

## Exercise 8.1

KEI PLAN WAS SIE MIT SET  $S$  MEINT...

- (a) Let  $M_1$  be a TM that is undefined on input 0 and  $M_2$  be another TM that is defined on input 0. Let both TM's compute an unary function over  $\mathbb{N}$ . The Rice theorem tells us that it's impossible to write a general algorithm that decides whether a given TM has a property or not except if the property is always true or false. Since  $M_1$  and  $M_2$  are two TM's that behave differently, the property of 0 being undefined is not always true or false and therefore the language  $L$  is undecidable.
- (b) Let  $M_1$  be a TM that rejects every input and  $M_2$  be another TM that accepts every input. Both TM's obviously behave differently and are therefore undecidable using the Rice theorem
- (c)  $L$  is decidable because if it halts on an even number of steps for input 0, it always halts on an even number of steps. (KEI PLAN WIESO... CHATGPT)
- (d)  $L$  is decidable because we can "simply" compute possible pairs of input and compare them with the output. (EBEFALLS CHATGPT. MACHT IRGENDWIE HALBWEGS SINN??)

## Exercise 8.2

## Exercise 8.3

- (a)
- (b)

## Exercise 8.4

- (a)  $n^3$  will dominate as  $n$  gets larger. Therefore the runtime of  $X$  is bound to a polynomial function and we can conclude that problem  $X$  is part of  $P$ .
- (b) The runtime of  $n^{\log_2(n)}$  grows faster than any polynomial function because the exponent is not constant and grows with the size of  $n$ . We cannot conclude that  $X$  is in  $P$  for that reason.