## PARTIAL EVALUATION TRANSFORMATION

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## Non-Recursive Function

old function:  $f(x,y) \triangleq e(x,y)$  } y static, x dynamic  $y \triangleq \tilde{y}$  —  $\tilde{y}$  ground term

new function:  $f'(x) \triangleq e(x, \tilde{y})$ 

$$+ff'$$
  $y=\tilde{y}=>f(x,\tilde{y})=f'(x)$  — trivial, by  $S_f$  and  $S_{f'}$ 

then optimize f' via further transformations

 $x \rightarrow x_1,...,x_n$   $y \rightarrow y_1,...,y_m$   $y \rightarrow \widetilde{y}_1,...,\widetilde{y}_m$   $y \rightarrow \widetilde{y}_1,...,\widetilde{y}_m$   $y \rightarrow \widetilde{y}_1,...,\widetilde{y}_m$ 

## Recursive Function - Default Treatment

old function:  $f(x,y) \triangleq ... f...$  } y static, x dynamic  $y \triangleq \widetilde{y}$  —  $\widetilde{y}$  ground term

new function:  $f'(x) \triangleq f(x, \tilde{y})$  — non-recursive — preliminary simple approach

 $+\left[ff'\right]$   $y=\widetilde{y} \Longrightarrow f(x,\widetilde{y})=f'(x)$  — trivial, by  $S_{f'}$ 

optimize f' via successive transformations, which may unfold the recursion completely if driven by y

$$\begin{array}{c} (\sqrt{f}) \quad \chi_{8f}(\times, y) \wedge \dots \\ \chi_{f'}(\times) \stackrel{\triangle}{=} \quad \chi_{f}(\times, \tilde{y}) \\ (-\sqrt{f'}) \quad \omega_{f'}(\times) \\ \omega_{f'}(\times) = \chi_{7f}(\times, \tilde{y}) \wedge \left(\chi_{f}(\times, \tilde{y}) \stackrel{\triangle}{=} \right) \chi_{f}(\times, \tilde{y}) \end{array}$$

$$\begin{array}{c} \chi_{f'}(\times) \stackrel{\triangle}{=} \quad \chi_{f}(\times, \tilde{y}) \\ \chi_{f'}(\times) \stackrel{\triangle}{=} \quad \chi_{f}(\times, \tilde{y}) \\ \chi_{f'}(\times) \stackrel{\triangle}{=} \quad \chi_{f}(\times, \tilde{y}) \end{array}$$

generalizes to more parameters os in non-recursive case

## Recursive Function with Unchanging Static Argument old function: $f(x,y) \leq if a(x,y)$ then b(x,y) else c(x,y,f(d(x,y),y)) $[T_f] 7a(x,y) \Rightarrow \mu_f(d(x,y)) \leq_f \mu_f(x) - x$ measured argument, y not measured argument y ≜ y - y ground term new function: $f'(x) \triangleq if a(x, \tilde{y})$ then $b(x, \tilde{y})$ else $c(x, \tilde{y}, f'(d(x, \tilde{y})))$ $M_{f'}(x) \triangleq M_{f}(x) \qquad \langle f' \triangleq \langle f \rangle$ $\vdash [T_{f'}] \neg \alpha(x, \tilde{y}) \Rightarrow M_{f'}(d(x, \tilde{y})) \prec_{f'} M_{f'}(x)$ $\vdash [ff'] y = y \Rightarrow f(x,y) = f'(x)$ $a(x,\tilde{y}) \xrightarrow{\delta f'} f'(x) = b(x,\tilde{y})$ $y := g' f(x,\tilde{y}) = b(x,\tilde{y})$ induct f' $y := g' f(x,\tilde{y}) = b(x,\tilde{y})$

L QED

 $\begin{array}{c} \boxed{\forall f} \quad \chi_{\mathcal{S}_{f}}(\varkappa, y) \wedge \left[\chi_{f}(\varkappa, y)\right] \Rightarrow \chi_{a}(\varkappa, y) \wedge \left[\alpha(\varkappa, y)\right] \Rightarrow \chi_{b}(\varkappa, y) \right] \wedge \left[\gamma_{a}(\varkappa, y)\right] \wedge \chi_{f}(d(\varkappa, y), y) \wedge \chi_{c}(\varkappa, y) + (d(\varkappa, y), y))\right] \\ \chi_{f}(\varkappa) \triangleq \chi_{f}(\varkappa, \mathfrak{F}) \end{array}$ 

generalizes to more parameters as in non-recursive case - y1,..., ym all unchanging