

# Lectures notes on knot theory

Andrew Berger

February 18, 2016

# Contents

<b>1</b>	<b>Disclaimer</b>	<b>3</b>
<b>2</b>	<b>1/19/16: Introduction + Motivation</b>	<b>4</b>
<b>3</b>	<b>1/21/16: 2nd class</b>	<b>5</b>
3.1	Logistical things . . . . .	5
3.2	Minimal introduction to point-set topology . . . . .	5
3.3	Equivalence of knots . . . . .	6
3.4	Reidemeister moves . . . . .	6
<b>4</b>	<b>1/26/16: recap of the last lecture</b>	<b>9</b>
4.1	Recap of last lecture . . . . .	9
4.2	Intro to knot complement . . . . .	9
4.3	Hard Unknots . . . . .	9
<b>5</b>	<b>1/28/16</b>	<b>11</b>
5.1	Logistical things . . . . .	11
5.2	Question from last time . . . . .	11
5.3	Connect sum operation, knot cancelling, prime knots . . . . .	11
<b>6</b>	<b>2/2/16</b>	<b>13</b>
6.1	Orientations . . . . .	13
6.2	Linking number . . . . .	13
<b>7</b>	<b>2/4/16</b>	<b>14</b>
7.1	Logistical things . . . . .	14
7.2	Seifert Surfaces . . . . .	14
7.3	Intro to research . . . . .	15
<b>8</b>	<b>2/9/16 – The trefoil is knotted</b>	<b>16</b>
8.1	The trefoil is not the unknot . . . . .	16
8.2	Braids . . . . .	16
8.2.1	The braid group . . . . .	16
<b>9</b>	<b>2/11: Coloring</b>	<b>17</b>
9.1	Logistical happenings . . . . .	17
9.2	(Tri)Colorings . . . . .	17
<b>10</b>	<b>2/16: <math>\pi_1</math></b>	<b>18</b>
10.1	Logistical things . . . . .	18
10.2	Crash course on the fundamental group . . . . .	18
<b>11</b>	<b>2/18: Wirtinger presentation</b>	<b>20</b>
<b>12</b>	<b>2/23: POLYNOMIALS</b>	<b>21</b>

## 1 Disclaimer

Caution - these lecture notes have not been proofread and may contain errors, due to either the lecturer or the scribe. Please send any corrections or suggestions to [andbberger@berkeley.edu](mailto:andbberger@berkeley.edu)

## 2 1/19/16: Introduction + Motivation

Start with a joke: “Applied Math majors may have misread the course title”

- hand out syllabus & table of knots
- administration & logistics
- “idea” of class structure
- find out about class (their backgrounds & interests)
  
- introduce some material:
  - 1) knots come from real life: tie ends of string together
    - unknots
    - simplest possible knot = trefoil
  - 2) not everything is a circle
    - want to distinguish them
    - knot complements
    - techniques?
  - 3) links ? can be complicated while component knots simple
  - 4) crossings / deformations
    - swap crossings of trefoil to get its mirror (interesting fact: they’re not equivalent!)
  - 5) usefulness:
    - want to know when objects are knotted
    - get all 3-dimensional spaces
    - applications to physics

### 3 1/21/16: 2nd class

#### 3.1 Logistical things

Talk to Chris if you're uncomfortable with group theory. There are going to be two projects (long and short) - we'll start breaking up in to groups to work on those starting the second week of February.

#### 3.2 Minimal introduction to point-set topology

Just to set terms and notation for future reference.

**Definition 1.** Some miscellaneous definitions:

- $\mathbb{R}^n := \mathbb{R} \oplus \dots \oplus \mathbb{R}$
- $B^n = \{x \in \mathbb{R}^n \mid |x| \leq 1\}$  note that in general this includes the boundary (which we call  $\partial B^n = S^{n-1}$ , the  $(n-1)$ -sphere)
- There are many ways to get the inclusion  $\mathbb{R}^m \subseteq \mathbb{R}^n$  for  $m \leq n$ , as a convention just take the first  $m$  elements from  $\mathbb{R}^n$

**Remark 1.** The (closed) line and circle are the only 1-dimensional topological spaces (once we have defined appropriate notions of equivalence). More on this later

**Definition 2** (Homeomorphism). We say that  $f$  is a *homeomorphism* if it is a bicontinuous bijection (both  $f$  and its inverse are continuous). Two topological spaces  $X, Y$  are homeomorphic if there exists  $f : X \rightarrow Y$  a homeomorphism, denoted  $X \cong Y$

**Remark 2.** Given  $X, Y$  topological spaces and  $f : X \rightarrow Y$  a homeomorphism we can 'pull' the notion of openness in  $X$  into a notion of openness in  $Y$ :  $A \subseteq Y$  is open if  $f^{-1}(A)$  is open in  $X$

**Example 1.** Examples of homeomorphisms

- $(0, 1) \cong \mathbb{R}$ . As an exercise find the homeomorphism that is witness to this fact (hint: use an arctangent)
- A square in the plane is homeomorphic to a circle in the plane
- $\mathbb{R}^n \not\cong \mathbb{R}^m$  for  $n \neq m$ : dimension is invariant under homeomorphism. This is a deep result that is supposed to be hard to prove

**Definition 3** (Knot). A knot is a one-dimensional subset of  $\mathbb{R}^3$  that is homeomorphic to  $S^1$ .

We can specify a knot  $K$  by specifying an embedding (smooth injective)  $f : S^1 \rightarrow \mathbb{R}^3$  so that  $K = f(S^1)$ . For  $f$  to be smooth, all of its derivatives must exist.

**Example 2.** Examples of embeddings specifying knots

- $f = \mathbb{1}$  (abuse of notation here) specifies a circle
- The infinite non-knot example we looked at yesterday fails to be a knot because the derivative does not exist at the limit point

**Definition 4** (Link). Given a collection of knots  $\{K_i\}$ , we define a link  $L = \bigcup_i K_i$  (disjoint union)

The study of links is different from the study of knots, due to “linking behavior”. Roughly speaking: knots can be very complicated as well their disjoint unions, but moreover, links can get very complicated while their connected components may all be unknots.

### 3.3 Equivalence of knots

Our word for equivalence of knots is ambient isotopy. This refers to the fact that the homeomorphism that is witness to the equivalence of the knots acts on the ambient space the knots lives and not only on the knot itself.

**Definition 5** (Equivalence of knots). For  $K_1, K_2$  knots, we say that  $K_1 \cong_{isotopic} K_2$  if there exists a (orientation-preserving) homeomorphism  $f : \mathbb{R}^3 \rightarrow \mathbb{R}^3$  such that  $f(K_1) = K_2$ .

More precisely we require that there exists a 1-parameter family  $\{f_t\}_{0 \leq t \leq 1}$  of smooth homeomorphisms such that  $f_0 = \mathbb{1}$  and  $f_1 = f$ . In particular, we cannot have an isotopy that shrinks a knot to a point.

This aligns pleasingly for the intuitive notion that knots are equal when they can be deformed to each other:  $f$  here is just the global deformation.

**Remark 3.** Careful about the notation here:  $K_1 \cong K_2 \cong S^1$  holds for any knots  $K_1, K_2$

**Definition 6** (Knotted). We say that a knot  $K$  is *knotted* if  $K \not\cong_{isotopic} S^1$

**Definition 7** (Planar diagram of  $K$ ). When we visualize knots we make some projection  $\mathbb{R}^3 \rightarrow \mathbb{R}^2$  (with defined coordinate system). The naïve projection that discards information about the crossings is called the universe and is not necessarily unique. The planar diagram is this projection with crossing information captured.

**Example 3.** Examples of invariants (under equivalence)

- Dimension is an invariant of isotopy
- The crossing number is the minimal number of crossings in a given diagram

### 3.4 Reidemeister moves

What does an isotopy mean for planar diagrams? The famous Reidemeister moves

**Definition 8** (Reidemeister moves). Refer to figure 1

1. R0: You can straighten wiggly lines
2. R1: You can undo twists
3. R2: You can separate underpasses (that don’t intersect)
4. R3: You can move a line behind an intersection across the intersection

**Theorem 1.**  $K_1 \cong_{isotopic} K_2$  iff their diagrams can be obtained from each other using the Reidemeister moves

Wherein Chris struggles to draw a trefoil.

- Exercise 1.**
1. Show that the figure eight knot is equivalent to its mirror.
  2. Take your favorite knot diagram and show that it is equivalent to the diagram obtained as follows: Take an arc on the left-most side of your diagram, put a twist in it and then pull it all over to the right side of the diagram. (The point of this exercise is to give an abstract proof of equivalence using Reidemeister moves without worrying about the explicit diagram.)
  3. Show that an unknot with a couple of loops is equivalent to the unknot *without* using R1

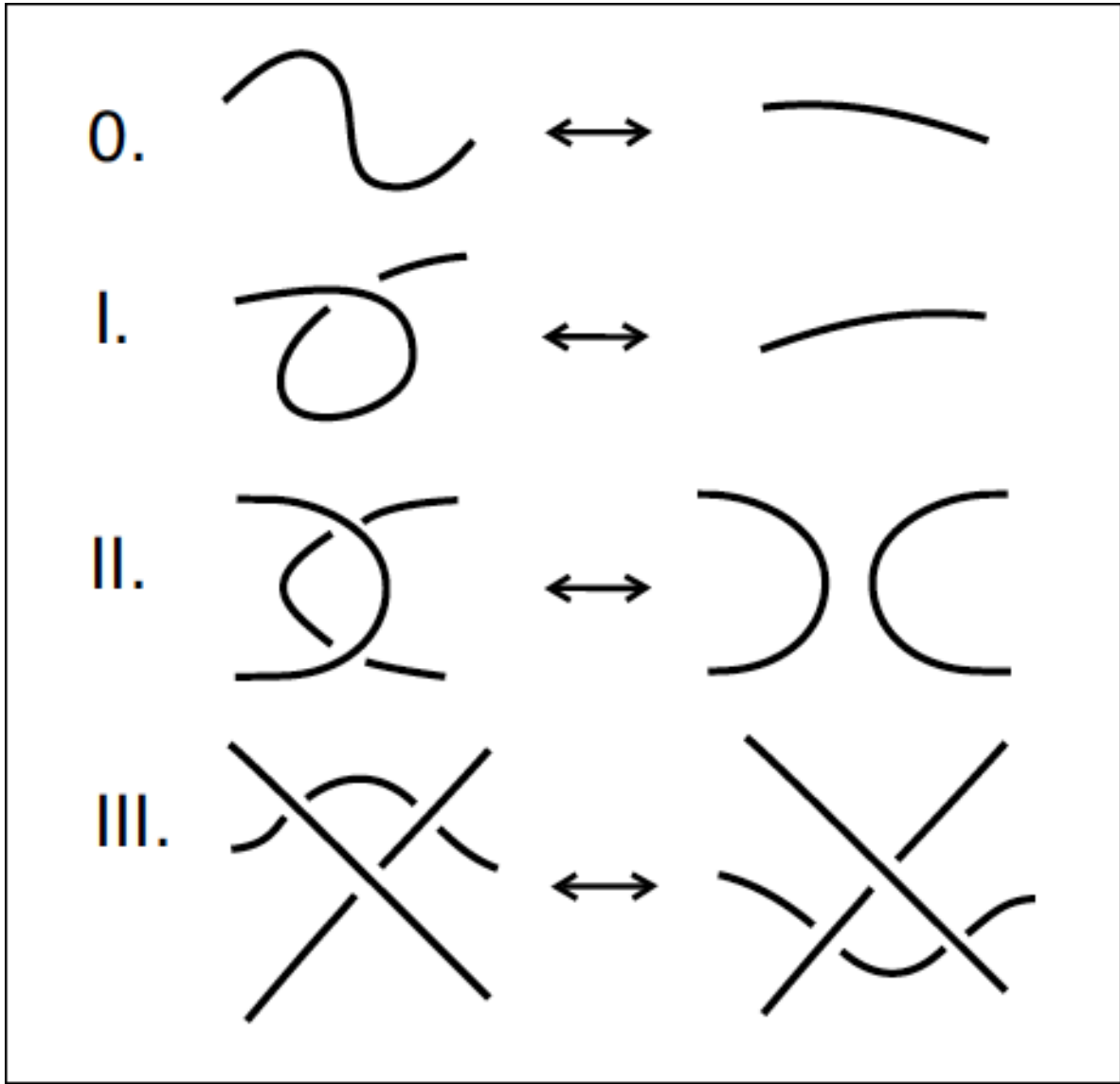


Figure 1: Sketch of allowed Reidemeister moves



## 4 1/26/16: recap of the last lecture

### 4.1 Recap of last lecture

What does it mean to be a knot? To be knotted?

**Recall 1.** Recall that a knot  $K$  is a subset of  $\mathbb{R}^3$  that is homeomorphic to  $S^1$ .

We could equally as well have defined a knot to be the image  $f(S^1)$  of a smooth embedding  $f : S^1 \rightarrow \mathbb{R}^3$ .

**Notation 1.** Let's fix our notation for ambient isotopy (the kind that captures a notion of knottedness) and homeomorphism (under which all knots are equivalent), being very careful to distinguish between them.

- For  $K_1, K_2$  knots that are equivalently knotted (or isotopic) and write  $K_1 \cong K_2$
- We write homeomorphism with  $\approx$ . For instance, all knots  $K$  are homeomorphic to  $S^1$ ,  $K \approx S^1$

### 4.2 Intro to knot complement

We expect that the knot complement  $\mathbb{R}^3 \setminus K$  should somehow capture the notion of knottedness. Can we formalize that?

'It's not the space itself - it's what the knot is doing in 3-dimensional space that matters, and this theorem captures that precisely'

**Theorem 2** (Gordon-Luecke). *If  $\mathbb{R}^3 \setminus K_1 \approx \mathbb{R}^3 \setminus K_2$  then  $K_1 \cong K_2$*

**Remark 4.** The fundamental group sometimes allows us to answer the question ' $\mathbb{R}^3 \setminus K_1 \approx \mathbb{R}^3 \setminus K_2$ ?' and therefore turn this into an algebraic problem.

Francesca pointed out that the converse of the preceding theorem is easily proved, directly from the definition of equivalence. Do it as an exercise!

**Question 1.** Marissa asked: when is a knot equal to its mirror? Chris isn't aware of any general classification results - maybe when the Jones polynomial has palindromic coefficients. Another question: if  $K_1 \cong K_2$  is it true that  $K_1^* \cong K_2^*$  (here  $K^*$  indicates the knot mirror) Yet another question: if  $K_1 \cong K_2$  and  $K_1 \cong K_1^*$  is it true that  $K_2 \cong K_2^*$ ? Someone pointed out that if you can prove the previous statement, then this statement is easily proved to be true.

### 4.3 Hard Unknots

Chris struggles again to draw a trefoil.

Let's fix a notion of complexity for knot diagrams - we say that the complexity of a diagram is the crossing number of that diagram.

Recall that a knot is equivalent to the unknot if it *has a diagram* with crossing number zero. A natural direction to head in towards proving that the trefoil is not the unknot might be to show that any Reidemeister move can not decrease the complexity of the trefoil diagram - but one needs to be very careful here as there may be a sequence that increases in complexity for a while before decreasing.

This motivates the concept of hard unknots: unknots that require a sequence of Reidemeister moves increasing the complexity of the diagram before the complexity finally falls to zero.

To begin with let's study a hard unknot called 'the culprit' from the paper *Hard Unknots and Collapsing Tangles* (Kauffman, Lambropoulou). Can you see how to untangle this knot?

## 5 1/28/16

Agenda: what sort of invariants; solved problem from last class; oriented knots, linking numbers, fundamental group of a knot

### 5.1 Logistical things

Start thinking about what you want to do for a project.

### 5.2 Question from last time

If  $K \cong J$  does  $K^* \cong J^*$ ?

One idea is that if  $D_K$  and  $D_J$  are diagrams for  $K, J$ , then we can find a sequence of R-moves  $\{R_i\}$  transforming  $D_K$  to  $D_J$ . It's easily imagined that if we just 'flip' the moves in  $\{R_i\}$  then that new flipped sequence transforms  $D_{K^*}$  to  $D_{J^*}$ . This idea can be used to prove that  $K \cong J \implies K^* \cong J^*$ .

Chris came up with a different proof:

First a false proof:

Fact:  $\mathbb{R}^3 \setminus K \approx \mathbb{R}^3 \setminus K^*$ . So if one naïvely applies the Gordon-Luecke theorem, they conclude  $K \cong K^*$ , which is obviously false.

What gives here is that the Gordon-Luecke theorem requires that the homeomorphism is orientation-preserving. The map that's witness to the fact that  $\mathbb{R}^3 \setminus K \approx \mathbb{R}^3 \setminus K^*$  is orientation reversing. Thus the Gordon-Luecke theorem does not apply.

Let's do a real proof:

**Lemma 1.** If  $K \cong J$  then  $K^* \cong J^*$

*Proof.* Suppose that  $K, J$  are knots such that  $K \cong J$ . By the Gordon-Luecke theorem,  $\mathbb{R}^3 \setminus K \approx \mathbb{R}^3 \setminus J$  in an orientation preserving fashion.

$\mathbb{R}^3 - K^* \approx \mathbb{R}^3 - K$  with an orientation reversing homeomorphism; same for  $J$ . Now we can easily obtain the desired orientation preserving map that's witness to  $\mathbb{R}^3 - K^* \approx \mathbb{R}^3 - J^*$  by the appropriate composition (note here, that the composition of an even number of orientation-reversing maps is orientation-preserving).  $\square$

### 5.3 Connect sum operation, knot cancelling, prime knots

**Definition 9** (connect sum). This is a general concept from algebraic topology, but for our purpose the connect sum of knots  $K$  and  $J$  (written  $K \# J$ ) is given by cutting knots  $K$  and  $J$  at some point and pasting them together.

**Remark 5.** Is this operation well-defined? In particular, does it matter where we choose to cut and paste the knots together?

Yes it is, and no it does not. A pictorial proof: if we connected the knots in one place and we want to connect in another we can just slide the knots along each other until we arrive at the desired location.

One thing that does matter is orientation - if we don't assign an orientation there are two ways to join the arcs together at one location. We can resolve this by giving knots an orientation and making  $\#$  preserve that.

Chris refined his technique for drawing the trefoil; it is now foolproof.

**Question 2.** Let  $\mathcal{K}$  be the collection of knots (up to  $\cong$ ). Is  $(K, \#)$  a group?

- $K \# 0 \cong K$ , so there is an identity element. (0 indicates the unknot)
- $K_1 \# K_2 \cong K_2 \# K_1$ , as can be seen by a pictorial proof, so this operation is commutative.
- It's associative too (we can attach anywhere and at any time).
- Okay, do knot inverses exist? No. Knots cannot be cancelled.

**Remark 6.** In particular, it is not possible that  $0 \# 0 \cong K$  with  $K$  knotted, despite how “complex” the knot diagrams for 0 are.

**Theorem 3.** *Knots cannot be cancelled under connect sum*

*Proof.* This is known as the Mazur-Swindle.

To prove this, we have to expand our notion of knots to include wild knots (those that shrink to a point).

Suppose that  $K \# J \cong 0$ . Now construct the infinite connect sum  $K \# J \# K \# J \dots$  letting the knots shrink in space as the sum goes along so that they occupy a closed topological space<sup>1</sup>.

We have that  $(K \# J) \# (K \# J) \dots \cong 0$  but  $K \# (J \# K) \# (J \# K) \cong K \# 0 \cong K$ , and thus our knots must be unknotted to begin with!  $\square$

**Question 3** (Unsolved). Does  $c(K \# J) = c(K) + c(J)$  where  $c$  is the minimal crossing number operation.

**Definition 10** (Prime, composite knots).  $K$  is called composite if it can be given as  $K_1 \# K_2 \cong K$  for  $K_1 \not\cong 0$  and  $K_2 \not\cong 0$ .

Otherwise,  $K$  is called prime.

**Exercise 2.** Show that the trefoil and figure 8 knots are prime.

---

<sup>1</sup>If we let the knots go off to infinity without shrinking them, then the result does not form a loop and is not captured by our notion of knots. By shrinking them they become wild, but are still homeomorphic to  $S^1$

## 6 2/2/16

### 6.1 Orientations

Put more structure on our knots:

- 1) Orient it by labelling the diagram with an arrow.
- 2) To each crossing we can associate a  $\pm 1$ . To make this well-defined: rotate the over-strand CCW (counterclockwise) with respect to the under-strand, and assign  $+1$  if the orientations (i.e. arrows) on those strands match up.

An *invariant* of a knot is one which doesn't depend on the diagram, which means it should be independent of the Reidemeister moves!

The crossing number (number of crossings) is not an invariant. Although R3 preserves it, R2 decreases (or increases) it by 2, and R1 decreases (or increases) it by 1.

OK, so what if we also consider the signs of the crossings? That is, maybe the *writhe* (number of crossings counted with their sign) is a better choice.... No, it is also not an invariant. Although R3 preserves it, and R2 preserves it (**exercise!**), but R1 changes it by  $\pm 1$ .

**Remark 7.** Let  $K$  denote the unknot diagram obtained from the standard  $O$  by two R1-moves of opposite orientation, i.e. the unit circle but with two “opposing twists.” *Remember our first exercise assigned at the beginning of the semester:* prove  $K \cong O$  via R2+R3 moves only. Now, this is consistent with the fact that the writhe is invariant under R2+R3 moves, noting that the writhe of the standard diagram for the unknot is 0. But the writhe changes by  $\pm 1$  for the R1-move, so it is impossible to do the analogous exercise above where the construction of  $K$  involves two “twists” with the same orientation (or even with just a single twist!).

### 6.2 Linking number

So far, our different counts of crossings (with signs) doesn't get us an invariant. Well, we do get something for links!

The *linking number* of a link  $L = K_1 \cup K_2$  is  $lk(L) \equiv lk(K_1, K_2) = \frac{1}{2} \sum_c \text{sign}(c)$ , where  $c$  sums over the crossings between  $K_1$  and  $K_2$  (not the self-crossings of  $K_i$ !). The factor  $\frac{1}{2}$  is to be consistent with our intuition that  $K_1$  “links”  $K_2$  while  $K_2$  “also links”  $K_1$ , so we shouldn't double-count.

**Exercise 3.** This is an invariant. R1 is obvious, it introduces self-intersections.

**Example 4.**  $lk(\text{unlink}) = 0$  independent of orientation.

$lk(\text{Hopf link}) = \pm 1$  depending on orientation.

$lk(\text{Whitehead link}) = 0...$  but it's “visually linked”! An invariant does not always provide a complete classification.

## 7 2/4/16

### 7.1 Logistical things

Introduce each other and discuss our interests for potential short projects (this will help us decide who to work with). Projects should begin by end of next week.

### 7.2 Seifert Surfaces

A *surface* (2-manifold)  $\Sigma \subset \mathbb{R}^3$  can be roughly regarded as a collection of points (a “space”), such that for each point  $x \in \Sigma$  there is a neighborhood of points which is homeomorphic either to  $\mathbb{R}^2$  or  $\mathbb{R}_{\geq 0}^2$  (the upper half-plane)... so the space is 2-dimensional. The points whose neighborhood is  $\mathbb{R}_{\geq 0}^2$  form the *boundary*  $\partial\Sigma$  of the surface, and is 1-dimensional.

Likewise, knots are 1-manifolds without boundary: they can be roughly regarded as a collection of points, such that for each point  $x \in K$  there is a neighborhood of points which is homeomorphic to  $\mathbb{R}$ .

**Remark 8.** Manifolds can be either *compact* or *noncompact*. **Assume all of our spaces are compact.** For the non-topologists: think of a compact surface as a surface which “cannot be arbitrarily stretched to infinity” in the ambient space  $\mathbb{R}^3$ ... the open interval  $(0, 1)$  can be arbitrarily stretched, as it is homeomorphic to  $\mathbb{R}$ , while the closed interval  $[0, 1]$  cannot.

Surfaces can be orientable or not. The Mobius strip is not orientable.

(Oriented) Surfaces without boundary are determined completely by the number of *holes* they have – this number is the *genus*. Sphere (0 holes), torus (1 hole), and  $n$ -holed tori.

**Definition 11.** *Seifert surface of a knot  $K$*  is an orientable surface  $\Sigma$  with boundary, such that  $\partial\Sigma = K$ .

**Example 5.** The unit circle (unknot) in the plane has the *disk* as a Seifert surface. The unlink (of two components) has the *cylinder* (or *annulus*) as a Seifert surface... and the same is true for the Hopf link (the annulus has 2 twists in this case).

Seifert surfaces always exist. Look up **Seifert’s Algorithm**.

But Seifert surfaces are not unique – we can arbitrarily add holes to them! So the genus of a Seifert surface is not a well-defined number for a knot. Immediately you should question, what can we do to get something well-defined?

**Definition 12.** *Genus of a knot  $g(K)$*  is the smallest possible genus of all Seifert surfaces of the knot.

**Theorem 4.**  $g(K) = 0$  iff  $K \cong O$ .

This theorem hinges on “Seifert surface” being orientable, to talk about “genus”. That is, we have a reason to require orientability in the definition of Seifert surfaces: The theorem  $g(K) = 0$  iff  $K \cong O$ , and the notion of genus makes sense for orientable surfaces.

Now, there is a notion of genus for nonorientable surfaces. Does the theorem extend if we allow Seifert surfaces to be nonorientable? Note that the standard picture of a Mobius strip (a strip with 1 twist) is a surface which bounds the unknot, but what would be its genus? And the trefoil is knotted but its standard diagram is bounded by a Mobius strip with 3 twists.

**Exercise 4.** The trefoil has the punctured torus as a Seifert surface. *Hint: First find a diagram for it which isn't the standard one (the "standard one" is the diagram which bounds the Mobius strip with 3 twists).*

We have a very limited amount of knowledge up to now in this course, so we can "easily" take all combinations of the things we learned to ask questions, such as what happens to the genus under the connect-sum operation. Well,

**Theorem 5.**  $g(K \# J) = g(K) + g(J)$ .

The inequality  $g(K \# J) \leq g(K) + g(J)$  is obvious: the  $\#$ -operation for 1-dimensional spaces also works for  $n$ -dimensional spaces, so wherever we "cut" the knots we can also "cut" the Seifert surfaces and then glue them. NOTE: This only gives us an inequality, because the resulting Seifert surface for  $K \# J$  might not be the one having minimal possible genus (even though the surfaces for  $K$  and  $J$  had minimal genus)!

**Corollary 1.** Knots cannot be cancelled, i.e. there are no knot-inverses under  $\#$ .

Our previous proof of this result involved passing to the realm of wild knots... here we don't have to!

### 7.3 Intro to research

Just as a tip: Question everything. See where proofs of theorems hinge on all of the hypotheses. Sometimes, the hypotheses are there to make life easier, rather than being the most general statement.

Typically, "orientability" is not necessary, and this is where a lot of statements you read in the literature can be generalized to handle non-orientable spaces.

When a proof of some theorem requires you to first move outside of a given framework and then come back into the given framework, then there is probably a way to prove the theorem without doing this. The "knots cannot be cancelled" is such an example (the given framework is "our study of tame (i.e. not wild) knots").

## 8 2/9/16 – The trefoil is knotted

### 8.1 The trefoil is not the unknot

Up until now has been the intro to knot theory. From now on we're going to pursue self studies and focus our attention on different topics to get depth. We'll get to polynomials next week as they're a good branching off point.

Believe it or not, we have yet to show that any given knot is knotted! We will now show (in a couple of ways) that

**Theorem 6.** *The trefoil  $T \not\cong O$ .*

Some methods of proof: braiding, (tri)colorability, fundamental group, (Jones) polynomial.

### 8.2 Braids

**Definition 13.** A braid  $\sigma$  of  $n$  strings consists of  $n$  oriented arcs with fixed ordered endpoints. Take  $n$  unlinks passing in parallel through a region, cut all of the strings in some boxed region, and permute where they attach to on the other side (via  $\sigma$ ). That gets you some link  $\hat{\sigma}$ .

**Theorem 7** (Alexander). *Any oriented link is given by  $\hat{\sigma}$  for some  $\sigma$ .*

*Every link can be obtained by a braid. Note that this includes knots, as they are 1 links.*

**Remark 9.** What things can we do with this theorem (as researchers)? It says this map  $\sigma \mapsto \hat{\sigma}$  is surjective - is it injective? No

**Theorem 8.** *Markov There exists moves  $\sigma_1 \mapsto \sigma_2$  which don't affect  $\hat{\sigma}_1$  (the produced link)*

**Example 6.** If  $\sigma = 1_n$  (the identity permutation) then  $\hat{\sigma}$  is the  $n$  unlink. But if  $\hat{\sigma}$  is the  $n$  unlink does  $\sigma = 1_n$ ? Yes (This is a theorem of Birman-Menasco) Note that this does not exclude permutations that produce less than  $n$  unlinks. There are certainly permutations that do this, for instance a simple twist of two strands.

#### 8.2.1 The braid group

Are braids a group under concatenation? We discovered that connected sum on knots forms the structure of a monoid as knots cannot be canceled.

Yes,  $(B_n = \{\sigma_i\}), \cdot$  is in fact a group! It can be represented as  $B_n = \langle \sigma_1, \dots, \sigma_{n-1} \mid \text{relations below} \rangle$

The generators have the form  $\sigma_i = (i[i+1])$  (they swap with their neighbor). If  $|i-j| \geq 2$  then  $\sigma_i \sigma_j = \sigma_j \sigma_i$  - when generators are far apart they commute. Else  $\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}$

**Example 7.** • Show that  $(\sigma_1 \dots \sigma_{n-1})^m$  produces the  $(n, m)$ -torus link.

• Show that the trefoil is not the unknot using braid theory



## 9 2/11: Coloring

### 9.1 Logistical happenings

- Does everyone have a partner? Subjects?
- Chris away March 15 - 17. Maybe we'll just meet anyway and do class presentations

We still don't know that the trefoil is knotted!

**Example 8.** Show with braid theory that the trefoil is knotted.

### 9.2 (Tri)Colorings

Assume for this section that our knots are connected (so ignore links).

**Definition 14** (Strand). A *strand* is a line in a diagram that goes from an undercross to an undercross (whereas an *arc* is a line in the diagram that goes from a crossing to crossing).

For instance the trefoil is made up of 3 strands.

**Definition 15** (Color). A “coloring” of a knot diagram is a choice of color for each strand.

**Definition 16** (Tri-colorability). A knot diagram has a “tri-coloring” if each strand can be colored one of 3 colors such that at each crossing either all three colors or only one color are present, and all 3 colors are present in the diagram.

Is (tri)colorability invariant under R-moves? Well-defined? Yes. It's an easy exercise to show.

Does every knot have a tri-coloring? No. The unknot is not tri-colorable. But the trefoil is! Thus  $T \not\cong O$ .

**Notation 2.**  $\tau(K)$  denotes the number of colorings (not just tri-colorings!) of knot  $K$

**Question 4.** Questions that come up about colorings

- Relation between colorability and coloring Seifert surfaces?
- Relation between prime knots and colorability?
- Relation to 4-Color Theorem of maps?

**Example 9.** •  $\tau(0) = 3 = 3^1$ . The unknot has 3 colorings

- $\tau(T) = 3 + 3 \cdot 2 = 3^2$ . The trefoil has 9 colorings (hence  $T \not\cong O$  again)
- The trefoil is tri-colorable, but the figure 8 knot isn't. So now our problem is to see that the figure 8 knot is knotted... What is  $\tau$  of it?
- In general, combinatorics says  $\tau(K) \leq 3^{\text{number of strands of } K}$

Tuesday is going to be a crash course on homotopy and groups. We'll learn algebraically that knots are knotted because of their complements. That will segue into polynomials and all sorts of other fun stuff

## 10 2/16: $\pi_1$

### 10.1 Logistical things

March 15 + 17: the week before spring week, during which Chris will be absent and we should be emailing projects in. Know and/or begin final project

March 8 + 10: will be class presentations.

March 1 + 3: email Chris progress/thoughts 'report'

### 10.2 Crash course on the fundamental group

There are tons of different kinds of mappings we could define from the set of knots to the integers, and to other sets such as integer polynomials. We have seen enough so far to believe that knots have a very rich structure - perhaps we'd also believe that the integers don't really have enough structure for mappings into them to be very useful, as far as revealing the structure of knots goes.

This motivates studying maps from knots into the polynomials, which obviously have much more structure than the integers.

Using the fundamental group for knots is motivated by:  $K \cong J \iff \mathbb{R}^3 \setminus J \approx \mathbb{R}^3 \setminus K$ .

1.  $\pi_1$  takes n-dimensional 'spaces' and gives groups
2.  $\pi_1$  is invariant under homeomorphism
3. For any knot  $K$ ,  $\pi_1(K) \cong \pi_1(S^1) \cong \mathbb{Z}$ , which is why we look at the knot complement
4.  $\pi_1(\mathbb{R}^3 - K)$  is called the **knot group**

**Definition 17.** For some 'space'  $\Sigma$ ,  $\pi(\Sigma) = \{ \text{loops (up to reparametrization)} \gamma(t) : S^1 \rightarrow \Sigma \text{ with respect to some basepoint} \}$ . The group operation is composition of loops  $\sigma_1 \cdot \sigma_2 : S^1 \rightarrow \Sigma$ . The inverse  $\bar{\gamma}$  is just the time reversed version of  $\gamma$ . The zero element is the constant map  $S^1 \rightarrow * \subseteq \Sigma$  (where  $*$  is the basepoint).

So far we're working up to reparametrization equivalence - but with that notion  $\gamma \cdot \bar{\gamma}$  goes around  $\gamma$  once in the forwards direction and once in the reverse direction, it most certainly does not say this is a constant map as claimed. We want to work up to homotopic equivalence

Crash course in homotopy: a path  $\gamma$  is a map  $[0, 1] \rightarrow \Sigma$ , a loop is a path with  $\gamma(0) = \gamma(1)$ . We say that  $\gamma_1$  is homotopically equivalent to  $\gamma_2$  if there exists a continuous family  $F_\tau$  (with  $\tau \in [0, 1]$ ) of "deformations" of the paths such that  $F_0 = \gamma_1$  and  $F_1 = \gamma_2$ . This is written  $\gamma_1 \simeq \gamma_2$

Do you see how the fundamental group captures the holes in the 'space'?

**Example 10.** •  $\pi_1(\mathbb{R}^n) = 0$  (the trivial group) all loops contract to basepoint

- $\pi_1(\text{annulus}) = \mathbb{Z}$ .  $\gamma_1 \not\simeq *$  is the loop that circles the whole once,  $\gamma^2 \not\simeq *$  winds twice (and is the same up to reparametrization because we require that reparametrizations are bijections  $S^1 \rightarrow S^1$ )
- $\pi_1(S^1) = \mathbb{Z}$ . The annulus is nothing more than a thickened circle.

- What about  $\pi_1(T^2)$ ? Is it  $\mathbb{Z} \oplus \mathbb{Z}$ ? If  $\alpha$  is the loop around one of the holes and  $\beta$  is the loops around the other hole, does  $\alpha\beta = \beta\alpha$ ? It turns out that  $\pi_1(T^2) = \mathbb{Z} \oplus \mathbb{Z}$ . But  $\pi_1(\text{double torus})$  is nonabelian!
- $\pi_1(\mathbb{R}^3 - O) = \mathbb{Z} = \langle g_1 | \emptyset \rangle$  the loops link around the unknot with any linking number

**Remark 10.** Computing  $\pi_1$  is hard in general.

## 11 2/18: Wirtinger presentation

We showed that the knot group of the unknot is the integers,  $\pi_1(\mathbb{R}^3 - O) \cong \mathbb{Z}$ . Thus to show a knot is knotted, it suffices to show its knot group is not the integers. *Exercise: If the knot group surjects onto a group  $G$  such that  $G \not\cong \mathbb{Z}_n$  nor  $\mathbb{Z}$ , then the knot is knotted.*

$\pi_1(\mathbb{R}^3 - O) = \mathbb{Z}$  in some sense tells us that nontrivial groups arise from knots due to crossings.

Knot group  $\pi_1(\mathbb{R}^3 - K)$  can be computed using the *Wirtinger presentation* for any knot diagram of  $K$ . Given a diagram with  $n \geq 1$  crossings, the presentation has  $n$  generators (plus specific relations). These generators correspond to loops that wrap around arcs of the diagram, and the relations correspond to the “obstruction” to homotoping the loops past the crossings.

### Algorithm:

1. Orient the knot diagram.
2. Label the 3 arcs at each crossing (1 arc for the overcrossing and 2 arcs for the undercrossing),  $g_i$  and  $g_j$  and  $g_k$ . Here  $g_j$  represents the loop (with respect to same fixed basepoint in  $\mathbb{R}^3 - K$ ) that encircles the  $j^{\text{th}}$  arc with linking number  $+1$ .
3. Stare hard enough to see the relations: Assuming  $g_k$  is the overcrossing-arc and the crossing is positive,  $g_j = g_k g_i g_k^{-1}$ . Basically, conjugation by the overcrossing-arc allows you to “pass a loop through the crossing.”

**Example 11.** *The trefoil is knotted:*  $\pi_1(\mathbb{R}^3 - T) \cong B_3$  the Braid group on 3 strands. Chris’ proof in class doesn’t establish this group-isomorphism, but he constructs a surjective homomorphism  $\pi_1(\mathbb{R}^3 - T) \twoheadrightarrow S_3$  onto the symmetric group of 3 letters (and since we know there is no surjection  $\mathbb{Z} \rightarrow S_3$ , we’re done). Chris’ map is determined by the 3 generators of the Wirtinger presentation:  $g_1 \mapsto (1\ 2)$  and  $g_2 \mapsto (2\ 3)$  and  $g_3 \mapsto (3\ 1)$

*Exercise:* The figure 8 knot is not equivalent to the trefoil.

*Exercise:* The Hopf link is linked.

**Remark 11.** Note that the Wirtinger presentation is not usually optimal in terms of number of generators. The unknot can have  $n$  twists in it, but  $\mathbb{Z}$  has 1 generator.

If  $\pi_1(\mathbb{R}^3 - K) \cong \pi_1(\mathbb{R}^3 - J)$  then does  $K \cong J$ ? *No, look at the mirror of any knot!*

## 12 2/23: POLYNOMIALS

$\mathcal{K} := \{\text{knots}\}$ . Our invariants so far took the form  $\mathcal{K} \rightarrow \mathbb{Z}$  (such as the genus and minimal crossing number), or  $\mathcal{K} \rightarrow \{\text{groups}\}$  via  $\pi_1$  of the knot complement. We will now expand to polynomial rings, which should be more sensitive to the richness of  $\mathcal{K}$  because  $\mathbb{Z} \subset \mathbb{Z}[x]$ .