

Knowledge Representation and Reasoning

Solutions to Review Exercises on First Order Logic

1 Alpine Club

Formulate the following pieces of knowledge as sentences of first-order logic:

Tony, Mike and John belong to the Alpine Club. Every member of the Alpine Club who is not a skier is a mountain climber. Mountain climbers do not like rain, and anyone who does not like snow is not a skier. Mike dislikes whatever Tony likes, and likes whatever Tony dislikes. Tony likes rain and snow.

Answer:

The answer uses the following predicates and constants:

- *Member*: unary predicate meaning a member of the Alpine Club;
- *Skier*: unary predicate meaning a skier;
- *Climber*: unary predicate meaning a climber;
- *Likes*: binary predicate where *Likes*(*x*, *y*) means that *x* likes *y*;
- constants *tony*, *mike*, *john*, *rain*, *snow*.

In the translation, we name sentences so that it is easy to refer to them later.

- Tony, Mike and John belong to the Alpine Club.

$$S1 : \text{Member}(\text{tony})$$

$$S2 : \text{Member}(\text{mike})$$

$$S3 : \text{Member}(\text{john})$$

Alternatively, we may use one conjunction $\text{Member}(\text{tony}) \wedge \text{Member}(\text{mike}) \wedge \text{Member}(\text{john})$.

- Every member of the Alpine Club who is not a skier is a mountain climber.

$$S4 : \forall x (\text{Member}(x) \wedge \neg \text{Skier}(x) \supset \text{Climber}(x))$$

- Mountain climbers do not like rain

$$S5 : \forall x (\text{Climber}(x) \supset \neg \text{Likes}(x, \text{rain}))$$

- and anyone who does not like snow is not a skier.

$$S6 : \forall x (\neg \text{Likes}(x, \text{snow}) \supset \neg \text{Skier}(x))$$

- Mike dislikes whatever Tony likes

$$S7 : \forall x (\text{Likes}(\text{tony}, x) \supset \neg \text{Likes}(\text{mike}, x))$$

- and likes whatever Tony dislikes.

$$S8 : \forall x (\neg \text{Likes}(\text{tony}, x) \supset \text{Likes}(\text{mike}, x))$$

- Tony likes rain and snow.

$$S9 : \text{Likes}(\text{tony}, \text{rain})$$

$$S10 : \text{Likes}(\text{tony}, \text{snow})$$

2 Reduction to CNF

Rewrite all sentences in $KB = \{(p \vee q) \supset r, r \supset s, p\}$ in conjunctive normal form, and present KB in clausal form.

Answer:

- $(p \vee q) \supset r$ is, by definition of \supset equivalent to $\neg(p \vee q) \vee r$
 $\neg(p \vee q) \vee r$ is by de Morgan's law equivalent to $(\neg p \wedge \neg q) \vee r$
By distributivity, $(\neg p \wedge \neg q) \vee r$ is equivalent to $(\neg p \vee r) \wedge (\neg q \vee r)$
 $(\neg p \vee r) \wedge (\neg q \vee r)$ is in CNF and corresponds to two clauses $[\neg p, r] [\neg q, r]$.
- $r \supset s$ is, by definition of \supset equivalent to $\neg r \vee s$
 $\neg r \vee s$ is in CNF and corresponds to the clause $[\neg r, s]$.

The KB written in clausal form is $KB = \{[\neg p, r], [\neg q, r], [\neg r, s], [p]\}$.

3 Propositional Resolution

a) Show by resolution that the following set of clauses is inconsistent (derive empty clause from it):

$$[A, B, C], [A, B, \neg C], [A, \neg B, C], [A, \neg B, \neg C] \\ [\neg A, B, C], [\neg A, B, \neg C], [\neg A, \neg B, C], [\neg A, \neg B, \neg C]$$

b) Show by resolution that the following sentence is inconsistent:

$$\neg\neg A \wedge (\neg A \vee ((\neg B \vee C) \wedge B)) \wedge \neg C$$

Answer:

a) We can apply resolution as follows:

1. $[A, B]$ from $[A, B, C], [A, B, \neg C]$.
2. $[A, \neg B]$ from $[A, \neg B, C], [A, \neg B, \neg C]$.
3. $[A]$ from 1. and 2.
4. $[\neg A, B]$ from $[\neg A, B, C], [\neg A, B, \neg C]$.
5. $[\neg A, \neg B]$ from $[\neg A, \neg B, C], [\neg A, \neg B, \neg C]$.
6. $[\neg A]$ from 4. and 5.
7. \square from 3. and 6.

b) We first need to transform into conjunctive normal form to obtain the clauses:

- $\neg\neg A \wedge (\neg A \vee ((\neg B \vee C) \wedge B)) \wedge \neg C$ is equivalent to
- $A \wedge (\neg A \vee \neg B \vee C) \wedge (\neg A \vee B) \wedge \neg C$ which corresponds to the clauses
- $[A], [\neg A, \neg B, C], [\neg A, B], [\neg C]$; Then:
 1. $[\neg B, C]$ from $[A], [\neg A, \neg B, C]$
 2. $[B]$ from $[A], [\neg A, B]$
 3. $[C]$ from 1. and 2.
 4. \square from 3. and $[\neg C]$

4 First-Order Resolution

Determine whether the following sentences are valid using resolution:

- a) $\exists x \forall y \forall z ((P(y) \supset Q(z)) \supset (P(x) \supset Q(x)))$
- b) $\exists x (P(x) \supset \forall y (P(y)))$
- c) $\neg \exists x \forall y (E(x, y) \leftrightarrow \neg E(y, y))$

Answer:

To do this we need to check if from the negation of the sentence we can derive an empty clause (a contradiction).

a) First transform the negation into clausal form:

$$\begin{aligned} & \neg \exists x \forall y \forall z ((P(y) \supset Q(z)) \supset (P(x) \supset Q(x))) \\ & \neg \exists x \forall y \forall z (\neg (\neg P(y) \vee Q(z)) \vee (\neg P(x) \vee Q(x))) \\ & \forall x \exists y \exists z \neg (\neg (\neg P(y) \vee Q(z)) \vee (\neg P(x) \vee Q(x))) \\ & \forall x \exists y \exists z (\neg \neg (\neg P(y) \vee Q(z)) \wedge \neg (\neg P(x) \vee Q(x))) \\ & \forall x \exists y \exists z ((\neg P(y) \vee Q(z)) \wedge (\neg \neg P(x) \wedge \neg Q(x))) \\ & \forall x \exists y \exists z ((\neg P(y) \vee Q(z)) \wedge P(x) \wedge \neg Q(x)) \\ & \forall x ((\neg P(f(x)) \vee Q(g(x))) \wedge P(x) \wedge \neg Q(x)) \end{aligned}$$

Clauses:

$$\begin{aligned} C1 & : [\neg P(f(x)) \vee Q(g(x))] \\ C2 & : [P(x_1)] \\ C3 & : [\neg Q(x_2)] \end{aligned}$$

Proof:

- 1. $[Q(g(x))]$ from $C1$ and $C2$, $x_1/f(x)$.
- 2. \square from (1) and $C3$, $x_2/g(x)$.

b) Transform the negation into clausal form:

$$\begin{aligned} & \neg \exists x (P(x) \supset \forall y (P(y))) \\ & \forall x \neg (\neg P(x) \vee \forall y (P(y))) \\ & \forall x (P(x) \wedge \exists y \neg P(y)) \\ & \forall x (P(x) \wedge \neg P(f(x))) \end{aligned}$$

Clauses:

$$\begin{aligned} C1 & : [P(x)] \\ C2 & : [\neg P(f(x_1))] \end{aligned}$$

Proof:

- 1. \square from $C1$ and $C2$, $x/f(x_1)$

c) Transform the negation into clausal form:

$$\begin{aligned} & \neg \neg \exists x \forall y (E(x, y) \leftrightarrow \neg E(y, y)) \\ & \exists x \forall y ((\neg E(x, y) \vee \neg E(y, y)) \wedge (E(x, y) \vee E(y, y))) \\ & \forall y ((\neg E(a, y) \vee \neg E(y, y)) \wedge (E(a, y) \vee E(y, y))) \end{aligned}$$

Clauses:

$$C1 : [\neg E(a, y), \neg E(y, y)]$$

$$C2 : [E(a, y_1), E(y_1, y_1)]$$

Proof:

1. $[\neg E(a, a)]$ factorization $C1, y/a$
2. $[E(a, a)]$ factorization $C2, y_1/a$
3. \square from 1. and 2.

5 Alpine Club and First-Order Resolution

As a follow-up to the Alpine Club Exercise, use resolution to prove that there exists a member of the Alpine club who is a climber but not a skier.

Answer:

Translation into first-order logic as given in the solution for Exercise 1 (S1 – S10) together with:

$$S11 : \exists x (Member(x) \wedge Climber(x) \wedge \neg Skier(x))$$

Now in clausal form (with S11 negated):

$$C1 : [Member(tony)]$$

$$C2 : [Member(mike)]$$

$$C3 : [Member(john)]$$

$$C4 : [\neg Member(x), Skier(x), Climber(x)]$$

$$C5 : [\neg Climber(x_1), \neg Likes(x_1, rain)]$$

$$C6 : [Likes(x_2, snow), \neg Skier(x_2)]$$

$$C7 : [\neg Likes(tony, x_3), \neg Likes(mike, x_3)]$$

$$C8 : [Likes(tony, x_4), Likes(mike, x_4)]$$

$$C9 : [Likes(tony, rain)]$$

$$C10 : [Likes(tony, snow)]$$

$$C11 : [\neg Member(x_5), \neg Climber(x_5), Skier(x_5)]$$

Prove that, together, $C1 - C11$ are inconsistent:

1. $[\neg Likes(mike, snow)]$ from $C10$ and $C7, x_3/snow$
2. $[\neg Skier(mike)]$ from 1. and $C6, x_2/mike$
3. $[\neg Member(mike), Climber(mike)]$ from 2. and $C4, x/mike$
4. $[Climber(mike)]$ from 3. and $C2$
5. $[\neg Member(mike), Skier(mike)]$ from (4) and $C11, x_5/mike$
6. $[Skier(mike)]$ from 5. and $C2$
7. \square from 6. and 2.

Can you find an even shorter proof?