## Lógica Quântica Lecture notes and exercise sheet 4

## Natural transformations

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**Definition 1.** Let  $\mathbf{C}$  and  $\mathbf{D}$  be categories, and let  $F: \mathbf{C} \longrightarrow \mathbf{D}$  and  $G: \mathbf{C} \longrightarrow \mathbf{D}$  be functors. A natural transformation  $\eta$  from F to G, written  $\eta: F \Longrightarrow G$ , is given by an arrow  $\eta_A: FA \longrightarrow GA$  of  $\mathbf{D}$  for each object  $A \in \mathsf{Ob}(\mathbf{C})$  satisfying the following property, called naturality: for every arrow  $h: A \longrightarrow B$  in  $\mathbf{C}$ ,

$$\eta_B \circ Fh = Gh \circ \eta_A$$
,

i.e. the following diagram commutes:

$$FA \xrightarrow{\eta_A} GA$$

$$\downarrow^{Ff} \qquad \qquad \downarrow^{Gf}$$

$$FB \xrightarrow{\eta_B} GB$$

The natural transformation  $\eta$  is called a *natural isomorphism* if each  $\eta_A$  is an iso.

**Exercise 1.** Recall the functor List: Set  $\longrightarrow$  Set from exercise 3.5. Show that the following are natural transformations:

(a) reverse: List  $\Longrightarrow$  List where

$$\mathsf{reverse}_X \colon \mathsf{List}X \longrightarrow \mathsf{List}X :: [x_1, \dots, x_n] \longmapsto [x_n, \dots, x_1] ;$$

(b) unit:  $Id \Longrightarrow List where$ 

$$\mathsf{unix}_X \colon X \longrightarrow \mathsf{List}X :: x \longmapsto [x] ;$$

(c) flatten: List  $\circ$  List  $\Longrightarrow$  List where

$$\mathsf{flatten}_X \colon \mathsf{List}(\mathsf{List}X) \longrightarrow \mathsf{List}X \\ \coloneqq [[x_{1,1}, \dots, x_{1,k_1}], \dots, [x_{n,1}, \dots, x_{n,k_n}]] \\ \longmapsto [x_{1,1}, \dots, x_{1,k_1}, \dots, x_{n,1}, \dots, x_{n,k_n}] \\ \ \vdots \\ \$$

 $\mathrm{(d)} \ \mathsf{concat} \colon \mathsf{List} \times \mathsf{List} \Longrightarrow \mathsf{List} \ \mathrm{where}$ 

$$\operatorname{concat}_X : \operatorname{List}X \times \operatorname{List}X \longrightarrow \operatorname{List}X :: ([x_1, \dots, x_n], [y_1, \dots, y_m],) \longmapsto [x_1, \dots, x_n, y_1, \dots, y_m]$$

**Exercise 2.** Let  $D: \mathbf{Set} \longrightarrow \mathbf{Set}$  be the functor  $D = \times \circ \langle \mathsf{Id}, \mathsf{Id} \rangle$  mapping a set X to  $X \times X$  and a function f to  $f \times f$ , and let  $P_1: \mathbf{Set} \times \mathbf{Set} \longrightarrow \mathbf{Set}$  be the functor mapping (X,Y) to X and (f,g) to f.

- (a) Show that  $\Delta_X : X \longrightarrow X \times X :: x \longmapsto (x,x)$  defines a natural transformation  $\mathsf{Id} \Longrightarrow D$ .
- (b) Show that  $p_{X,Y}: X \times Y \longrightarrow X :: (x,y) \longmapsto x$  defines a natural transformation  $x \Longrightarrow P_1$ .
- (c) Show that the  $\Delta$  and p are the only natural transformations between these functors.

**Exercise 3.** Define analogous transformations  $\Delta$  and p for any category C with binary products.

**Exercise 4.** Let  $\mathbb{C}$  be a category with a terminal object T. Let  $K_T \colon \mathbb{C} \longrightarrow \mathbb{C}$  be the constant functor with value T, i.e.  $K_T$  maps any object to T and any arrow to  $\mathrm{id}_T$ . Show that the canonical arrows to the terminal object,  $\tau_A \colon A \longrightarrow T$  define a natural transformation  $\mathrm{Id}_{\mathbb{C}} \Longrightarrow K_T$ .

**Exercise 5** (Uniform deleting). Let  $\mathbb{C}$  be a category and T be any object of  $\mathbb{C}$ . A category  $\mathbb{C}$  has uniform deleting to T if there is a natural transformation  $e: \mathsf{Id}_{\mathbb{C}} \Longrightarrow K_T$  with  $e_T = \mathsf{id}_T$ . Show that  $\mathbb{C}$  has uniform deleting to T if and only if T is terminal.

**Exercise 6.** Recall the definition of dual of a vector space from exercise 3.6. Let V be a finite-dimensional vector space.

- (a) Show that V is isomorphic to its dual  $V^*$  and second dual  $V^{**}$ .
- (b) Show that the isomorphism  $V \cong V^{**}$  is natural, while there is no natural isomorphism  $V \cong V^{*}$ . Note how this is related to *basis independence*.

**Exercise 7.** Let **C** be a category with binary products and a terminal object **1**. Show that there are natural isomorphisms:

- (a)  $a_{A,B,C}: A \times (B \times C) \xrightarrow{\cong} (A \times B) \times C$  (Hint:  $\langle \langle \pi_1, \pi_1 \circ \pi_2 \rangle, \pi_2 \circ \pi_2 \rangle$ .
- (b)  $s_{A,B} \colon A \times B \xrightarrow{\cong} B \times A$
- (c)  $l_A \colon \mathbf{1} \times A \xrightarrow{\cong} A$
- (d)  $r_A : A \times \mathbf{1} \xrightarrow{\cong} A$

**Exercise 8.** Let P,Q be posets (seen as categories) and  $f,g:P\longrightarrow Q$  be functors, i.e. monotone functions. When is there a natural transformation  $f\Longrightarrow g$ ?

## Yoneda lemma