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Ans: to Q: No: 1

$$M(x) = x^7 + x^9 + x^3 + x^1 = 10011010$$

$$P(x) = x^3 + x^2 + 1 = 1101$$

Since, P consists 4 bits, append, $k = 3$ bits

$k = 000$; at the end of M

$$\therefore S = 10011010000$$

At sender,

$$P \rightarrow 1101 \mid 10011010000 \longleftrightarrow S$$

$$\begin{array}{r} 1101 \\ \hline 1001 \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1000 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1011 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline \end{array}$$

$$\begin{array}{r} 1101 \\ \hline 1100 \end{array}$$

At receiver,

Transmit sequence or
codeword $T = 10011010101$

$$\begin{array}{r} 1000 \\ 1101 \\ \hline 1010 \end{array} \leftarrow R$$

Anw. to Q. No' 2

(a)

Suppose that there is no error in the received sequence, hence, $r = c = 1101010$

Hence $(n, k) = (7, 4)$

$n = 7 = \text{code length} = q + k \text{ bits}$; $n = k + q$

$$\begin{aligned}q &= n - k \\&= 7 - 4 \\&= 3 \text{ bits}\end{aligned}$$

$m = 4$; k bits

so, $c = \begin{bmatrix} 1 & 1 & 0 & 1 & 0 & 1 & 0 \\ 2 & & M \end{bmatrix}$

message, $M = 1010$

$$\begin{aligned}\therefore \text{codeword, } c &= M \times G \\&= [1010] \times \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 & 1 \end{bmatrix} \\&= [0 \ 0 \ 1 \ 1 \ 0 \ 1 \ 0]\end{aligned}$$

$$1 \times 1 \oplus 0 \times 0 \oplus 1 \times 1 \oplus 1 \times 0 = 1 \oplus 0 \oplus 1 \oplus 0 = 0$$

$$1 \times 1 \oplus 0 \times 1 \oplus 1 \times 1 \oplus 0 \times 0 = 0$$

$$1 \times 1 \oplus 0 \times 1 \oplus 1 \times 0 \oplus 0 \times 1 = 1$$

$$1 \times 1 \oplus 0 \times 0 \oplus 1 \times 0 \oplus 0 \times 0 = 1$$

$$1 \times 0 \oplus 0 \times 1 \oplus 1 \times 0 \oplus 0 \times 0 = 0$$

$$1 \times 0 \oplus 0 \times 0 \oplus 1 \times 1 \oplus 0 \times 0 = 1$$

$$1 \times 0 \oplus 0 \times 0 \oplus 1 \times 0 \oplus 0 \times 1 = 0$$

⑥

We know,

Parity - check matrix, $H = T[I_2, P_{k \times 2}]$

$$P_{k \times 2} = P_{3 \times 2} = \begin{bmatrix} 1 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}; I_2 = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$= P_{3 \times 2}^T = \begin{bmatrix} 1 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \end{bmatrix}$$

$$H = \begin{bmatrix} 1 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{bmatrix}$$

⑦

Syndrome $\hat{s} = rH^T$

$$\hat{s} = [1 \ 1 \ 0 \ 1 \ 0 \ 1 \ 0] \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{bmatrix}$$

$$= [1 \ 1 \ 1]$$

Here, the non-zero syndrome is an erroneous reception. Here, $\hat{s} = [1 \ 1 \ 1]$ which is in fourth element in receive sequence. So, in the fourth element from 1 to 0 sequence will be after from 1 to 0 so, the valid codeword is $[1 \ 1 \ 0 \ 0 \ 0 \ 0 \ 0]$

Ans: Q.No. 3

IPv4 address: 189, 93, 57, 31
convert it into IPv6

Presenting each octet with 8 bits binary

$$189 = \underbrace{1011}_B \underbrace{1000}_8 = B8$$

$$93 = \underbrace{0101}_5 \underbrace{1101}_0 = 5D$$

$$57 = \underbrace{0011}_3 \underbrace{1001}_9 = 39$$

$$31 = \underbrace{0001}_1 \underbrace{1111}_F = 1F$$

IPv6 address will be: :: B85D:391F

Ans: Q.No. 9

MAC address: 99 E7 F2 5A 03 61

Convert local address:

$$\begin{aligned}9957 &= 01001001 \ 1110 \ 0111 \\&= 01001011 \ 1110 \ 0111 \\&= 9BE7\end{aligned}$$

So, the local address is,

FE80::9B57:F2FF:FE5A:0361