

# 1 Minima of convex functions

Consider the constrained problem

$$\min_{x \in \mathcal{X}} f(x). \quad (P)$$

Recall definition of convex functions.

**Theorem 1.** Consider the problem (P). Suppose  $f$  is convex, and  $\mathcal{X}$  is convex, closed and non-empty. Then:

1. Any local solution to (P) is also a global solution.

2. The set of global solutions to (P) is convex.

*pf.* Part I: Suppose for contra,  $x^*$  is a local sol. is not a global sol.

$$\exists \bar{x} \text{ s.t. } f(\bar{x}) < f(x^*). \quad \mathcal{X} \text{ is cvx} \Rightarrow (1-\alpha)x^* + \alpha\bar{x} \in \mathcal{X}.$$

$$\because f \text{ is cvx} \quad \therefore f((1-\alpha)x^* + \alpha\bar{x}) \leq (1-\alpha)f(x^*) + \alpha f(\bar{x}) < (1-\alpha + \alpha)f(x^*) = f(x^*).$$

$$\text{Set } \alpha \rightarrow 1. \Rightarrow \nexists B(x^*, \varepsilon) \text{ s.t. } \forall x \in B(x^*, \varepsilon), f(x) \geq f(x^*).$$

Contradicts with  $x^*$  is a local min.

Part 2.

Suppose  $x^*, y^* \in \mathcal{X}$   
 $f(x^*) = f(y^*)$ , both global sol.  $\forall x \in \mathcal{X}, f(x) \geq f(x^*) = f(y^*)$ .

$$\text{Consider } (1-\alpha)x^* + \alpha y^* \in \mathcal{X}. \quad \text{Since } f \text{ is cvx, } f((1-\alpha)x^* + \alpha y^*) \leq (1-\alpha)f(x^*) + \alpha f(y^*) \\ = f(x^*)$$

$$\because x^* \text{ is a global min} \quad \therefore f((1-\alpha)x^* + \alpha y^*) \geq f(x^*)$$

$$\therefore f((1-\alpha)x^* + \alpha y^*) = f(x^*) \Rightarrow (1-\alpha)x^* + \alpha y^* \text{ is a global sol.}$$

$\therefore$  The set of all global sol. is cvx.

## 1.1 Continuously differentiable convex functions

**Theorem 2** (Equivalent characterization of convexity). The following are true.

1. Let  $f: \mathbb{R}^d \rightarrow \mathbb{R}$  be continuously differentiable. The function  $f$  is convex if and only if

$$\forall x, y: f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle. \quad (1)$$

(A picture. From local to global.)

2. Let  $f: \mathbb{R}^d \rightarrow \mathbb{R}$  be twice continuously differentiable. The function  $f$  is convex if and only if

$$\forall x: \nabla^2 f(x) \succeq 0.$$

Recall: Def of cvx function:  $\text{epi}(f)$  is a cvx set.

pf: part I:

$$f(\alpha x + (1-\alpha)y) \leq \alpha f(x) + (1-\alpha)f(y).$$

( $\Rightarrow$ ) Suppose  $f$  is cvx.  $f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y)$ .

$$\begin{aligned} f(y) - f(x) &\geq \frac{f((1-\alpha)x + \alpha y) - f(x)}{\alpha} = \frac{\langle \nabla f(x), \alpha(y-x) \rangle + o(\|\alpha\|)}{\alpha} \\ &= \langle \nabla f(x), y-x \rangle + \frac{o(\|\alpha\|)}{\alpha} \end{aligned}$$

Set  $\alpha \geq 0$ .  $\square$ .

( $\Leftarrow$ ). Set  $z = (1-\alpha)x + \alpha y$ .

$$f(x) \geq f(z) + \langle \nabla f(z), \alpha(x-y) \rangle$$

$$f(y) \geq f(z) + \langle \nabla f(z), (1-\alpha)(y-x) \rangle$$

$$\begin{aligned} \text{Hence } f(z) &= (1-\alpha)f(x) + \alpha f(y) \leq (1-\alpha)(f(x) - \langle \nabla f(z), \alpha(x-y) \rangle) \\ &\quad + \alpha(f(y) - \langle \nabla f(z), (1-\alpha)(y-x) \rangle) \\ &= (1-\alpha)f(x) + \alpha f(y). \end{aligned}$$

$$\therefore f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y).$$

Part 2.

By Taylor's thm,  $\alpha > 0$ ,  $\exists \alpha \in (0, 1)$

$$f(x + \alpha u) = f(x) + \langle \nabla f(x), \alpha u \rangle + \frac{\alpha^2}{2} u^T \nabla^2 f(x + \alpha u) u.$$

$$\begin{aligned} (\Rightarrow) \nabla^2 f(x) \succeq 0 &\Rightarrow u^T \nabla^2 f(x + \alpha u) u \geq 0, \forall u \Rightarrow f(x + \alpha u) \geq f(x) + \langle \nabla f(x), \alpha u \rangle \\ &\Rightarrow f(y) \geq f(x) + \langle \nabla f(x), y-x \rangle. \end{aligned}$$

$$(\Leftarrow) \text{ If } f \text{ is cvx. } f(x + \alpha u) \geq f(x) + \langle \nabla f(x), \alpha u \rangle \Rightarrow u^T \nabla^2 f(x) u \geq 0, \forall u.$$

$$\Rightarrow \nabla^2 f(x) \succeq 0.$$

Convex function: 4 Equiv statement:

$$f(y) \geq f(x) + \langle \nabla f(x), y-x \rangle, \forall x, y.$$

$$\Leftrightarrow \nabla^2 f(x) \succeq 0, \forall x$$

$$\Leftrightarrow f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y)$$

$$\Leftrightarrow \text{epi}(f) \text{ is a convex set.}$$

**Theorem 3** (Sufficient condition for global optimality). Consider the problem (P), where  $f$  is continuously differentiable and convex. If  $x^* \in \mathcal{X}$  and  $\nabla f(x^*) = 0$ , then  $x^*$  is a global minimizer of  $f$ .

Recall: For smooth only functions, check local min is NP-hard.  
But for convex function, it's easy to check global min.

Pf:  $\nabla f(x^*) = 0$ .  $x = x^* + p$ , arbitrary  $x \in \mathcal{X}$ .

$$f(x) - f(x^*) \geq \langle \nabla f(x^*), p \rangle = 0. \Rightarrow \forall x, f(x) \geq f(x^*). \Rightarrow x^* \text{ is a global min.}$$

**Remark 1.** Theorem 3 holds for both unconstrained (i.e.,  $\mathcal{X} = \mathbb{R}^d$ ) and constrained problems. Using terminology from last time,  $x^*$  being a stationary point is sufficient for global optimality. For unconstrained problem, this is also necessary (Lecture 4, Theorem 1). For constrained problem, this may not be necessary (example).

## 2 Strongly convex functions

We use Euclidean norm  $\|\cdot\|_2$  in this section.

**Definition 1** (Strong convexity). Given  $m > 0$ , we say that  $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$  is strongly convex with modulus/parameter  $m$  (or  $m$ -strongly convex for short), if

$$\forall x, y \in \mathbb{R}^d : f((1-\alpha)x + \alpha y) \leq (1-\alpha)f(x) + \alpha f(y) - \frac{m}{2}(1-\alpha)\alpha \|y-x\|_2^2.$$

**Remark 2.** Verify yourself that the above is equivalent to convexity of the function  $f(x) - \frac{m}{2}\|x\|_2^2$ .

↓  
Verification.

$$g(x) = f(x) - \frac{m}{2}\|x\|_2^2.$$

( $\Rightarrow$ ) Assume  $f(x)$  is  $m$ -strongly-conv.

$$\text{Then } g((1-\alpha)x + \alpha y) = f((1-\alpha)x + \alpha y) - \frac{m}{2}\|(1-\alpha)x + \alpha y\|_2^2$$

$$\leq (1-\alpha)f(x) + \alpha f(y) - \frac{m}{2}(1-\alpha)\alpha \|y-x\|_2^2 - \frac{m}{2}\|(1-\alpha)x + \alpha y\|_2^2$$

$$= (1-\alpha)(f(x) - \frac{m}{2}\|x\|_2^2) + \alpha(f(y) - \frac{m}{2}\|y\|_2^2) = (1-\alpha)g(x) + \alpha g(y).$$

( $\Leftarrow$ ) Assume  $g(x)$  is CVX.

$$f((1-\alpha)x + \alpha y) - \frac{m}{2} \|(1-\alpha)x + \alpha y\|_2^2 \leq (1-\alpha) \left( f(x) - \frac{m}{2} \|x\|_2^2 \right) + \alpha \left( f(y) - \frac{m}{2} \|y\|_2^2 \right)$$

整理可得.

**Theorem 4** (Equivalent characterization of strong convexity). *The following hold.*

1. Suppose  $f$  is continuously differentiable. Then  $f$  is  $m$ -strong convexity if and only if

$$f(y) \geq f(x) + \langle \nabla f(x), y - x \rangle + \frac{m}{2} \|y - x\|_2^2.$$

(A picture. Compare with convexity only. Complements  $L$ -smoothness.)

2. Suppose  $f$  is twice continuously differentiable. Then  $f$  is  $m$ -strong convexity if and only if

$$\forall x : \nabla^2 f(x) \succeq mI.$$

(Compare with  $L$ -smoothness)

pf. Apply Thm 2 to  $g(x) = f(x) - \frac{m}{2} \|x\|_2^2$ .

**Theorem 5.** Suppose that  $f : \mathbb{R}^d \rightarrow \bar{\mathbb{R}}$  is continuously differentiable and  $m$ -strongly convex for some  $m > 0$ . If  $x^* \in \mathcal{X}$  satisfies  $\nabla f(x^*) = 0$ , then  $x^*$  is the unique global minimizer of  $f$ .

Proof. By Part 1 of Theorem 4:

$$f(x) \geq f(x^*) + \langle \nabla f(x^*), x - x^* \rangle + \underbrace{\frac{m}{2} \|x - x^*\|_2^2}_{>0 \text{ unless } x=x^*}.$$

强凸性保证 global min 的唯一性.

### 3 Algorithmic setup

1. First-order oracle:

$$x \rightarrow \text{oracle} \rightarrow f(x), \nabla f(x)$$

2. Second-order oracle:

$$x \rightarrow \text{oracle} \rightarrow f(x), \nabla f(x), \nabla^2 f(x)$$

All algorithms we consider in this course are iterative:

- start with some  $x_0$
- at iteration  $k = 0, 1, 2, \dots$ 
  - get oracle answers for  $x_k$ , choose  $x_{k+1}$

### 4 Basic descent methods

Take the form

$$x_{k+1} = x_k + \alpha_k p_k, \quad k = 0, 1, \dots$$

**Definition 2.**  $p \in \mathbb{R}^d$  is a descent direction for  $f$  at  $x$  if

$$f(x + tp) < f(x)$$

for all sufficiently small  $t > 0$ .

**Proposition 1.** If  $f$  is continuously differentiable (in a neighborhood of  $x$ ), then any  $p$  such that  $\langle -\nabla f(x), p \rangle > 0$  is a descent direction.

$$\langle \nabla f(x), p \rangle < 0.$$

pf: By Taylor's thm,

$$f(x+tp) = f(x) + \langle \nabla f(x), tp \rangle + o(\|t\|).$$

Set  $t \rightarrow 0$ . For  $p$  s.t.  $\langle \nabla f(x), p \rangle < 0$ ,  $f(x+tp) < f(x)$ .  $\square$ .

## 5 Gradient descent

Any  $p$  with  $\langle -\nabla f(x), p \rangle > 0$  is a descent direction. What would be a good choice? One that maximizes  $\langle -\nabla f(x), p \rangle$  over some set of  $p$ 's.

For example, look at all  $p$  with  $\|p\|_2 = 1$ . Then

$$\sup_{\|p\|_2=1} \langle -\nabla f(x), p \rangle = \|\nabla f(x)\|_2$$

$\Downarrow$   
maximize this!

attained for  $p = -\frac{\nabla f(x)}{\|\nabla f(x)\|_2}$ .

✓.  $-\nabla f(x)$  is a unit vec.

That is, try to move in the direction of the negative gradient,  $-\nabla f(x)$ .

"Simplest" descent algorithm:

$$x_{k+1} = x_k - \alpha_k \nabla f(x_k),$$

where  $\alpha_k$  is the step size. Ideally, choose  $\alpha_k$  small enough so that

$$f(x_{k+1}) < f(x_k)$$

for large  $\alpha_k$ ,

$\alpha_k \|\nabla f(x_k)\|$  need to be considered.

when  $\nabla f(x_k) \neq 0$ .

Known as "gradient method", "gradient descent", "steepest descent" (w.r.t. the  $\ell_2$  norm).