Dynamic Software Updates: A VM-centric Approach

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The only thing that is constant is change. Heraclitus of Ephesus

Motivation

- Software applications change all the time
- Deployed systems must be updated with bug fixes, new features
- Updating typically involves: stop, apply patch, restart

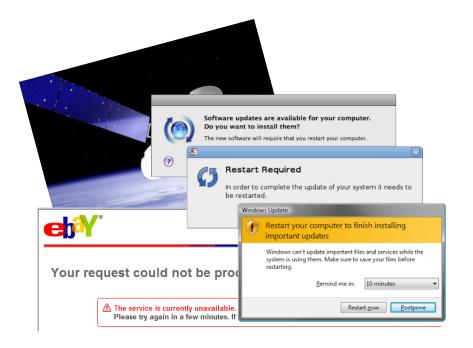






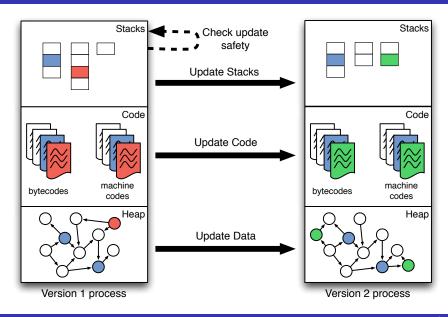
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The fundamental problem is losing state because of downtime.

Dynamic software updating



Overheard on programming.reddit







Dynamic updating systems

- Special-purpose architectures, application-specific solutions exist
- General-purpose solutions gaining strength
 - K42, Ksplice for OS updates
 - Polus, Ginseng for C applications
- Not for managed languages

DSU opportunity for managed languages

DSU Solutions for C/C++ typically

- Require special compilation
- Statically/dynamically insert indirection for function calls
- Restrict structure updates, require extra allocation
- Impose space/time overheads on normal execution
- Make type-safety for updates difficult
- Not multi-threaded

Our solution

- JVOLVE a Java Virtual Machine with DSU support
- Key insight: Extend existing VM services
- No DSU-related overhead during normal execution
- Support updates to real world applications

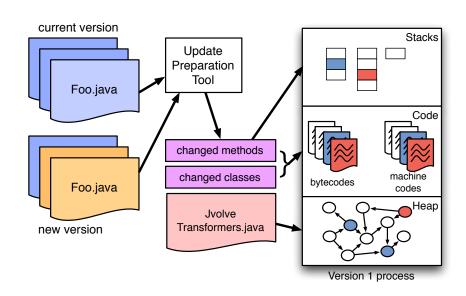
Dynamic software updating in managed languages can be achieved in a safe, flexible and efficient manner by naturally extending existing VM services.

DSU support should be a standard feature of future VMs.

Outline

- Introduction
- JVOLVE
 - VM Implementation
 - Evaluation
- Future Work
- Conclusion

JVOLVE - System overview



Supported updates

Changes within the body of a method

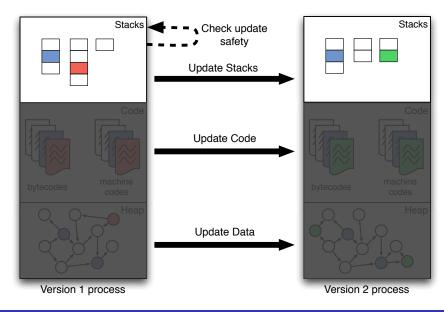
```
public static void main(String args[]) {
    System.out.println("Hello, World.");
    System.out.println("Hello again, World.");
}
```

- Class signature updates
 - Add, remove, change the type signature of fields and methods

```
public class Line {
- private final Point2D p1, p2;
+ private final Point3D p1, p2;
...
}
```

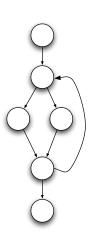
Signature updates require an object transformer function

Check for update safety



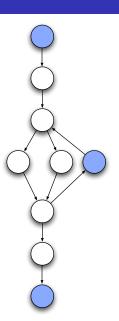
Safe point for the update

- Update must be atomic
- Updates happen at "safe points"
- Safe points are VM yield points, and restrict what methods can be on stack
- Extend the thread scheduler to suspend all application threads
- If any stack has a restricted method, delay the update



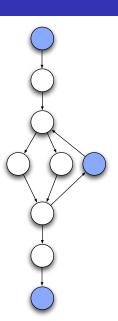
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Restricted methods

- (1) Methods changed by the update
- (2) Methods identified by the user as unsafe based on semantic information about the application

Install return barriers that trigger DSU upon unsafe method's return

- (3) Methods whose bytecode is unchanged, but compiled representation is changed by the update
 - Offsets of fields and methods hard-coded in machine code
 - Inlined callees may have changed

Utilize on-stack replacement to recompile base-compiled methods

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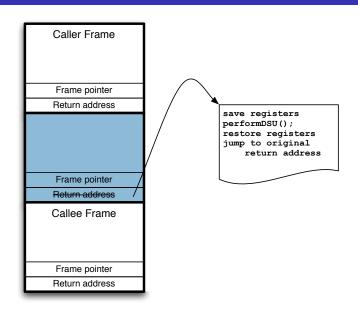
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Utilize on-stack replacement to recompile base-compiled methods

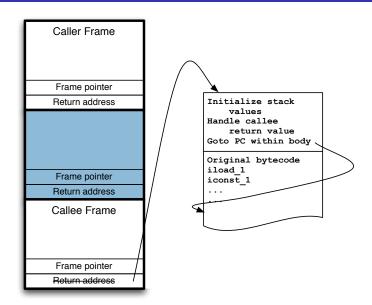
Handling restricted methods

Caller Frame Frame pointer Return address Frame pointer Return address Callee Frame Frame pointer Return address

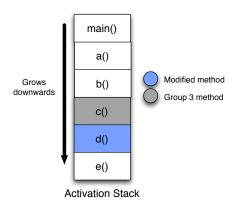
Installing return barrier for DSU



Performing On-stack replacement

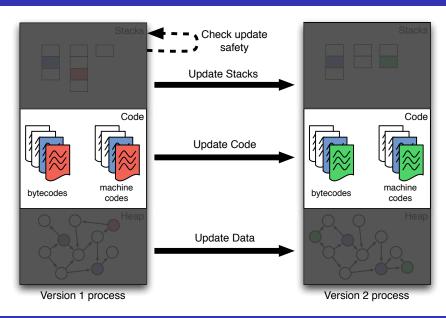


Reaching a safe point



Install a return barrier for d(). Wait till it returns. On-stack replace new machine code for c().

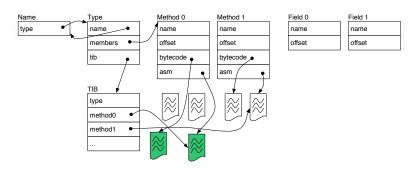
Update code



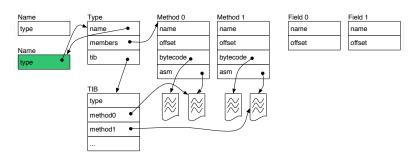
Update code

- Modify class loader to recognize new versions of classes
- Install new versions of classes and methods
- Rely on Just-in-time Compiler to compile new versions of methods on demand
- Extend On-stack replacement to update active methods

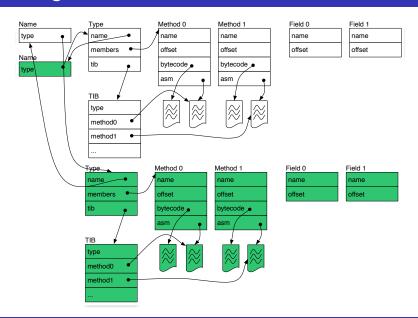
Updating a method



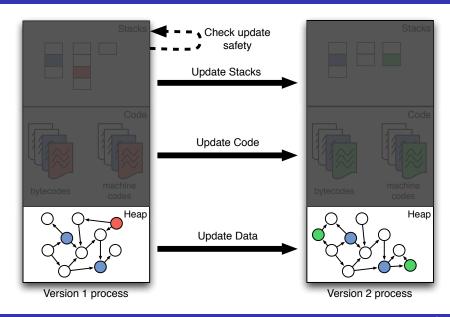
Updating a class



Updating a class



Update data



Example of an update (JavaEmailServer)

```
public class User {
  private final String username, domain, password;
 private String[] forwardAddresses;
 private EmailAddress[] forwardAddresses;
  public User(...) {...}
  public String[] getForwardedAddresses() {...}
  public void setForwardedAddresses(String[] f) {...}
}
public class ConfigurationManager {
  private User loadUser(...) {
     User user = new User(...);
     String[] f = ...;
     user.setForwardedAddresses(f):
     return user;
```

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     EmailAddress[] f = ...;
     user.setForwardedAddresses(f):
     return user;
```

Example of an update (JavaEmailServer)

```
public class v131_User {
  private final String username, domain, password;
 private String[] forwardAddresses;
                                              Stub generated by UPT
                                              for the old version
public class JvolveTransformers {
public static void jvolveClass(User unused) {}
public static void jvolveObject(User to, v131_User from) {
    to.username = from.username:
    to.domain = from.domain:
                                              Default transformer copies
    to.password = from.password;
                                              old fields, initializes new
    // to.forwardAddresses = null;
                                              ones to null
    to.forwardAddresses = new EmailAddress[len...
111
```

Example of an update (JavaEmailServer)

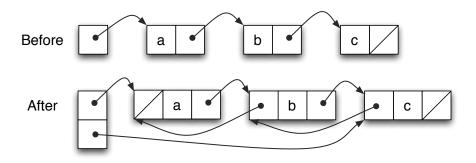
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public class JvolveTransformers {
public static void jvolveClass(User unused) {}
public static void jvolveObject(User to, v131_User from) {
    to.username = from.username:
    to.domain = from.domain:
    to.password = from.password;
    // to.forwardAddresses = null:
    int len = from.forwardAddresses.length;
    to.forwardAddresses = new EmailAddress[len]:
    for (int i = 0; i < len; i++) {
      to.forwardAddresses[i] =
       new EmailAddress(from.forwardAddresses[i]);
111
```

Compiling transformation functions

```
public static void jvolveObject(User to, v131_User from) {
   to.username = from.username;
   ...
}
```

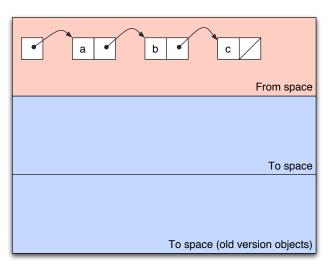
- Very close to Java semantics
- Field username is private and final
- Functions compiled specially by a JastAddJ extension to the Java language
- Ignores access protection and allows assigning to final fields

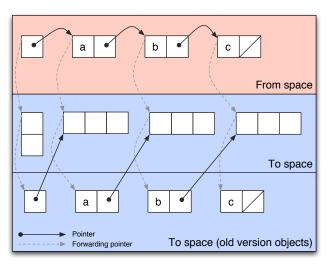
Transforming objects in the GC

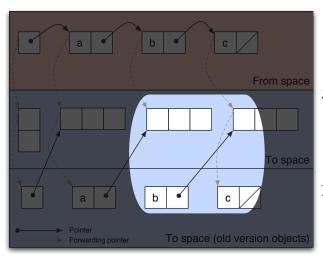


Happens in two steps

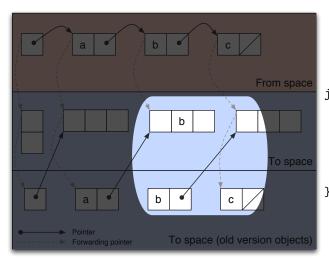
- Garbage collector creates an additional empty copy for updated objects
- Walk through and transform all these objects



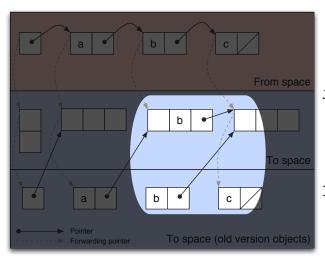




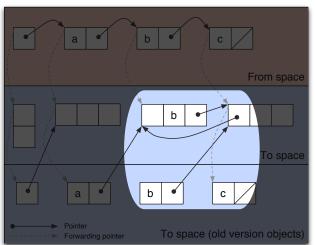
```
jvolveObject(Node to,
        old_Node from) {
  to.data = from.data;
  to.next = from.next;
  if (to.next != null)
        to.next.prev = to;
}
```



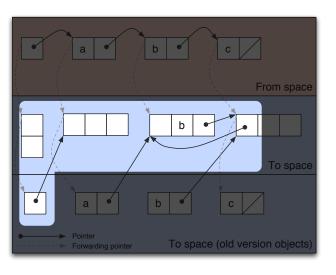
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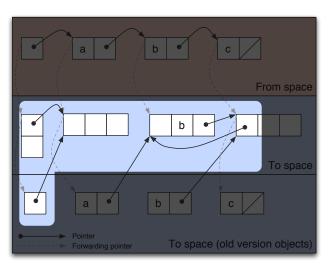


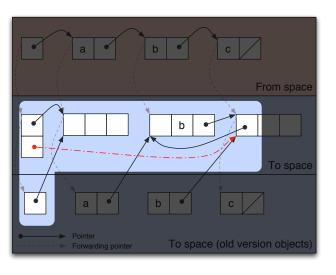
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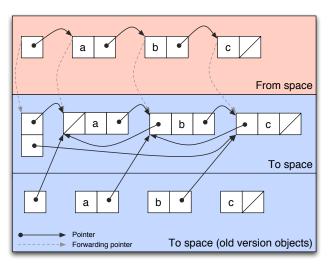


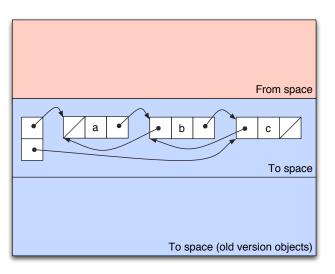
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Application Experience

- Jetty webserver
 - 11 versions, 5.1.0 through 5.1.10, 1.5 years
 - 45 KLOC
- JavaEmailServer
 - 10 versions, 1.2.1 through 1.4, 2 years
 - 4 KLOC
- CrossFTP server
 - 4 versions, 1.05 through 1.08, more than a year
 - 18 KLOC

What works

Support 20 of 22 updates

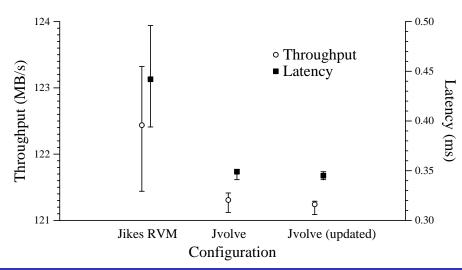
- 13 updates change class signature by adding new fields
- Several updates require On-stack replacement support
- Two versions update an infinite loop, postponing the update indefinitely

Unsupported updates

- JavaEmailServer 1.2.4 to 1.3
 - Update reworks the configuration framework of the server
 - Many classes are modified to refer to the configuration system
 - Including infinite loops in SMTP and POP threads
- Jetty 5.1.2 to 5.1.3
 - The application would never reach a safe point
 - Modified method ThreadedServer.acceptSocket() that waits for connections is nearly always on stack
 - Return barrier not sufficient since the main method in other threads PoolThread.run() is itself modified

${ m JVOLVE}$ performance

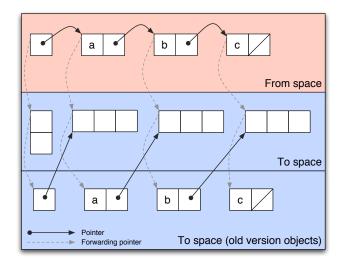
No overhead during steady-state execution



Update pause time

- No apriori overhead during normal execution (before and after the update)
- Only effect on execution time is the update pause time
 - Comparable to GC pause time

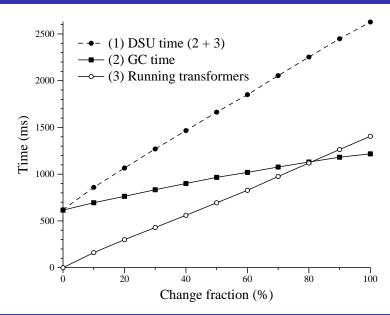
Update pause time



Update pause time

```
DSU Pause Time \cong Regular GC Time +
    Time to allocate upd. objects +
    Time to transform objects
    \propto Upd. objects fraction
    Heap size
```

DSU pause times (microbenchmark)



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Looking ahead

What's in the path towards mainstream adoption?

Safety

- Currently, we guarantee type safety. Can we do more?
- Safe behavior in a general context is undecidable [Gupta et al., 1996]
- Divide program into transactions, ensure they are version consistent [Neamtiu et al., 2008]
- Exhaustively testing updates

Flexibility

- Supporting every possible update is not a goal
- Support most of the updates that occur in practice
- Large and complex applications and updates
- Understand update semantics from refactorings

Efficiency

- Update time roughly proportional to GC time
- Semi-space GC is not practical
- We have two requirements
 - Copying objects
 - Full-heap collection
- Other collectors? Concurrent collectors?

Conclusion

- JVOLVE, a Java VM with support for Dynamic Software Updating
- Most-featured, best-performing DSU system for Java
- Naturally extends existing VM services
- Supports about two years worth of updates

Dynamic software updating in managed languages can be achieved in a safe, flexible and efficient manner.

Source code and other information: http://www.cs.utexas.edu/~suriya/jvolve

References I

- Gupta, D., Jalote, P., and Barua, G. (1996).

 A formal framework for on-line software version change.

 IEEE Trans. Softw. Eng., 22(2):120–131.
- Neamtiu, I., Hicks, M., Foster, J. S., and Pratikakis, P. (2008). Contextual effects for version-consistent dynamic software updating and safe concurrent programming. In *Proc. POPL*.