

Q1 What is the difference between *basic* regular expressions and *extended* regular expressions?

Basic regular expressions are those defined by the grammar of $c, 1, 0, +, \dots$ - they present the formal concept of regular expressions.

Extended regular expressions are those found in most modern program languages, with extended features for matching patterns in strings, such as $a\{n\}$, $a+$, etc. Some of these features are simply syntactic sugar over features found in basic regular expressions. Extended expression can build up on basic expression as an expression.

Q2 What is the language recognised by the regular expressions $(0^*)^*$.

$$L((0^*)^*) = \{\epsilon\}$$

Remember $*$ always introduces the empty string

$$L(0) = \{\epsilon\}, \text{ but because of } * \text{ we get } L(0^{**}) = \{\epsilon\}$$

Q3 Review the first handout about sets of strings and read the second handout. Assuming the alphabet is the set $\{a, b\}$, decide which of the following equations are true in general for arbitrary languages A, B and C :

$$(A \cup B) @ C \stackrel{?}{=} A @ C \cup B @ C$$

$$\text{equal} \rightarrow (A+B)C = A \cdot C + B \cdot C$$

$$A^* \cup B^* \stackrel{?}{=} (A \cup B)^*$$

$$\text{not equal} \rightarrow \text{left regex doesn't support ABABABABA}$$

$$A^* @ A^* \stackrel{?}{=} A^*$$

$$\text{equal} \rightarrow A^* @ A^* = A^*(^*+^*) = A^*$$

$$(A \cap B) @ C \stackrel{?}{=} (A @ C) \cap (B @ C)$$

$$\text{not equal, try } A=\{[a]\}, B=\{\epsilon\} \text{ and } C=\{[a], \epsilon\}$$

Q4 Given the regular expressions $r_1 = 1$ and $r_2 = 0$ and $r_3 = a$. How many strings can the regular expressions r_1^* , r_2^* and r_3^* each match?

r_1^* : empty string

r_2^* : empty string

r_3^* : infinite

Q5 Give regular expressions for (a) decimal numbers and for (b) binary numbers. Hint: Observe that the empty string is not a number. Also observe that leading 0s are normally not written—for example the JSON format for numbers explicitly forbids this. So 007 is not a number according to JSON.

Decimal: $0 + (1+2+3+4+5+6+7+8+9) \cdot (1+2+3+4+5+6+7+8+9+0)^*$

Binary: $0 + 1 \cdot (1+0)^*$

Q6 Decide whether the following two regular expressions are equivalent $(1+a)^* \equiv^? a^*$ and $(a \cdot b)^* \cdot a \equiv^? a \cdot (b \cdot a)^*$.

$(1+a)^* \equiv^? a^*$ Yes, they are equal. Both match the empty string.

$(a \cdot b)^* \cdot a \equiv^? a \cdot (b \cdot a)^*$ Yes, they are equal.

Q7 Given the regular expression $r = (a \cdot b + b)^*$. Compute what the derivative of r is with respect to a , b and c . Is r nullable?

r is nullable

$\text{der } a(r) = (b) \cdot (a \cdot b + b)^*$; $\text{der } b(r) = (a \cdot b + b)^*$; $\text{der } c(r) = 0$

Q8 Give an argument for why the following holds: if r is nullable then $r^{\{n\}} \equiv r^{\{..n\}}$.

8/

$$[] \in L(r)$$

$$L(r^{\{n\}}) = \bigcup_{0 \leq n} L(r)^n$$

$$L(r^{\{..n\}}) = L(r)^0 \cup L(r)^1 \cup \dots \cup L(r)^n$$

if $[] \in A$, then $A^n \subset A^{n+m}$ for every $0 \leq m$

Eg. $A = \{a, b, c, []\}$ from HW1

Therefore, $L(r)^0 \subset L(r)^n, \dots, L(r)^{n-1} \subset L(r)^n$

$$\begin{aligned} L(r^{\{..n\}}) &= L(r)^0 \cup L(r)^1 \cup \dots \cup L(r)^n \\ &= L(r)^n \\ &= L(r^{\{n\}}) \end{aligned}$$

$$\therefore r^{\{n\}} \equiv r^{\{..n\}}$$

$$A = \{a, b, []\}$$

$$A^2 = 2^0 + 2^1 + 2^2$$

$$= 1 + 2 + 4$$

$$= 7$$

$$\{a, b, []\}$$

$$= \{a, b, aa, ab, ba, bb, []\}$$

Q9 Define what is meant by the derivative of a regular expression with respect to a character.
(Hint: The derivative is defined recursively.)

$$\begin{aligned}
 \text{der } c(0) &\stackrel{\text{def}}{=} 0 \\
 \text{der } c(1) &\stackrel{\text{def}}{=} 0 \\
 \text{der } c(d) &\stackrel{\text{def}}{=} \text{if } c = d \text{ then } 1 \text{ else } 0 \\
 \text{der } c(r_1 + r_2) &\stackrel{\text{def}}{=} \text{der } c r_1 + \text{der } c r_2 \\
 \text{der } c(r_1 \cdot r_2) &\stackrel{\text{def}}{=} \text{if } \text{nullable}(r_1) \\
 &\quad \text{then } (\text{der } c r_1) \cdot r_2 + \text{der } c r_2 \\
 &\quad \text{else } (\text{der } c r_1) \cdot r_2 \\
 \text{der } c(r^*) &\stackrel{\text{def}}{=} (\text{der } c r) \cdot (r^*)
 \end{aligned}$$

Q10 Assume the set *Der* is defined as $\text{Der } c A = \{s \mid c :: s \in A\}$. What is the relation between *Der* and the notion of derivative of regular expressions?

$\text{Der } c A$ is for a set of world, and $\text{der } c(a)$ is for the regular expressions

$$L(\text{der } c r) = \text{Der } c(L(r))$$

Q11 Give a regular expression over the alphabet $\{a, b\}$ recognising all strings that do not contain any substring *bb* and end in *a*.

$$r = (a^*ba)^* \cdot a + (a^*ab)^* \cdot a$$

Q12 Do $(a+b)^+ \cdot b^+$ and $(a^+ \cdot b^+) + (b^+ \cdot b^+)$ define the same language?

No. For example, abbab

Q13 Define the function *zeroable* by recursion over regular expressions.

$$\begin{aligned}
 \text{zeroable}(0) &\stackrel{\text{def}}{=} \text{true} \\
 \text{zeroable}(1) &\stackrel{\text{def}}{=} \text{false} \\
 \text{zeroable}(c) &\stackrel{\text{def}}{=} \text{false} \\
 \text{zeroable}(r_1 + r_2) &\stackrel{\text{def}}{=} \text{zeroable}(r_1) \wedge \text{zeroable}(r_2) \\
 \text{zeroable}(r_1 \cdot r_2) &\stackrel{\text{def}}{=} \text{zeroable}(r_1) \vee \text{zeroable}(r_2) \\
 \text{zeroable}(r^*) &\stackrel{\text{def}}{=} \text{true}
 \end{aligned}$$

Q14 Give a regular expression that can recognise all strings from the language $\{a^n \mid \exists k. n=3k+1\}$.

$$r = a \cdot (a^3)^*$$

Q15 Give a regular expression that can recognise an odd number of *as* or an even number of *bs*.

$$r = a \cdot (aa)^* + bb \cdot (bb)^*$$