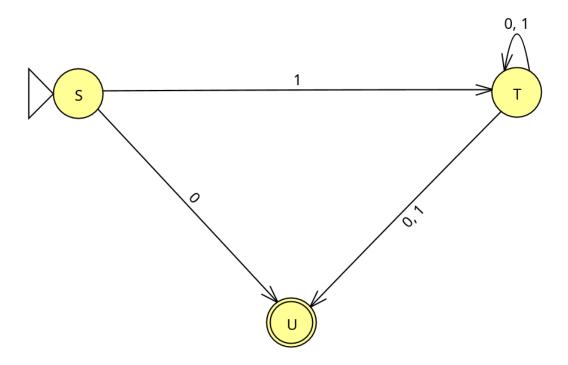


Figure 1: NFA for G