## ARBITRARY RECURSION

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One Recursive Call
arbitrary, possibly non-terminating, recursive 'definition' (e.g. a program-mode function):
      f(\overline{x}) \stackrel{.}{=} if a(\overline{x}) \stackrel{.}{\underline{then}} b(\overline{x}) \stackrel{.}{\underline{else}} c(\overline{x}, f(\overline{d}(\overline{x}))) \qquad \overline{x} = (x_1, ..., x_n) \qquad \overline{d}(\overline{x}) = (d_1(\overline{x}), ..., d_n(\overline{x}))
         a \in \mathcal{U}' b: \mathcal{U}' \rightarrow \mathcal{U}'' d_i: \mathcal{U}'' \rightarrow \mathcal{U} c: \mathcal{U}'' \times \mathcal{U}'' \rightarrow \mathcal{U}'' m > 0 (number of results)
  t(\bar{z}) \triangleq |\exists k \in \mathbb{N}. \ \alpha(\bar{d}^k(\bar{z}))| - f \text{ terminates on } \bar{z}
 always-terminating "version" of f (a logic-mode function, if a,b,c,d are logic-mode):
      f(\bar{z}) \triangleq if t(\bar{z}) then [if a(\bar{z})] then b(\bar{z}) else c(\bar{z}, f(\bar{d}(\bar{z})))] else ... any value (irrelevant)
       \mu_{\hat{\tau}}(\bar{z}) \triangleq if t(\bar{z}) then min \{k \in \mathbb{N} \mid a(\bar{d}^k(\bar{z}))\} else ... any natural number (irrelevant) \mu_{\hat{\tau}}: \mathcal{U}^n \to \mathbb{N}
       \vdash \boxed{\mu\text{-end}} \ \ t(\vec{z}) \Rightarrow \alpha(\vec{d}^{M_{\hat{f}}(\vec{z})}(\vec{z}))

\begin{array}{ccc}
 & \xrightarrow{S_{N_{\overline{a}}}} \mu_{\widehat{x}}(\overline{x}) = \min \left\{ k \in \mathbb{N} \left[ a(\overline{d}^{k}(\overline{x})) \right] \longrightarrow \mu_{\widehat{x}}(x) \in \left\{ k \in \mathbb{N} \left[ a(\overline{d}^{k}(\overline{x})) \right] \longrightarrow a(\overline{d}^{\mu_{\widehat{x}}(\overline{x})}(\overline{x})) \right\} \\
\end{array}

  + \left[\mu - \min\right] \alpha(\overline{d}(\overline{z})) \Longrightarrow l \gg \mu_{\widehat{f}}(\overline{z})
     a(\overline{d}(\overline{z})) \xrightarrow{S_t} t(\overline{z}) \xrightarrow{S_{M_f^f}} \mu_f(\overline{x}) = \min \left\{ k \in \mathbb{N} \mid a(\overline{d}^k(\overline{z})) \right\}
\Rightarrow l \in \left\{ k \in \mathbb{N} \mid a(\overline{d}^k(\overline{x})) \right\}
     QED
  \vdash \mid \tau_{\hat{\mathfrak{f}}} \mid t(\bar{\mathfrak{z}}) \wedge \neg \alpha(\bar{\mathfrak{z}}) \Longrightarrow \mu_{\hat{\mathfrak{f}}}(\bar{d}(\bar{\mathfrak{z}})) \prec_{\hat{\mathfrak{f}}} \mu_{\hat{\mathfrak{f}}}(\bar{\mathfrak{z}})
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alternative recursive definition of the measure function;
 φ(u) = if u ∈ N then u else O - fixing function for N
 \mathcal{E}_{t}(\bar{z}) \triangleq \mathcal{E}_{k}. \alpha(\bar{d}^{\varphi(k)}(\bar{z})) - \text{witness of } t, for slightly madified definition t(\bar{z}) \triangleq [\exists k. \alpha(\bar{d}^{\varphi(k)}(\bar{z}))]
 M_{\mathcal{N}}(\overline{z},k) \triangleq \mathcal{Y}\left(\varepsilon_{\mathsf{t}}(\overline{z}) - \mathcal{Y}(k)\right) \forall_{\mathcal{N}} \triangleq \mathcal{Z} \subseteq N_{\mathcal{N}}N
  \vdash [\tau_{v}] \neg a(\overline{d}^{g(k)}(\overline{z})) \wedge g(k) \langle g(\varepsilon_{t}(\overline{z})) \Rightarrow \mu_{v}(\overline{z}, \varphi(k) + 1) \prec_{v} \mu_{v}(\overline{z}, k)
                  \mu_{\nu}(\bar{z}, \varphi(k)+1) = \varphi(\varepsilon_{+}(\bar{z}) - \varphi(\varphi(k)+1)) = \varphi(\varepsilon_{+}(\bar{z}) - \varphi(k)-1) = \varepsilon_{+}(\bar{z}) - \varphi(k)-1
                  \varphi(k) < \varphi\left(\varepsilon_{t}(\bar{z})\right) \longrightarrow \varphi\left(\varepsilon_{t}(\bar{z})\right) \neq 0 \xrightarrow{\delta_{\varphi}} \varepsilon_{t}(\bar{z}) \in \mathbb{N} \longrightarrow \varphi(k) < \varepsilon_{t}(\bar{z}) \longrightarrow \varepsilon_{t}(\bar{z}) - \varphi(k) > 0
\mu_{v}(\bar{z}, k) \stackrel{\delta_{r^{v}}}{=} \varphi\left(\varepsilon_{t}(\bar{z}) - \varphi(k)\right) \stackrel{\varepsilon}{=} \varepsilon_{t}(\bar{z}) - \varphi(k) \longrightarrow \mu_{v}(\bar{z}, \varphi(k) + 1) < \nu_{v} \mu_{v}(\bar{z}, k) \stackrel{\varepsilon}{=} \varepsilon_{t}(\bar{z}) - \varphi(k)
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$$| V | V(\bar{x}, k) \in N$$

$$| base \rangle = \delta_{v} \rightarrow v(\bar{x}, k) = \varphi(k)$$

$$| \delta_{\varphi} \rightarrow \varphi(k) \in N$$

$$| step \rangle = \delta_{v} \rightarrow v(\bar{x}, k) = v(\bar{x}, \varphi(k) + 1)$$

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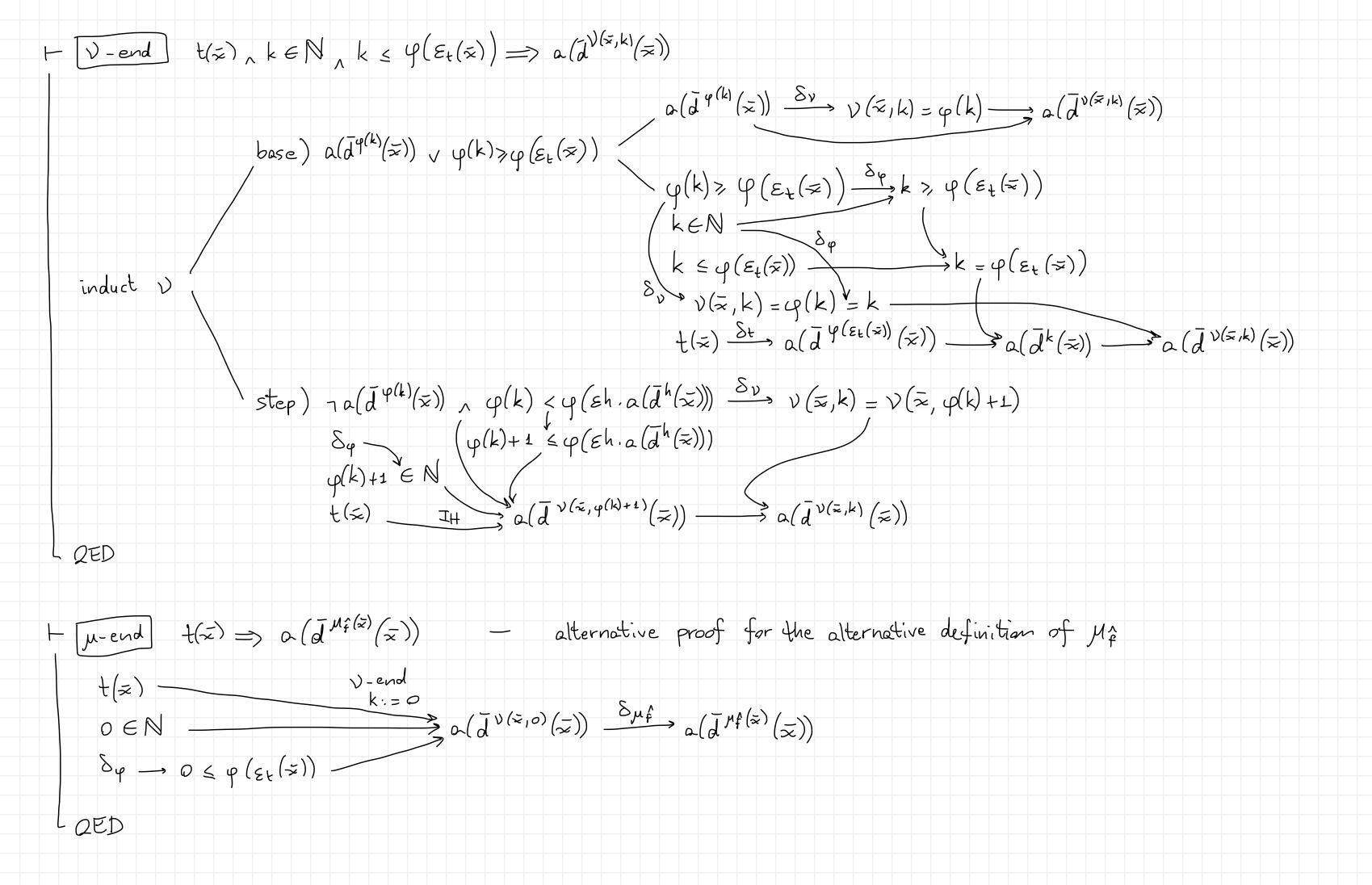
$$| c | c | c | c | c | c |$$

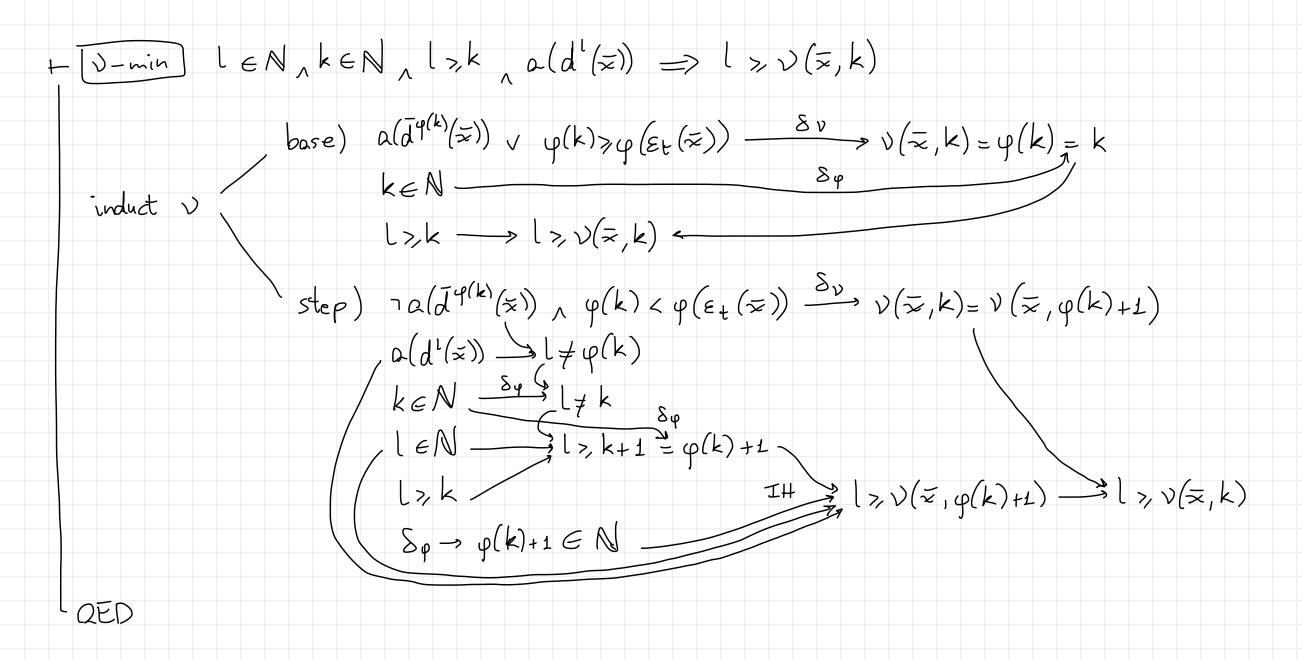
$$| c | c | c | c | c |$$

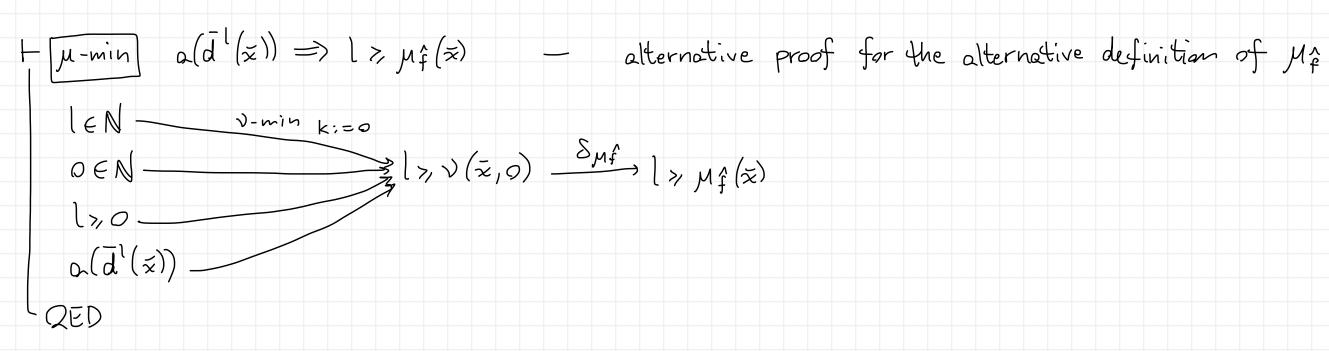
$$| c | c | c | c | c | c |$$

$$\mu_{\hat{x}}(\bar{z}) \triangleq \nu(\bar{z}, 0)$$
 - alternative definition of  $\mu_{\hat{x}}$ 

$$\mu_{\hat{f}}: \mathcal{U}^n \rightarrow \mathcal{N} \quad (\Leftarrow \nu \mathcal{N})$$







Ti proved as before using u-end and u-min