#### Lecture 2.1

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## Algorithms: Efficiency & Analysis

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#### Algorithms

- An algorithm is a step-by-step procedure used to solve a problem. But making sure the developer is using the most efficient algorithm is very crucial no matter how fast computers become or how cheap memory gets.
- In order to determine how efficiently an algorithm solves a problem, we need to analyze the algorithm.
- Order helps group algorithms according to their eventual behavior.

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#### Algorithms (cont.)

- A computer program is composed of individual modules, understandable by a computer, that solve specific tasks (such as sorting).
- The concentration here is not on the design of entire programs, but rather on the design of the individual modules that accomplish the specific tasks.
- These specific tasks are called **problems**. A problem is a question to which we seek an answer. 

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#### Algorithms (cont.)

- A problem may contain variables that are not assigned specific values in the statement of the problem. These variables are called parameters to the problem.
- Because a problem contains **parameters**, it represents a class of problems, one for each assignment of values to the parameters. Each specific assignment of values to the parameters is called an **instance** of the problem.

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#### Algorithms (cont.)

■ To produce a computer program that can solve all instances of a problem, we must specify a general step-by-step procedure for producing the solution to each instance. This step-by-step procedure is called an algorithm. We say that the algorithm solves the problem.

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# The Importance of Developing Efficient Algorithms

- Efficiency is an important consideration when working with algorithms. A solution is said to be efficient if it solves the problem within the required resource constraints.
- Since one problem can be solved by many different algorithms, it may be important to determine which one is the most efficient. Eg. Searching with sequential search and binary search.

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#### Sequential Search vs Binary Search

- The sequential search algorithm begins at the first position in the array and looks at each value in turn until the item is found.
- The binary search algorithm first compares x with the middle item of the array. If they are equal, the algorithm is done. If x is smaller than the middle item, then x must be in the first half of the array (if it is present at all), and the algorithm repeats the searching procedure on the first half of the array

#### Binary Search (cont.)

■ (That is, x is compared with the middle item of the first half of the array. If they are equal, the algorithm is done, etc.) If x is larger than the middle item of the array, the search is repeated on the second half of the array. This procedure is repeated until x is found or it is determined that x is not in the array.

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#### Binary Search (cont.)



Figure 1.1: The array items that Binary Search compares with x when x is large than all the items in an array of size 32. The items are numbered according to the order in which they are compared.

#### Sequential Search

- Write the algorithm of sequential search for array of integers.
- Solution will be discussed in the lecture

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#### Sequential Search in Java

```
static int sequentialSearch(final Comparable key,
final ArrayList<? extends Comparable> in) {
   int loc = 0;
   while(loc<in.size() && in.get(loc).compareTo(key) != 0) {
      loc++;
   }
   if(loc>=in.size())
      loc = -1;
   return loc;
}
```

#### Binary Search

- Write the algorithm of binary search for array of integers.
- Solution will be discussed in the lecture

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#### Binary Search in Java

```
static int binarySearch
final Comparable key,
final ArrayList<? extends Comparable> in) {
   int low, mid, high; //indices
   int loc = -1;

   low = 0; high = in.size()-1;

   while(low<=high && loc == -1) {
      mid = (int)Math.floor((low+high)/2.0);
      if(in.get(mid).compareTo(key) == 0)
            loc = mid;
      else if(in.get(mid).compareTo(key) > 0)
            high = mid-1;
      else
            low = mid+1;
   }

   return loc;
}
```

#### Efficiency comparison

- To compare Sequential Search to Binary Search, Sequential Search does *n* comparisons to determine that *x* is not in the array of size *n*. If *x* is in the array, the number of comparisons is no greater than *n*.
- A Binary Search is an algorithm for locating the position of an element in a sorted list. The method reduces the number of elements that need to be examined by two each time.
- Binary Search appears to be much more efficient than Sequential Search

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#### Efficiency comparison (cont.)

Compare the number of comparisons for the worst case:

Array Size	Sequential Search	Binary Search
32	32	6
64	64	7
128	128	8
n	n	log <sub>2</sub> n + 1
1,024	1,024	11
1,048,576	1,048,576	21

■ Therefore Binary Search appears to be much more efficient than Sequential Search, based on the number of comparisons done.

#### Recursive vs iterative Fibonacci

- The Fibonacci sequence is given as:
- Fib(n) = Fib(n-1) + Fib(n-2) for  $n \ge 2$ ; Fib(0) = 0, Fib(1) = 1 Example: 0, 1, 1, 2, 3, 5, 8, ...
- Recursive Algorithm (*n*<sup>th</sup> Fibonacci term):

```
int fib(int n)
{
  if (n <= 1) return n;
  else
  return fib(n-1) + fib(n-2);
}</pre>
```

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## Iterative (n<sup>th</sup> Fibonacci term)

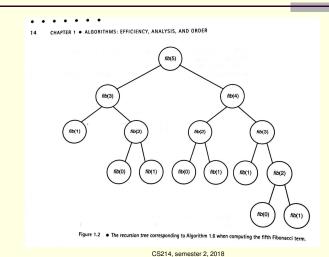
```
int fib2(int n)
{
  index i;
  int f[0..n]; //array f[0] = 0;
  if (n > 0) {
    f[1] = 1;

for (i = 2; i <= n; i++)
  f[i] = f[i-1] + f[i-2];
    }
  return f[n];
}</pre>
```

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#### Recursive Fibonacci computations



### Efficiency comparison

■ Recursive vs iterative Fibonacci.

n	Recursive	Iterative
0	1	1
1	1	2
2	3	3
3	5	4
4	9	5
n	> 2 <sup>n/2</sup>	n + 1
40	> 1,048,576	41

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