

$$\boxed{\sigma \mapsto \sigma'}$$

Machine takes one step from state  $\sigma$  to state  $\sigma'$

$s; \rho \triangleright x$	$\mapsto s; \rho' \triangleleft V \ (\rho[x] = (V, \rho'))$
$s; \rho \triangleright (\text{con } tn \ cn)$	$\mapsto s; \rho \triangleleft (\text{con } tn \ cn)$
$s; \rho \triangleright (\text{lam } x \ M)$	$\mapsto s; \rho \triangleleft (\text{lam } x \ M)$
$s; \rho \triangleright (\text{delay } M)$	$\mapsto s; \rho \triangleleft (\text{delay } M)$
$s; \rho \triangleright (\text{force } M)$	$\mapsto s, (\text{force } \_); \rho \triangleright M$
$s; \rho \triangleright [M \ N]$	$\mapsto s, [\_ \ N]; \rho \triangleright M$
$s; \rho \triangleright (\text{builtin } bn)$	$\mapsto s; \rho \triangleright M \ (bn \text{ computes to } M)$
$s; \rho \triangleright (\text{builtin } bn \ \overline{M\overline{M}})$	$\mapsto s, ((\text{builtin } bn \ \_ \overline{M}), \rho); \rho \triangleright M$
$s; \rho \triangleright (\text{error})$	$\mapsto \blacklozenge$
$\cdot \triangleleft V$	$\mapsto \Box(V, \rho)$
$s, [\_ \ N]; \rho \triangleleft V$	$\mapsto s, [V \ \_]; \rho \triangleright N$
$s, [(\text{lam } x \ A \ M), \rho'] \_]; \rho \triangleleft V$	$\mapsto s; \rho'[x \mapsto (V, \rho)] \triangleright M$
$s, (\text{force } \_); \rho \triangleleft (\text{delay } M)$	$\mapsto s; \rho \triangleright M$
$s, ((\text{builtin } bn \ \overline{C} \_ \overline{M\overline{M}}), \rho'); \rho \triangleleft V$	$\mapsto s, ((\text{builtin } bn \ \overline{C}(V, \rho) \_ \overline{M}), \rho'); \rho' \triangleright M$
$s, ((\text{builtin } bn \ \overline{C} \_), \rho'); \rho \triangleleft V$	$\mapsto s; \rho' \triangleright M \ (bn \text{ computes on } \overline{C}(V, \rho) \text{ to } M)$

(c) CEK machine transitions for type-erased Plutus Core

Figure 23 (Continued): A CEK machine for type-erased Plutus Core