

Queues

This lesson is borrowed from the following:

Reference

CS 367 – Introduction to Data Structures

<http://pages.cs.wisc.edu/~mattmcc/cs367/notes/Queues.ppt>

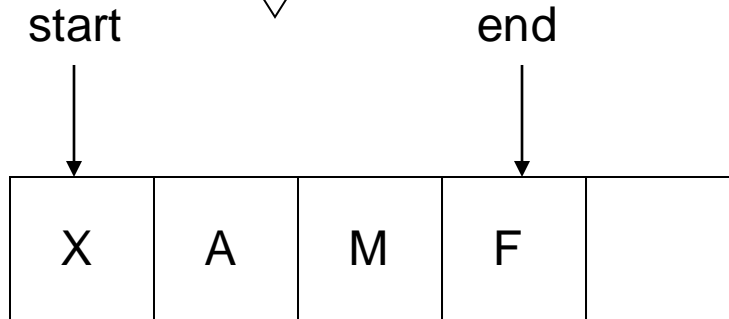
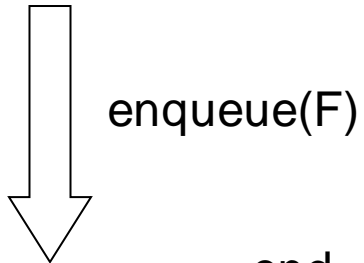
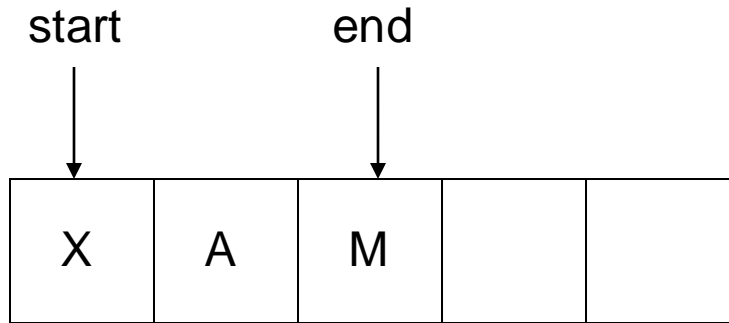
Queue

- A queue is a data structure that stores data in such a way that the last piece of data stored, is the last one retrieved
 - also called First-In, First-Out (FIFO)
- Only access to the stack is the first and last element
 - consider people standing in line
 - they get service in the order that they arrive

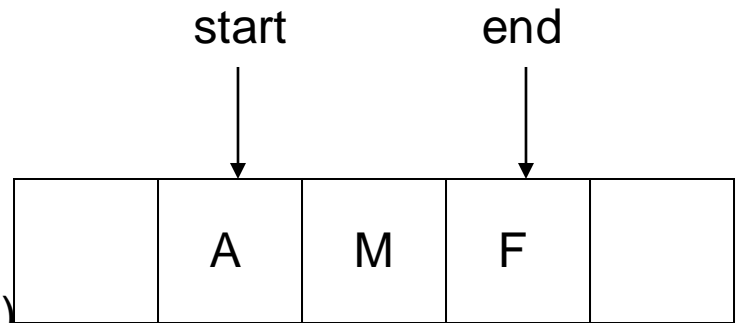
Queues

- *Enque*
 - operation to place a new item at the tail of the queue
- *Dequeue*
 - operation to remove the next item from the head of the queue

Queue



item = dequeue()
item = X



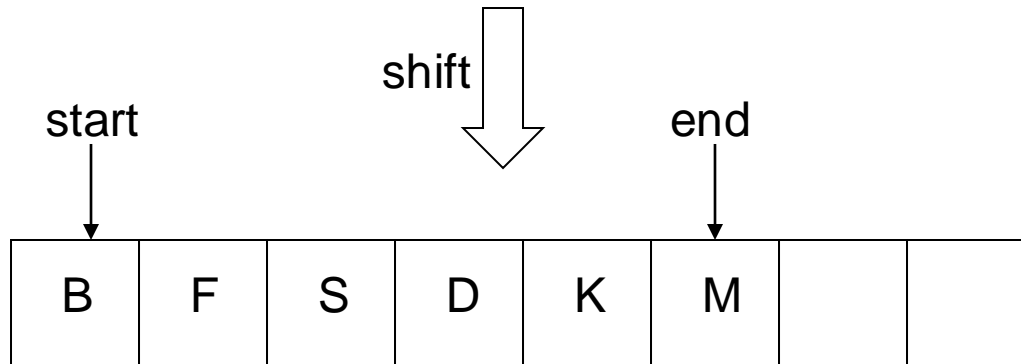
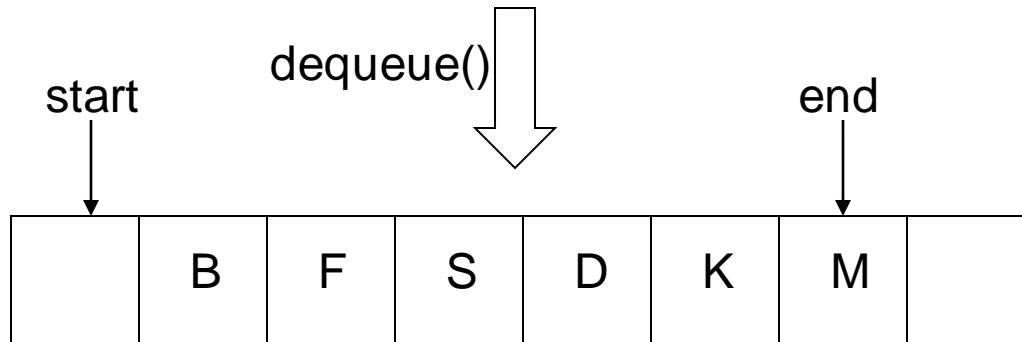
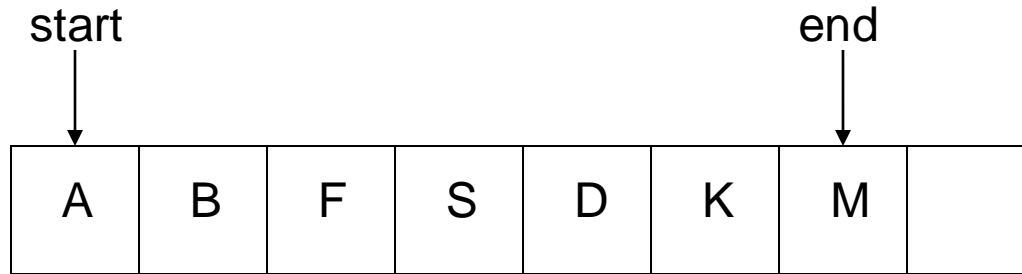
Implementing a Queue

- At least two common methods for implementing a queue
 - array
 - linked list
- Which method to use depends on the application
 - advantages? disadvantages?

Regular Linear Array

- In a standard linear array
 - 0 is the first element
 - `array.length - 1` is the last element
- All objects removed would come from element 0
- All objects added would go one past the last currently occupied slot
- To implement this, when an object is removed, all elements must be shifted down by one spot

Regular Linear Array



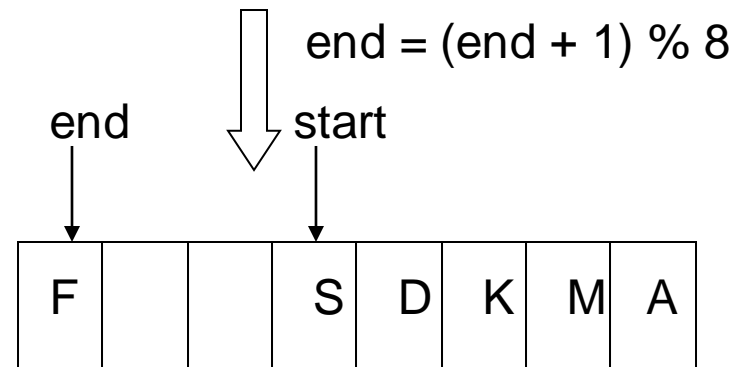
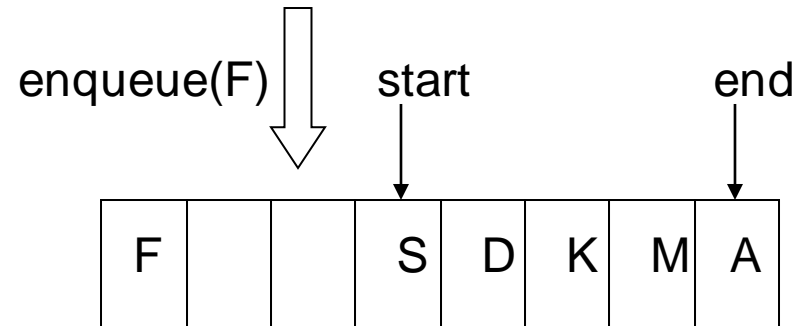
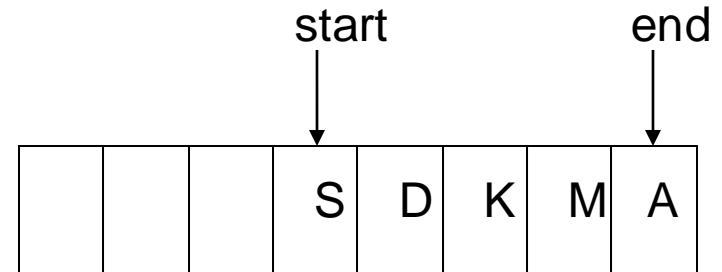
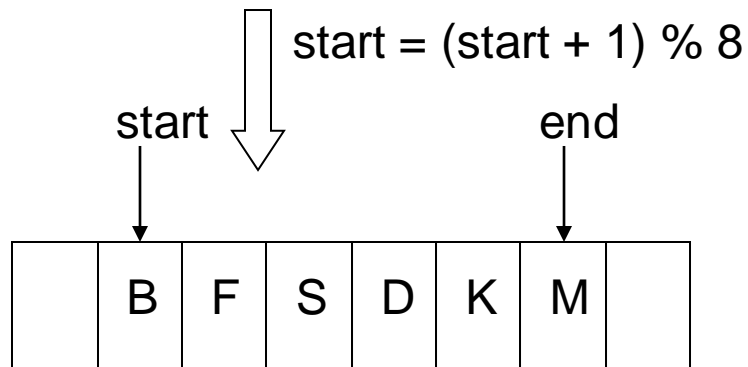
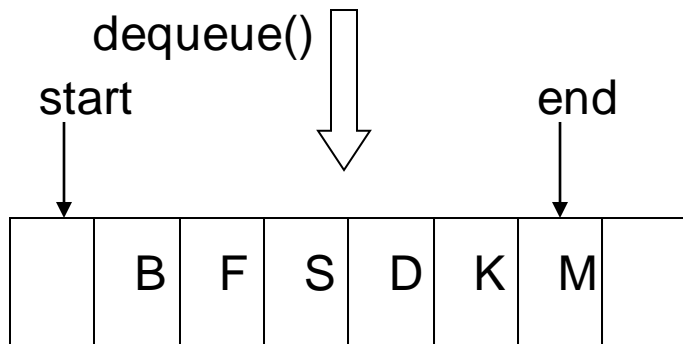
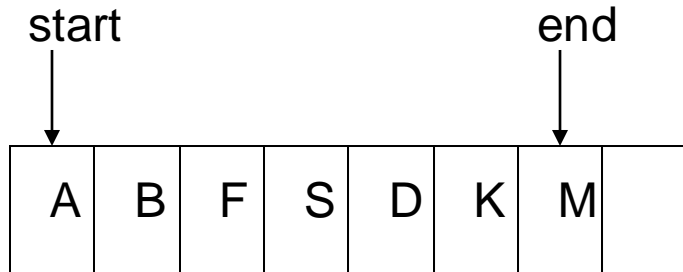
Regular Linear Array

- Very expensive structure to use for a queue
 - shifting all the data down by one is very time consuming
- Would prefer to let the data “wrap-around”
 - the start would not always be zero
 - it would be the first occupied cell
 - the last item may actually appear in the array before the first item
 - this is called a *circular array*

Circular Array

- Need to keep track of the index that holds the first item
- Need to keep track of the index that holds the last item
- The “wrap-around” is accomplished through the use of the *mod* operator (%)
 - $\text{index} = (\text{end} + 1) \% \text{array.length}$
 - assume $\text{array.length} = 5$ and $\text{end} = 4$
 - then: $\text{index} = (4 + 1) \% 5 = 0$

Circular Array



Circular List

- If the list is empty, *start* and *end* would refer to the same spot
- If the list is full, *start* and *end* would refer to the same spot
- How do you tell the difference between an empty and a full list?
 - if the list is empty, make *start* and *end* refer to -1
 - then if *start* and *end* both refer to an element greater than -1, the list is full

Implementing Queues: Array

- Advantage:
 - best performance
- Disadvantage:
 - fixed size
- Basic implementation
 - initially empty circular array
 - two fields: *start* and *end*
 - where the next data goes in, and the next data comes out
 - if array is full, *enqueue()* returns false
 - otherwise the data is added to the queue
 - if array is empty, *dequeue()* returns null
 - otherwise removes the start and returns it

Queue Class (array based)

```
class QueueArray {  
    private Object[ ] queue;  
    private int start, end;  
    public QueueArray(int size) {  
        queue = new Object[size];  
        start = end = -1;  
    }  
    public boolean enqueue(Object data);  
    public Object dequeue();  
    public void clear();  
    public boolean isEmpty();  
    public boolean isFull();  
}
```

enqueue() Method (array based)

```
public boolean enqueue(Object data) {  
    if(((end + 1) % queue.length) == start)  
        return false; // queue is full  
  
    // move the end of the queue and add the element  
    end = (end + 1) % queue.length;  
    queue[end] = data;  
    if(start == -1) { start = 0; }  
    return true;  
}
```

dequeue() Method (array based)

```
public Object dequeue() {  
    if(start == -1)  
        return null; // empty list  
  
    // get the object, update the start, and return the object  
    Object tmp = queue[start];  
    if(start == end)  
        start = end = -1;  
    else  
        start = (start + 1) % queue.length;  
    return tmp;  
}
```

Notes on *enqueue()* and *dequeue()*

- Just implementing a circular list
 - if start and end equal -1, the list is empty
 - if start and end are the same and not equal to -1, there is only one item in the list
 - if the end of the list is one spot “behind” the start of the list, the list is full
 - `if(((end + 1) % queue.length) == start) { ... }`
 - always remove from the start and add to the end
 - make sure to move *end* before adding
 - make sure to move *start* after removing

Remaining Methods (array based)

```
public void clear() {  
    start = end = -1;  
}
```

```
public boolean isEmpty() {  
    return start == -1;  
}
```

```
public boolean isFull() {  
    return ((end + 1) % queue.length) == start;  
}
```

Implementing a Stack: Linked List

- Advantages:
 - can grow to an infinite size
 - lists are well suited for implementing a queue
 - makes things very easy
- Disadvantage
 - potentially slower than an array based queue
 - can grow to an infinite size
- Basic implementation
 - add new node to the tail of the list
 - remove a node from the head of the list

Queue Class (list based)

```
class QueueList {  
    private LinkedList queue;  
    public QueueList() {  
        queue = new LinkedList();  
    }  
    public boolean enqueue(Object data) { list.addTail(data); }  
    public Object dequeue() { return list.deleteHead(); }  
    public void clear() { list.clear(); }  
    public boolean isEmpty() { return list.isEmpty(); }  
}
```

Additional Notes

- It should appear obvious that linked lists are very well suited for queues
 - *addTail()* and *deleteHead()* are basically the *enqueue()* and *dequeue()* methods, respectively
- Our original list implementation did not have a *clear()* method
 - all it has to do is set the *head* and *tail* to null
- Again, no need for the *isFull()* method
 - list can grow to an infinite size

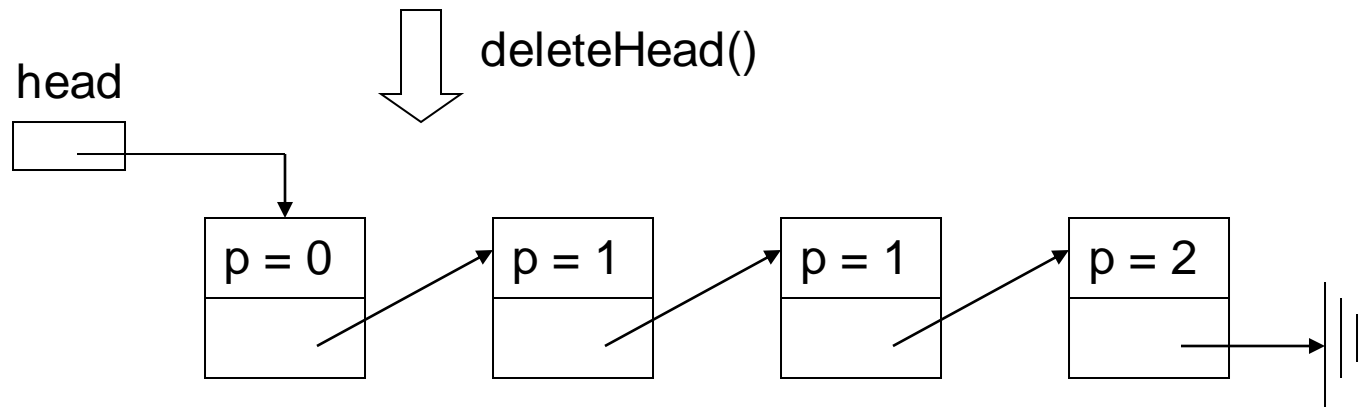
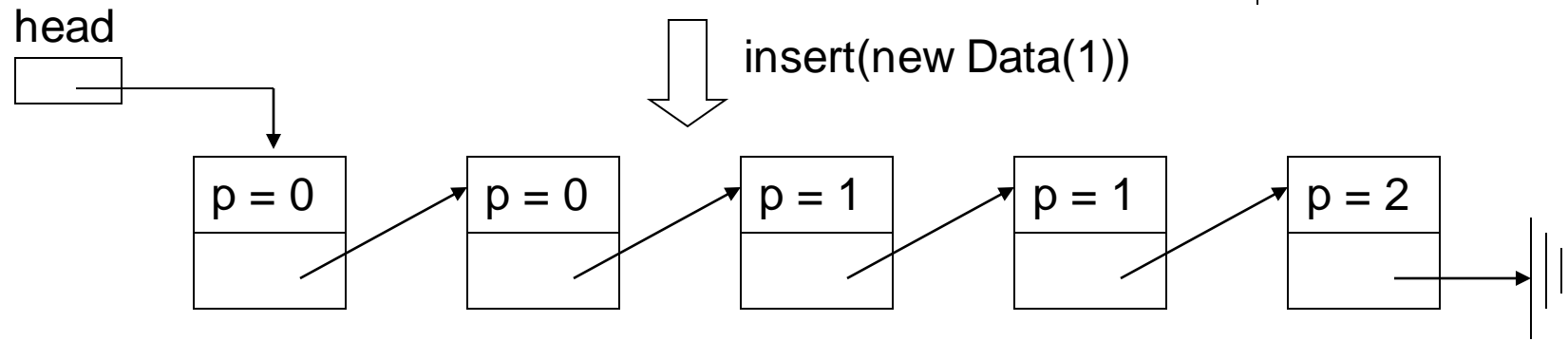
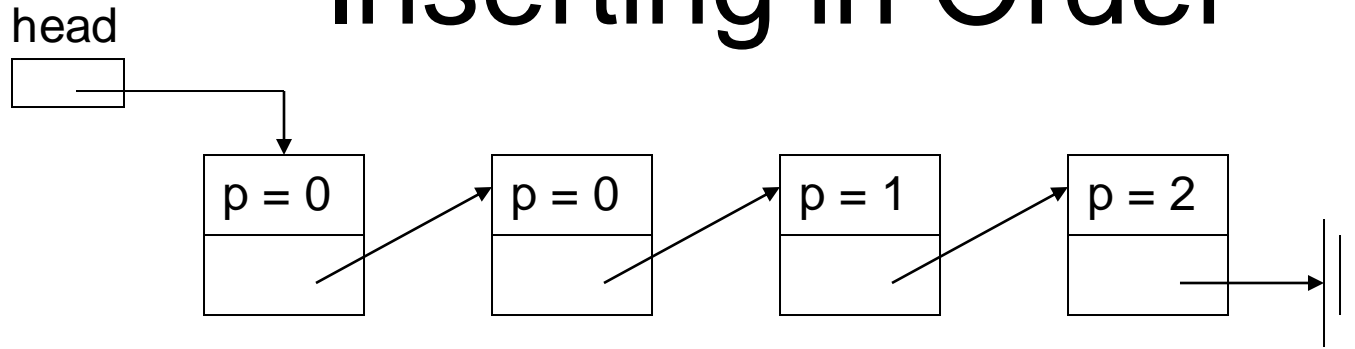
Priority Queue

- Sometimes it is not enough just do FIFO ordering
 - may want to give some items a higher priority than other items
 - these should be serviced before lower priority even if they arrived later
- Two major ways to implement a priority queue
 - insert items in a sorted order
 - always remove the head
 - insert in unordered order and search list on remove
 - always add to the tail
 - either way, time is $O(n)$
 - either adding data takes time and removing is quick, or
 - adding data is quick and removing takes time

Inserting in Order

- Use the very first linked list class shown
 - only need to use the *add()* and *removeHead()* methods
 - *add()* method puts things in the list in order
 - the *compareTo()* method (implemented by your data class) should return a value based on priority
 - usually consider lower number a higher priority
 - Performance
 - $O(n)$ to add
 - $O(1)$ to remove

Inserting in Order



Queue Applications

- As with stacks, queues are very common
 - networking
 - routers queue packets before sending them out
 - operating systems
 - disk scheduling, pipes, sockets, etc.
 - system modeling and simulation
 - queueing theory
- Any of these queues can be done as a priority queue
 - consider a disk scheduler
 - higher priority is given to a job closer to the current position of the disk head
 - next request done is that closest to the current position

Disk Scheduler

- Requests for disk sectors arrive randomly
- Disk requests are completed at a much slower rate than disk requests arrive
 - need to place waiting jobs in a queue
- All requests should be placed in a priority queue
 - jobs closest to the current position get placed closer to the front of the queue
- When the current job finishes, the next job is removed from the head of the queue

Code for a Priority Queue Class

```
class QueuePriority {  
    private LinkedList queue;  
    public Queue() { queue = new LinkedList(); }  
    public void enqueue(Object data) { queue.add(data); }  
    public Object dequeue() { return queue.deleteHead(); }  
    public void clear() { queue.clear(); }  
    public boolean isEmpty() { return queue.isEmpty(); }  
}
```

Code for a Disk Request Class

```
class DiskRequest implements Comparable{
    private int cylinder;
    private int head;
    private int sector;
    public DiskRequest(int cylinder, int head, int sector) {
        this.cylinder = cylinder;
        this.head = head;
        this.sector = sector;
    }
    public int getCylinder() { return cylinder; }
    public int getHead() { return head; }
    public int getSector() { return sector; }
    public int compareTo(Object obj) {
        DiskRequest req = (DiskRequest)obj;
        return cylinder - req.getCylinder();
    }
    public String toString() {
        String msg = new String("Cylinder(" + cylinder + ")\tHead(" +
                                head + ")\tSector(" + sector + ")");
        return msg;
    }
}
```

```
public static void main(String[] args) {  
    QueuePriority diskQueue = new QueuePriority();  
    int option = getOption();  
    DiskRequest req;  
    while(option != 3) {  
        switch(option) {  
            case 1:  
                req = getRequest();  
                diskQueue.enqueue(req);  
                break;  
            case 2:  
                if(diskQueue.isEmpty())  
                    System.out.println("Disk queue is empty.");  
                else {  
                    req = (DiskRequest)diskQueue.dequeue();  
                    System.out.println("Removed request: " + req.toString());  
                }  
                break;  
            case 3:  
                break;  
            default:  
                System.out.println("Error: invalid entry.");  
        }  
        option = getOption();  
    }  
}
```