## Interpretação e Compilação de Linguagens—2016-2017 Interpretation and Compilation of Programming Languages

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Notes: The test is open book. Students can use any (individual) printed material that each student

brings along. The test has a duration of 1h30.

Q-1 [7 val.] This question is about the definition of an abstract syntax and the operational semantics for a programming language. Consider the programming language, called microIt, presented in class with the concrete syntax given by the following grammar:

```
\begin{split} E ::= num \mid bool \mid E_1 + E_2 \\ \mid x \mid \mathsf{decl} \ x = E_1 \text{ in } E_2 \\ \mid \mathsf{var}(E) \mid !E \mid E := E \\ \mid E_1 ? E_2 : E_3 \\ \mid \text{ for } x \text{ in } E_1 .. E_2 \text{ do } E_3 \end{split}
```

The language comprises the base constructs for: **integer literals** (num), **boolean literals** (bool), and their usual operations, represented here by operation E+E; **identifier** use (x) and **declaration** decl  $x = E_1$  in  $E_2$ . The language includes the expression for creation of a variable var(E), the dereferencing of a variable !E, and the assignment expression E:=E. The semantics of the presented language follows the semantics presented in the course lectures. The language also includes a conditional expression that yields the result of  $E_2$  if  $E_1$  denotes true, and of  $E_3$  if  $E_1$  denotes false; and an iterator operation (for x in  $E_1...E_2$  do  $E_3$ ) over a range given by two integer values (denoted by  $E_2$  and  $E_3$ ), evaluating expression  $E_3$  for all possible values of identifier x, that range from the integer value denoted by expression  $E_1$  to the integer value of expression  $E_2$  (inclusive). Consider the example written in the programming language microIt:

```
decl s = var(0) in
decl x = for x in 1..10 do (x\%2==0)?(s:=!s+1):0 in !s
```

Notice that operators !=, ==, and % are generally represented in the language by the operator +.

- a) [1 val.] Define the abstract syntax of the iteration operation in language microIt by means of an abstract data type, defined by set of (abbreviated) Java classes and interfaces.
- b) [2 val.] Define the set of values of language microIt by means of an abstract data type, defined by a set of (abbreviated) Java classes and interfaces.
- c) [2 val.] Define the operational semantics of language microIt by means of a method eval for the iteration expression.
- d) [1 val.] State the denotation (value) of the example above according to the semantics defined in the previous question, and the expected semantics for the remaining operators.
- e) [1 val.] Explain why the iteration expression above only makes sense in the context of an imperative language. (max 50 words.)

Q-2 [7 val.] This question is about the definition of the type system for language microIt. To answer the following questions you may use abstract data types, defined by a set of Java classes and interfaces, and the corresponding methods using Java Code. You may also use the notation used in the lecture slides.

- a) [2 val.] Define the set of types used to type programs of language microIt.
- b) [2 val.] Define the type system of language microIt by means of a typecheck function, for the iteration expression.
- c) [1 val.] State the type denotation of the example expression in question Q-1, according to the type semantics defined in question Q-2a.
- d) [1 val.] Enumerate the execution errors that may occur during the execution of a program written in language microlt, according to the semantics defined in question Q-1.
- e) [1 val.] Indicate and justify which execution errors may be prevented by the type system, and those that cannot.

Q-3 [6 val.] This question is about the evaluation and compilation of programs using mutable environments. Consider the following program written in the microIt language.

- a) [1 val.] Indicate what are the evaluation environments for the subexpressions w + y, and x + y + v + z, highlighted in the listing above with (\*).
- b) [1 val.] Indicate what are the compilation environments for the subexpressions w + y, and x + y + v + z, highlighted in the listing above with (\*).
- c) [4 val.] List the set of instructions that results from translating expression x + y + v + z to the Jasmin assembly language.

Clearly State the compilation addresses (result of function find) of all 4 identifiers used.

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