

Written examination, May 29, 2020

Course: Modelling and Analysis of Real-Time Systems

Course no. 02224

Aids allowed: All

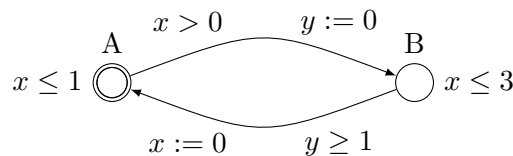
Exam duration: 2 hours

Weighting: PROBLEM 1: approx. 25 % PROBLEM 3: approx. 40 %
 PROBLEM 2: approx. 35 %

You are encouraged to briefly justify your answers. If you happen to be in doubt about the precise meaning of a question, you should write down how you choose to understand it.

PROBLEM 1 (approx. 25 %)

In this problem we consider the following timed automaton P with two locations: A and B.



The timed automaton P

Question 1.1:

- Give two states of the transition system underlying the timed automaton P .
- Give two transitions of the transition system underlying the timed automaton P , where one transition should be a discrete transition and the other a progress of time.

Question 1.2:

A region automaton is a finite automaton constructed on the basis of a timed automaton.

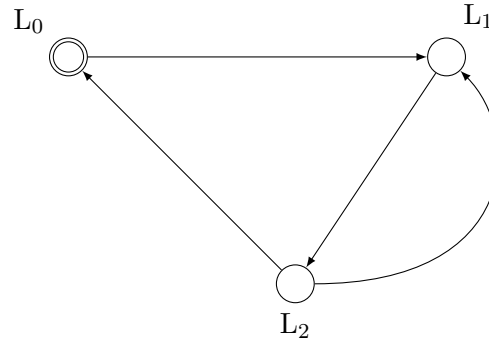
- A *clock region* may be of one of the following types: A *point*, a bounded *line segment*, an unbounded *half line (ray)*, a *bounded area*, or an *unbounded area*. Present an example of each region type for P .
- How many clock regions are there of the different types and how many in total?
- Present two states of the region automaton for P .
- Provide three consecutive transitions of the region automaton for P , where at least one transition should be a discrete step. [Transitions are consecutive if the end state of one transition is the start state of the next one.]

Question 1.3:

Let the automaton P' be like P with the A-invariant $x \leq 1$ removed. Explain why P' may not be considered a sound model.

PROBLEM 2 (approx. 35 %)

Consider the following timed automaton Q with three locations L_0 , L_1 and L_2 .



The timed automaton Q

Question 2.1: Extend this automaton Q so that

- a stay in L_1 lasts for at most 2 seconds (use a clock x to express this property),
- a stay in location L_2 lasts for at least 3 seconds and at most 5 seconds (the clock x should be reused to express this property), and
- if L_0 is left, then L_0 must be reentered between 15 and 20 seconds later (use a clock y to express this property).

Question 2.2: For each of the following Uppaal queries, determine whether it is satisfied by your extension of Q . The answers must be accompanied with brief justifications.

- $E \lt \!> Q.L_2$
- $A[] (Q.L_2 \text{ imply } y > 0)$
- $Q.L_2 \dashrightarrow Q.L_1$
- $A[] (y > 17 \text{ imply } (Q.L_0 \text{ or } Q.L_2))$

Question 2.3: Formulate each of the following properties as Uppaal queries:

- It is possible that L_2 is never reached.
- A stay in L_1 will always lead to L_2 .
- It is possible that the value of y will never exceed 15.
- In L_2 , the value of y cannot exceed the value of x by more than 17.

Question 2.4:

- For your extension of Q , determine the minimum and maximum number of times the transition from L_2 to L_1 may be taken from the moment where L_0 is left till L_0 is reentered again.
- How would you use Uppaal to verify these minimum and maximum numbers?
Hint: Further extension of Q is allowed.

PROBLEM 3 (approx. 40 %)

The questions in this problem can be solved independently of each other.

A real-time system to be run on a single-processor computer has three periodic tasks **a**, **b** and **c** with the following parameters:

	T	C
a	8	3
b	20	7
c	40	3

The deadline of each task is equal to its period. Initially, the tasks are assumed to be independent.

Question 3.1:

- Calculate the load of each task and the total load of the system.
- Discuss whether schedulability of the task set using fixed-priority scheduling (FPS) can be guaranteed based on the load figures.

In the given system, the set of task is to be scheduled by a preemptive fixed-priority scheduler (FPS) using rate monotonic priority assignment.

Question 3.2:

Calculate the response time for each of the three tasks. Show all intermediate response time estimates.

Now, a protected, shared resource M is to be used by all three tasks. Each task is always going to use the resource M for exactly 1 of its computation time units C . Furthermore, the usage of M will take place at the end of the computation for each of the three tasks.

Recall that the *blocking time* B_t of a task t is the maximum time the task t may experience being suspended while *lower priority tasks* execute.

Question 3.3:

- Explain why the blocking times of tasks **b** and **c** are 1 and 0 respectively.
- Determine the blocking time B_a of task **a** and justify your answer.
[Notice that blocking of task **a** may be subject to *priority inversion*.]
- Explain why the task set using M is not schedulable for the given priority assignment.
- Show that the task set using M may become schedulable by choosing another priority assignment:

State the new priority assignment, determine the blocking time for each task and calculate the response time for each task taking its blocking time into account.