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\*\*\*PLEASE USE BEING-EVE.PDF TO GRADE\*\*\*

\*\*I've included this txt in the submission because the pdf is harder to read than this txt.\*\*

YOUR JOB FOR DIFFIE HELLMAN

Figure out the shared secret agreed upon by Alice and Bob. This will be an integer.

The secret message is 6.

Show your work. Exactly how did you figure out the shared secret?

On paper scan

Show precisely where in your process you would have failed if the integers involved were much larger.

Computing the value of  $x$  and  $y$  in  $g^x \bmod p$  and  $g^y \bmod p$  respectively gets slow if  $g$  and  $p$  are large integers. It's also possible that there can be multiple values of  $x$  and  $y$ . This would take more time to figure out the secret message ( $g^{(xy)} \bmod p$ ).

YOUR JOB FOR RSA

Figure out the encrypted message sent from Alice to Bob.

Hi Bob. I'm walking from now on. Your pal, Alice.

<https://foundation.mozilla.org/en/privacynotincluded/articles/its-official-cars-are-the-worst-product-category-we-have-ever-reviewed-for-privacy/>

Show your work. Exactly how did you figure out the message? (You should include an explanation of how the message from Alice to Bob is encoded. That is, how does Alice's intended message (whatever manner of message it may be) correspond to the integers in the plaintext that you end up with after decrypting the encrypted message?)

To begin, I found  $p$  and  $q$  using brute force (lines 44–47). By definition,  $p * q = n$ .  $n$  is given in Bob's public key, so I brute forced the values of  $p$  and  $q$ .

Once I found  $p$  and  $q$ , I could brute force the value of  $d$  (Bob's private key). By definition  $ed = 1 \bmod (p - 1)(q - 1)$ .  $e$  is in the public key, and I solved for  $p$  and  $q$ . I then used

function:

$$m(c) = c^d \bmod n$$

Found on wikipedia:

[https://en.wikipedia.org/wiki/RSA\\_\(cryptosystem\)](https://en.wikipedia.org/wiki/RSA_(cryptosystem))

This gives an integer. Turning this integer into binary reveals 2  
ascii characters smushed  
back-to-back. Converting the ascii into English letters reveals  
the message.

Alice encrypted her message by taking 2 characters at a time,  
concatenating their binary  
ascii values together, then taking that binary number as an  
integer. That integer is then  
encoded using Alice's public key as per RSA.  
$$c(m) = m^e \bmod n$$

Show precisely where in your process you would have failed if the  
integers involved were large  
(e.g., on the order of a 2000-bit value for n).

Brute forcing for d would have failed (taken way too much time) if  
n were large (lines 53–56).

If n were large, it would also have taken much longer to find p  
and q (lines 43–47).

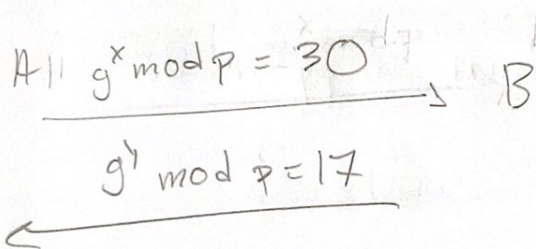
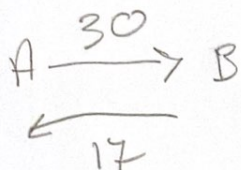
Explain, briefly, why the message encoding Alice used would be  
insecure even if Bob's keys  
involved very large integers.

Alice's message would have been insecure because the 2 character  
strings that Alice encoded  
can be matched to English. From there, it's possible to figure out  
the message without  
knowing the encoding scheme.

$$g=7, P=61$$

DH

Robert  
Baron



x=Alice  
y=Bob

$$7^x \bmod 61 = 30$$

$$7^y \bmod 61 = 17$$



$$\boxed{x=4, y=23} \text{ (in PY)}$$

shared secret =  $g^{yx} \bmod p = 7^{23 \cdot 4} \bmod 61 = \boxed{6}$

where would my process fail if the ints were larger?

computing  $g^x \bmod p$  and  $g^y \bmod p$  gets very slow if  $g$  and  $p$  are large integers

- there can also be multiple answers for  $x$  and  $y$ .  
This would take more time ( ) to figure out the secret message (bc the secret message uses both  $x$  and  $y$ )



## RSA

Bob's public key

$$(n_{\text{bob}}, e_{\text{bob}}) = (170171, 17)$$

$$n_B = \text{modulus} = 170171$$

$$e_B = \text{encryption exponent} = 17$$

LCM of  $p$  and  $q$

$$P_B = 449 \quad \left. \begin{array}{l} \\ \end{array} \right\} \text{found in py code}$$

$$P_B \cdot L_{n_B}, q_B \cdot L_{n_B}$$

$$q_B = 379$$

45

$$m(c) = c \cdot \text{mod } n$$

18537  $\rightarrow$  2 ascii chars

$$100100001101001001001001$$

H

i

8 bits

8 bits

$$m(c) = c \cdot \text{mod } n$$

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