

Assignment 1

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October 25, 2020

My code consists of three parts.

1. Get the whole input data from 'hw1_input.txt',
and map key attribute string to integer values between 0 and 63.
2. Sort only key attribute values using optimized radix sort.
(Optimized by adding two mapping lists)
3. Write 'hw1_output.txt' in order according to map_list.

Below are detailed explanations and performance analysis of each part.

1. Get the whole input data from 'hw1_input.txt',
and map key attribute string to integer values between 0 and 63.

```
53     input_state = fgets(input_buffer,BUFFER_SIZE,fd_input);
54     //n is stored in input buffer
55     while(input_buffer[idx]!='\n')
56     {
57         data_size*=10;
58         data_size += input_buffer[idx]-'0';
59         idx++;
60     }
61     // gets $
62     input_state = fgets(input_buffer,BUFFER_SIZE,fd_input);
63     // gets list of attributes
64     input_state = fgets(input_buffer,BUFFER_SIZE,fd_input);
65     idx=0;
66     while(input_buffer[idx]!='\n')
67     {
68         if(input_buffer[idx]==DELIMITER)
69             attribute_num++;
70         else if(input_buffer[idx]==OPEN_BR)
71             key_attribute_idx=attribute_num-1;
72         idx++;
73     }
74     // gets $
75     input_state = fgets(input_buffer,BUFFER_SIZE,fd_input);
```

It gets the value of n and t, and gets what is the key attribute.

```

89 for(int data_idx=0;data_idx<data_size;data_idx++)
90 {
91     input_state = fgetc(input_buffer,BUFFER_SIZE,fd_input);
92     int deli_cnt=0;
93     int buff_idx=0;
94     int data_m_idx=0;
95     while(deli_cnt!=key_attribute_idx)
96     {
97         if(input_buffer[buff_idx]==DELIMITER)
98         {
99             deli_cnt++;
100             buff_idx++;
101         }
102         else
103             buff_idx++;
104     }
105     // printf("deli_cnt : %d key_attribute_idx : %d \n",deli_cnt,key_attribute_idx);
106     while(input_buffer[buff_idx]!=DELIMITER && input_buffer[buff_idx]!=NEW_LINE && input_buffer[buff_idx]!=EOF && input_buffer[buff_idx]!=0 )
107     {
108         unsigned char data=0;
109         char b_data = input_buffer[buff_idx];
110         if(b_data<NUM_0) //space
111             data_block[data_idx][data_m_idx]=1;
112         else if(b_data<AL_A) //numerical
113             data_block[data_idx][data_m_idx]=b_data-NUM_0+2;
114         else if(b_data<AL_a) //Upper letters
115             data_block[data_idx][data_m_idx]=2*(b_data-AL_A)+13;
116         else //Lower letters
117             data_block[data_idx][data_m_idx]=2*(b_data-AL_a)+12;
118         // printf("%c -> %d \n",b_data,data_block[data_idx][data_m_idx]);
119         buff_idx++;
120         data_m_idx++;
121     }
122     buff_idx=-1;
123     do
124     {
125         buff_idx++;
126         // printf("buff_idx : %d\n",buff_idx);
127         org_data[data_idx][buff_idx]=input_buffer[buff_idx];
128     } while(input_buffer[buff_idx]!=NEW_LINE&&input_buffer[buff_idx]!=0);
129     org_data_len[data_idx]=buff_idx+1;
130 }

```

And gets each input line. Here, it maps each key attribute character into a integer value as below.

Blank → 0 (Blank is not space. It means there is no character.)

Space → 1

‘0’ → 2

‘1’ → 3

.

.

.

‘a’ → 12

‘A’ → 13

.

.

.

‘z’ → 62

‘Z’ → 63

It stores these mapped integer values into ‘data_block’ and ‘temp_block’.

And it stores the whole attributes to ‘org_data’.

```

20 static char org_data[MAX_BLOCK_N][ (MAX_M)*(MAX_T) ];
21 static int org_data_len[MAX_BLOCK_N];
22 static unsigned char data_block[MAX_BLOCK_N][MAX_M];
23 static unsigned char temp_block[MAX_BLOCK_N][MAX_M];

```

I declared global static arrays, ‘data_block’, ‘temp_block’, and ‘org_data’.

This code block takes $O(nmt)$, where n is the number of objects, m is the maximum length of each attribute value, and t is the number of attributes for each object. Since each maximum values of n , m , and t are 1,000,000, 15, 10, the execution count is as similar as 150,000,000. (not sec, just representation of counts the code should run.) It does $O(nmt)$ twice because it first maps into integer

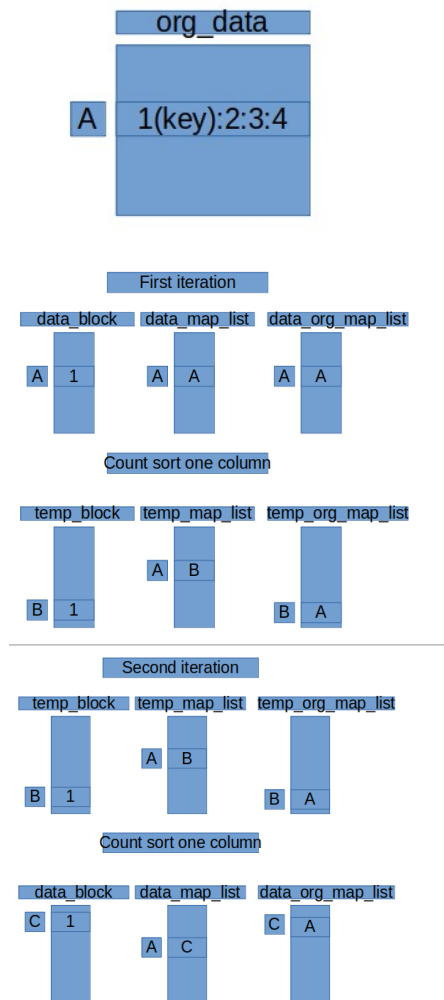
values and then store the whole attributes. I implemented like this for some optimization for storing whole attributes.

Anyway, this code blocks costs $O(nmt)$

2. Sort only key attribute values using optimized radix sort. (Optimized by adding two mapping lists)

```
152 void cnt_sort(int data_size, int col_idx) //count sort the directed index
153 {
154     for(int i=0;i<MAX_R;i++)
155         cnt_list[i]=0;
156     if(now_block_num==0)
157     {
158         for(int data_idx=0;data_idx<data_size;data_idx++)
159         {
160             unsigned char d_data = data_block[data_idx][col_idx];
161             cnt_list[d_data]++;
162         }
163         for(int i=1;i<MAX_R;i++) //getting prefix sum
164             cnt_list[i]+=cnt_list[i-1];
165         for(int data_idx=data_size-1;data_idx>=0;data_idx--)
166         {
167             unsigned char d_data = data_block[data_idx][col_idx];
168             int next_idx = cnt_list[d_data]-1;
169             for(int m_idx=0;m_idx<MAX_M;m_idx++)
170             {
171                 temp_block[next_idx][m_idx]=data_block[data_idx][m_idx];
172             }
173             int org_idx = data_org_map_list[data_idx];
174             temp_map_list[org_idx]=next_idx;
175             temp_org_map_list[next_idx]=org_idx;
176             cnt_list[d_data]-=1;
177         }
178         now_block_num=1;
179     }
180     else
181     {
182         for(int data_idx=0;data_idx<data_size;data_idx++)
183         {
184             unsigned char d_data = temp_block[data_idx][col_idx];
185             cnt_list[d_data]++;
186         }
187         for(int i=1;i<MAX_R;i++) //getting prefix sum
188             cnt_list[i]+=cnt_list[i-1];
189         for(int data_idx=data_size-1;data_idx>=0;data_idx--)
190         {
191             unsigned char d_data = temp_block[data_idx][col_idx];
192             int next_idx = cnt_list[d_data]-1;
193             for(int m_idx=0;m_idx<MAX_M;m_idx++)
194             {
195                 data_block[next_idx][m_idx]=temp_block[data_idx][m_idx];
196             }
197             int org_idx = temp_org_map_list[data_idx];
198             data_map_list[org_idx]=next_idx;
199             data_org_map_list[next_idx]=org_idx;
200             cnt_list[d_data]-=1;
201         }
202         now_block_num=0;
203     }
204 }
205
206 return;
207 }
208
```

I implemented count sorting function. The column index which should be sorted is given to the function. As I said above, there is only 64 integer values. So, it takes $O(n+64)$ to count sort one column. As count sorting doesn't sort in place, I implemented one more data block called 'temp_block'. The code uses both blocks alternatively. And two mapping lists are given to each block. One is 'map_list' and another one is 'org_map_list'. 'map_list' records where the original data is in block. And, 'org_map_list' records which original data is stored in block. It is illustrated below.



It iterates this for 15 times. So after 15 iterations, the sorted information is stored in 'temp_block', 'temp_map_list', and 'temp_org_map_list'.

This count sorting has some overhead. This one needs to manage two map_list, and it takes $2*n$ for one iteration. So, the time complexity for count sorting once is $O(n+64) + O(2*n) = O(n)$.

And this radix sort does count sorting for 15 times. So, it takes $O(15*n) = O(n)$.

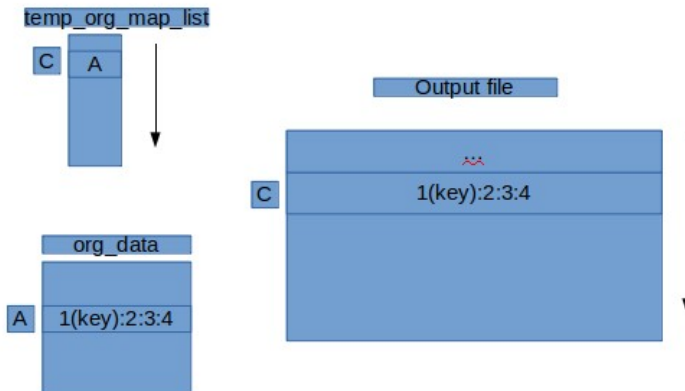
In conclusion, this sorting is done in $O(n)$.

3. Write 'hw1_output.txt' in order according to map_list.

```

143     for(int i=0;i<data_size-1;i++)
144     {
145         org_idx = temp_org_map_list[i];
146         fwrite(org_data[org_idx],1,org_data_len[org_idx]-1,fd_output);
147         fputc(NEW_LINE,fd_output);
148     }
149     org_idx = temp_org_map_list[data_size-1];
150     fwrite(org_data[org_idx],1,org_data_len[org_idx]-1,fd_output);

```



In 'temp_org_map_list', the original object index is stored. So, just writing the proper original object in order is enough to make 'hw1_output.txt'. It can decide to write which object in the following place. It just goes through the 'temp_org_map_list' from 0~(n-1) index, writing the corresponding original object to output file. **This code block takes $O(nmt)$.**

In summary,

first part takes $O(nmt)$, second part takes $O(n)$, last part takes $O(nmt)$.

So, the total complexity is $O(nmt) + O(n) + O(nmt) = \underline{O(nmt)}$

Thank you for reading.