Feedback — Problem Set-6

You submitted this quiz on **Sat 28 Nov 2015 8:50 AM PST**. You got a score of **4.00** out of **5.00**. You can attempt again in 1 minutes.

Question 1

Suppose we use a hash function h to hash n distinct keys into an array T of length m. Assuming simple uniform hashing --- that is, with each key mapped independently and uniformly to a random bucket --- what is the expected number of keys that get mapped to the first bucket? More precisely, what is the expected cardinality of the set $\{k: h(k) = 1\}$.

| Your Answer | Score | Explanation |
|------------------|----------------|--|
| ○ 1/m | | |
| O m/(2n) | | |
| 0 nl(2m) | | |
| ○ 1/n | | |
| \bigcirc m/n | | |
| ● n/m | 1.00 | Use linearity of expectation, with one indicator variable for each key. The probability that one key hashes to the first bucket is $1/m$, and by linearity of expectation the total expected number of keys that hash to the first bucket is just n/m . |
| Total | 1.00 / 1.00 | |

Question 2

You are given a binary tree (via a pointer to its root) with n nodes, which may or may not be a

binary search tree. How much time is necessary and sufficient to check whether or not the tree satisfies the search tree property?

| Your Answer | Score | Explanation |
|--------------------|--------|--|
| $\Theta(n)$ | 1.00 | For the lower bound, if there is a violation of the search tree property, you might need to examine all of the nodes to find it (in the worst case). |
| 0 | | |
| $\Theta(n \log n)$ | | |
| 0 | | |
| $\Theta(height)$ | | |
| 0 | | |
| $\Theta(\log n)$ | | |
| Total | 1.00 / | |
| | 1.00 | |

Question 3

You are given a binary tree (via a pointer to its root) with n nodes. As in lecture, let size(x) denote the number of nodes in the subtree rooted at the node x. How much time is necessary and sufficient to compute size(x) for every node x of the tree?

| Score | Explanation |
|----------------|---|
| × 0.00 | How many distinct quantities are you responsible for computing? |
| | |
| | |
| | |
| 0.00 / 1.00 | |
| | 0.00 / |

Question 4

Which of the following is not a property that you expect a well-designed hash function to have?

| Your Answer | Score | Explanation |
|---|---------------|---|
| The hash function should "spread out" most (i.e., "non-pathological") data sets (across the buckets/slots of the hash table). | | |
| • The hash function should "spread out" every data set (across the buckets/slots of the hash table). | ✓ 1.00 | As discussed in lecture, unfortunately, there is no such hash function. |
| The hash function should be easy to compute (constant time or close to it). | | |
| The hash function should be easy to store (constant space or close to it). | | |
| Total | 1.00 / | |
| | 1.00 | |

Question 5

Suppose we relax the third invariant of red-black trees to the property that there are no *three* reds in a row. That is, if a node and its parent are both red, then both of its children must be black. Call these *relaxed* red-black trees. Which of the following statements is *not* true?

| Your Answer | | Score | Explanation |
|--|----------|-------|--|
| There is a relaxed red-black tree that is not also a red-black tree. | | | |
| • Every binary search tree can be turned into a relaxed red-black tree (via some coloring of the nodes as black or red). | ~ | 1.00 | A chain with four nodes is a counterexample. |
| Every red-black tree is also a relaxed red-black tree. | | | |
| O The height of every relaxed red-black tree with n nodes is $O(\log n)$. | | | |

Total 1.00 / 1.00