# Simply-Typed Lambda Calculus with Extensions

CS231

# 1 Language of Booleans and Integers

# 1.1 Syntax

## 1.2 Small-Step Operational Semantics

$$\frac{\phantom{a}}{\text{if true then } \mathsf{t}_2 \text{ else } \mathsf{t}_3 \longrightarrow \mathsf{t}_2} \tag{E-IFTRUE}$$

$$\frac{\texttt{t}_1 \longrightarrow \texttt{t}_1'}{\text{if } \texttt{t}_1 \text{ then } \texttt{t}_2 \text{ else } \texttt{t}_3 \longrightarrow \text{if } \texttt{t}_1' \text{ then } \texttt{t}_2 \text{ else } \texttt{t}_3} \tag{E-IF)}$$

$$\frac{\texttt{t}_1 \longrightarrow \texttt{t}_1'}{\texttt{t}_1 + \texttt{t}_2 \longrightarrow \texttt{t}_1' + \texttt{t}_2} \tag{E-PLUS1}$$

$$\frac{\mathtt{t}_2 \longrightarrow \mathtt{t}_2'}{\mathtt{v}_1 + \mathtt{t}_2 \longrightarrow \mathtt{v}_1 + \mathtt{t}_2'} \tag{E-PLUS2}$$

$$\frac{\mathbf{n} = \mathbf{n}_1 \, \llbracket + \rrbracket \, \mathbf{n}_2}{\mathbf{n}_1 + \mathbf{n}_2 \longrightarrow \mathbf{n}} \tag{E-PLUSRED}$$

$$\frac{\mathtt{t}_1 \longrightarrow \mathtt{t}_1'}{\mathtt{t}_1 \ > \ \mathtt{t}_2 \longrightarrow \mathtt{t}_1' \ > \ \mathtt{t}_2} \tag{E-GT1}$$

$$\frac{\mathsf{t}_2 \longrightarrow \mathsf{t}_2'}{\mathsf{v}_1 > \mathsf{t}_2 \longrightarrow \mathsf{v}_1 > \mathsf{t}_2'} \tag{E-GT2}$$

$$\frac{\mathbf{v} = \mathbf{n}_1 \ [\![ > \!]\!] \ \mathbf{n}_2}{\mathbf{n}_1 \ > \ \mathbf{n}_2 \longrightarrow \mathbf{v}}$$
 (E-GTRED)

## 1.3 Static Type System

$$\frac{}{\text{true: Bool}} \qquad \qquad \text{(T-True)} \qquad \qquad \frac{}{\text{false: Bool}} \qquad \qquad \text{(T-False)}$$

$$\frac{\mathsf{t}_1 \colon \mathsf{Int} \qquad \mathsf{t}_2 \colon \mathsf{Int}}{\mathsf{t}_1 \, + \, \mathsf{t}_2 \colon \mathsf{Int}} \tag{T-PLUS}$$

$$\frac{\mathsf{t}_1 \colon \mathsf{Int} \qquad \mathsf{t}_2 \colon \mathsf{Int}}{\mathsf{t}_1 \, > \, \mathsf{t}_2 \colon \mathsf{Bool}} \tag{T-GT}$$

# 2 Simply-Typed Lambda Calculus

# 2.1 Syntax

```
t ::= x | function x:T \rightarrow t | t t v ::= function x:T \rightarrow t T ::= T_1 \rightarrow T_2
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#### 2.2 Substitution

$$\begin{split} & [x \mapsto v] \, x = v \\ & [x \mapsto v] \, x' = x', \text{where } x \neq x' \\ & [x \mapsto v] \, \text{function } x \colon T \to t_0 = \text{function } x \colon T \to t_0 \\ & [x \mapsto v] \, \text{function } x_0 \colon T \to t_0 = \text{function } x_0 \colon T \to [x \mapsto v] t_0, \text{where } x \neq x_0 \\ & [x \mapsto v] t_1 \, t_2 = [x \mapsto v] t_1 \, [x \mapsto v] t_2 \end{split}$$

# 2.3 Small-Step Operational Semantics

$$\frac{}{((function x:T \rightarrow t) v) \longrightarrow [x \mapsto v]t}$$
 (E-APPBETA)

$$\frac{\mathtt{t}_1 \longrightarrow \mathtt{t}_1'}{\mathtt{t}_1 \ \mathtt{t}_2 \longrightarrow \mathtt{t}_1' \ \mathtt{t}_2} \qquad \text{(E-APP1)} \qquad \frac{\mathtt{t}_2 \longrightarrow \mathtt{t}_2'}{\mathtt{v}_1 \ \mathtt{t}_2 \longrightarrow \mathtt{v}_1 \ \mathtt{t}_2'} \qquad \text{(E-APP2)}$$

# 2.4 Static Type System

 $\Gamma$  is a finite function from variable names to types.

$$\frac{\Gamma(x) = T}{\Gamma \vdash x : T}$$
 (T-VAR)

$$\frac{\Gamma, \text{x:} T_1 \vdash \text{t:} T_2}{\Gamma \vdash \text{function x:} T_1 \rightarrow \text{t:} T_1 \rightarrow T_2} \tag{T-Fun}$$

$$\frac{\Gamma \vdash \mathsf{t}_1 : \mathsf{T}_2 \to \mathsf{T} \qquad \Gamma \vdash \mathsf{t}_2 : \mathsf{T}_2}{\Gamma \vdash \mathsf{t}_1 \ \mathsf{t}_2 : \mathsf{T}} \tag{T-APP}$$

#### 3 Extensions

We augment our language with a unit value, pairs, tagged unions, let, letrec.

#### 3.1 Syntax

t ::= () | (t,t) | fst t | snd t| left t | right t | (match t with left x -> t | right x -> t) | let x=t in t | letrec x=v in t v ::= () | (v,v) | left v | right vT ::= True | False | T  $\wedge$  T | T  $\vee$  T

#### 3.2 Small-Step Operational Semantics

$$\frac{t_1 \longrightarrow t_1'}{(t_1, t_2) \longrightarrow (t_1', t_2)} \qquad \text{(E-PAIR1)} \qquad \frac{t_2 \longrightarrow t_2'}{(v_1, t_2) \longrightarrow (v_1, t_2')} \qquad \text{(E-PAIR2)}$$

$$\frac{t \longrightarrow t'}{\text{fst } t \longrightarrow \text{fst } t'} \qquad \text{(E-FST)} \qquad \frac{\text{fst } (v_1, v_2) \longrightarrow v_1}{\text{fst } (v_1, v_2) \longrightarrow v_2} \qquad \text{(E-SNDRED)}$$

$$\frac{t \longrightarrow t'}{\text{snd } t \longrightarrow \text{snd } t'} \qquad \text{(E-SND)} \qquad \frac{t \longrightarrow t'}{\text{right } t \longrightarrow \text{right } t'} \qquad \text{(E-RIGHT)}$$

 $\frac{t\longrightarrow t'}{\text{match t with left } x_1 \ -> \ t_1 \ | \ \text{right } x_2 \ -> \ t_2 \longrightarrow \text{match } t' \ \text{with left } x_1 \ -> \ t_1 \ | \ \text{right } x_2 \ -> \ t_2}$ 

 $\frac{}{\text{match left v with left } x_1 \ \text{->} \ t_1 \ | \ \text{right } x_2 \ \text{->} \ t_2 \longrightarrow [x_1 \ \mapsto \ \text{v}]t_1} \, (\text{E-MATCHLEFT})}$ 

 $\frac{}{\text{match right v with left } x_1 \text{ -> } t_1 \text{ | right } x_2 \text{ -> } t_2 \longrightarrow [x_2 \mapsto \text{vlt}_2} \text{ (E-MATCHRIGHT)}$ 

$$\frac{\texttt{t}_1 \longrightarrow \texttt{t}_1'}{\texttt{let x=t}_1 \text{ in } \texttt{t}_2 \longrightarrow \texttt{let x=t}_1' \text{ in } \texttt{t}_2} (\texttt{E-LET}) \qquad \qquad \frac{\texttt{let x=v in t} \longrightarrow [\texttt{x} \mapsto \texttt{v}] \texttt{t}}{\texttt{let x=v in t} \longrightarrow [\texttt{x} \mapsto \texttt{v}] \texttt{t}} (\texttt{E-LETRED})$$

(E-LETREC) letrec x=v in t  $\longrightarrow$  let x=[x  $\mapsto$  letrec x=v in x]v in t

#### 3.3 Static Type System

$$\Gamma \vdash ()$$
 : True (T-TRUE)

$$\frac{\Gamma \vdash t : False}{\Gamma \vdash t : T}$$
 (T-False)

$$\frac{\Gamma \vdash \mathsf{t}_1 \,:\, \mathsf{T}_1 \qquad \Gamma \vdash \mathsf{t}_2 \,:\, \mathsf{T}_2}{\Gamma \vdash (\mathsf{t}_1, \mathsf{t}_2) \,:\, \mathsf{T}_1 \,\wedge\, \mathsf{T}_2} \tag{T-PAIR}$$

$$\frac{\Gamma \vdash \texttt{t} : \texttt{T}_1 \ \land \ \texttt{T}_2}{\Gamma \vdash \texttt{fst} \ \texttt{t} : \texttt{T}_1} \qquad \qquad (\texttt{T-SND}) \qquad \qquad \frac{\Gamma \vdash \texttt{t} : \texttt{T}_1 \ \land \ \texttt{T}_2}{\Gamma \vdash \texttt{snd} \ \texttt{t} : \texttt{T}_2} \qquad \qquad (\texttt{T-SND})$$

$$\frac{\Gamma \vdash \texttt{t} : \texttt{T}_1}{\Gamma \vdash \texttt{left t} : \texttt{T}_1 \ \lor \ \texttt{T}_2} \qquad (\texttt{T-LEFT}) \qquad \qquad \frac{\Gamma \vdash \texttt{t} : \texttt{T}_2}{\Gamma \vdash \texttt{right t} : \texttt{T}_1 \ \lor \ \texttt{T}_2} \qquad (\texttt{T-RiGHT})$$

$$\frac{\Gamma \vdash \texttt{t} : \texttt{T}_1 \ \lor \ \texttt{T}_2 \qquad \Gamma, \texttt{x}_1 : \texttt{T}_1 \vdash \texttt{t}_1 : \texttt{T} \qquad \Gamma, \texttt{x}_2 : \texttt{T}_2 \vdash \texttt{t}_2 : \texttt{T}}{\Gamma \vdash \texttt{match} \ \texttt{t} \ \texttt{with} \ \texttt{left} \ \texttt{x}_1 \ -> \ \texttt{t}_1 \ | \ \texttt{right} \ \texttt{x}_2 \ -> \ \texttt{t}_2 : \texttt{T}} \tag{T-MATCH)}$$

$$\frac{\Gamma \vdash \mathsf{t}_1 \,:\, \mathsf{T}_1 \qquad \Gamma, \mathsf{x} \colon \mathsf{T}_1 \vdash \mathsf{t}_2 \,:\, \mathsf{T}}{\Gamma \vdash \mathsf{let} \ \mathsf{x} = \mathsf{t}_1 \ \mathsf{in} \ \mathsf{t}_2 \,:\, \mathsf{T}} \quad (\mathsf{T}\text{-}\mathsf{LET}) \qquad \frac{\Gamma, \mathsf{x} \colon \mathsf{T}_1 \vdash \mathsf{v}_1 \,:\, \mathsf{T}_1 \qquad \Gamma, \mathsf{x} \colon \mathsf{T}_1 \vdash \mathsf{t}_2 \,:\, \mathsf{T}}{\Gamma \vdash \mathsf{let}\mathsf{rec} \ \mathsf{x} = \mathsf{v}_1 \ \mathsf{in} \ \mathsf{t}_2 \,:\, \mathsf{T}} \, (\mathsf{T}\text{-}\mathsf{LETREC})$$