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**Topic : Huffman Coding**

**Algorithm:**

```
1 text = "message" [T.C: O(1)]
2 freq = [] [T.C: O(1)]
3 struct Node [T.C: O(1)]
4   3.1 char data
5   3.2 unsigned freq
6   3.3 Node *left, *right
7   3.4 Node(data, freq)
8     3.4.1 this->data = data
9     3.4.2 this->freq = freq
10    3.4.3 left = right = NULL
11
12 4 calcFrequencies(string text) [T.C: O(n)]
13    4.1 for char in text O(n)
14      4.2 freq[char]++
15
16 5 buildHuffmanTree(unordered_map(char, int) freq) [T.C: O(n+mlogm)]
17    5.1 priority_queue(Node) minHeap = [] O(1)
18    5.2 for char,freqValue in freq O(m)
19      5.3   minHeap.push(new Node(char, freqValue)) O(log(m))
20    5.4 while minHeap.size > 1 O(m-1)
21      5.5   Node left = minHeap.top()
22      5.6   minHeap.pop() O(log(m))
23      5.7   Node right = minHeap.top()
24      5.8   minHeap.pop() O(log(m))
25      5.9   Node internal = new Node('$',left.freq + right.freq)
26      5.10  internal.left = left
27      5.11  internal.right = right
28      5.12  minHeap.push(internal) O(log(m))
29
30 5.13 return minHeap.top()
31
32 6 storeCodes(
33   Node rootNode,
34   string code,
35   unordered_map(char, int) codes
36 )
37 [T.C: O(m*L)]
38
39 6.1 if rootNode == NULL
40 6.2   return
41 6.3 if rootNode.left == NULL and rootNode.right == NULL
42 6.4   codes[rootNode.data] = code
43 6.5 storeCodes(rootNode.left, code+"0", codes) O(L)
44 6.6 storeCodes(rootNode.right, code+"1", codes) O(L)
```

## Time Complexity:

### **IMPORTANT CLARIFICATION**

- A Balanced tree can have  $m \neq n$  OR  $m = n$ 
  - 'abcdefg' → balanced tree w/  $m = n$  (special case, all unique)
  - 'aaaabbbbcccc' → balanced tree with  $m \neq n$
- A Skewed tree can have  $m \neq n$ 
  - 'bbbbbcda' → generates skewed tree with  $m \neq n$
- It is not possible to get a Skewed tree with  $m = n$

#### 1. buildHuffmanTree()

A temporary data structure 'minHeap' used to create the Huffman Tree.

##### Time Complexity:

$$O(m\log(m)) + O((m-1)*3\log(m)) = O(m\log(m))$$

where  $m$  = no. of unique characters

- Repeated chars ( $m \neq n$ ) =  $O(m\log(m))$ 
  - Same for skewed and balanced tree
- When all characters unique ( $m = n$ ) =  $O(n\log(n))$

#### 2. storeCodes()

- Total nodes visited :  $2m - 1$       ' $m$ ' original + ' $m-1$ ' internal
- At each node :  $O(L)$       String concatenation
- $m \neq n$  Case  $\rightarrow O((2m-1)*L) : O(m*L)$        $L$  = Huffman Tree Height
  - Balanced Tree  $\rightarrow O(m\log(m))$
  - Skewed Tree  $\rightarrow O(m^2)$
- $m = n$  Case  $\rightarrow O((2n-1)*L) : O(n*L)$

Where,

Huffman Tree Height ( $L$ ) is determined by  
no. of unique characters ( $m$ )

NOTE: calcFrequency() is  $O(n)$  and not dependent on ' $m$ ' or ' $L$ '

So, including calcFrequency() the total time complexity will be:

$$\text{Total Time Complexity} = O(n) + O(m\log(m)) + O(m*L)$$

- Balanced Tree (Typical/Expected Structure):  
 $O(n) + O(m\log(m)) + O(m\log(m))$

##### Total Time Complexity

$$\text{Repeated chars } (m \neq n) = O(n + m\log m)$$

$$\text{All chars unique } (m = n) = O(n\log n) \quad (\text{special case})$$

- Skewed Tree (Pathological Structure):

$$O(n) + O(m\log(m)) + O(m^2)$$

##### Only one case

$$\text{Repeated chars} = O(m^2)$$

## Source Code:

```
#include <iostream>
#include <queue>
#include <unordered_map>
#include <vector>
#include <string>
using namespace std;

struct Node {
    char data;
    unsigned freq;
    Node *left, *right;

    Node(char data, unsigned freq) {
        this->data = data;
        this->freq = freq;
        left = right = nullptr;
    }
};

struct Compare {
    bool operator()(Node* l, Node* r) {
        return l->freq > r->freq;
    }
};

// Non-Recursive Function
Node* buildHuffmanTree(unordered_map<char, int>& freq) {
    // Define minHeap
    priority_queue<Node*, vector<Node*>, Compare> minHeap;

    // Populate minHeap
    for (auto pair: freq) {
        minHeap.push(new Node(pair.first, pair.second));
    }
    // Build Huffman Tree
    while (minHeap.size() > 1){
        // Select the two elements from the minHeap
        // having least frequencies
        Node* left = minHeap.top();
        minHeap.pop();
        Node* right = minHeap.top();
        minHeap.pop();

        // Generate non-leaf node
        // Generated node character -> '$'
        Node* internal = new Node('$', left->freq + right->freq);
        internal->left = left;
        internal->right = right;

        // Push generated node to tree
        minHeap.push(internal);
    }
    return minHeap.top();
}

// Recursive Function
```

```

// Tree is traversed top-to-bottom
void storeCodes(Node* root, string code, unordered_map<char, string>& codes) {
    if (!root) return;
    if (!root->left && !root->right) {
        codes[root->data] = code;
        return;
    } // If at leaf node, character code is complete
    storeCodes(root->left, code+"0", codes);           // If left edge, append 0
    storeCodes(root->right, code+"1", codes);          // If right edge, append 1
}

int main() {
    // 1. Get text input
    string text;
    cout << "Enter text to encode: ";
    getline(cin, text);

    // 2. Calculate frequencies
    cout << endl << "Calculating frequencies..." << endl;
    unordered_map<char, int> freq;
    for (char c: text) {
        freq[c]++;
    }
    // 3. Build the tree
    cout << "Building Huffman tree..." << endl;
    Node* root = buildHuffmanTree(freq);

    // 4. Generate the huffman codes
    cout << "Generating Huffman encoding table..." << endl;
    unordered_map<char, string> codes;
    storeCodes(root, "", codes);
    cout << endl;

    // 4. Print out encoding table
    cout << "Huffman Encoding Table" << endl;
    for (auto pair: codes) {
        cout << pair.first << ": " << pair.second << endl;
    }
    cout << endl;

    // 5. Print out compression results
    int originalBits = 8*text.length();
    int compressedBits = 0;
    string compressedMessage = "";

    for (char c: text) {
        compressedBits += codes[c].length();
        compressedMessage += codes[c];
    } // Calculating new length of compressed message

    cout << "Compressed Message: " << compressedMessage << endl << endl;
    cout << "No. of bits required to store original message: " << originalBits << endl;
    cout << "No. of bits required to store compressed version: " << compressedBits <<
endl;
    cout << "Size reduction: " << (originalBits-compressedBits)*100/originalBits << "%"
<< endl;

    return 0;
}

```

## Sample Output:

```
rug@Oxide [Algorithm-Code]>> ./a.out
Enter text to encode: Hello World

Calculating frequencies...
Building Huffman tree...
Generating Huffman encoding table...

Huffman Encoding Table
l: 10
: 011
W: 1111
d: 010
o: 110
H: 001
e: 1110
r: 000

Compressed Message: 00111101010110011111110000100010

No. of bits required to store original message: 88
No. of bits required to store compressed version: 32
Size reduction: 63%
rug@Oxide [Algorithm-Code]>> ./a.out
Enter text to encode: She sells sea shells by the sea shore

Calculating frequencies...
Building Huffman tree...
Generating Huffman encoding table...

Huffman Encoding Table
t: 10110
: 111
a: 1010
e: 110
r: 01111
l: 100
y: 01110
b: 10111
o: 01101
s: 01100
n: 010
h: 00

Compressed Message: 01100010110110011010010000111001101011000101101001000011101101110111011001011011001101011000100011010111110

No. of bits required to store original message: 296
No. of bits required to store compressed version: 118
Size reduction: 60%
```