Week 1 - Assignment

1. List Operations - Part 2

Implement the functions below. Note that many of these functions are available in the standard library, but the goal of this exercise is to practice by implementing them from scratch. When writing a recursive function involving lists, put some thought into choosing the right base case.

Recall that the syntax of pattern-matching on a list is as follows (where x is the head of the list and xs is the tail):

```
function :: [...] \rightarrow ...
function [] = ...
function (x:xs) = ...
```

Note that HLint (on E-Systant) may generate warnings, you can ignore these.

• Write a function myProduct:: [Integer] -> Integer, which takes a list of integers and computes their product.

Note that we use Integer here instead of Int. The former can represent numbers of arbitrary size, whereas the latter will overflow.

```
Main> myProduct [1,2,3]
6

Main> myProduct []
1

Main> myProduct [-2,3,-4,5,-6]
-720
```

• Write a function insert:: Int -> [Int] -> [Int], which takes an integer and a list of integers and inserts the integer into the list at the first position where it is less than or equal to the next element.

```
Main> insert 0 [1,2,3] [0,1,2,3]
```

 $^{^1}$ For example, see module Data.List, which can be found at http://downloads.haskell.org/~ghc/7.6.3/docs/html/libraries/base/.

```
Main> insert 2 [1,0,3] [1,0,2,3]

Main> insert 4 [1,2,3] [1,2,3,4]
```

• Write a function myLast:: [Int] -> Int which returns the last element of a list. Assume that the input lists are non-empty²

```
Main> myLast [1,2,3,4,5]
```

2. Rock - Paper - Scissors

In the Rock, Paper and Scissors game, two players choose one of the following gestures after counting to three:

- \bullet A clenched fist which represents a rock.
- A flat hand representing a piece of paper.
- Index and middle finger extended which represents a pair of scissors.

The result of a round is decided this way:

- Rock defeats scissors, because a rock will blunt a pair of scissors.
- Paper defeats rock, because a paper can wrap up a rock.
- Scissors defeat paper, because scissors cut paper.
- Otherwise, the players chose the same gesture and it's a draw.

1. Moves

- Define a datatype Move with three choices Rock, Paper and Scissors that represent the valid moves. Note that the "deriving (Eq, Show)" should not be removed from the declaration otherwise the testing framework won't work.
- Write a function beat :: Move -> Move such that beat m is the move that beats move m.
- Write a function lose:: Move -> Move such that lose m is the move that will lose against move m.

²You may return an error using the error function in case the list is empty.

2. Playing The Game

- Define a datatype Result that represents the outcome of a round of Rock Paper Scissors. As we explained above, a player can either Win, Lose or the game may end up in a Draw. Like before, the "deriving (Eq, Show)" should not be removed from the declaration otherwise the testing framework won't work.
- Write a function outcome:: Move -> Move -> Result that takes as arguments two moves (the first argument is the move of the first player and the second is the move of the second player) and calculates the outcome for the first player.

NOTE: Since defining the data types is part of the exercise and their definition is thus not known to the testing framework (E-Systant), the tests for this exercise **only partially check the correctness of your solutions, less accurately than for other exercises.** Hence, be extra careful when defining functions beat, lose and outcome.

3. Lists, Continued

In contrast to the previous exercise on lists, try to implement these functions using list comprehensions (e.g., [a+1 | a <- as]) as well as list ranges: [1..n], [1,5,..,n]. Do not use explicit recursion!

Note that many of these functions are available in the standard library,³ but the goal of this exercise is to practice by implementing them from scratch.

• In mathematics, the factorial of a non-negative integer number n is defined recursively as follows:

$$factorial\ (n) = \begin{cases} 1 & \text{, if } n = 0\\ n * factorial\ (n-1) & \text{, if } n > 0 \end{cases}$$

Alternatively, we can also express it as: factorial (n) = 1 * 2 * ... * (n-1) * n

Write a function factorial:: Integer -> Integer, such that factorial n is the factorial of n. If n is a negative number, factorial n should result in 1. **Hint:** You can use the standard library function product, to compute the product of every element in a list.

```
Main> factorial 5
120

Main> factorial 0
1

Main> factorial (-10)
```

³For example, see module Data.List, which can be found on http://downloads.haskell.org/~ghc/7.6.3/docs/html/libraries/base/.

• Write a function myRepeat:: Int -> Int -> [Int] such that myRepeat n x returns a list with n times the number x. If n is less than zero return the empty list.

```
Main> myRepeat 4 5
[5,5,5,5]

Main> myRepeat (-1) 5
[]

Main> myRepeat 0 5
[]
```

• Write a function flatten:: [[Int]] -> [Int] which converts a list of lists to a single list.

```
Main> flatten [[1,2],[3,4],[5,6]]
[1,2,3,4,5,6]

Main> flatten []
[]
```

• Write a function range:: Int -> Int -> [Int] which returns a list of the consecutive numbers between the two given numbers, both numbers included. If the first number is greater than the second you should return the empty list.

```
Main> range 1 10
[1,2,3,4,5,6,7,8,9,10]

Main> range (-10) (-5)
[-10,-9,-8,-7,-6,-5]

Main> range 10 1
[]
```

• Write a function sumInts:: Int -> Int, which takes an integer low and an integer high and computes the sum:

$$\mathtt{sumInts\ low\ high} = \mathtt{low} + (\mathtt{low} + 1) + (\mathtt{low} + 2) + \ldots + (\mathtt{high} - 1) + \mathtt{high}$$

If low > high then the sum should be zero. Hint: You can use the standard library function sum, to compute the sum of every element in a list.

```
Main> sumInts 3 5
12
```

```
Main> sumInts 5 3

0

Main> sumInts 5 5
```

• Write a function removeMultiples:: Int -> [Int] -> [Int] which removes all multiples of a number from the list. Use n `mod` d or mod n d. Assume that the first argument will never be zero.

```
Main> removeMultiples 2 (range 1 10)
[1,3,5,7,9]
Main> removeMultiples 5 []
[]
```

4. Arithmetic Expressions

The following Exp datatype encodes the abstract syntax for arithmetic expressions. Note that it is recursively defined: Exp occurs in the right-hand side of its own definition.

1. Interpreter

• Write an interpreter eval:: Exp -> Int that evaluates arithmetic expressions.

Examples

```
Main> eval (Add (Mul (Const 2) (Const 4)) (Const 3))

11

Main> eval (Sub (Const 42) (Mul (Const 6) (Const 7)))

0
```

2. Compiler

Instead of evaluating an arithmetic expression directly, we can also compile it to a program of a simple stack machine and subsequently execute the program. We represent a program as a list of instructions. The instructions IAdd, ISub, IMul take the two topmost elements from the stack, perform the corresponding

operation, and push the result onto the stack. The IPush instruction pushes the given value onto the stack. A stack is modelled by a list of integers.

```
data Inst = IPush Int | IAdd | ISub | IMul
  deriving (Show, Eq)

type Prog = [Inst]

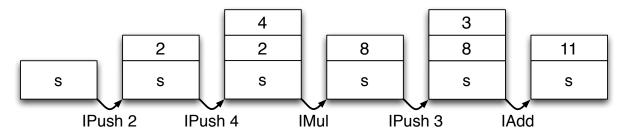
type Stack = [Int]
```

For example, the following arithmetic expression

```
Add (Mul (Const 2) (Const 4)) (Const 3)
```

is equivalent to the stack program

The stack program leaves the result on top of the stack. The stack machine performs the following steps when executing the program on an initial stack s



• Write an execution function execute:: Inst -> Stack -> Stack that executes a single instruction. Since there are cases where the function can crash (e.g. a stack overflow in cases where we want to execute an IAdd instruction but the stack contains fewer than two elements), you can use the following exception-raising function where needed:

```
runtimeError :: Stack
runtimeError = error "Runtime error."
```

- Write a function run:: Prog -> Stack -> Stack that runs a whole program on a given initial stack.
- Write a compiler compile:: Exp -> Prog that compiles arithmetic expressions to stack machine programs. The compiled program should leave the result of the computation as the top element on the stack. Make sure that your compiler uses a left-to-right evaluation order and that it produces results equivalent to the interpreter, i.e. the following identity holds

```
forall (s :: Stack). run (compile e) s == (eval e) : s
```

Examples

```
Main> execute IAdd [4,5,6]
[9,6]

Main> execute ISub [4,5,6]
[1,6]

Main> execute (IPush 2) [4,5,6]
[2,4,5,6]

Main> run [IAdd, ISub] [4,5,6]
[-3]

Main> run [IAdd, ISub, IPush 7, IMul] [4,5,6,8]
[-21,8]

Main> run [IPush 1,IPush 2,IPush 3,IMul,ISub] []
[-5]

Main> compile (Sub (Const 1) (Mul (Const 2) (Const 3)))
[IPush 1,IPush 2,IPush 3,IMul,ISub]
```

5. EXTRA: Caesar Cipher

Note: For this assignment you may find some functions from the Data. Char library useful.

Julius Caesar before sending his messages often encoded them, by replacing each letter by the letter three places further down in the alphabet (wrapping around at the end of the alphabet). For example, the string

```
"haskell is fun"
would be encoded as
"kdvnhoo lv ixq"
```

In general, we can do even more than Caesar did, and encode our strings using any integer between 1 and 25 (since the alphabet has 26 letters), having 25 different ways of encoding a string. For example, with a shift factor of 10, the original string would be encoded as:

```
"rkcuovv sc pex"
```

Encoding and Decoding

For simplicity, in this exercise we will only encode lowercase letters, leaving all other letters unchanged. Function let2int converts a lowercase letter between 'a' and 'z' to an integer from 0 to 25, and function

int2let does the inverse:

```
let2int :: Char -> Int
let2int c = ord c - ord 'a'
int2let :: Int -> Char
int2let n = chr (ord 'a' + n)
```

(The library functions ord::Char -> Int and chr::Int -> Char convert a character to its unicode representation and vice-versa) For example:

```
Main> let2int 'a'
0
Main> int2let 0
'a'
```

- Define function shift:: Int -> Char -> Char, which applies a shift factor to a lowercase letter and leaves any other character unchanged (**Hint:** Use the above functions, as well as function mod to ensure that the resulting integer representation does not exceed 26).
- Define function encode:: Int -> String -> String by means of function shift, which, given a shift factor, encodes a whole string.

Examples

```
Main> shift 3 'a'
'd'

Main> shift 3 'z'
'c'

Main> shift (-3) 'c'
'z'

Main> shift 3 ' '
' ' '

Main> encode 3 "haskell is fun"
"kdvnhoo lv ixq"

Main> encode (-3) "kdvnhoo lv ixq"
"haskell is fun"
```

Note that there is no need for a "decode" function, since if a string is encoded using a shift factor n, we can always take it back be re-encoding it using (-n) as a shift factor.

Frequency Tables

The key to crack the Caesar cipher is the observation that some letters of the English alphabet appear more often than others. In fact, by analyzing a large volume of text, one can derive the following table of approximate percentage frequencies of the 26 letters of the alphabet:

```
table :: [Float]
table = [ 8.2, 1.5, 2.8, 4.3, 12.7, 2.2, 2.0, 6.1, 7.0, 0.2, 0.8, 4.0, 2.4
, 6.7, 7.5, 1.9, 0.1, 6.0, 6.3, 9.1, 2.8, 1.0, 2.4, 0.2, 2.0, 0.1 ]
```

For example, letter 'e' occurs most often, with a frequency of 12.7%, while 'q' and 'z' appear least often, with a frequency of 0.1% each.

• Define function percent::Int -> Int -> Float which computes the percentage of an integer with respect to another (**Hint:** Use library function fromIntegral::(Integral a, Num b) => a -> b to convert the arguments to Float before dividing them). For example:

```
Main> percent 6 12 50.0 Main> percent 3 15 20.0
```

• Define function freqs::String -> [Float] which computes the frequencies of the 26 letters of the alphabet for a given string. Assume that the given string will contain at least one lowercase letter. For example:

That is, letter 'a' appears with frequency 6.7, letter 'b' with frequency 13.3 and so on.

Cracking The Cipher

Now that we have laid the foundations, it is time to crack Caesar's Cipher.

A standard method for comparing a list of observed frequencies o with a list of expected frequencies e is the *chi-square* statistic, defined as follows:

$$\text{chisqr o } \mathbf{e} = \sum_{i=0}^{n-1} \frac{(\mathbf{o}_i - \mathbf{e}_i)^2}{\mathbf{e}_i}$$

The details of the chi-square method are not important to us, only the fact that the smaller the result of chisqr o e, the better the match between frequency tables o and e.

• Implement function chisqr:: [Float] -> [Float] -> Float. Hint: this exercise can be easily solved using the zip function and list comprehensions.

• Implement function rotate:: Int -> [a] -> [a], which rotates the elements of a list a given number of times to the left. For example:

```
Main> rotate 3 [1,2,3,4,5] [4,5,1,2,3]
```

You can assume that the integer argument is always between 0 and the length of the list. **Hint:** Use functions take and drop to implement this exercise.

Now, if we are given an encoded string but not the shift factor used for the encoding, we can find the shift factor as follows:

- 1. We produce the frequency table of the encoded string
- 2. We calculate the chi-square statistic for each possible rotation of this table with respect to the expected frequencies (value table)
- 3. The position of the minimum chi-square value is the most probable shift-factor used to encode the string.

```
For example, if table' = freqs "kdvnhoo lv ixq", then
```

```
[ chisqr (rotate n table') table | n <- [0..25] ]
```

gives the result

```
[1408.8, 640.3, 612.4, 202.6, 1439.8, 4247.2, 651.3, \cdots, 626.7]
```

, in which the minimum value is 202.6, appearing in position 3 in this list (counting from 0). Hence, we can conclude that the shift factor used to encode the string was 3, and to retrieve the original string, we just need to encode it again using -3.

• Define function crack::String -> String which takes an encoded string and, using the above method, computes the original string. Hint: In addition to all the functions you have already defined, you will also find functions minimum::Ord a => [a] -> a and elemIndex::Eq a => a -> [a] -> Maybe Int useful for solving this exercise (function elemIndex is defined in the Data.List library).

Examples

```
Main> crack "kdvnhoo lv ixq"

"haskell is fun"

Main> crack "vscd mywzboroxcsyxc kbo ecopev"

"list comprehensions are useful"
```

Note that cracking is not always accurate, especially in cases where the encoded word is too short, or has an unusual distribution of letters.

```
Main> crack (encode 3 "haskell")
"piasmtt"

Main> crack (encode 3 "boxing wizards jump quickly")
"wjsdib rduvmyn ephk lpdxfgt"
```

6. EXTRA : Approximating π

The goal of this exercise is to compute π using known mathematical formulas. Use both types of list ranges ([1..n], [1,5,..,n]) and list comprehensions [a+1 | a <- as] to solve this exercise. Do not use explicit recursion!

Some Basic Functions

The following functions will come in handy when implementing the π -approximations, so it would be better to implement them first:

- Function sumf:: [Float] -> Float, which computes the sum of the elements of a list.⁴
- Function productf:: [Float] -> Float, which computes the product of the elements of a list.

Examples

```
Main> sumf []
0.0

Main> sumf [1, 5.0, 6.32]
12.32

Main> productf []
1.0

Main> productf [1, 5.0, 6.32]
31.6
```

Approximation 1

One way to compute π analytically is by using the following formula:

$$\pi(n) \approx 8 * \left(\frac{1}{1*3} + \frac{1}{5*7} + \frac{1}{9*11} + \dots + \frac{1}{(4n+1)*(4n+3)}\right)$$

⁴Float is a single-precision floating-point number, Double is a double-precision floating-point number.

The higher the value of n, the closer the value of $\pi(n)$ to the actual value of π :

$$\pi(0) = 8 * \left(\frac{1}{1*3}\right) = 2.6666667$$

$$\pi(1) = 8 * \left(\frac{1}{1*3} + \frac{1}{5*7}\right) = 2.8952382$$

$$\pi(2) = 8 * \left(\frac{1}{1*3} + \frac{1}{5*7} + \frac{1}{9*11}\right) = 2.9760463$$

$$\dots$$

$$\pi(100) = 8 * \left(\frac{1}{1*3} + \frac{1}{5*7} + \dots + \frac{1}{401*403}\right) = 3.1366422$$

• Implement function piSum:: Float \rightarrow Float that, given an n, approximates π using the above formula. You can assume that n is a natural number, e.g., 0.0, 1.0, 2.0, ...

Approximation 2

Similarly, we can also implement π using the following formula:

$$\pi(n) \approx 4 * \frac{2 * 4}{3 * 3} * \frac{4 * 6}{5 * 5} * \frac{6 * 8}{7 * 7} * \frac{8 * 10}{9 * 9} * \dots * \frac{(2n+2) * (2n+4)}{(2n+3)^2}$$

• Implement function piProd::Float -> Float that, given an n, approximates π using the above formula. You can assume that n is a natural number, e.g., 0.0, 1.0, 2.0, ...

7. EXTRA: Prime numbers

Write a function sieve::Int -> [Int] which returns all prime numbers smaller than the given number. Implement your function following Eratosthenes' algorithm.⁵ You can actually stop filtering the moment you've reached the square root of the input number. Ignore this optimization in your first implementation.

Examples

```
Main> sieve 20
[2,3,5,7,11,13,17,19]

Main> sieve 49
[2,3,5,7,11,13,17,19,23,29,31,37,41,43,47]
```

Extension

Haskell provides the sqrt function to find the square root of a number. However, this function requires an argument of type Double, whereas your argument is an Int. To convert this Int to a Double, use the fromIntegral function. The result of sqrt will also be a Double. To convert this Double back to an Int, use the floor function.⁶ Because these functions work with type classes, we give you versions of these

⁵See http://en.wikipedia.org/wiki/Sieve_of_Eratosthenes.

⁶See http://en.wikipedia.org/wiki/Floor_and_ceiling_functions#Examples for more information about the floor func-

functions with the right types: $\mathtt{sqrtMono}$, $\mathtt{i2d}$, and $\mathtt{floorMono}$.

```
sqrtMono :: Double -> Double
sqrtMono = sqrt

i2d :: Int -> Double
i2d = fromIntegral

floorMono :: Double -> Int
floorMono = floor
```

Using these functions, try to write the function floorSquare which *floors* the square root of the given Int argument. Use this floorSquare function to make sieve more efficient.