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Course: Computer  
Networks

Ans: to Q. No. 1

$$M(x) = x^7 + x^6 + x^3 + x^1 = 10011010$$

$$P(x) = x^3 + x^2 + 1 = 1101$$

Since,  $P$  consists 4 bits, appends,  $k = 3$  bits

$k = 000$ ; at the end of  $M$

$$\therefore S = 10011010000$$

At sender,

$$P \rightarrow 1101 \mid 10011010000 \longleftrightarrow S$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1001 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1000 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1011 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1100 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

At receiver,

Transmit sequence or  
codeword  $T = 10011010101$

$$\begin{array}{r} 1000 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1010 \leftarrow R \end{array}$$

## Ans to Q. No. 2

(a)

Suppose that there is no error in the received sequence, Hence,  $r = c = 1101010$

Hence  $(n, k) = (7, 4)$

$n = 7 = \text{code length} = r + k \text{ bits}; n = k + r$

$$\begin{aligned} r &= n - k \\ &= 7 - 4 \\ &= 3 \text{ bits} \end{aligned}$$

$m = 4; k \text{ bits}$

So,  $c = [\underbrace{110}_M \underbrace{1010}_M]$

message,  $M = 1010$

$$\begin{aligned} \therefore \text{Code word, } c &= M \times G \\ &= [1010] \times \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix} \\ &= [0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0] \end{aligned}$$

$$1 \times 1 \oplus 0 \times 0 \oplus 1 \times 1 \oplus 1 \times 0 = 1 \oplus 0 \oplus 1 \oplus 0 = 0$$

$$1 \times 1 \oplus 0 \times 1 \oplus 1 \times 1 \oplus 0 \times 0 = 0$$

$$1 \times 1 \oplus 0 \times 1 \oplus 1 \times 0 \oplus 0 \times 1 = 1$$

$$1 \times 1 \oplus 0 \times 0 \oplus 1 \times 0 \oplus 0 \times 0 = 1$$

$$1 \times 0 \oplus 0 \times 1 \oplus 1 \times 0 \oplus 0 \times 0 = 0$$

$$1 \times 0 \oplus 0 \times 0 \oplus 1 \times 1 \oplus 0 \times 0 = 1$$

$$1 \times 0 \oplus 0 \times 0 \oplus 1 \times 0 \oplus 0 \times 1 = 0$$

(b)

We know,

Parity-check matrix,  $H = [I_3, P_{k \times 2}]$

$$P_{k \times 2} = P_{3 \times 4} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}; I_3 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$= P_{3 \times 4}^T = \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & 0 \\ 1 & 1 & 0 & 1 \end{bmatrix}$$

$$H = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{bmatrix}$$

(c)

Syndrome  $\vec{s} = rHT^T$

$$r = [1 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0]$$

$$H^T = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

$$= [1 \ 1 \ 1]$$

Here, the non-zero syndrome an error occurred reception. Here,  $s = [1 \ 1 \ 1]$  which is in fourth row. So, in the fourth element in receive sequence will be after from 1 to 0. So, the valid codeword is  $= [1 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0]$

Ans: to: Q: No: 3

IPv4 address: 184, 93, 57, 31

Convert it into IPv6

Representing each octet with 8 bits binary

$$184 = \underbrace{1011}_B \underbrace{1000}_8 = B8$$

$$93 = \underbrace{0101}_5 \underbrace{1101}_D = 5D$$

$$57 = \underbrace{0011}_3 \underbrace{1001}_9 = 39$$

$$31 = \underbrace{0001}_1 \underbrace{1111}_F = 1F$$

IPv6 address will be: 0::B85D:391F

Ans: to: Q: No: 4

MAC address: 99 E7: F2 5A: 03 61

Convert local address

$$9957 = 0100100111100111$$

$$= 0100101111100111$$

$$= 9BE7$$

So, the local address is,

FE80:: 9B57: F2FF: FE5A: 0361