

Introduction to Optimization

Improving convergence rates.

Lecture 07: Quadratic functions and finite termination.

Lecture 08: Inertial algorithms. Stochastic gradient.



Systems of linear equations

Let $\mathcal{A} \in \mathbb{R}^{N \times N}$ be invertible (to simplify), and let $\beta \in \mathbb{R}^N$. The problem

$$\mathcal{A}x = \beta$$

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$$f(x) = \frac{1}{2} \|\mathcal{A}x - \beta\|^2,$$

and also the unique minimizer of

$$\phi(x) = \frac{1}{2} x \cdot Ax - b \cdot x,$$

where $A = \mathcal{A}^T \mathcal{A}$ (symmetric and positive definite) and $b = \mathcal{A}^T \beta$.

Conjugate gradient methods

Conjugate gradient methods iterate


$$x_{n+1} = x_n + \alpha_n d_n$$

for convenient choices of d_n and α_n in such a way that

$$\{d_0, \dots, d_{N-1}\}$$

is a basis of \mathbb{R}^N and x_n minimizes ϕ on the affine subspace

$$x_0 + \text{span}\{d_0, \dots, d_{n-1}\}.$$

As a consequence, the solution is found in at most N steps for every initial point x_0 . 

Conjugate directions

Vectors $\{d_0, \dots, d_{N-1}\}$ are **conjugate with respect to A** if

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is an inner product in \mathbb{R}^N . Being conjugate means being orthogonal with respect to this inner product. Conjugate vectors are linearly independent.

An abstract conjugate gradient method

Theorem

Let $\{d_0, \dots, d_{N-1}\}$ be conjugate, and let $x_{n+1} = x_n + \alpha_n d_n$, where

$$\alpha_n = \operatorname{Argmin}_{\alpha > 0} \phi(x_n + \alpha d_n) = -\frac{d_n \cdot \nabla \phi(x_n)}{\|d_n\|_A^2} = -\frac{d_n \cdot (Ax_n - b)}{\|d_n\|_A^2},$$

for $n = 0, \dots, N-1$. Then x_N minimizes ϕ , whence $Ax_N = b$.

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for $n = 0, \dots, N-1$. Then x_N minimizes ϕ , whence $Ax_N = b$.

Moreover, $\nabla \phi(x_n) \cdot d_j = 0$ for $j = 0, \dots, n-1$, and x_n minimizes ϕ over

$$x_0 + \operatorname{span}\{d_0, \dots, d_{n-1}\}.$$

This is known as *expanding subspace minimization*.

Implementation of a conjugate gradient method

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If $g_n = 0$, we stop; otherwise, we compute $\alpha_n = \frac{\|g_n\|^2}{\|d_n\|_A^2}$, and then update

$$x_{n+1} = x_n + \alpha_n d_n$$

$$g_{n+1} = g_n + \alpha_n A d_n$$

$$d_{n+1} = -g_{n+1} + \beta_n d_n, \quad \text{with} \quad \beta_n = \frac{\|g_{n+1}\|^2}{\|g_n\|^2}.$$

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The procedure described above produces a conjugate set $\{d_0, \dots, d_n\}$, with $n \leq N$, and terminates at x_n , where $Ax_n = b$.

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If A has exactly k distinct eigenvalues, the algorithm terminates in at most k steps.

Nonlinear extensions

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At step n , we know x_n , $g_n = \nabla f(x_n)$ and d_n .

If $g_n = 0$, we stop; otherwise, we compute (backtracking) α_n satisfying

$$\begin{aligned} f(x_n + \alpha_n d_n) &\leq f(x_n) + c_1 \alpha_n g_n \cdot d_n \\ |\nabla f(x_n + \alpha_n d_n) \cdot d_n| &\leq -c_2 g_n \cdot d_n, \end{aligned}$$

with $0 < c_1 < c_2 < 1/2$ (**strong Wolfe conditions**).

Nonlinear extensions, continued

Then, we update

$$\begin{aligned}x_{n+1} &= x_n + \alpha_n d_n \\ d_{n+1} &= -g_{n+1} + \beta_n d_n,\end{aligned}$$

where several choices for β_n are possible, such as:

- Fletcher-Reeves: $\frac{\|g_{n+1}\|^2}{\|g_n\|^2}$
- Polak-Ribière: $\frac{g_{n+1} \cdot (g_{n+1} - g_n)}{\|g_n\|^2}$
- Hestenes-Stiefel: $\frac{g_{n+1} \cdot (g_{n+1} - g_n)}{d_n \cdot (g_{n+1} - g_n)}$
- Dai-Yuan: $\frac{\|g_{n+1}\|^2}{d_n \cdot (g_{n+1} - g_n)}$

Break

Newton's method

Quadratic model of f given by Taylor's expansion at x_n :

$$f(x) \simeq f(x_n) + \nabla f(x_n) \cdot (x - x_n) + \frac{1}{2}(x - x_n) \cdot \nabla^2 f(x_n)(x - x_n).$$

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Example: $f(x) = \frac{1}{2} \|Ax - b\|^2$.



Newton's method

$$x_{n+1} = x_n - \nabla^2 f(x_n)^{-1} \nabla f(x_n)$$

Theorem

Consider $f : \mathbb{R}^N \rightarrow \mathbb{R}$ and $\hat{x} \in \mathbb{R}^N$ such that $\nabla f(\hat{x}) = 0$.

Suppose $\|\nabla^2 f(x) - \nabla^2 f(y)\| \leq L\|x - y\|$ and $\|\nabla^2 f(x)^{-1}\| \leq M$ for all x, y in a neighborhood of \hat{x} .

Then, there is $\delta > 0$ such that if $\|x_0 - \hat{x}\| < \delta$, then

$$\|x_n - \hat{x}\| \leq r^{2^n}$$

for some $r \in (0, 1)$ and all $n \geq 0$.

Avoid the cost of $\nabla^2 f(x_n)^{-1}$

Define $x_{n+1} = x_n - \alpha_n D_n \nabla f(x_n)$, where

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One simple heuristic is **periodic evaluation**: Choose $p \in \mathbb{N}$ and define

$$D_n = \nabla^2 f(x_{kp})^{-1} \quad \text{for } n = kp, kp + 1, \dots, (k + 1)p - 1.$$

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Another idea is to define D_{n+1} as a function of D_n (not $\nabla^2 f(x_n)$ or $\nabla^2 f(x_{n+1})$) while keeping the **essence** of Newton's method.

Quasi-Newton methods

The key equality defining Newton's method is the **secant condition**:

$$\nabla^2 f(x_n)^{-1} [\nabla f(x_{n+1}) - \nabla f(x_n)] = x_{n+1} - x_n.$$

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Let us construct D_{n+1} so that it is symmetric (like the Hessian), **no too different** from D_n , and satisfies

$$D_{n+1}g_n = s_n,$$

where $g_n = \nabla f(x_{n+1}) - \nabla f(x_n)$ and $s_n = x_{n+1} - x_n$.

Popular instances

- DFP: Davidon (1959), Fletcher and Powell (1987)

$$D_{n+1} = D_n - \frac{(D_n g_n)(D_n g_n)^T}{g_n \cdot D_n g_n} + \rho_n (s_n s_n^T), \quad \rho_n = \frac{1}{g_n \cdot s_n}.$$

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- BFGS: Broyden, Fletcher, Goldfarb and Shanno (1970)

$$D_{n+1} = (I - \rho_n s_n g_n^T) D_n (I - \rho_n s_n g_n^T)^T + \rho_n (s_n s_n^T).$$

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Theorem

Let $f : \mathbb{R}^N \rightarrow \mathbb{R}$ be μ -strongly convex and L -smooth, and let D_0 be positive definite. Then, x_n converges to the minimizer of f .

It does so in at most N steps if f is quadratic.