### Exercise 7.1

(a)

$$x_1 = \operatorname{dec}(101) = 5, \qquad x_2 = \operatorname{dec}(0) = 0, \qquad x_3 = \operatorname{dec}(11) = 3, \qquad y = \operatorname{dec}(1001) = 9$$

(b) 
$$x_1=\mathrm{dec}(11)=3, \qquad x_2=\mathrm{dec}(100)=4, \qquad x_3=\mathrm{dec}(1)=1, \qquad y=\mathit{undefined}$$

### Exercise 7.2

Step 1. Check if  $n_2 = 0$ : if **true** write  $n_1 = 1$ , remove  $n_2, n_3$  from tape and terminate.

Step 2. Check if  $n_2 = 1$ : if **true** remove  $n_2, n_3$  from tape and terminate.

Step 3. Copy  $n_1$  to  $n_3$ .

Step 4.  $M_{pred_1}(n_2)$ .

Step 5.  $M_{mul}(n_1, n_3)$  and write result to  $n_1$ .

Step 6. Check if  $n_2 = 1$ : if **true** remove  $n_2, n_3$  from tape and terminate.

Step 7. Jump to Step 4.

#### Exercise 7.3

To show that the composition  $(f \circ g) : \Sigma_1^* \to \Sigma_2^*$  is Turing-computable, we can build a TM that simulates the computations of f and g sequentially. This TM would work as follows:

- 1. Start the TM with the input x.
- 2. Use a Turing machine that simulates the computation of g on x. If g(x) is undefined, then halt and output undefined.
- 3. Use a Turing machine that simulates the computation of f on g(x). If f(g(x)) is undefined, then halt and output undefined.
- 4. If f(g(x)) is defined, output it as the result of the composition.

Since we assume that f and g are Turing-computable, we know that there exist Turing machines that can compute f and g respectively. Therefore, we can construct a Turing machine that combines the computations of these two machines, and this new machine can compute the composition of f and g for any input. Therefore, we have shown that the composition of two Turing-computable partial functions is also Turing-computable.

# Exercise 7.4

Let w be an arbitrary string in A. Since A is a subset of B, we know that w is also in B. Similarly, since B is a subset of C, we know that w is also in C. Therefore, w is in A and in C, which implies that A is a subset of C, or  $A \leq C$ .

Thus, we have shown that if  $A \leq B$  and  $B \leq C$ , then  $A \leq C$  for any languages A, B and C. This is a result of the transitive property of set inclusion.

# Exercise 7.5