

# Semantic Scaffolds for Pseudocode-to-Code Generation



Ruiqi Zhong, Mitchell Stern, Dan Klein



## Input-output Test Cases



# Task

Generate

## # Pseudocode

```
1 declare constant integer numOfAlphabets = 26
2
3 create string s
4 let Count = 0 be an integer
5 read s
6 Set Ch to be a
7 for i = 0 to i less than the length of s
8     set Count to Count + min(abs(s[i] - Ch),
9                             numOfAlphabets - abs(s[i] - Ch))
9     set s[i] to Ch
10
11 print Count
12
13
```

## Code

```
const int numOfAlphabets = 26;
int main() {
    string s;
    int Count = 0;
    cin >> s;
    char Ch = 'a';
    for (int i = 0; i < s.length(); i++) {
        Count += min(abs(s[i] - Ch),
                    numOfAlphabets - abs(s[i] - Ch));
        Ch = s[i];
    }
    cout << Count << endl;
    return 0;
}
```

Code: solutions to the coding challenge.  
Pseudocode: human annotated instructions based on the code.



# Language Ambiguities

## # Pseudocode

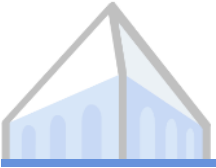
```
1 declare constant integer numOfAlphabets = 26
2
3 create string s
4 let Count = 0 be an integer
5 read s
6 Set Ch to be a
7 for i = 0 to i less than the length of s
8     set Count to Count + min(abs(s[i] - Ch),
9                             numOfAlphabets - abs(s[i] - Ch))
9     set s[i] to Ch
10
11 print Count
12
13
```

Several  
Possibilities

```
char Ch = 'a';
char Ch = a;
Ch = 'a';
Ch = a;
```

## Code

```
const int numOfAlphabets = 26;
int main() {
    string s;
    int Count = 0;
    cin >> s;
    char Ch = 'a';
    for (int i = 0; i < s.length(); i++) {
        Count += min(abs(s[i] - Ch),
                    numOfAlphabets - abs(s[i] - Ch));
        Ch = s[i];
    }
    cout << Count << endl;
    return 0;
}
```



# Evaluation

## # Pseudocode

```
1 declare constant integer numOfAlphabets = 26
2
3 create string s
4 let Count = 0 be an integer
5 read s
6 Set Ch to be a
7 for i = 0 to i less than the length of s
8     set Count to Count + min(abs(s[i] - Ch),
9                             numOfAlphabets - abs(s[i] - Ch))
9     set s[i] to Ch
10
11 print Count
12
13
```

## Top-k Code Candidate

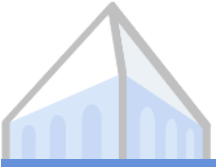
```
const int numOfAlphabets = 26;
int main() {
    string s;
    int Count = 0;
    cin >> s;
    char Ch = 'a';
    for (int i = 0; i < s.length(); i++) {
        Count += min(abs(s[i] - Ch),
                    numOfAlphabets - abs(s[i] - Ch));
        Ch = s[i];
    }
    cout << Count << endl;
    return 0;
}
```

Evaluation: execute the top-k (k=1, 10, 100, 1000) program candidates on the inputs and check outputs. Inputs/outputs come from the online judges.

Input: “zeus” -> Output: 18

Input: “map” -> Output: 35

Input: “ares” -> Output: 34



# Task Properties

---

## ▶ Length

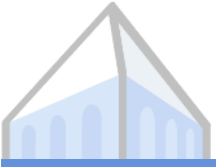
- ▶ 15 lines per-program on average, 457 lines max.
- ▶ Need to generate line-by-line (instead of learning end-to-end).

## ▶ Semantic Evaluation

- ▶ Execution correctness (instead of BLEU score/other surface form metrics).
- ▶ “`i += 1;`” is equivalent to “`i = i + 1;`”

## ▶ Search

- ▶ Top-1 solution is not enough.
- ▶ Quality of top-1000 solution might matter.
- ▶ We care about efficiency.



# Task Properties

---

## ▶ Length

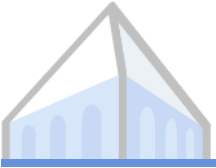
- ▶ 15 lines per-program on average, 457 lines max.
- ▶ Need to generate line-by-line (instead of learning end-to-end).

## ▶ Semantic Evaluation

- ▶ Execution correctness (instead of BLEU score/other surface form metrics).
- ▶ “`i += 1;`” is equivalent to “`i = i + 1;`”

## ▶ Search

- ▶ Top-1 solution is not enough.
- ▶ Quality of top-1000 solution might matter.
- ▶ We care about efficiency.



# Task Properties

---

## ▶ Length

- ▶ 15 lines per-program on average, 457 lines max.
- ▶ Need to generate line-by-line (instead of learning end-to-end).

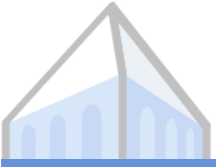
## ▶ Semantic Evaluation

- ▶ Execution correctness (instead of BLEU score/other surface form metrics).
- ▶ “`i += 1;`” is equivalent to “`i = i + 1;`”

## ▶ Search

- ▶ Top-1 solution is not enough.
- ▶ Quality of top-1000 solution might matter.
- ▶ We care about efficiency.





# Task Properties

---

## ▶ Length

- ▶ 15 lines per-program on average, 457 lines max.
- ▶ Need to generate line-by-line (instead of learning end-to-end).

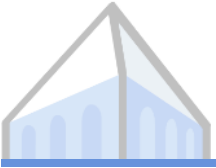
## ▶ Semantic Evaluation

- ▶ Execution correctness (instead of BLEU score/other surface form metrics).
- ▶ “`i += 1;`” is equivalent to “`i = i + 1;`”

## ▶ Search

- ▶ Top-1 solution is not enough.
- ▶ Quality of top-1000 solution might matter.
- ▶ We care about efficiency.

We propose a method to **efficiently** generate **1000** candidates for **long** programs and search for a **semantically correct** solution.



# Baseline

## ▶ Notation.

- ▶  $L$ : # of lines.
- ▶  $l$ : line index.
- ▶  $x_l$ : the pseudocode input at line  $l$ .

## $l$ Pseudocode $x_l$

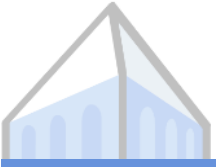
```
3 create string s
4 let Count = 0 be an integer
5 read s
X6 6 Set Ch to be a
```

## ▶ Translate each line independently.

- ▶ Generate 100 *code fragment* candidates  $y_{lc}$  for each  $x_l$ .
- ▶  $y_{lc}$  is assigned a probability score  $p_{lc}$ .

## ▶ Generate full program $y$ .

- ▶ For each line  $l$  we select one code fragment  $c_l$ , then concatenate the code fragments.
- ▶ We score a full program by taking the independent scoring for each line.



# Baseline

## ► Notation.

- $L$ : # of lines.
- $l$ : line index.
- $x_l$ : the pseudocode input at line  $l$ .

## ► Translate each line independently.

- Generate 100 *code fragment* candidates  $y_{lc}$  for each  $x_l$ .
- $y_{lc}$  is assigned a probability score  $p_{lc}$ .

## ► Generate full program $y$ .

- For each line  $l$  we select one code fragment  $c_l$ , then concatenate the code fragments.
- We score a full program by taking the independent scoring for each line.

### $l$ Pseudocode $x_l$

```
3 create string s
4 let Count = 0 be an integer
5 read s
```

**X6** 6 Set Ch to be a

$p_{61} = 0.70$ ,  $y_{61}$  `char Ch = 'a';`

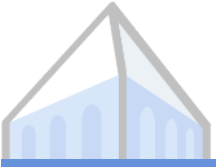
$p_{62} = 0.15$ ,  $y_{62}$  `char Ch = a;`

$p_{63} = 0.05$ ,  $y_{63}$  `Ch = 'a';`

$p_{64} = 0.02$ ,  $y_{64}$  `Ch = a;`

• • •

Others omitted,  
altogether 100 of candidates.



# Baseline

## ► Notation.

- $L$ : # of lines.
- $l$ : line index.
- $x_l$ : the pseudocode input at line  $l$ .

## ► Translate each line independently.

- Generate 100 *code fragment* candidates  $y_{lc}$  for each  $x_l$ .
- $y_{lc}$  is assigned a probability score  $p_{lc}$ .

## ► Generate **full program $y$** .

- For each line  $l$  we select one code fragment  $c_l$ , then concatenate the code fragments.
- We score a full program by taking the independent scoring for each line.

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

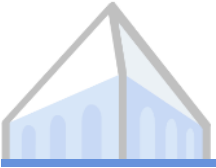
### $l$ Pseudocode $x_l$

```
3 create string s
4 let Count = 0 be an integer
5 read s
```

**X6** 6 Set Ch to be a

$p_{61} = 0.70$ ,  $y_{61}$  `char Ch = 'a';`  
 $p_{62} = 0.15$ ,  $y_{62}$  `char Ch = a;`  
 $p_{63} = 0.05$ ,  $y_{63}$  `Ch = 'a';`  
 $p_{64} = 0.02$ ,  $y_{64}$  `Ch = a;`

• • •  
Others omitted,  
altogether 100 of candidates.

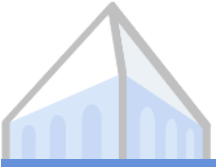


# An Abstract View

Select one column (fragment) for each row (line) to form a full program.

Line number $l$	Best Candidate	2nd Candidate	3rd Candidate	4th Candidate	[Other]
1	$y_{1,1}$	$y_{1,2}$	$y_{1,3}$	$y_{1,4}$	...
2	$y_{2,1}$	$y_{2,2}$	$y_{2,3}$	$y_{2,4}$	...
3	$y_{3,1}$	$y_{3,2}$	$y_{3,3}$	$y_{3,4}$	...
4	$y_{4,1}$	$y_{4,2}$	$y_{4,3}$	$y_{4,4}$	...
5	$y_{5,1}$	$y_{5,2}$	$y_{5,3}$	$y_{5,4}$	...
6	$y_{6,1}$	$y_{6,2}$	$y_{6,3}$	$y_{6,4}$	...
7	$y_{7,1}$	$y_{7,2}$	$y_{7,3}$	$y_{7,4}$	...
8	$y_{8,1}$	$y_{8,2}$	$y_{8,3}$	$y_{8,4}$	...
9	$y_{9,1}$	$y_{9,2}$	$y_{9,3}$	$y_{9,4}$	...
10	$y_{10,1}$	$y_{10,2}$	$y_{10,3}$	$y_{10,4}$	...
...	...	...	...	...	...

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$



# Baseline Top-1 Combination

Exact top-K solutions can be generated in  $O(KL \log K)$  time.

Line number $l$	Best Candidate	2nd Candidate	3rd Candidate	4th Candidate	[Other]
1	$y_{1,1}$	$y_{1,2}$	$y_{1,3}$	$y_{1,4}$	...
2	$y_{2,1}$	$y_{2,2}$	$y_{2,3}$	$y_{2,4}$	...
3	$y_{3,1}$	$y_{3,2}$	$y_{3,3}$	$y_{3,4}$	...
4	$y_{4,1}$	$y_{4,2}$	$y_{4,3}$	$y_{4,4}$	...
5	$y_{5,1}$	$y_{5,2}$	$y_{5,3}$	$y_{5,4}$	...
6	$y_{6,1}$	$y_{6,2}$	$y_{6,3}$	$y_{6,4}$	...
7	$y_{7,1}$	$y_{7,2}$	$y_{7,3}$	$y_{7,4}$	...
8	$y_{8,1}$	$y_{8,2}$	$y_{8,3}$	$y_{8,4}$	...
9	$y_{9,1}$	$y_{9,2}$	$y_{9,3}$	$y_{9,4}$	...
10	$y_{10,1}$	$y_{10,2}$	$y_{10,3}$	$y_{10,4}$	...
...	...	...	...	...	...

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$



# Performance

---

Method	Top 1	Top 10	Top 100	Top 1000
Baseline	0.0%	8.1%	29.2%	44.3%



# Baseline Ignores Dependencies

## ▶ Syntactic dependency and constraints between lines.

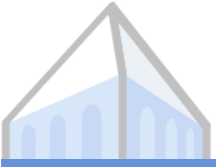
- ▶ Result code needs to be a legal derivation of the Abstract Syntax Tree.
- ▶ Pseudocode does not specify this detail whether to delay scope opening “ { ” to the next line.
- ▶ “Enumerate index i of the array s” can be translated to:
  - ▶ `for (int i = 0; i < s.length(); i++) {`
  - ▶ `for (int i = 0; i < s.length(); i++)`

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

## ▶ Semantic dependency and constraints between lines.

- ▶ Cannot use an undeclared variable or redeclare a declared variable in the same scope.
- ▶ Pseudocode does not explicitly state variable usage and declaration.
- ▶ “Set Ch to be a”
  - ▶ `char Ch = ‘a’;`
  - ▶ `char Ch = a;`
  - ▶ `Ch = ‘a’;`





# Baseline Ignores Dependencies

## ▶ Syntactic dependency and constraints between lines.

- ▶ The full program needs to be a **legal derivation** of the Abstract Syntax Tree.
- ▶ Pseudocode does not specify whether to delay scope opening “{” to the next line.
- ▶ “Enumerate index i of the array s” can be translated to:

- ▶ `for (int i = 0; i < s.length(); i++) {`
- ▶ `for (int i = 0; i < s.length(); i++)`

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

## ▶ Semantic dependency and constraints between lines.

- ▶ Cannot use an undeclared variable or redeclare a declared variable in the same scope.
- ▶ Pseudocode does not explicitly state variable usage and declaration.
- ▶ “Set Ch to be a”
  - ▶ `char Ch = ‘a’;`
  - ▶ `char Ch = a;`
  - ▶ `Ch = ‘a’;`



# Baseline Ignores Dependencies

## ▶ Syntactic dependency and constraints between lines.

- ▶ The full program needs to be a legal derivation of AST.
- ▶ Pseudocode does not specify whether to delay scope opening “ { ” to the next line.
- ▶ “Enumerate index i of the array s” can be translated to:

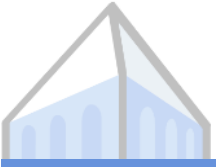
- ▶ **for** (int i = 0; i < s.length(); i++) {

- ▶ **for** (int i = 0; i < s.length(); i++)

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

## ▶ Semantic dependency and constraints between lines.

- ▶ Cannot use an undeclared variable or redeclare a declared variable in the same scope.
- ▶ Pseudocode does not explicitly state variable usage and declaration.
- ▶ “Set Ch to be a”
  - ▶ char Ch = ‘a’;
  - ▶ char Ch = a;
  - ▶ Ch = ‘a’;



# Syntactic Constraints

**Fragments**

Parse each fragment



```
int main () {  
    int n, ans = 1;  
    for (int i = 1; i <= n / 2; i++) {  
        . . .  
    }
```

**“High level symbols”**

```
[return type] [function name] ( ) {  
    [statement] ;  
    for [for conditions] {  
        . . .  
    }
```



# Syntactic Constraints

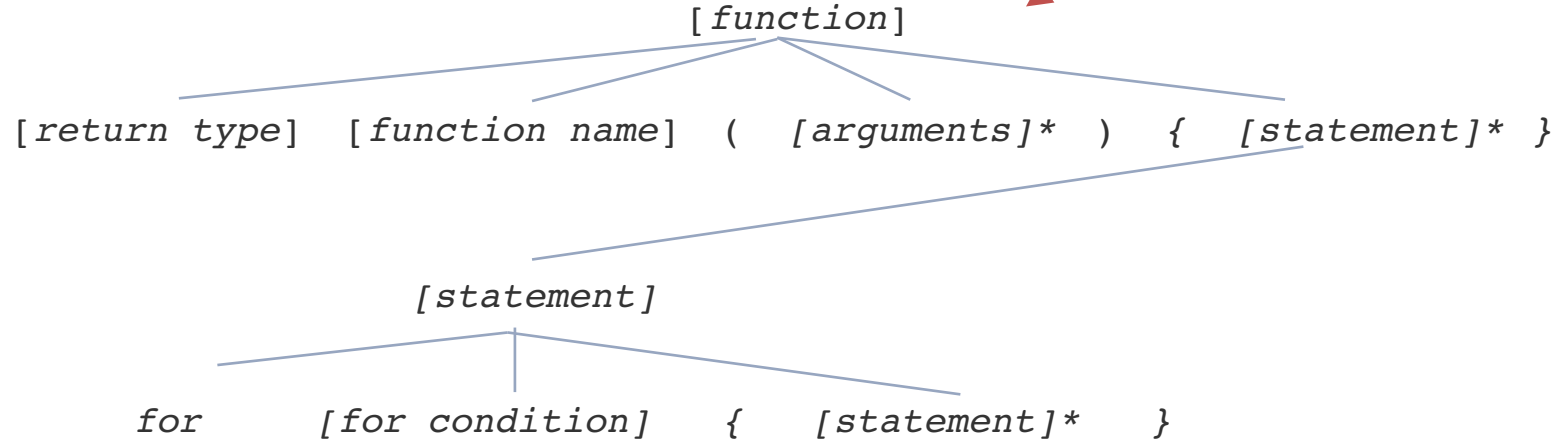
## Fragments

```
int main () {  
    int n, ans = 1;  
    for (int i = 1; i <= n / 2; i++) {  
        . . .  
    }
```

Parse each fragment

## “High level symbols”

```
[return type] [function name] ( ) {  
    [statement] ;  
    for [for conditions] {  
        . . .  
    }
```



We use an incremental parser to check whether the high level symbol combination is a possible derivation of the AST grammar.



# Baseline Ignores Dependencies

## ▶ Syntactic dependency and constraints between lines.

- ▶ The full program needs to be a legal derivation of AST.
- ▶ Pseudocode does not specify whether to delay scope opening “ { ” to the next line.
- ▶ “Enumerate index i of the array s” can be translated to:

- ▶ `for (int i = 0; i < s.length(); i++) {`

- ▶ `for (int i = 0; i < s.length(); i++)`

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

## ▶ Semantic dependency and constraints between lines.

- ▶ Cannot **use an undeclared variable** or **redeclare a declared variable** in the same scope.
- ▶ Pseudocode does not explicitly state variable usage and declaration.
- ▶ “Set Ch to be a”
  - ▶ `char Ch = ‘a’;`
  - ▶ `char Ch = a;`
  - ▶ `Ch = ‘a’;`



# Baseline Ignores Dependencies

## ▶ Syntactic dependency and constraints between lines.

- ▶ The full program needs to be a legal derivation of AST.
- ▶ Pseudocode does not specify whether to delay scope opening “ { ” to the next line.
- ▶ “Enumerate index i of the array s” can be translated to:

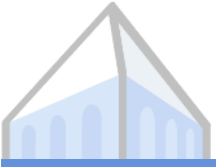
- ▶ `for (int i = 0; i < s.length(); i++) {`

- ▶ `for (int i = 0; i < s.length(); i++)`

$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$

## ▶ Semantic dependency and constraints between lines.

- ▶ Cannot use an undeclared variable or redeclare a declared variable in the same scope.
- ▶ Pseudocode does not explicitly state variable usage and declaration.
- ▶ “Set Ch to be a”
  - ▶ `char Ch = 'a';`
  - ▶ `char Ch = a;`
  - ▶ `Ch = 'a';`



# Semantic Constraints

Parse each fragment

## Fragments

```
int main () {  
    int n, ans = 1;  
    for (int i = 1; i <= n / 2; i++) {  
        ...  
    }
```

## Variables Used and Declared

declared: main	used:
declared: n, ans	used:
declared: i	used: n

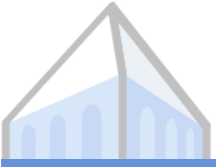


# Semantic Constraints

Parse each fragment							
Fragments	Variables Used and Declared						
<pre>int main () {     int n, ans = 1;     for (int i = 1; i &lt;= n / 2; i++) {         ...     }</pre>	<table><tr><td>declared: main</td><td>used:</td></tr><tr><td>declared: n, ans</td><td>used:</td></tr><tr><td>declared: i</td><td>used: n</td></tr></table>	declared: main	used:	declared: n, ans	used:	declared: i	used: n
declared: main	used:						
declared: n, ans	used:						
declared: i	used: n						

We keep a variable table to check whether any fragment re-declares a declared variable, or uses an undeclared variable.



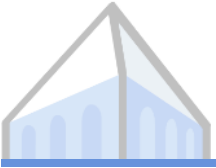


# Beam Search

Line number $l$	Best Candidate	2nd Candidate	3rd Candidate	4th Candidate	[Other]
1	$y_{1,1}$	$y_{1,2}$	$y_{1,3}$	$y_{1,4}$	...
2	$y_{2,1}$	$y_{2,2}$	$y_{2,3}$	$y_{2,4}$	...
3	$y_{3,1}$	$y_{3,2}$	$y_{3,3}$	$y_{3,4}$	...
4	$y_{4,1}$	$y_{4,2}$	$y_{4,3}$	$y_{4,4}$	...
5	$y_{5,1}$	$y_{5,2}$	$y_{5,3}$	$y_{5,4}$	...
6	$y_{6,1}$	$y_{6,2}$	$y_{6,3}$	$y_{6,4}$	...
7	$y_{7,1}$	$y_{7,2}$	$y_{7,3}$	$y_{7,4}$	...
8	$y_{8,1}$	$y_{8,2}$	$y_{8,3}$	$y_{8,4}$	...
9	$y_{9,1}$	$y_{9,2}$	$y_{9,3}$	$y_{9,4}$	...
10	$y_{10,1}$	$y_{10,2}$	$y_{10,3}$	$y_{10,4}$	...
...	...	...	...	...	...

1. To form one program candidate, we select one column for each row; the score is given by the equation below.
2. We want to find the top-K scoring candidates that satisfy the syntactic and semantic constraints.
3. We can efficiently check whether the first  $l$  fragment combination is valid under these constraints.

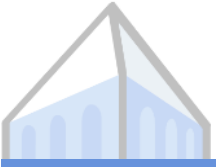
$$y = \text{Concat}_{l=1}^L y_{lc_l} \quad \text{score}(y) = \prod_{l=1}^L p_{lc_l}$$



# Performance

Method	Top 1	Top 10	Top 100	Top 1000
Baseline	0.0%	8.1%	29.2%	44.3%
Beam Syntactic	42.4% <sup>↑</sup>	51.3% <sup>↑</sup>	58.2% <sup>↑</sup>	N/A
Beam Semantic	45.6% <sup>↑</sup>	54.9% <sup>↑</sup>	61.9% <sup>↑</sup>	N/A

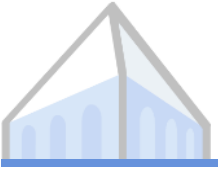
Adding constraints significantly improve over the baseline; semantic constraints also improve over syntactic constraints.



# Performance

Method	Top 1	Top 10	Top 100	Top 1000
Baseline	0.0%	8.1%	29.2%	44.3%
Beam Syntactic	42.4% <sup>↑</sup>	51.3% <sup>↑</sup>	58.2% <sup>↑</sup>	N/A
Beam Semantic	45.6% <sup>↑</sup>	54.9% <sup>↑</sup>	61.9% <sup>↑</sup>	N/A

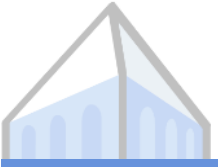
Time complexity of beam search grows **quadratically** with beam width ( $\geq K$ ), which becomes intractable if top-1000 candidate is considered.



# Motivation for Scaffold

---

- ▶ Baseline assumes independence between lines, which makes top-K generation **fast and exact**.
- ▶ Beam search deals with the inherent dependence between lines and **satisfies the constraint**.
- ▶ Scaffold search: first search the line dependency structure, then independently generate each line.



# Motivation for Scaffold

---

- ▶ Baseline assumes independence between lines, which makes top-K generation fast and exact.
- ▶ Beam search deals with the inherent dependence between lines and satisfies the constraint.
- ▶ Scaffold search: first search the line dependency structure, then independently generate each line.



# Configurations and Scaffold

Parse each fragment →	
Fragments	Syntactic Configurations $\phi$
int main () {	[return type] [function name] ( ) {
int n, ans = 1;	[statement] ;
for (int i = 1; i <= n / 2; i++) {	for [for conditions] {
...	

Parse each fragment →	
Fragments	Semantic Configurations $\phi$
int main () {	declared: main      used:
int n, ans = 1;	declared: n, ans    used:
for (int i = 1; i <= n / 2; i++) {	declared: i          used: n
...	



# Configurations and Scaffold

Fragments Parse each fragment

```
int main () {
    int n, ans = 1;
    for (int i = 1; i <= n / 2; i++) {
        ...
    }
}
```

## Syntactic Configurations $\phi$

```
[return type] [function name] ( ) {
    [statement] ;
    for [for conditions] {
        ...
    }
}
```

Fragments Parse each fragment

```
int main () {
    int n, ans = 1;
    for (int i = 1; i <= n / 2; i++) {
        ...
    }
}
```

## Semantic Configurations $\phi$

declared: main	used:
declared: n, ans	used:
declared: i	used: n

[return type] [function name] ( ) {	declared: main	used:
[statement] ;	declared: n, ans	used:
for [for conditions] {	declared: i	used: n
. . .	. . .	

Combine configuration from each line to form a program scaffold.



# Example Scaffolds

Valid Scaffold

<code>[return type] [function name] ( ) {</code>	declared: main	used:
<code>[statement] ;</code>	declared: n, ans	used:
<code>for [for conditions] {</code>	declared: i	used: n
<code>. . .</code>	<code>. . .</code>	

Invalid Scaffold

<code>[return type] [function name] ( ) {</code>	declared: main	used:
<code>[statement] ;</code>	declared: ans	used:
<code>for [for conditions] {</code>	declared: i	used: <b>n</b>
<code>. . .</code>	<code>. . .</code>	

Use of undeclared variable *n*

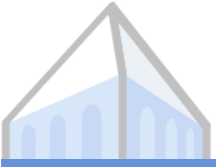
Invalid Scaffold

<code>[return type] [function name] ( )</code> <input type="checkbox"/>	declared: main	used:
<code>[statement] ;</code>	declared: ans	used:
<code>for [for conditions] {</code>	declared: i	used: n
<code>. . .</code>	<code>. . .</code>	

Missing “{”

We can efficiently check whether the combination of the first  $l$  configuration satisfies the constraint.





# Scoring Configs and Scaffolds

$p_{1,1}: 0.8$	$p_{1,1}: 0.1$	$p_{1,1}: 0.03$	$p_{1,1}: 0.01$	$\dots$
$y_{1,1}: i += 1;$	$y_{1,2}: \text{int } i = 1;$	$y_{1,3}: \text{int } i;$	$y_{1,4}: i -= 1;$	$\dots$

Marginalize code fragments to assign scores to configurations

$p = 0.81, \phi_{1,1}$

statement; variable used: None; declared  $i$ .

$p = 0.13, \phi_{1,2}$

statement; variable used:  $i$ ; declared: None.

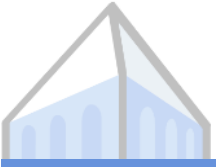
The score of a scaffold is the product of the probability scores for each configurations that form the scaffold.



# Beam Search for Scaffold

Line number $l$	Best Config	2nd Config	3rd Config	4th Config	[Other]
1	$\phi_{1,1}$	$\phi_{1,2}$	$\phi_{1,3}$	$\phi_{1,4}$	...
2	$\phi_{2,1}$	$\phi_{2,2}$	$\phi_{2,3}$	$\phi_{2,4}$	...
3	$\phi_{3,1}$	$\phi_{3,2}$	$\phi_{3,3}$	$\phi_{3,4}$	...
4	$\phi_{4,1}$	$\phi_{4,2}$	$\phi_{4,3}$	$\phi_{4,4}$	...
5	$\phi_{5,1}$	$\phi_{5,2}$	$\phi_{5,3}$	$\phi_{5,4}$	...
6	$\phi_{6,1}$	$\phi_{6,2}$	$\phi_{6,3}$	$\phi_{6,4}$	...
7	$\phi_{7,1}$	$\phi_{7,2}$	$\phi_{7,3}$	$\phi_{7,4}$	...
8	$\phi_{8,1}$	$\phi_{8,2}$	$\phi_{8,3}$	$\phi_{8,4}$	...
9	$\phi_{9,1}$	$\phi_{9,2}$	$\phi_{9,3}$	$\phi_{9,4}$	...
10	$\phi_{10,1}$	$\phi_{10,2}$	$\phi_{10,3}$	$\phi_{10,4}$	...
...	...	...	...	...	...

1. To form one scaffold, we select one config for each row; the score is the product of the config scores.
2. We want to find the top-50 scoring scaffolds that satisfy the syntactic and semantic constraints.
3. We can efficiently check whether the first  $l$  config combination is valid under these constraints.



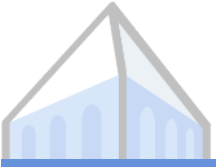
# Scaffold to Code

<code>[return type] [function name] ( ) {</code>	declared: main	used:
<code>[statement] ;</code>	declared: n, ans	used:
<code>for [for conditions] {</code>	declared: i	used: n
<code>. . .</code>	<code>. . .</code>	

`int n, ans;  
int n, ans = 0;`

`for (int i = 0; i < n; i++) {  
 for (int i = 0; i <= n - 1; i++) {  
 for (int i = n; i > 1; i--) {`

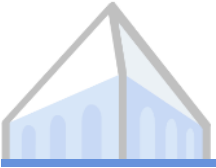
1. Given a scaffold, we only consider code fragment candidates that agree with the configuration for each line.
2. Any code fragment combination satisfies the constraint, if the scaffold is valid.
3. We can efficiently generate top-K candidates given a scaffold.



# Performance

	Method	Top 1	Top 10	Top 100	Top 1000
Beam size 200 Beam size 50	Baseline	0.0%	8.1%	29.2%	44.3%
	Beam Syntactic	42.4%	51.3%	58.2%	N/A
	Beam Semantic	45.6%	54.9%	61.9%	N/A
	Scaffold Semantic	45.8% <sup>↑</sup>	55.3% <sup>↑</sup>	62.8% <sup>↑</sup>	67.6%

1. With scaffold search, we obtain better performance with ~16x less compute.
2. Scaffold search allows us to calculate the performance of top-1000 candidates.



# Conclusion

---

- ▶ We propose a method to efficiently generate 1000 candidates for long programs and search for a semantically correct solution.
- ▶ We need semantic constraints as well to improve performance.
- ▶ Scaffold search improves both search qualities and efficiency.

# Thank you!

Paper: <https://arxiv.org/abs/2005.05927>

Code: <https://github.com/ruiqi-zhong/SemanticScaffold>

