EXPERIMENT NO. 1

Study of Basic Commands in Linux Operating System

File and Directory Management

ls

- Description: Lists files and directories in the current directory.
- Usage: ls [options] [directory]
- Examples:
 - 1s List files and directories.
 - ∘ ls -1 List with detailed information.
 - ∘ ls -a List all files, including hidden ones.

cd

- Description: Changes the current directory.
- Usage: cd [directory]
- Examples:
 - cd /home/user Change to /home/user directory.
 - ∘ cd .. Move up one directory.
 - ∘ cd ~ Move to the home directory.

pwd

- Description: Prints the current working directory.
- Usage: pwd
- Examples:
 - pwd Display the full path of the current directory.

mkdir

- Description: Creates a new directory.
- Usage: mkdir [options] directory_name

- Examples:
 - mkdir new_folder Create a directory namednew_folder.
 - mkdir -p parent_folder/child_folder Create parent and child directories.

rmdir

- Description: Removes an empty directory.
- Usage: rmdir directory_name
- Examples:
 - rmdir old_folder Remove the empty directoryold_folder.

rm

- Description: Removes files or directories.
- Usage: rm [options] file_name
- Examples:
 - orm file.txt Remove file.txt.
 - rm -r folder_name Remove folder_name and its contents recursively.

ср

- Description: Copies files or directories.
- Usage: cp [options] source destination
- Examples:
 - o cp file1.txt file2.txt Copy file1.txt to file2.txt.
 - o cp -r dir1 dir2 Copy dir1 to dir2 recursively.

mν

- Description: Moves or renames files or directories.
- Usage: mv [options] source destination
- Examples:
 - o mv file1.txt /home/user/ Move file1.txt to /home/user/.
 - mv old_name.txt new_name.txt Rename old_name.txt to new_name.txt.

File Viewing and Editing

cat

- Description: Concatenates and displays file content.
- Usage: cat [options] file_name
- Examples:
 - o cat file.txt Display the content offile.txt.
 - o cat file1.txt file2.txt Concatenate and display both files.

less

- Description: Views file content with backward and forward navigation.
- Usage: less file_name
- Examples:
 - less file.txt View and navigate throughfile.txt.

free

- Description: Displays memory usage.
- Usage: free [options]
- Examples:
 - free Show memory usage.
 - free -h Show memory usage in human-readable format.

Process Management

ps

- Description: Displays information about active processes.
- Usage: ps [options]
- Examples:
 - ps Display processes for the current user.
 - ps aux Display detailed information about all processes.

Networking

ping

- Description: Tests network connectivity to a host.
- Usage: ping [options] destination
- Examples:
 - ping google.com Send ICMP packets to google.com.

ifconfig

- Description: Displays or configures network interfaces.
- Usage: ifconfig [interface] [options]
- Examples:
 - ifconfig Display network interfaces and their details.

netstat

- Description: Displays network connections and routing tables.
- Usage: netstat [options]
- Examples:
 - netstat -tuln Display listening ports and network connections.

User Management

whoami

- Description: Displays the currently logged-in user.
- Usage: whoami
- Examples:
 - whoami Show the username of the current user.

adduser

- Description: Adds a new user to the system.
- Usage: adduser username
- Examples:
 - o adduser newuser Add a new user named newuser.

File Permissions

chmod

- Description: Changes file permissions.
- Usage: chmod [options] mode file_name
- Examples:
 - $^{\circ}$ chmod 755 file.txt Set permissions of file.txt to rwxr-xr-x .

Aim:

Write a program to implement Thread and Process

Theory:

A thread is a lightweight process sharing memory with other threads in the same process, while a process is an independent program with its own memory space. Threads enable concurrent execution within a process, and process management involves handling separate program instances

```
#include <pthread.h>
#include <stdio.h>
#include <unistd.h>
void *thread function(void *arg) {
 printf("Thread: Running in thread\n");
  return NULL;
void create_thread() {
 pthread_t thread;
 if (pthread_create(&thread, NULL, thread_function, NULL)) {
   printf("Error creating thread\n");
   return;
 pthread_join(thread, NULL);
 printf("Thread: Finished execution\n");
void create_process() {
 pid_t pid = fork();
 if (pid == 0) {
   printf("Process: This is the child process\n");
 } else if (pid > 0) {
   printf("Process: This is the parent process\n");
 } else {
   printf("Process: Fork failed\n");
 }
int main() {
 printf("Main: Starting thread and process examples\n");
 create_thread();
 create_thread();
 create_process();
  return 0;
}
```

rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code\$./cmprun.sh lab2.c

Name: Bishnu Chalise

Main: Starting thread and process examples

Thread: Running in thread Thread: Running in thread
Thread: Finished execution
Thread: Running in thread
Thread: Finished execution
Process: This is the parent process
Process: This is the child process

Aim:

Write a C program to implement the producer-consumer problem using Semaphores

Theory:

The producer-consumer problem involves two types of processes: producers that generate data and consumers that process it. Semaphores are used to synchronize access to a shared buffer, preventing race conditions and ensuring that producers and consumers operate efficiently without conflicts.

```
#include <pthread.h>
#include <semaphore.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>
#define BUFFER SIZE 10
#define PRODUCER COUNT 1
#define CONSUMER COUNT 1
#define PRODUCE_COUNT 20
int buffer[BUFFER SIZE];
int in = 0;
int out = 0;
sem t empty;
sem t full;
pthread mutex t mutex;
void *producer(void *arg) {
  for (int i = 0; i < PRODUCE COUNT; i++) {</pre>
    sem wait(&empty);
    pthread mutex lock(&mutex);
    buffer[in] = i;
    printf("Produced: %d ", buffer[in]);
    in = (in + 1) % BUFFER SIZE;
    pthread mutex unlock(&mutex);
    sem post(&full);
    usleep(rand() % 100000); // Simulate variable production time
  }
  return NULL;
void *consumer(void *arg) {
```

```
for (int i = 0; i < PRODUCE COUNT; i++) {</pre>
    sem wait(&full);
    pthread_mutex_lock(&mutex);
    int item = buffer[out];
    printf("Consumed: %d\n", item);
    out = (out + 1) % BUFFER_SIZE;
    pthread_mutex_unlock(&mutex);
    sem_post(&empty);
    usleep(rand() % 1000000); // Simulate variable consumption time
  }
  return NULL;
}
int main() {
  pthread_t producers[PRODUCER_COUNT];
  pthread_t consumers[CONSUMER_COUNT];
  sem_init(&empty, 0, BUFFER_SIZE);
  sem_init(&full, 0, 0);
  pthread_mutex_init(&mutex, NULL);
  for (int i = 0; i < PRODUCER_COUNT; i++) {</pre>
    pthread_create(&producers[i], NULL, producer, NULL);
  for (int i = 0; i < CONSUMER_COUNT; i++) {</pre>
    pthread_create(&consumers[i], NULL, consumer, NULL);
  for (int i = 0; i < PRODUCER_COUNT; i++) {</pre>
    pthread_join(producers[i], NULL);
  for (int i = 0; i < CONSUMER_COUNT; i++) {</pre>
    pthread_join(consumers[i], NULL);
  sem_destroy(&empty);
  sem_destroy(&full);
  pthread_mutex_destroy(&mutex);
  return 0;
}
  rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab3.c
   Name: Bishnu Chalise
   Produced: 0 Consumed: 0
   Produced: 1 Consumed: 1
   Produced: 2 Consumed: 2 Produced: 3 Consumed: 3
   Produced: 4 Consumed: 4
   Produced: 5 Produced: 6 Consumed: 5
   Consumed: 6
   Produced: 7 Consumed: 7
   Produced: 8 Consumed: 8
   Produced: 9 Consumed: 9
   Produced: 10 Consumed: 10
   Produced: 11 Produced: 12 Consumed: 11
   Consumed: 12
Produced: 13 Consumed: 13
Produced: 14 Consumed: 14
   Produced: 15 Consumed: 15
   Produced: 16 Consumed: 16
   Produced: 17 Consumed: 17
   Produced: 18 Consumed: 18
Produced: 19 Consumed: 19
```

Aim:

Write a C program to simulate the concept of the Dining Philosophers problem

Theory:

The Dining Philosophers problem is a classic synchronization problem where philosophers alternate between thinking and eating from a shared set of forks. The challenge is to avoid deadlock and ensure that each philosopher gets a chance to eat without resource starvation.

```
#include <pthread.h>
#include <semaphore.h>
#include <stdio.h>
#include <unistd.h>
#define N 2
sem t forks[N];
pthread_mutex_t mutex;
void *philosopher(void *num) {
 int id = *(int *)num;
 while (1) {
    printf("Philosopher %d is thinking...\n", id);
    sleep(1);
    pthread mutex lock(&mutex);
    sem wait(&forks[id]);
    sem wait(&forks[(id + 1) % N]);
    pthread mutex unlock(&mutex);
    printf("Philosopher %d is eating...\n", id);
    sleep(2);
    sem post(&forks[id]);
    sem post(&forks[(id + 1) % N]);
   printf("Philosopher %d finished eating and is thinking again...\n", id);
 }
}
int main() {
 pthread t philosophers[N];
 int ids[N];
 pthread_mutex_init(&mutex, NULL);
 for (int i = 0; i < N; i++) {
```

```
sem_init(&forks[i], 0, 1);
ids[i] = i;
}

for (int i = 0; i < N; i++) {
   pthread_create(&philosophers[i], NULL, philosopher, &ids[i]);
}

for (int i = 0; i < N; i++) {
   pthread_join(philosophers[i], NULL);
}

for (int i = 0; i < N; i++) {
   sem_destroy(&forks[i]);
}
pthread_mutex_destroy(&mutex);
return 0;
}</pre>
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab4.c
Name: Bishnu Chalise
Philosopher 1 is thinking...
Philosopher 0 is thinking...
Philosopher 1 is eating...
Philosopher 1 finished eating and is thinking again...
Philosopher 1 is thinking...
Philosopher 0 is eating...
Philosopher 0 finished eating and is thinking again...
Philosopher 0 is thinking...
Philosopher 1 is eating...
Philosopher 1 finished eating and is thinking again...
Philosopher 1 is thinking...
Philosopher 0 is eating...
Philosopher 0 finished eating and is thinking again...
Philosopher 0 is thinking...
Philosopher 1 is eating...
Philosopher 1 finished eating and is thinking again...
Philosopher 1 is thinking...
Philosopher 0 is eating...
Philosopher 0 finished eating and is thinking again...
Philosopher 0 is thinking...
Philosopher 1 is eating...
```

Aim:

To simulate and calculate the turnaround time and waiting time for the following non-preemptive CPU scheduling algorithms:

- a) FCFS (First-Come, First-Served)
- b) SJF (Shortest Job First)
- c) Priority Scheduling

Theory:

FCFS (First-Come, First-Served):

Description: Processes are executed in the order they arrive. The process that arrives first gets executed first.

Turnaround Time: The total time taken from arrival to completion

of the process. Waiting Time: The total time a process spends waiting in the queue before getting executed.

SJF (Shortest Job First):

Description: The process with the shortest burst time (execution time) is executed next. This algorithm can be either preemptive or non-preemptive, but here we'll consider the non-preemptive version.

Turnaround Time: Time from the arrival of the process to its completion.

Waiting Time: The total time a process spends waiting in the queue before it starts executing.

Priority Scheduling:

Description: Each process is assigned a priority. The process with the highest priority (typically the smallest number) is executed first. In case of a tie, FCFS is used.

Turnaround Time: Time from the arrival of the process to its completion.

Waiting Time: The total time a process spends waiting in the queue before it starts executing.

Algorithm:

FCFS Scheduling:

- 1.1 Read the number of processes.
- 1.2 Read the arrival and burst times for each process.
- 1.3 Calculate the completion, turnaround, and waiting times:

Completion Time (CT) = Arrival Time + Burst Time Turnaround Time (TAT) = Completion Time - Arrival Time Waiting Time (WT) = Turnaround Time - Burst Time

SJF Scheduling:

- 2.1 Read the number of processes.
- 2.2 Read the arrival and burst times for each process.
- 2.3 Sort processes by burst time.
- 2.4 Calculate completion, turnaround, and waiting times:
 Completion Time (CT) is computed based on the sorted order.
 Turnaround Time (TAT) = Completion Time Arrival Time
 Waiting Time (WT) = Turnaround Time Burst Time

Priority Scheduling:

- 3.1 Read the number of processes.
- 3.2 Read the arrival time, burst time, and priority for each process.
- 3.3 Sort processes by priority (lower number indicates higher priority).
- 3.4 Calculate completion, turnaround, and waiting times:
 Completion Time (CT) is computed based on the priority order.
 Turnaround Time (TAT) = Completion Time Arrival Time
 Waiting Time (WT) = Turnaround Time Burst Time

```
#include <stdio.h>
#define MAX 10
typedef struct {
 int id;
 int arrival;
 int burst;
 int priority;
 int waiting;
 int turnaround;
} Process;
void swap(Process *a, Process *b) {
 Process temp = *a;
 *a = *b:
 *b = temp;
void sortByArrival(Process proc[], int n) {
  for (int i = 0; i < n - 1; i++) {
    for (int j = i + 1; j < n; j++) {
      if (proc[i].arrival > proc[j].arrival) {
        swap(&proc[i], &proc[j]);
     }
   }
 }
}
```

```
void sortByBurst(Process proc[], int n) {
  for (int i = 0; i < n - 1; i++) {
    for (int j = i + 1; j < n; j++) {
      if (proc[i].burst > proc[j].burst) {
        swap(&proc[i], &proc[j]);
      }
   }
 }
}
void sortByPriority(Process proc[], int n) {
 for (int i = 0; i < n - 1; i++) {
    for (int j = i + 1; j < n; j++) {
      if (proc[i].priority > proc[j].priority) {
        swap(&proc[i], &proc[j]);
      }
   }
 }
}
void FCFS(Process proc[], int n) {
 int total_wt = 0, total_tat = 0;
 proc[0].waiting = 0;
 proc[0].turnaround = proc[0].burst;
 for (int i = 1; i < n; i++) {
    proc[i].waiting = proc[i - 1].waiting + proc[i - 1].burst - proc[i].arrival;
    if (proc[i].waiting < 0)</pre>
      proc[i].waiting = 0;
    proc[i].turnaround = proc[i].waiting + proc[i].burst;
    total_wt += proc[i].waiting;
   total_tat += proc[i].turnaround;
 printf("\nFCFS Scheduling:\n");
 for (int i = 0; i < n; i++) {
    printf("Process %d: Waiting Time = %d, Turnaround Time = %d\n", proc[i].id,
           proc[i].waiting, proc[i].turnaround);
 printf("Average Waiting Time = %.2f\n", (float)total_wt / n);
 printf("Average Turnaround Time = %.2f\n", (float)total_tat / n);
void SJF(Process proc[], int n) {
 int total_wt = 0, total_tat = 0;
 int time = 0;
 proc[0].waiting = 0;
 proc[0].turnaround = proc[0].burst;
 for (int i = 1; i < n; i++) {
    proc[i].waiting = time - proc[i].arrival;
    if (proc[i].waiting < 0)</pre>
      proc[i].waiting = 0;
    proc[i].turnaround = proc[i].waiting + proc[i].burst;
    total_wt += proc[i].waiting;
    total_tat += proc[i].turnaround;
    time += proc[i].burst;
 }
 printf("\nSJF Scheduling:\n");
 for (int i = 0; i < n; i++) {
    printf("Process %d: Waiting Time = %d, Turnaround Time = %d\n", proc[i].id,
           proc[i].waiting, proc[i].turnaround);
 }
```

```
printf("Average Waiting Time = %.2f\n", (float)total wt / n);
 printf("Average Turnaround Time = %.2f\n", (float)total_tat / n);
void PriorityScheduling(Process proc[], int n) {
  int total_wt = 0, total_tat = 0;
 int time = 0;
  proc[0].waiting = 0;
 proc[0].turnaround = proc[0].burst;
 for (int i = 1; i < n; i++) {
    proc[i].waiting = time - proc[i].arrival;
    if (proc[i].waiting < 0)</pre>
      proc[i].waiting = 0;
    proc[i].turnaround = proc[i].waiting + proc[i].burst;
    total_wt += proc[i].waiting;
    total_tat += proc[i].turnaround;
    time += proc[i].burst;
 printf("\nPriority Scheduling:\n");
 for (int i = 0; i < n; i++) {
    printf("Process %d: Waiting Time = %d, Turnaround Time = %d\n", proc[i].id,
           proc[i].waiting, proc[i].turnaround);
 printf("Average Waiting Time = %.2f\n", (float)total_wt / n);
 printf("Average Turnaround Time = %.2f\n", (float)total_tat / n);
int main() {
 Process proc[MAX];
 int n;
 printf("Enter number of processes: ");
 scanf("%d", &n);
 for (int i = 0; i < n; i++) {
    printf("Enter arrival time, burst time, and priority for process %d: ",
           i + 1);
    proc[i].id = i + 1;
    scanf("%d %d %d", &proc[i].arrival, &proc[i].burst, &proc[i].priority);
 sortByArrival(proc, n);
 FCFS(proc, n);
 sortByArrival(proc, n);
 sortByBurst(proc, n);
 SJF(proc, n);
 sortByArrival(proc, n);
 sortByPriority(proc, n);
 PriorityScheduling(proc, n);
 return 0;
}
```

```
• rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab5.c
 Name: Bishnu Chalise
 Enter number of processes: 3
 Enter arrival time, burst time, and priority for process 1: 0 5 1
 Enter arrival time, burst time, and priority for process 2: 2 6 2
 Enter arrival time, burst time, and priority for process 3: 1 4 3
 FCFS Scheduling:
 Process 1: Waiting Time = 0, Turnaround Time = 5
Process 3: Waiting Time = 4, Turnaround Time = 8
 Process 2: Waiting Time = 6, Turnaround Time = 12
 Average Waiting Time = 3.33
 Average Turnaround Time = 6.67
 SJF Scheduling:
 Process 3: Waiting Time = 0, Turnaround Time = 4
 Process 1: Waiting Time = 0, Turnaround Time = 5
 Process 2: Waiting Time = 3, Turnaround Time = 9
 Average Waiting Time = 1.00
 Average Turnaround Time = 4.67
 Priority Scheduling:
 Process 1: Waiting Time = 0, Turnaround Time = 5
 Process 2: Waiting Time = 0, Turnaround Time = 6
 Process 3: Waiting Time = 5, Turnaround Time = 9
 Average Waiting Time = 1.67
 Average Turnaround Time = 5.00
```

Aim:

To simulate the following preemptive CPU scheduling algorithms to find turnaround time and waiting time. a) Round Robin To simulate and calculate the turnaround time and waiting time for the following preemptive CPU scheduling algorithms:

- a) Round Robin
- b) Priority
- b) Priority

Theory:

Round Robin:

Description: Each process is assigned a fixed time quantum (or time slice). Processes are executed in a circular order, with each process getting a chance to execute for the duration of the time quantum. If a process does not complete within its time quantum, it is placed at the end of the queue to wait for its next turn.

Turnaround Time: Time from the arrival of the process to its completion.

Waiting Time: The total time a process spends waiting in the queue before it gets executed.

Priority Scheduling:

Description: Each process is assigned a priority. The process with the highest priority (typically the smallest number) is executed first. In preemptive priority scheduling, if a new process arrives with a higher priority than the currently running process, the current process is preempted, and the new process is executed.

Turnaround Time: Time from the arrival of the process to its completion.

Waiting Time: The total time a process spends waiting in the queue before it starts executing.

Algorithm:

Round Robin Scheduling:

- 1.1 Read the number of processes and the time quantum.
- 1.2 Read the arrival and burst times for each process.
- 1.3 Initialize the queue and start processing each process for

the duration of the time quantum.

- 1.4 If a process's burst time is greater than the time quantum, update the burst time and move it to the end of the queue.
- 1.5 Calculate completion, turnaround, and waiting times: Completion Time (CT) is updated as the process finishes execution.

Turnaround Time (TAT) = Completion Time - Arrival Time Waiting Time (WT) = Turnaround Time - Burst Time

Priority Scheduling:

- 2.1 Read the number of processes.
- 2.2 Read the arrival time, burst time, and priority for each process.
- 2.3 Sort processes by priority (lower number indicates higher priority).
- 2.4 Implement the preemptive scheduling by maintaining a priority queue and preempting the current process if a new process with higher priority arrives.
- 2.5 Calculate completion, turnaround, and waiting times: Completion Time (CT) is updated based on process execution and preemption.

Turnaround Time (TAT) = Completion Time - Arrival Time Waiting Time (WT) = Turnaround Time - Burst Time

```
#include <limits.h>
#include <stdio.h>
#define MAX 10
typedef struct {
 int id;
 int arrival time;
 int burst time;
 int priority;
 int remaining time;
 int turnaround time;
 int waiting time;
} Process;
void round robin(Process processes[], int n, int quantum) {
  int time = 0, i, flag;
  int remaining_processes = n;
 while (remaining_processes > 0) {
    flaq = 0:
    for (i = 0; i < n; i++) {
      if (processes[i].arrival_time <= time &&</pre>
          processes[i].remaining_time > 0) {
        flag = 1;
        if (processes[i].remaining_time <= quantum) {</pre>
          time += processes[i].remaining_time;
          processes[i].turnaround_time = time - processes[i].arrival_time;
          processes[i].waiting_time =
```

```
processes[i].turnaround time - processes[i].burst time;
          processes[i].remaining_time = 0;
          remaining_processes--;
        } else {
          time += quantum;
          processes[i].remaining_time -= quantum;
        }
      }
    if (flag == 0) {
      time++;
    }
 }
}
void priority_scheduling(Process processes[], int n) {
  int time = 0, i, min_priority_index;
 int completed = 0;
 while (completed < n) {</pre>
    min_priority_index = -1;
    for (i = 0; i < n; i++) {
      if (processes[i].arrival_time <= time &&</pre>
          processes[i].remaining_time > 0) {
        if (min_priority_index == -1 ||
            processes[i].priority < processes[min_priority_index].priority) {</pre>
          min_priority_index = i;
        }
      }
    if (min_priority_index != -1) {
      processes[min_priority_index].remaining_time--;
      if (processes[min_priority_index].remaining_time == 0) {
        processes[min_priority_index].turnaround_time =
            time - processes[min_priority_index].arrival_time;
        processes[min_priority_index].waiting_time =
            processes[min_priority_index].turnaround_time -
            processes[min_priority_index].burst_time;
        completed++;
      }
    } else {
     time++;
    }
 }
}
void print results(Process processes[], int n, const char *algorithm) {
  int i;
  float total_turnaround_time = 0, total_waiting_time = 0;
  printf("%s Scheduling Results:\n", algorithm);
  printf("ID\tArrival\tBurst\tPriority\tTurnaround\tWaiting\n");
  for (i = 0; i < n; i++) {
    total_turnaround_time += processes[i].turnaround_time;
    total_waiting_time += processes[i].waiting_time;
    printf("%d\t%d\t%d\t\t%d\t\t%d\n", processes[i].id,
           processes[i].arrival_time, processes[i].burst_time,
           processes[i].priority, processes[i].turnaround_time,
           processes[i].waiting_time);
  }
 printf("Average Turnaround Time: %.2f\n", total_turnaround_time / n);
 printf("Average Waiting Time: %.2f\n", total_waiting_time / n);
```

```
int main() {
  int n, i, quantum;
 Process processes[MAX];
  printf("Enter number of processes: ");
 scanf("%d", &n);
 for (i = 0; i < n; i++) {
    printf("Enter arrival time, burst time, and priority for process %d: ",
           i + 1);
    processes[i].id = i + 1;
    scanf("%d %d %d", &processes[i].arrival_time, &processes[i].burst_time,
          &processes[i].priority);
    processes[i].remaining_time = processes[i].burst_time;
    processes[i].turnaround_time = 0;
   processes[i].waiting_time = 0;
 printf("Enter time quantum for Round Robin: ");
 scanf("%d", &quantum);
 Process rr_processes[MAX];
 Process pr_processes[MAX];
  for (i = 0; i < n; i++) {
    rr_processes[i] = processes[i];
    pr_processes[i] = processes[i];
  round_robin(rr_processes, n, quantum);
 print_results(rr_processes, n, "Round Robin");
 priority_scheduling(pr_processes, n);
 print_results(pr_processes, n, "Priority");
  return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab6.c
 Name: Bishnu Chalise
 Enter number of processes: 3
 Enter arrival time, burst time, and priority for process 1: 0 5 2
 Enter arrival time, burst time, and priority for process 2: 1 3 1 Enter arrival time, burst time, and priority for process 3: 2 8 3
 Enter time quantum for Round Robin: 2
 Round Robin Scheduling Results:
 ID
          Arrival Burst
                            Priority
                                              Turnaround
                                                                Waiting
          0
                            2
                                              12
                                              8
                   8
                                              14
                                                                6
 Average Turnaround Time: 11.33
 Average Waiting Time: 6.00
 Priority Scheduling Results:
 TD
          Arrival Burst Priority
                                              Turnaround
                                                                Waiting
          0
                                              8
 2
                   8
                                              14
                                                                6
 Average Turnaround Time: 8.33
 Average Waiting Time: 3.00
```

Aim:

To simulate the Banker's Algorithm to ensure safe allocation of resources and avoid deadlocks in a system with multiple processes and resource types.

Theory:

Banker's Algorithm:

Description: The Banker's Algorithm is used to allocate resources to processes in a way that ensures the system is always in a safe state. It checks for safe states using a request matrix and allocation matrix, which track resource allocations and requests. The algorithm calculates if resource allocation to processes can lead to a safe sequence, where all processes can complete without causing a deadlock.

Safe State: A state is considered safe if there exists a sequence of processes that allows all of them to complete successfully without causing a deadlock. The Banker's Algorithm ensures that every resource request can be granted while maintaining this safe state.

Request Matrix: Represents the maximum demand of each process for each resource.

Allocation Matrix: Represents the current allocation of resources to each process.

Available Resources: Represents the resources available in the system that are not currently allocated.

Algorithm:

Banker's Algorithm:

- 1.1 Initialize matrices for available resources, maximum demands, and current allocations.
- 1.2 For each resource request by a process:
- 1.2.1 Check if the request is less than or equal to the process's maximum demand.
- 1.2.2 Check if the request is less than or equal to the available resources.
- 1.2.3 Temporarily allocate the requested resources and check if the system is in a safe state by simulating the execution of processes with the new allocation.
- 1.2.4 If the system is in a safe state, grant the request.

Otherwise, revert to the previous state and deny the request. 1.3 Repeat the above steps for all resource requests and process executions.

```
#include <stdio.h>
#define MAX 10
#define RESOURCE_TYPES 3
void calculate_need(int need[MAX][RESOURCE_TYPES], int max[MAX][RESOURCE_TYPES],
                    int allot[MAX][RESOURCE TYPES], int n) {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < RESOURCE_TYPES; j++) {</pre>
      need[i][j] = max[i][j] - allot[i][j];
    }
 }
}
int is_safe(int processes[], int n, int m, int available[],
            int max[][RESOURCE_TYPES], int allot[][RESOURCE_TYPES]) {
  int need[MAX][RESOURCE_TYPES];
  calculate_need(need, max, allot, n);
  int work[RESOURCE_TYPES];
  int finish[MAX];
  for (int i = 0; i < m; i++) {
   work[i] = available[i];
  for (int i = 0; i < n; i++) {
    finish[i] = 0;
  }
  int safeSeq[MAX];
  int count = 0;
 while (count < n) {</pre>
    int found = 0;
    for (int p = 0; p < n; p++) {
      if (!finish[p]) {
        int j;
        for (j = 0; j < m; j++) {
          if (need[p][j] > work[j]) {
            break;
          }
        if (j == m) {
          for (int k = 0; k < m; k++) {
            work[k] += allot[p][k];
          safeSeq[count++] = p;
          finish[p] = 1;
          found = 1;
        }
      }
    if (!found) {
      printf("System is not in a safe state.\n");
      return 0;
```

```
}
 printf("System is in a safe state.\nSafe sequence is: ");
 for (int i = 0; i < n; i++) {
   printf("P%d ", safeSeq[i]);
 printf("\n");
 return 1;
int main() {
 int n, m;
 int processes[MAX];
 int available[RESOURCE_TYPES];
 int max[MAX][RESOURCE_TYPES];
 int allot[MAX][RESOURCE_TYPES];
 printf("Enter number of processes: ");
 scanf("%d", &n);
 printf("Enter number of resource types: ");
 scanf("%d", &m);
 printf("Enter the number of available instances for each resource:\n");
 for (int i = 0; i < m; i++) {
   printf("Resource %d: ", i + 1);
   scanf("%d", &available[i]);
 }
 printf("Enter maximum matrix:\n");
 for (int i = 0; i < n; i++) {
   processes[i] = i;
   printf("Process %d:\n", i + 1);
   for (int j = 0; j < m; j++) {
     printf("Maximum resource %d: ", j + 1);
     scanf("%d", &max[i][j]);
   }
 }
 printf("Enter allocation matrix:\n");
 for (int i = 0; i < n; i++) {
   printf("Process %d:\n", i + 1);
    for (int j = 0; j < m; j++) {
      printf("Allocated resource %d: ", j + 1);
      scanf("%d", &allot[i][j]);
   }
 }
 is_safe(processes, n, m, available, max, allot);
 return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab7.c
Name: Bishnu Chalise
Enter number of processes: 3
Enter number of resource types: 3
Enter the number of available instances for each resource:
Resource 1: 1
Resource 2: 0
Resource 3: 1
Enter maximum matrix:
Process 1:
Maximum resource 1: 7
Maximum resource 2: 5
Maximum resource 3: 3
Process 2:
Maximum resource 1: 3
Maximum resource 2: 2
Maximum resource 3: 2
Process 3:
Maximum resource 1: 9
Maximum resource 2: 0
Maximum resource 3: 2
Enter allocation matrix:
Process 1:
Allocated resource 1: 0
Allocated resource 2: 1
Allocated resource 3: 0
Process 2:
Allocated resource 1: 2
Allocated resource 2: 0
Allocated resource 3: 0
Process 3:
Allocated resource 1: 3
Allocated resource 2: 0
Allocated resource 3: 1
System is not in a safe state.
```

Aim:

Aim:

To simulate the Banker's Algorithm for deadlock prevention, ensuring that resources are allocated in a way that prevents the system from entering an unsafe state that could lead to a deadlock.

Theory:

Banker's Algorithm:

Description: The Banker's Algorithm is used to prevent deadlocks by ensuring that resource allocations are made in such a way that the system always remains in a safe state. This involves checking whether a resource allocation will result in a safe or unsafe state before granting the request. If granting a request results in a safe state, the request is granted; otherwise, it is denied. Safe State: A system is in a safe state if there is a sequence of processes that allows all processes to complete without causing a deadlock. The algorithm ensures that resource allocations do not lead to a state where no such sequence exists.

Request Matrix: Represents the maximum resources each process may need.

Allocation Matrix: Represents the current allocation of resources to each process.

Available Resources: Represents the resources currently available in the system.

Algorithm:

Banker's Algorithm for Deadlock Prevention:

- 1.1 Initialize the matrices for available resources, maximum demands, and current allocations.
- 1.2 For each resource request by a process:
- 1.2.1 Check if the request is less than or equal to the process's maximum demand.
- 1.2.2 Check if the request is less than or equal to the available resources.
- 1.2.3 Temporarily allocate the requested resources and check if the system remains in a safe state using a safety algorithm:
- 1.2.3.1 Calculate the new available resources, allocation matrix, and need matrix.

- 1.2.3.2 Use a safety algorithm to determine if there exists a safe sequence.
- 1.2.4 If the system is in a safe state, grant the request.0therwise, revert to the previous state and deny the request.1.3 Repeat the above steps for all resource requests and process
- executions.

```
#include <stdbool.h>
#include <stdio.h>
#define MAX 10
#define RESOURCE_TYPES 3
void calculate_need(int need[MAX][RESOURCE_TYPES], int max[MAX][RESOURCE_TYPES],
                     int allot[MAX][RESOURCE_TYPES], int n) {
  for (int i = 0; i < n; i++) {
    for (int j = 0; j < RESOURCE_TYPES; j++) {</pre>
      need[i][j] = max[i][j] - allot[i][j];
    }
 }
}
bool is_safe(int processes[], int n, int m, int available[],
              int max[][RESOURCE_TYPES], int allot[][RESOURCE_TYPES],
              int need[MAX][RESOURCE_TYPES]) {
  int work[RESOURCE_TYPES];
  bool finish[MAX];
  for (int i = 0; i < m; i++) {
    work[i] = available[i];
  for (int i = 0; i < n; i++) {
    finish[i] = false;
  int safeSeq[MAX];
  int count = 0;
 while (count < n) {</pre>
    bool found = false;
    for (int p = 0; p < n; p++) {
      if (!finish[p]) {
        int j;
        for (j = 0; j < m; j++) {
          if (need[p][j] > work[j]) {
            break;
          }
        \textbf{if} \ (\texttt{j} == \texttt{m}) \ \{
          for (int k = 0; k < m; k++) {
            work[k] += allot[p][k];
          safeSeq[count++] = p;
          finish[p] = true;
          found = true;
        }
     }
    }
```

```
if (!found) {
      return false:
 }
  printf("System is in a safe state.\nSafe sequence is: ");
  for (int i = 0; i < n; i++) {
   printf("P%d ", safeSeq[i]);
 printf("\n");
  return true;
void request_resources(int processes[], int n, int m, int available[],
                       int max[][RESOURCE_TYPES], int allot[][RESOURCE_TYPES],
                       int need[MAX][RESOURCE_TYPES]) {
 int request[MAX][RESOURCE_TYPES];
 int i, j, process_id;
 printf("Enter process number making request: ");
 scanf("%d", &process_id);
 if (process_id < 0 || process_id >= n) {
    printf("Invalid process number.\n");
    return;
 }
 printf("Enter request for resources:\n");
 for (i = 0; i < m; i++) {
    printf("Resource %d: ", i + 1);
    scanf("%d", &request[process_id][i]);
 }
 for (i = 0; i < m; i++) {
    if (request[process_id][i] > need[process_id][i]) {
      printf("Request exceeds maximum claim.\n");
      return;
    }
 }
 for (i = 0; i < m; i++) {
   if (request[process_id][i] > available[i]) {
      printf("Resources not available.\n");
      return;
    }
 }
 for (i = 0; i < m; i++) {
    available[i] -= request[process_id][i];
    allot[process_id][i] += request[process_id][i];
    need[process_id][i] -= request[process_id][i];
 }
 if (is_safe(processes, n, m, available, max, allot, need)) {
    printf("Request granted.\n");
 } else {
    printf("Request cannot be granted due to unsafe state. Rolling back.\n");
    for (i = 0; i < m; i++) {
      available[i] += request[process_id][i];
      allot[process_id][i] -= request[process_id][i];
      need[process_id][i] += request[process_id][i];
    }
 }
```

```
}
int main() {
 int n, m;
 int processes[MAX];
 int available[RESOURCE TYPES];
 int max[MAX][RESOURCE_TYPES];
 int allot[MAX][RESOURCE_TYPES];
 int need[MAX][RESOURCE_TYPES];
 printf("Enter number of processes: ");
 scanf("%d", &n);
 printf("Enter number of resource types: ");
 scanf("%d", &m);
 printf("Enter the number of available instances for each resource:\n");
 for (int i = 0; i < m; i++) {
   printf("Resource %d: ", i + 1);
   scanf("%d", &available[i]);
 }
 printf("Enter maximum matrix:\n");
 for (int i = 0; i < n; i++) {
   processes[i] = i;
   printf("Process %d:\n", i + 1);
   for (int j = 0; j < m; j++) {
     printf("Maximum resource %d: ", j + 1);
     scanf("%d", &max[i][j]);
   }
 }
 printf("Enter allocation matrix:\n");
 for (int i = 0; i < n; i++) {
   printf("Process %d:\n", i + 1);
    for (int j = 0; j < m; j++) {
      printf("Allocated resource %d: ", j + 1);
      scanf("%d", &allot[i][j]);
   }
 }
 calculate_need(need, max, allot, n);
  request_resources(processes, n, m, available, max, allot, need);
 return 0;
}
```

```
Name: Bishnu Chalise
Enter number of processes: 2
Enter number of resource types: 3
Enter the number of available instances for each resource:
Resource 1: 3
Resource 2: 2
Resource 3: 2
Enter maximum matrix:
Process 1:
Maximum resource 1: 7
Maximum resource 2: 5
Maximum resource 3: 3
Process 2:
Maximum resource 1: 3
Maximum resource 2: 2
Maximum resource 3: 2
Enter allocation matrix:
Process 1:
Allocated resource 1: 0
Allocated resource 2: 1
Allocated resource 3: 0
Process 2:
Allocated resource 1: 2
Allocated resource 2: 0
Allocated resource 3: 0
Enter process number making request: 1
Enter request for resources:
Resource 1: 1
Resource 2: 0
Resource 3: 2
System is in a safe state.
Safe sequence is: P1 P0
```

rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code\$./cmprun.sh lab8.c

Aim:

Program to simulate the MVT and MFT memory management techniques

Theory:

MVT (Multi-Programming with Variable Partitions): MVT is a memory management technique where the memory is divided into variable-sized partitions. Each partition is allocated to a process based on its needs. The partitions can grow or shrink dynamically as processes are loaded or terminated, allowing more efficient use of memory and accommodating varying process sizes.

MFT (Multi-Programming with Fixed Partitions): MFT divides memory into fixed-size partitions. Each partition is allocated to a process, and processes are limited to the partition size. This method simplifies memory management but can lead to internal fragmentation if a process does not fully use its allocated partition.

```
#include <stdbool.h>
#include <stdio.h>
#define TOTAL MEMORY 100
#define MFT_BLOCKS 50
void printMemory(int mft[], int mvt[], int mft_size, int mvt_size) {
  printf("MFT Memory: ");
  for (int i = 0; i < mft_size; i++) {</pre>
    printf("%d ", mft[i]);
  printf("\n");
  printf("MVT Memory: ");
  for (int i = 0; i < mvt_size; i++) {</pre>
    if (mvt[i] == -1) {
      printf(". ");
    } else {
      printf("%d ", mvt[i]);
  printf("\n");
int allocateMFT(int mft[], int task size, int mft size) {
  for (int i = 0; i <= mft_size - task_size; i++) {</pre>
    bool fit = true;
    for (int j = i; j < i + task_size; j++) {</pre>
```

```
if (mft[j] != 0) {
        fit = false;
        break:
      }
    if (fit) {
      for (int j = i; j < i + task_size; j++) {</pre>
       mft[j] = 1;
      return i;
    }
  }
  return -1;
void deallocateMFT(int mft[], int start, int task_size) {
  for (int i = start; i < start + task_size; i++) {</pre>
   \mathsf{mft[i]} = 0;
  }
}
int allocateMVT(int mvt[], int task_size, int mvt_size) {
  for (int i = 0; i <= mvt_size - task_size; i++) {</pre>
    bool fit = true;
    for (int j = i; j < i + task\_size; j++) {
      if (mvt[j] != -1) {
        fit = false;
        break;
      }
    }
    if (fit) {
      for (int j = i; j < i + task_size; j++) {</pre>
        mvt[j] = 1;
      }
      return i;
  }
  return -1;
void deallocateMVT(int mvt[], int start, int task_size) {
  for (int i = start; i < start + task_size; i++) {</pre>
    mvt[i] = -1;
  }
}
int main() {
  int mft[MFT_BLOCKS] = {0};
  int mvt[TOTAL_MEMORY] = {-1};
  printf("Initial State:\n");
  printMemory(mft, mvt, MFT_BLOCKS, TOTAL_MEMORY);
  printf("\nAllocating tasks:\n");
  int task_size = 10;
  printf("Allocating task of size %d in MFT\n", task_size);
  int start = allocateMFT(mft, task_size, MFT_BLOCKS);
  if (start != -1) {
    printf("Task allocated at index %d\n", start);
  } else {
    printf("Failed to allocate task in MFT\n");
  printf("Allocating task of size %d in MVT\n", task_size);
```

```
start = allocateMVT(mvt, task size, TOTAL MEMORY);
  if (start != -1) {
    printf("Task allocated at index %d\n", start);
  } else {
    printf("Failed to allocate task in MVT\n");
  printf("\nCurrent State:\n");
  printMemory(mft, mvt, MFT_BLOCKS, TOTAL_MEMORY);
  printf("\nDeallocating tasks:\n");
  printf("Deallocating task of size %d from MFT starting at index %d\n",
         task_size, 0);
  deallocateMFT(mft, 0, task_size);
  printf("Deallocating task of size %d from MVT starting at index %d\n",
         task_size, 0);
  deallocateMVT(mvt, 0, task_size);
  printf("\nFinal State:\n");
  printMemory(mft, mvt, MFT_BLOCKS, TOTAL_MEMORY);
  return 0;
}
rudy@rudy:-/Desktop/untitled/github/d-tfn/operating-system/codes ./cmprun.sh lab9.c
Name: Bishnu Chalise
Allocating tasks:
Allocating task of size 10 in MFT
Allocating task of size 10 in MVT Failed to allocate task in MVT
Deallocating tasks:
Deallocating task of size 10 from MFT starting at index 0
```

Final State:

Aim:

To simulate the paging technique of memory management

Theory:

Paging: Paging is a memory management technique that divides memory into fixed-size blocks called pages. Processes are also divided into pages of the same size. When a process is executed, its pages are loaded into any available memory frames. This allows for efficient use of memory and helps in managing processes that are larger than the available physical memory. Paging eliminates external fragmentation and simplifies memory allocation.

```
#include <stdio.h>
#include <stdlib.h>
#include <time.h>
#define PAGE SIZE 256
#define PHYSICAL MEM SIZE 1024
#define LOGICAL MEM SIZE 2048
#define NUM FRAMES (PHYSICAL MEM SIZE / PAGE SIZE)
#define NUM_PAGES (LOGICAL_MEM_SIZE / PAGE_SIZE)
int physical memory[PHYSICAL MEM SIZE];
int page_table[NUM_PAGES];
void initialize_page_table() {
  for (int i = 0; i < NUM_PAGES; i++) {</pre>
   page_table[i] = -1;
int allocate_frame(int page_number) {
 int frame_number = rand() % NUM_FRAMES;
 page_table[page_number] = frame_number;
  return frame_number;
int translate_address(int logical_address) {
  int page_number = logical_address / PAGE_SIZE;
 int offset = logical_address % PAGE_SIZE;
 int frame_number = page_table[page_number];
  if (frame_number == -1) {
    frame_number = allocate_frame(page_number);
```

```
return frame_number * PAGE_SIZE + offset;
}
void write_memory(int logical_address, int value) {
  int physical_address = translate_address(logical_address);
  physical_memory[physical_address] = value;
int read_memory(int logical_address) {
  int physical_address = translate_address(logical_address);
  return physical_memory[physical_address];
int main() {
  srand(time(0));
  initialize_page_table();
 write_memory(500, 42);
 write_memory(1200, 84);
  printf("Value at logical address 500: %d\n", read_memory(500));
  printf("Value at logical address 1200: %d\n", read_memory(1200));
  return 0;
}
```

```
    rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab10.c
Name: Bishnu Chalise
    Value at logical address 500: 42
    Value at logical address 1200: 84
```

Aim:

To simulate the FIRST-FIT contiguous memory allocation technique

Theory:

FIRST-FIT: The FIRST-FIT memory allocation technique allocates the first available block of memory that is large enough to satisfy the request. The memory is scanned from the beginning, and the first sufficiently sized block is chosen for allocation. This method is simple and fast but can lead to fragmentation over time.

```
#include <stdio.h>
#define MEMORY SIZE 1000
#define NUM BLOCKS 5
int memory[MEMORY SIZE];
int block size[NUM BLOCKS] = {100, 200, 300, 150, 250};
int block allocated[NUM BLOCKS] = {0};
void allocate memory(int process size, int process id) {
  for (int i = 0; i < NUM BLOCKS; i++) {
    if (!block_allocated[i] && block_size[i] >= process_size) {
      block allocated[i] = process id;
      printf("Process %d allocated to block %d\n", process_id, i);
      return;
   }
 }
 printf("Process %d cannot be allocated\n", process id);
void deallocate memory(int process id) {
 for (int i = 0; i < NUM BLOCKS; i++) {
    if (block allocated[i] == process id) {
      block allocated[i] = 0;
      printf("Process %d deallocated from block %d\n", process id, i);
      return;
   }
 printf("Process %d not found in memory\n", process id);
int main() {
 allocate memory(120, 1);
 allocate memory(180, 2);
 allocate_memory(90, 3);
 allocate memory(200, 4);
 allocate_memory(80, 5);
```

```
deallocate_memory(2);
allocate_memory(180, 6);
return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab11.c
Name: Bishnu Chalise

Process 1 allocated to block 1
Process 2 allocated to block 2
Process 3 allocated to block 0
Process 4 allocated to block 4
Process 5 allocated to block 3
Process 2 deallocated from block 2
Process 6 allocated to block 2
```

Aim:

To simulate the BEST-FIT contiguous memory allocation technique

Theory:

BEST-FIT: The BEST-FIT memory allocation technique allocates the smallest available block of memory that is large enough to satisfy the request. The memory is searched for the block that most closely matches the size of the request, aiming to minimize wasted space. This method can reduce fragmentation but may require more time to find the best fit.

```
#include <stdio.h>
#define MEMORY SIZE 1000
#define NUM BLOCKS 5
int memory[MEMORY SIZE];
int block size[NUM BLOCKS] = {100, 200, 300, 150, 250};
int block allocated[NUM BLOCKS] = {0};
void allocate memory(int process size, int process id) {
  int best index = -1;
 for (int i = 0; i < NUM BLOCKS; i++) {
    if (!block_allocated[i] && block_size[i] >= process_size) {
      if (best_index == -1 || block_size[i] < block_size[best_index]) {</pre>
        best index = i;
      }
   }
 }
 if (best index !=-1) {
    block allocated[best index] = process id;
    printf("Process %d allocated to block %d\n", process id, best index);
    printf("Process %d cannot be allocated\n", process id);
void deallocate memory(int process id) {
 for (int i = 0; i < NUM BLOCKS; i++) {
    if (block_allocated[i] == process_id) {
      block allocated[i] = 0;
      printf("Process %d deallocated from block %d\n", process id, i);
      return;
   }
 printf("Process %d not found in memory\n", process_id);
```

```
int main() {
  allocate_memory(120, 1);
  allocate_memory(180, 2);
  allocate_memory(90, 3);
  allocate_memory(200, 4);
  allocate_memory(80, 5);

  deallocate_memory(2);
  allocate_memory(180, 6);
  return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab12.c
Name: Bishnu Chalise

Process 1 allocated to block 3
Process 2 allocated to block 1
Process 3 allocated to block 0
Process 4 allocated to block 4
Process 5 allocated to block 2
Process 2 deallocated from block 1
Process 6 allocated to block 1
```

Aim:

To simulate the WORST-FIT contiguous memory allocation technique

Theory:

WORST-FIT: The WORST-FIT memory allocation technique allocates the largest available block of memory to a process. By choosing the biggest block, it aims to leave significant free space for future allocations. However, this can lead to higher fragmentation compared to other techniques like FIRST-FIT and BEST-FIT.

```
#include <stdio.h>
#define MEMORY SIZE 1000
#define NUM BLOCKS 5
int memory[MEMORY SIZE];
int block size[NUM BLOCKS] = {100, 200, 300, 150, 250};
int block allocated[NUM BLOCKS] = {0};
void allocate_memory(int process_size, int process_id) {
  int worst index = -1;
 for (int i = 0; i < NUM BLOCKS; i++) {
    if (!block allocated[i] && block size[i] >= process size) {
      if (worst_index == -1 || block_size[i] > block_size[worst_index]) {
        worst index = i;
     }
   }
 }
 if (worst index != -1) {
    block allocated[worst index] = process id;
    printf("Process %d allocated to block %d\n", process id, worst index);
    printf("Process %d cannot be allocated\n", process id);
void deallocate memory(int process id) {
 for (int i = 0; i < NUM BLOCKS; i++) {
    if (block_allocated[i] == process_id) {
      block allocated[i] = 0;
      printf("Process %d deallocated from block %d\n", process id, i);
      return;
   }
 printf("Process %d not found in memory\n", process_id);
```

```
int main() {
  allocate_memory(120, 1);
  allocate_memory(180, 2);
  allocate_memory(90, 3);
  allocate_memory(200, 4);
  allocate_memory(80, 5);

  deallocate_memory(2);
  allocate_memory(180, 6);
  return 0;
}
```

```
Process 1 allocated to block 2
Process 2 allocated to block 1
Process 3 allocated to block 1
Process 4 cannot be allocated
Process 5 allocated to block 3
Process 2 deallocated to block 4
Process 6 allocated from block 4
```

Aim:

To simulate the FIFO page replacement algorithm

Theory:

FIFO (**First-In**, **First-Out**): FIFO is a page replacement algorithm where the oldest page in memory is replaced when a new page needs to be loaded. Pages are placed in a queue, and the first page to enter is the first one removed when memory is full. This simple approach can lead to suboptimal performance, especially in cases where the oldest page is frequently used.

Algorithm:

- 1. Initialize: Maintain a queue to track the pages in memory.
- 2. Page Requests: For each page request, check if the page is already in memory.
- **3.** Page Replacement: If the page is not in memory and memory is full, remove the page at the front of the queue (oldest) and load the new page.
- 4. Repeat: Continue this process for all incoming page requests.

```
#include <stdio.h>
#define MAX FRAMES 3
#define MAX_PAGES 10
int frames[MAX_FRAMES], pages[MAX_PAGES];
int front = 0, rear = 0, num_frames_filled = 0;
int is_page_in_frames(int page) {
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    if (frames[i] == page) {
      return 1;
   }
 }
  return 0;
void print_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    if (frames[i] == -1) \overline{\{}
      printf("- ");
```

```
} else {
      printf("%d ", frames[i]);
 }
  printf("\n");
void fifo_page_replacement(int num_pages) {
  for (int i = 0; i < num_pages; i++) {
    if (!is_page_in_frames(pages[i])) {
      if (num_frames_filled < MAX_FRAMES) {</pre>
        frames[rear] = pages[i];
        rear = (rear + 1) % MAX_FRAMES;
        num_frames_filled++;
      } else {
        frames[front] = pages[i];
        front = (front + 1) \% MAX_FRAMES;
        rear = (rear + 1) % MAX_FRAMES;
     print_frames();
    }
 }
}
int main() {
 int num_pages;
  printf("Enter number of pages: ");
  scanf("%d", &num_pages);
  printf("Enter the page sequence: ");
  for (int i = 0; i < num_pages; i++) {</pre>
    scanf("%d", &pages[i]);
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    frames[i] = -1;
  fifo_page_replacement(num_pages);
  return 0;
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab14.c
Name: Bishnu Chalise

Enter number of pages: 7
Enter the page sequence: 1 2 3 2 1 4 5
1 - -
1 2 -
1 2 3
4 2 3
4 5 3
```

Aim:

To simulate the LRU page replacement algorithm.

Theory:

LRU (Least Recently Used): The LRU page replacement algorithm replaces the page that has not been used for the longest period of time. By tracking the order in which pages are accessed, the algorithm ensures that the least recently used page is removed when a new page needs to be loaded. This approach generally leads to better performance than FIFO, as it more closely mimics actual memory usage patterns.

Algorithm:

- 1. Initialize: Maintain a list or stack to track the order of page accesses.
- Page Requests: For each page request, check if the page is already in memory.
- **3.** Update List: If the page is in memory, move it to the most recent position in the list.
- **4.** Page Replacement: If the page is not in memory and memory is full, remove the least recently used page (the oldest in the list) and load the new page.
- 5. Repeat: Continue this process for all incoming page requests.

```
#include <stdio.h>
#define MAX_FRAMES 3
#define MAX_PAGES 10

int frames[MAX_FRAMES], pages[MAX_PAGES];
int page_age[MAX_FRAMES];
int num_pages;

void initialize_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {
    frames[i] = -1;
    page_age[i] = 0;
  }
}</pre>
```

```
void print_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {
    if (frames[i] == -1) {
      printf("- ");
    } else {
      printf("%d ", frames[i]);
 }
  printf("\n");
int find_lru_index() {
  int min_age = page_age[0];
  int lru_index = 0;
  for (int i = 1; i < MAX_FRAMES; i++) {</pre>
    if (page_age[i] < min_age) {</pre>
      min_age = page_age[i];
      lru_index = i;
  return lru_index;
void lru_page_replacement() {
  int time = 0;
  for (int i = 0; i < num_pages; i++) {
    int page = pages[i];
    int found = 0;
    for (int j = 0; j < MAX_FRAMES; j++) {
      if (frames[j] == page) {
        page_age[j] = time++;
        found = 1;
        break;
      }
    }
    if (!found) {
      int lru_index = find_lru_index();
      frames[lru_index] = page;
      page_age[lru_index] = time++;
    print_frames();
 }
int main() {
  printf("Enter number of pages: ");
  scanf("%d", &num_pages);
  printf("Enter the page sequence: ");
  for (int i = 0; i < num_pages; i++) {
    scanf("%d", &pages[i]);
  initialize_frames();
  lru_page_replacement();
  return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab15.c
Name: Bishnu Chalise

Enter number of pages: 7
Enter the page sequence: 1 2 3 2 1 4 5
1 - -
2 - -
2 3 -
2 3 -
2 3 1
2 4 1
5 4 1
```

Aim:

To simulate the LFU page replacement algorithm.

Theory:

LFU (Least Frequently Used): The LFU page replacement algorithm replaces the page that has been used the least number of times. It tracks the frequency of page accesses and evicts the page with the lowest usage count when a new page needs to be loaded. This technique assumes that pages accessed frequently in the past will continue to be accessed more often than others.

Algorithm:

- 1. Initialize: Maintain a count for each page that tracks the frequency of its access.
- Page Requests: For each page request, check if the page is already in memory.
- **3.** Update Frequency: If the page is in memory, increment its frequency count.
- **4.** Page Replacement: If the page is not in memory and memory is full, remove the page with the lowest frequency count and load the new page.
- 5. Repeat: Continue this process for all incoming page requests.

```
#include <limits.h>
#include <stdio.h>
#define MAX FRAMES 3
#define MAX PAGES 10
int frames[MAX FRAMES], pages[MAX PAGES];
int page_frequency[MAX_FRAMES];
int num_pages;
void initialize_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    frames[i] = -1;
    page_frequency[i] = 0;
  }
void print_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    if (frames[i] == -1) {
      printf("- ");
    } else {
      printf("%d ", frames[i]);
  printf("\n");
```

```
int find_lfu_index() {
  int min_freq = INT_MAX;
  int lfu_index = 0;
  for (int i = 0; i < MAX_FRAMES; i++) {
    if (page_frequency[i] < min_freq) {</pre>
      min_freq = page_frequency[i];
      lfu_index = i;
    }
  }
  return lfu_index;
void lfu_page_replacement() {
  for (int i = 0; i < num_pages; i++) {
    int page = pages[i];
    int found = 0;
    for (int j = 0; j < MAX_FRAMES; j++) {
      if (frames[j] == page) {
        page_frequency[j]++;
        found = 1;
        break;
    }
    if (!found) {
      int lfu_index = find_lfu_index();
      frames[lfu_index] = page;
      page_frequency[lfu_index] = 1;
    print_frames();
  }
}
int main() {
  printf("Enter number of pages: ");
  scanf("%d", &num_pages);
  printf("Enter the page sequence: ");
  for (int i = 0; i < num_pages; i++) {
    scanf("%d", &pages[i]);
  initialize_frames();
  lfu_page_replacement();
  return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab16.c
Name: Bishnu Chalise

Enter number of pages: 7
Enter the page sequence: 1 2 3 2 1 4 5
1 - -
1 2 -
1 2 3
1 2 3
1 2 3
1 2 4
1 2 5
```

Aim:

To simulate the Optimal page replacement algorithm.

Theory:

Optimal Page Replacement: The Optimal page replacement algorithm replaces the page that will not be used for the longest period in the future. It provides the lowest possible page-fault rate but is not practical for implementation, as it requires future knowledge of the reference string. This algorithm is often used as a benchmark to compare the performance of other page replacement algorithms.

Algorithm:

- 1. Initialize: Maintain a list of future page requests.
- **2.** Page Requests: For each page request, check if the page is already in memory.
- **3.** Page Replacement: If the page is not in memory and memory is full, look ahead in the future reference string and replace the page that is not needed for the longest period of time.
- **4.** Load the new page and continue processing the remaining page requests.

```
#include <limits.h>
#include <stdio.h>

#define MAX_FRAMES 3
#define MAX_PAGES 10

int frames[MAX_FRAMES], pages[MAX_PAGES];
int num_pages;

void initialize_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {
    frames[i] = -1;
  }
}

void print_frames() {
  for (int i = 0; i < MAX_FRAMES; i++) {
    if (frames[i] == -1) {</pre>
```

```
printf("- ");
    } else {
      printf("%d ", frames[i]);
 }
 printf("\n");
int find_optimal_index(int current_index) {
  int furthest_use = -1;
  int optimal_index = 0;
  for (int i = 0; i < MAX_FRAMES; i++) {</pre>
    \quad \text{int } j; \\
    for (j = current_index; j < num_pages; j++) {</pre>
      if (frames[i] == pages[j]) {
        if (j > furthest_use) {
          furthest_use = j;
          optimal_index = i;
        }
        break;
      }
    if (j == num_pages) {
      return i;
  return optimal_index;
void optimal_page_replacement() {
  int current_index = 0;
  for (int i = 0; i < num_pages; i++) {</pre>
    int page = pages[i];
    int found = 0;
    for (int j = 0; j < MAX_FRAMES; j++) {
      if (frames[j] == page) {
        found = 1;
        break;
      }
    }
    if (!found) {
      int optimal_index = find_optimal_index(i + 1);
      frames[optimal_index] = page;
    print_frames();
  }
int main() {
  printf("Enter number of pages: ");
  scanf("%d", &num_pages);
  printf("Enter the page sequence: ");
  for (int i = 0; i < num_pages; i++) {</pre>
    scanf("%d", &pages[i]);
  initialize_frames();
  optimal_page_replacement();
```

```
return 0;
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab17.c
Name: Bishnu Chalise

Enter number of pages: 7
Enter the page sequence: 1 2 3 2 1 4 5
1 - -
1 2 -
1 2 3
1 2 3
1 2 3
4 2 3
5 2 3
```

Aim:

To simulate the following file organization techniques

Theory:

Single-Level Directory: In this technique, all files are stored in a single directory. All users and processes share the same directory space, which can lead to difficulties in organizing and managing files as the number of files grows.

Two-Level Directory: In this method, each user has their own directory, separate from others. The system maintains a master directory that contains pointers to each user's directory. This technique improves file organization but may still become complex for large systems.

Hierarchical Directory: A hierarchical (or tree-structured) directory system allows for multiple levels of directories. This technique enables users to create subdirectories within their directories, forming a tree-like structure. It is highly scalable and efficient for managing a large number of files and directories.

```
#include <stdio.h>
#include <string.h>
#define MAX FILES 10
#define MAX SUBDIRS 5
#define MAX FILE NAME 20
#define MAX_DIR_NAME 20
typedef struct {
  char name[MAX_FILE_NAME];
} File;
typedef struct {
  char name[MAX_DIR_NAME];
  File files[MAX FILES];
  int num files;
} Directory;
typedef struct {
  char name[MAX DIR NAME];
  Directory subdirs[MAX SUBDIRS];
  int num subdirs;
} TwoLevelDirectory;
```

```
typedef struct HierarchicalDirectory HierarchicalDirectory;
struct HierarchicalDirectory {
  char name[MAX DIR NAME];
 HierarchicalDirectory *subdirs[MAX_SUBDIRS];
 int num subdirs;
 File files[MAX_FILES];
 int num_files;
};
void initialize_directory(Directory *dir) { dir->num_files = 0; }
void add_file(Directory *dir, const char *file_name) {
 if (dir->num_files < MAX_FILES) {</pre>
    strcpy(dir->files[dir->num_files].name, file_name);
    dir->num_files++;
 } else {
    printf("Directory is full.\n");
}
void print_single_level_directory(Directory *dir) {
 printf("Files in directory '%s':\n", dir->name);
  for (int i = 0; i < dir->num_files; i++) {
    printf(" %s\n", dir->files[i].name);
 }
}
void initialize_two_level_directory(TwoLevelDirectory *two_level_dir) {
 two_level_dir->num_subdirs = 0;
void add_subdir(TwoLevelDirectory *two_level_dir, const char *subdir_name) {
 if (two_level_dir->num_subdirs < MAX_SUBDIRS) {</pre>
    strcpy(two_level_dir->subdirs[two_level_dir->num_subdirs].name,
           subdir_name);
    initialize_directory(&two_level_dir->subdirs[two_level_dir->num_subdirs]);
    two_level_dir->num_subdirs++;
  } else {
    printf("Two-level directory is full.\n");
}
void print_two_level_directory(TwoLevelDirectory *two_level_dir) {
  printf("Two-Level Directory '%s':\n", two_level_dir->name);
  for (int i = 0; i < two_level_dir->num_subdirs; i++) {
    printf(" Subdirectory '%s':\n", two_level_dir->subdirs[i].name);
    for (int j = 0; j < two_level_dir->subdirs[i].num_files; j++) {
      printf("
                 %s\n", two_level_dir->subdirs[i].files[j].name);
    }
 }
void initialize_hierarchical_directory(HierarchicalDirectory *hier_dir,
                                        const char *name) {
  strcpy(hier_dir->name, name);
 hier_dir->num_files = 0;
 hier_dir->num_subdirs = 0;
void add_file_hierarchical(HierarchicalDirectory *hier_dir,
                           const char *file_name) {
  if (hier_dir->num_files < MAX_FILES) {</pre>
```

```
strcpy(hier dir->files[hier dir->num files].name, file name);
    hier_dir->num_files++;
  } else {
    printf("Hierarchical directory is full.\n");
}
void add_subdir_hierarchical(HierarchicalDirectory *hier_dir,
                               HierarchicalDirectory *subdir) {
  if (hier_dir->num_subdirs < MAX_SUBDIRS) {</pre>
    hier_dir->subdirs[hier_dir->num_subdirs] = subdir;
    hier_dir->num_subdirs++;
  } else {
    printf("Hierarchical directory has reached maximum subdirectories.\n");
}
void print_hierarchical_directory(HierarchicalDirectory *hier_dir, int level) {
  for (int i = 0; i < level; i++) {</pre>
    printf(" ");
  printf("Directory '%s':\n", hier_dir->name);
  for (int i = 0; i < hier_dir->num_files; i++) {
    for (int j = 0; j < level + 1; j++) {
                ");
      printf("
    printf("File: %s\n", hier_dir->files[i].name);
  for (int i = 0; i < hier_dir->num_subdirs; i++) {
    print_hierarchical_directory(hier_dir->subdirs[i], level + 1);
}
int main() {
  Directory single_level_dir;
  strcpy(single_level_dir.name, "Root");
  initialize_directory(&single_level_dir);
  add_file(&single_level_dir, "file1.txt");
add_file(&single_level_dir, "file2.txt");
  print_single_level_directory(&single_level_dir);
  TwoLevelDirectory two_level_dir;
  strcpy(two_level_dir.name, "MainDir");
  initialize_two_level_directory(&two_level_dir);
  add_subdir(&two_level_dir, "SubDir1");
  add_file(&two_level_dir.subdirs[0], "file3.txt");
add_file(&two_level_dir.subdirs[0], "file4.txt");
  add_subdir(&two_level_dir, "SubDir2");
  add_file(&two_level_dir.subdirs[1], "file5.txt");
  print_two_level_directory(&two_level_dir);
  HierarchicalDirectory root_dir;
  initialize_hierarchical_directory(&root_dir, "RootHier");
  HierarchicalDirectory subdir1;
  initialize_hierarchical_directory(&subdir1, "SubDirHier1");
  HierarchicalDirectory subdir2;
  initialize_hierarchical_directory(&subdir2, "SubDirHier2");
  add_file_hierarchical(&root_dir, "file6.txt");
  add_file_hierarchical(&root_dir, "file7.txt");
  add_file_hierarchical(&subdir1, "file8.txt");
  add_file_hierarchical(&subdir2, "file9.txt");
  add_subdir_hierarchical(&root_dir, &subdir1);
```

```
add_subdir_hierarchical(&root_dir, &subdir2);
print_hierarchical_directory(&root_dir, 0);
return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab18.c
Name: Bishnu Chalise
Files in directory 'Root':
  file1.txt
  file2.txt
Two-Level Directory 'MainDir':
  Subdirectory 'SubDir1':
    file3.txt
    file4.txt
  Subdirectory 'SubDir2':
    file5.txt
Directory 'RootHier':
File: file6.txt
  File: file7.txt
 Directory 'SubDirHier1':
File: file8.txt
  Directory 'SubDirHier2':
    File: file9.txt
```

Aim:

To simulate all file allocation strategies

Theory:

Sequential Allocation: In sequential allocation, files are stored in contiguous blocks on the disk. Each file occupies a set of consecutive disk blocks, making it easy to access sequentially. However, it can lead to external fragmentation and inefficient use of space for files that grow over time.

Indexed Allocation: In this method, an index block is created for each file. The index block contains pointers to the actual blocks where the file data is stored. This allows random access to file blocks, eliminating fragmentation but requiring extra space for index blocks.

Linked Allocation: Linked allocation stores file blocks in a linked list, where each block points to the next one. There is no need for contiguous allocation, reducing fragmentation. However, it can lead to slower access times due to the need to traverse the linked list to read or write a file.

```
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#define MAX BLOCKS 10
#define MAX FILES 5
#define MAX_FILE_NAME 20
int disk[MAX_BLOCKS];
typedef struct {
  char name[MAX_FILE_NAME];
  int start_block;
  int size;
} SequentialFile;
typedef struct {
  char name[MAX FILE NAME];
  int index blocks[MAX BLOCKS];
  int num blocks;
} IndexedFile;
typedef struct Node {
```

```
int block;
  struct Node *next:
} LinkedNode;
typedef struct {
  char name[MAX_FILE_NAME];
  LinkedNode *head;
} LinkedFile;
void initialize_disk() {
  for (int i = 0; i < MAX_BLOCKS; i++) {</pre>
    disk[i] = 0;
}
void print_disk() {
  printf("Disk blocks: ");
  for (int i = 0; i < MAX_BLOCKS; i++) {</pre>
    printf("%d ", disk[i]);
  printf("\n");
void allocate_sequential(SequentialFile *file, int start_block, int size) {
  if (start_block + size > MAX_BLOCKS) {
    printf("Error: Not enough space on the disk.\n");
    return;
  for (int i = start_block; i < start_block + size; i++) {</pre>
    if (disk[i] != 0) {
      printf("Error: Space already allocated.\n");
      return;
    }
  file->start_block = start_block;
  file->size = size;
  strcpy(file->name, "SequentialFile");
  for (int i = start_block; i < start_block + size; i++) {</pre>
    disk[i] = 1;
  printf("Sequential file '%s' allocated from block %d to %d.\n", file->name,
         start_block, start_block + size - 1);
}
void allocate_indexed(IndexedFile *file, int blocks[], int num_blocks) {
  if (num_blocks > MAX_BLOCKS) {
    printf("Error: Too many blocks requested.\n");
    return;
  for (int i = 0; i < num_blocks; i++) {</pre>
    if (blocks[i] >= MAX_BLOCKS || disk[blocks[i]] != 0) {
      printf("Error: Block %d is not available.\n", blocks[i]);
      return;
    }
  file->num_blocks = num_blocks;
  strcpy(file->name, "IndexedFile");
  for (int i = 0; i < num_blocks; i++) {</pre>
    disk[blocks[i]] = 1; // Mark block as allocated
    file->index_blocks[i] = blocks[i];
  printf("Indexed file '%s' allocated at blocks: ", file->name);
  for (int i = 0; i < num_blocks; i++) {</pre>
    printf("%d ", file->index_blocks[i]);
```

```
}
 printf("\n");
void allocate linked(LinkedFile *file, int blocks[], int num_blocks) {
 LinkedNode *prev = NULL;
 LinkedNode *head = NULL;
 for (int i = 0; i < num_blocks; i++) {</pre>
    if (blocks[i] >= MAX_BLOCKS || disk[blocks[i]] != 0) {
      printf("Error: Block %d is not available.\n", blocks[i]);
      while (head != NULL) {
        LinkedNode *temp = head;
        head = head->next;
        free(temp);
      }
      return;
    }
   disk[blocks[i]] = 1;
    LinkedNode *node = (LinkedNode *)malloc(sizeof(LinkedNode));
    node->block = blocks[i];
    node->next = NULL;
    if (prev == NULL) {
     head = node;
    } else {
     prev->next = node;
   prev = node;
 file->head = head;
 strcpy(file->name, "LinkedFile");
 printf("Linked file '%s' allocated with blocks: ", file->name);
 LinkedNode *current = head;
 while (current != NULL) {
    printf("%d ", current->block);
    current = current->next;
 printf("\n");
int main() {
 initialize_disk();
 print_disk();
 SequentialFile seq_file;
 int start_block = 2;
 int size = 4;
 allocate_sequential(&seq_file, start_block, size);
 print_disk();
 IndexedFile idx_file;
 int index_blocks[] = {6, 7, 8};
 allocate_indexed(&idx_file, index_blocks, 3);
 print_disk();
 LinkedFile lnk_file;
 int linked_blocks[] = {1, 3, 5};
 allocate_linked(&lnk_file, linked_blocks, 3);
 print_disk();
  return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab19.c
Name: Bishnu Chalise

Disk blocks: 0 0 0 0 0 0 0 0 0 0
Sequential file 'SequentialFile' allocated from block 2 to 5.
Disk blocks: 0 0 1 1 1 1 0 0 0 0
Indexed file 'IndexedFile' allocated at blocks: 6 7 8
Disk blocks: 0 0 1 1 1 1 1 1 1 0
Error: Block 3 is not available.
Disk blocks: 0 1 1 1 1 1 1 1 0
```

Aim:

To simulate disk scheduling algorithms

Theory:

FCFS (First-Come, First-Served): FCFS is the simplest disk scheduling algorithm. Requests are handled in the order they arrive, without considering their proximity to the current head position. This can lead to high seek times if requests are scattered across the disk.

SCAN (Elevator Algorithm): In the SCAN algorithm, the disk arm moves in one direction, servicing all requests until it reaches the end of the disk. It then reverses direction and continues servicing requests. This approach reduces the variability in seek times compared to FCFS.

LOOK: The LOOK algorithm is a variation of SCAN. Instead of going all the way to the end of the disk, the arm only moves as far as the last request in each direction before reversing. This reduces unnecessary movement, improving efficiency.

Algorithm:

- 1. FCFS:
- a. Initialize a queue of disk requests.
- **b.** Process each request in the order it arrives.
- **c.** Move the disk head to the requested track and process the request.

2. SCAN:

- **a.** Start by moving the disk head in one direction (either toward the lowest or highest track).
- **b.** Service all requests in that direction until the disk's end is reached.
- c. Reverse the direction and continue servicing requests.

3. L00K:

- a. Move the disk head in one direction.
- **b.** Service requests up to the last request in that direction (don't go to the disk's end).
- c. Reverse the direction and repeat the process.

```
#include <stdio.h>
#include <stdlib.h>
#define MAX REQUESTS 10
void fcfs(int requests[], int num_requests, int start) {
  int seek count = 0;
  int distance:
  int cur_track = start;
  printf("FCFS Disk Scheduling:\n");
  for (int i = 0; i < num_requests; i++) {</pre>
    distance = abs(requests[i] - cur_track);
    seek_count += distance;
    cur_track = requests[i];
    printf("Move to track %d\n", requests[i]);
  }
  printf("Total seek time: %d\n", seek_count);
void scan(int requests[], int num_requests, int start, int direction) {
  int seek_count = 0;
  int cur_track = start;
  int distance;
  printf("SCAN Disk Scheduling:\n");
  int direction_change = (direction == 1) ? 0 : 1;
  int max = 0, min = 0;
  for (int i = 0; i < num\_requests; i++) {
    if (requests[i] > max)
      max = requests[i];
    if (requests[i] < min)</pre>
      min = requests[i];
  if (direction == 1) {
    for (int i = 0; i < num\_requests; i++) {
      if (requests[i] >= cur_track) {
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
      }
    if (cur_track != max) {
      seek_count += abs(max - cur_track);
      cur_track = max;
      printf("Move to track %d\n", max);
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] < cur_track) {</pre>
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
    }
  } else {
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] <= cur_track) {</pre>
```

```
distance = abs(requests[i] - cur_track);
        seek count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
    if (cur_track != min) {
      seek_count += abs(min - cur_track);
      cur_track = min;
      printf("Move to track %d\n", min);
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] > cur_track) {
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
      }
    }
  printf("Total seek time: %d\n", seek_count);
void look(int requests[], int num_requests, int start, int direction) {
  int seek_count = 0;
  int cur_track = start;
  int distance;
  int max = 0, min = 0;
  printf("LOOK Disk Scheduling:\n");
  for (int i = 0; i < num_requests; i++) {</pre>
    if (requests[i] > max)
      max = requests[i];
    if (requests[i] < min)</pre>
      min = requests[i];
  if (direction == 1) {
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] >= cur_track) {
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
      }
    for (int i = 0; i < num\_requests; i++) {
      if (requests[i] < cur_track) {</pre>
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
      }
    }
  } else {
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] <= cur_track) {</pre>
        distance = abs(requests[i] - cur_track);
        seek_count += distance;
        cur_track = requests[i];
        printf("Move to track %d\n", requests[i]);
      }
    }
    for (int i = 0; i < num_requests; i++) {</pre>
      if (requests[i] > cur_track) {
```

```
distance = abs(requests[i] - cur track);
        seek count += distance;
        cur track = requests[i];
        printf("Move to track %d\n", requests[i]);
   }
 }
 printf("Total seek time: %d\n", seek_count);
int main() {
 int requests[] = {55, 58, 39, 18, 90, 160, 150};
  int num_requests = sizeof(requests) / sizeof(requests[0]);
 int start = 50;
 int direction = 1;
 printf("Disk Scheduling Algorithms Simulation\n");
 fcfs(requests, num_requests, start);
 printf("\n");
 scan(requests, num_requests, start, direction);
 printf("\n");
 look(requests, num_requests, start, direction);
 return 0;
}
```

```
rudy@rudy:~/Desktop/untitled/github/d-tfn/operating-system/code$ ./cmprun.sh lab20.c
Name: Bishnu Chalise
Disk Scheduling Algorithms Simulation
FCFS Disk Scheduling:
Move to track 55
Move to track 58
Move to track 39
Move to track 18
Move to track 90
Move to track 160
Move to track 150
Total seek time: 200
SCAN Disk Scheduling:
Move to track 55
Move to track 58
Move to track 90
Move to track 160
Move to track 55
Move to track 39
Move to track 18
Total seek time: 252
LOOK Disk Scheduling:
Move to track 55
Move to track 58
Move to track 90
Move to track 160
Move to track 55
Move to track 39
Move to track 18
```