

## The SSM Instruction Set

- Basic instructions are colour-coded green. Additional instructions that your compilers are likely to use (eventually) are colour coded blue.
- The notation for describing the effect of an instruction on the opstack is borrowed from the JVM specification. Example: **sub**: ...,  $x, y \rightarrow \dots, x-y$   
Items on the stack (4-byte words) are listed with the uppermost item to the right. An ellipsis (...) denotes the lower part of the opstack. So, in this example,  $y$  is on top with  $x$  below  $y$ . Both items are popped and the result ( $x-y$ ) is pushed. Any items which happen to be on the opstack below  $x$  and  $y$  are left unchanged.
- mem[a]** denotes the single byte of memory at address  $a$ , while **mem4[a]** denotes the four-byte word which starts at memory address  $a$ .
- u16(x)** denotes the unsigned 2-byte integer obtained by discarding the two high-order bytes from the 4-byte word  $x$ .
- byte(x)** denotes the single byte obtained by discarding the three high-order bytes from the 4-byte word  $x$ .
- pad(b)** denotes the 4-byte word obtained by high-padding the single byte  $b$  with three zero bytes.

### SSM Instructions (1-byte)

Opcode	Mnemonic	Description
0 (0x00)	<b>noop</b>	No effect.
1 (0x01)	<b>halt</b>	Halt the machine. The opstack should normally be empty (normal termination); otherwise $n$ is popped from the opstack and <b>u16(n)</b> is interpreted as an error code for abnormal termination (configurable <sup>1</sup> ).
2 (0x02)	<b>pop</b>	Discard the top item from the opstack: ..., $x \rightarrow \dots$
3 (0x03)	<b>dup</b>	Push a duplicate of the top item onto the opstack: ..., $x \rightarrow \dots, x, x$
4 (0x04)	<b>swap</b>	Swap the top two items on the opstack: ..., $x, y \rightarrow \dots, y, x$
5 (0x05)	<b>rot</b>	“Rotate” the top three items on the opstack (pull third item to the top): ..., $x, y, z \rightarrow \dots, y, z, x$
6 (0x06)	<b>add</b>	Twos-complement integer addition: ..., $x, y \rightarrow \dots, x+y$
7 (0x07)	<b>sub</b>	Twos-complement integer subtraction: ..., $x, y \rightarrow \dots, x-y$
8 (0x08)	<b>mul</b>	Twos-complement integer multiplication: ..., $x, y \rightarrow \dots, x*y$
9 (0x09)	<b>div</b>	Twos-complement integer division: ..., $x, y \rightarrow \dots, x/y$
10 (0x0a)	<b>test_z</b>	If top item is zero, replace by 1, otherwise replace by 0: ..., $0 \rightarrow \dots, 1$ if $x = 0$ or ..., $x \rightarrow \dots, 0$ if $x \neq 0$
11 (0x0b)	<b>test_n</b>	If top item is negative, replace by 1, otherwise replace by 0: ..., $x \rightarrow \dots, 1$ if $x < 0$ or ..., $x \rightarrow \dots, 0$ if $x \geq 0$

<sup>1</sup> Defaults: 0 = normal termination, 1 = “Null Pointer”, 2 = “Array Index Out of Range”, 3 = “Heap Exhausted”.

## SSM Instructions (1-byte) continued

Opcode	Mnemonic	Description
12 (0x0c)	<b>get_dp</b>	Push the value of the DP register onto the opstack: $\dots \rightarrow \dots, \text{DP}$
13 (0x0d)	<b>get_fp</b>	Push the value of the FP register onto the opstack: $\dots \rightarrow \dots, \text{FP}$
14 (0x0e)	<b>get_sp</b>	Push the value of the SP register onto the opstack: $\dots \rightarrow \dots, \text{SP}$
15 (0x0f)	<b>load</b>	Load a 4-byte word from memory: $\dots, x \rightarrow \dots, \text{mem4}[\text{u16}(x)]$
16 (0x10)	<b>loadb</b>	Load a single byte from memory: $\dots, x \rightarrow \dots, \text{pad}(\text{mem}[\text{u16}(x)])$
17 (0x11)	<b>store</b>	Store a 4-byte word in memory: $\dots, x, y \rightarrow \dots$ with effect: $\text{mem4}[\text{u16}(x)] := y$
18 (0x12)	<b>storeb</b>	Store a single byte in memory: $\dots, x, y \rightarrow \dots$ with effect: $\text{mem}[\text{u16}(x)] := \text{byte}(y)$
19 (0x13)	<b>jump</b>	Unconditional jump: $\dots, x \rightarrow \dots$ with effect: control jumps to the instruction at address $\text{u16}(x)$
20 (0x14)	<b>jump_z</b>	Conditional jump: $\dots, x, y \rightarrow \dots$ with effect: control jumps to the instruction at address $\text{u16}(y)$ if $x = 0$
21 (0x15)	<b>jump_n</b>	Conditional jump: $\dots, x, y \rightarrow \dots$ with effect: control jumps to the instruction at address $\text{u16}(y)$ if $x < 0$
22 (0x16)	<b>call</b>	Function call <sup>2</sup> : $\dots, f, x_1, \dots, x_n, n \rightarrow \dots, v$ with effect: call the function at address $\text{u16}(f)$ passing $n$ four-byte arguments on the call-stack; push the function's return value $v$ on the opstack on return
23 (0x17)	<b>ret</b>	Return from call <sup>3</sup> : $\dots, v, n \rightarrow \dots$ with effect: pop $n+1$ four-byte words off the call-stack, restore the previous opstack, push $v$ on the restored opstack, jump to the current call's return address

<sup>2</sup> The function code executes in the context of a new, initially empty, opstack. The old opstack is restored when the call returns.

<sup>3</sup>  $n$  should match the total number of four-byte words (parameters plus locals) allocated for the current call;  $n+1$  words will be popped in total since one additional word is always used for stack management (return address and saved FP).

## SSM Instructions (2-byte)

Opcode	Mnemonic	Data (1 byte)	Description
24 (0x18)	<b>pushb</b>	<b>b</b>	Push a single byte on the opstack: ... → ..., <b>pad(b)</b>
25 (0x19)	<b>sysc</b>	<b>n</b>	Execute system call number <b>n</b> .

## System Calls

Syscall number	Literal	Arity <sup>4</sup>	Description
0 (0x00)	<b>OUT_BYTE</b>	1	Pop <b>n</b> from the opstack. Write the low-order byte of <b>n</b> on the standard output stream.
1 (0x01)	<b>OUT_CHAR</b>	1	Pop <b>n</b> from the opstack. Print the ASCII character specified by the low-order byte of <b>n</b> to the standard output stream.
2 (0x02)	<b>OUT_LN</b>	0	Print a line-ending to the standard output stream.
3 (0x03)	<b>OUT_DEC</b>	1	Pop twos-complement integer <b>n</b> from the opstack. Print the decimal representation of <b>n</b> to the standard output stream.
4 (0x04)	<b>OUT_STR</b>	1	Pop <b>x</b> from the opstack. Print the string stored at address <b>u16(x)</b> to the standard output. The following memory layout is assumed for string data: the first two bytes provide the string length (an unsigned integer <b>n</b> in the range $0 \leq n < 65536$ ); the following <b>n</b> bytes contain the ASCII codes of the letters in the string.
5 (0x05)	<b>READ_BYTE</b>	0	Read a single byte from the standard input stream, high-pad with three zero bytes, and push on the opstack.
6 (0x06)	<b>READ_INT</b>	0	Attempt to read a valid decimal integer string from the standard input stream. If successful (ie a valid string was entered and parsed as integer <b>n</b> ) push <b>n</b> followed by a 1: ... → ..., <b>n</b> , 1. If an invalid string was entered, just push 0: ... → ..., 0.
7 (0x07)	<b>PUSH_ARGC</b>	0	Push <b>n</b> on the opstack where <b>n</b> is the number of command-line arguments that were provided to the SSM on start-up.
8 (0x08)	<b>PUSH_ARG</b>	1	Pop <b>n</b> from the opstack. Use <b>n</b> as an index into the command-line parameter array and push <b>a</b> on the opstack, where <b>a</b> is the memory address of the corresponding string data. The machine will halt with an error if <b>n</b> is out of range.
9(0x09)	<b>MALLOC</b>	1	Pop <b>n</b> from the opstack. Allocate a contiguous block of <b>n</b> bytes of memory in the heap and push <b>a</b> on the opstack, where <b>a</b> is the memory address of the start of the allocated block; push 0 if allocation fails (heap exhausted).
10(0x0a)	<b>CALLOC</b>	1	The same as MALLOC but the allocated memory is filled with zeroes.
11(0x0b)	<b>FREE</b>	1	Pop <b>x</b> from the opstack. Deallocate the previously allocated block of memory starting at address <b>u16(x)</b> . It is an error if <b>u16(x)</b> is not an address previously returned by MALLOC or CALLOC, or if already deallocated.

<sup>4</sup> The “arity” of a system call is the number of four-byte words that it consumes from the opstack.

## SSM Instructions (4-byte)

Opcode	Mnemonic	Data (2 bytes)	Description
26 (0x1a)	<b>loadi</b>	$a$	Load a four-byte word from memory: $\dots \rightarrow \dots, \text{mem4}[a]$
27 (0x1b)	<b>loadbi</b>	$a$	Load a single byte from memory: $\dots \rightarrow \dots, \text{pad}(\text{mem}[a])$
28 (0x1c)	<b>storei</b>	$a$	Store a four-byte word in memory: $\dots, x \rightarrow \dots$ with effect: $\text{mem4}[a] := x$
29 (0x1d)	<b>storebi</b>	$a$	Store a single byte in memory: $\dots, x \rightarrow \dots$ with effect: $\text{mem}[a] := \text{byte}(x)$
30 (0x1e)	<b>jumpi</b>	$a$	Unconditional jump: control jumps to the instruction at address $a$ (opstack is unchanged)
31 (0x1f)	<b>jumpi_z</b>	$a$	Conditional jump: $\dots, x \rightarrow \dots$ with effect: control jumps to the instruction at address $a$ if $x = 0$
32 (0x20)	<b>jumpi_n</b>	$a$	Conditional jump: $\dots, x \rightarrow \dots$ with effect: control jumps to the instruction at address $a$ if $x < 0$
33 (0x21)	<b>calli</b>	$a$	Function call <sup>5</sup> : $\dots, x_1, \dots, x_n, n \rightarrow \dots, v$ with effect: call the function at address $a$ passing $n$ four-byte arguments on the call-stack; push the function's return value $v$ on the opstack on return
34 (0x22)	<b>salloc</b>	$n$	Allocate space on the call-stack: $\dots, n \rightarrow \dots$ with effect: allocate $n$ four-byte words of space on the call-stack (decrement SP by $4*n$ )
35 (0x23)	<b>sfree</b>	$n$	De-allocate space from the call-stack: $\dots, n \rightarrow \dots$ with effect: de-allocate $n$ four-byte words of space from the call-stack (increment SP by $4*n$ )

## SSM Instructions (5-byte)

Opcode	Mnemonic	Data (4 bytes)	Description
36 (0x24)	<b>push</b>	$x$	Push a four-byte word on the opstack: $\dots \rightarrow \dots, x$

<sup>5</sup> Note that the top value on the opstack ( $n$ ) determines how many additional items (the function arguments) will be consumed from the opstack and pushed onto the call-stack. The function code executes in the context of a separate opstack of its own (initially empty).