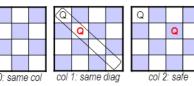
### Recursive Strategy for n-Queens

- Consider one row at a time. Within the row, consider one column at a time, looking for a "safe" column to place a queen.
- If we find one, place the queen, and *make a recursive call* to place a queen on the next row.
- If we can't find one, *backtrack* by returning from the recursive call, and try to find another safe column in the previous row.
- Example for n = 4:

row 0:

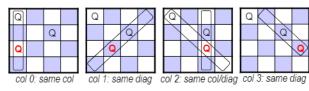


row 1:

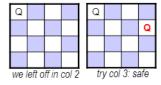


# 4-Queens Example (cont.)

row 2:

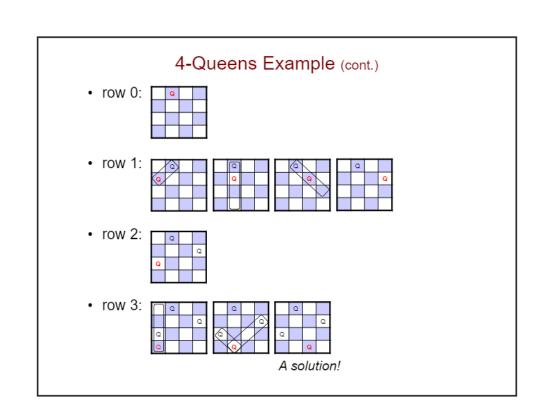


- We've run out of columns in row 2!
- Backtrack to row 1 by returning from the recursive call.
  - pick up where we left off
  - we had already tried columns 0-2, so now we try column 3:

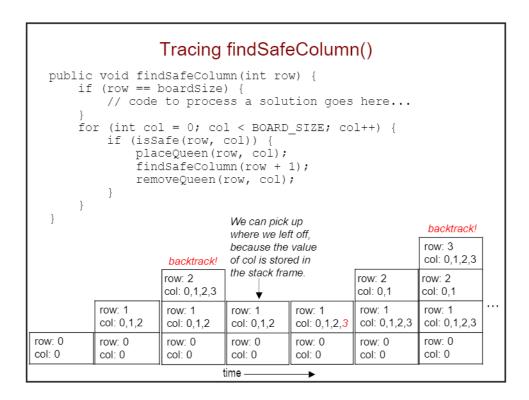


· Continue the recursion as before.

# 4-Queens Example (cont.) row 2: Out of the same color of the same colo



```
findSafeColumn() Method
public void findSafeColumn(int row) {
    if (row == boardSize) { // base case: a solution!
        solutionsFound++;
        displayBoard();
        if (solutionsFound >= solutionTarget)
             System.exit(0);
        return;
    }
    for (int col = 0; col < boardSize; col++) {
        if (isSafe(row, col)) {
             placeQueen(row, col);
                                              Note: neither row++
             // Move onto the next row.
                                             nor ++row will work
             findSafeColumn(row + 1) ;
                                              here.
             // If we get here, we've backtracked.
             removeQueen(row, col);
        }
    }
}
       (See ~cscie119/examples/recursion/Queens.java)
```



### Template for Recursive Backtracking

```
void findSolutions(n, other params) {
   if (found a solution) {
      solutionsFound++;
      displaySolution();
      if (solutionsFound >= solutionTarget)
            System.exit(0);
      return;
   }

   for (val = first to last) {
      if (isValid(val, n)) {
            applyValue(val, n);
            findSolutions(n + 1, other params);
            removeValue(val, n);
      }
   }
}
```

## Template for Finding a Single Solution

```
boolean findSolutions(n, other params) {
   if (found a solution) {
      displaySolution();
      return true;
   }

   for (val = first to last) {
      if (isValid(val, n)) {
         applyValue(val, n);
        if (findSolutions(n + 1, other params))
            return true;
      removeValue(val, n);
      }
   }

   return false;
}
```

### Data Structures for n-Queens

- Three key operations:
  - isSafe (row, col): check to see if a position is safe
  - placeQueen(row, col)
  - removeQueen(row, col)
- A two-dim. array of booleans would be sufficient:

```
public class Queens {
    private boolean[][] queenOnSquare;
```

Advantage: easy to place or remove a queen:

```
public void placeQueen(int row, int col) {
        queenOnSquare[row][col] = true;
}
public void removeQueen(int row, int col) {
        queenOnSquare[row][col] = false;
}
```

Problem: isSafe() takes a lot of steps. What matters more?

### Additional Data Structures for n-Queens

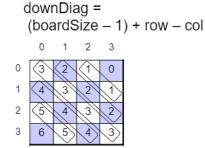
• To facilitate isSafe(), add three arrays of booleans:

```
private boolean[] colEmpty;
private boolean[] upDiagEmpty;
private boolean[] downDiagEmpty;
```

- An entry in one of these arrays is:
  - true if there are no queens in the column or diagonal
  - false otherwise

upDiag = row + col

· Numbering diagonals to get the indices into the arrays:



### Using the Additional Arrays

 Placing and removing a queen now involve updating four arrays instead of just one. For example:

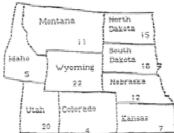
```
public void placeQueen(int row, int col) {
    queenOnSquare[row][col] = true;
    colEmpty[col] = false;
    upDiagEmpty[row + col] = false;
    downDiagEmpty[(boardSize - 1) + row - col] = false;
}
```

• However, checking if a square is safe is now more efficient:

```
public boolean isSafe(int row, int col) {
    return (colEmpty[col]
    && upDiagEmpty[row + col]
    && downDiagEmpty[(boardSize - 1) + row - col]);
}
```

# Recursive Backtracking II: Map Coloring

- Using just four colors (e.g., red, orange, green, and blue), we want color a map so that no two bordering states or countries have the same color.
- Sample map (numbers show alphabetical order in full list of state names):



This is another example of a problem that can be solved using recursive backtracking.