

# Fourier Method.

To compute  $\mathbb{E}(f(x)) := \mathbb{E}(f(\xi))$ :

$$\mathbb{E}(f(\xi)) = \int f \, dP_\xi \rightarrow \int T^*(Tf)(\omega) \, dP_\xi \rightarrow \int Tf(\omega) T(P_\xi)(\omega)$$

where  $T$  is suitable transf.

Remark: This method can be used when RHS is much simpler than LHS i.e. when  $Tf$ ,  $T(P_\xi)$  is well-known explicitly. E.g. Lévy's process.

(1) Preliminary:

$$\underline{\text{Def}}: \hat{f}(n) := \int_{\mathbb{R}^d} e^{inx} f(u) \, du, n \in \mathbb{C}.$$

Remark: Recall  $f \in L' \Rightarrow \hat{f} \in C_0$

$$\langle \hat{f}(n) \rangle = \frac{1}{2} \int_{\mathbb{R}^d} e^{inx} (f(x) - f(x - \frac{2}{n})) \, dx$$

Denote:  $L_{BC}^{(k)} = C_B^{(k)} \cap L^{(k)}$ .

Lemma: For  $f \in L'$ ,  $h \in L'$

$$i) g(x) = f(x) e^{inx} \Rightarrow \hat{g}(n) = \hat{f}(n + \omega)$$

$$g(x) = f(x-a) \Rightarrow \hat{g}(u) = e^{iau} \hat{f}(u)$$

$$\text{i)} g(x) = f(x/\lambda) \Rightarrow \hat{g}(u) = \lambda \hat{f}(\lambda u)$$

$$\text{ii)} g(x) = \overline{f(-x)} \Rightarrow \hat{g}(u) = \overline{\hat{f}(u)}$$

$$\text{iv)} \widehat{f \cdot h} = \hat{f} \cdot \hat{h}$$

$$\text{v)} g(x) = ix f(x) \in L' \Rightarrow \hat{f} \in C, (\hat{f})' = \hat{f}'$$

$$\text{vi)} f \cdot f' \in L_{loc}^1(\mathbb{R}) \Rightarrow \hat{f}'(u) = -iu \hat{f}(u).$$

$$\text{vii)} f \in C^2(\mathbb{R}), \text{ and } f \cdot f' \cdot f'' \in L_{loc}^1(\mathbb{R}) \\ \Rightarrow \hat{f} \in L_{loc}^1(\mathbb{R}).$$

Pf: Only prove vii):

$$\text{Note } \hat{f}''(u) = -u^2 \hat{f}(u) \in L_0 \text{ by vi).}$$

$$\Rightarrow (1+u^2) \hat{f}(u) \text{ is bad.}$$

$$\text{So: } \int \hat{f} \leq c \int \frac{1}{1+u^2} < \infty.$$

Thm.  $\leftarrow$  inversion formula

$$\tilde{f}(u) = \frac{1}{2\pi} \int_{-\infty}^{\infty} e^{-ixu} g(x) dx \text{ for } g \in L'$$

Thm: if  $f \in L_{loc}^1$ ,  $\hat{f} \in L'$   $\Rightarrow \forall x \in \mathbb{R}$ .

$$f(x) = (\hat{f})^*(x)$$

(2) Optional pricing:

Def: For  $f_K(x) = e^{-Rx} f(x)$ ,  $M_x(a) = \mathbb{E}[e^{ax}]$ .

$\mathcal{Z} := \{K \in \mathcal{K}' : f_K \in L^1_{loc}, \hat{f}_K \in L^1\}$ .

$\mathcal{I} := \{K \in \mathcal{K}' : M_x(K) < \infty\}$ .

Thm.  $R := \mathcal{Z} \cap \mathcal{I} \neq \emptyset$ . Let  $R \in \mathcal{R}$ . Then:

$$I(f, x) = \frac{1}{2\pi} \int_{\mathcal{K}'} M_x(R-iu) \hat{f}(u) du.$$

If: Note  $\hat{f}_K \in \mathcal{L}'$ . by inversion:

$$f_K(x) = \frac{1}{2\pi} \int_{\mathcal{K}'} e^{-ixu} \hat{f}_K(u) du.$$

$$\begin{aligned} \text{So: LHS} &= \int e^{Rx} f_K(x) dP_x \\ &= \int e^{Rx} \cdot \frac{1}{2\pi} \int e^{-ixu} \hat{f}_K(u) dP_x \\ &= \frac{1}{2\pi} \int M_x(R-iu) \hat{f}(u) du \end{aligned}$$

follows from Fubini: Thm.

$$\begin{aligned} \left| \int e^{Rx} / e^{inx} \hat{f}_K(u) dP_x \right| &\leq \int e^{Rx} \|\hat{f}\|_{\mathcal{L}'} dP_x \\ &\leq M_x(R) \|\hat{f}\|_{\mathcal{L}'} \end{aligned}$$

Rmk: i) Actually,  $R \neq \emptyset \Rightarrow f \in C_c(\mathcal{K}')$ .

For L1-convergence case:

Set  $\mathcal{I}' := \{K \in \mathcal{K}' : f_K \in L^1(\mathcal{K}')\}$ .

$\mathcal{J}' := \{K \in \mathcal{K}' : M_x(R) < \infty, M_x(R-i \cdot) \in L^1\}$

Then we can replace " $R \in R \neq \emptyset$ "

by " $R \in R' := \mathcal{J} \cap \mathcal{J}' \neq \emptyset$ " but it

implies  $\ell^{R_n}(\mu_x \llcorner \Omega)$  admits a conti.

by Lebesgue Integrality  $\ell, \mu'$ , a  
trading on conti. of  $f$  as hist.

of  $X$  / integrability of  $\hat{f}$  and  $\mu_X$ )

ii) Set  $Z_{\min} := \{R \in \mathcal{R} : f_R \in L^1(\mathbb{R})\}$ .

$J_{\min} := \{R \in \mathcal{R} : \mu_X(R) < \infty\}$ .

The minimal assumption of them  
above is:  $R \in Z_{\min} \cap J_{\min} \neq \emptyset$ .

Then the formula exists as a  
pointwise limit.

iii) The only hard part to check is  
the condition  $f_R \in L^1(\mathbb{R})$ . As for  
" $\hat{f}_R \in L^1(\mathbb{R})$ ". We have:

Lem.  $f \in H^1(\mathbb{R}) \Rightarrow \widehat{f'}(n) = -i n \widehat{f}(n), \widehat{f}, \widehat{f'} \in L^2$ .

Lem.  $f \in H^1(\mathbb{R}) \Rightarrow \widehat{f} \in L^1(\mathbb{R})$ .

$$\underline{\text{Pf:}} \int |\tilde{f}|^2 (1+|m|^2) = \int |\tilde{f}|^2 + |\tilde{f}^*|^2 < \infty$$

$$\int_0^T |\tilde{f}| \leq (\int |\tilde{f}|^2 (1+|m|^2))^{\frac{1}{2}}.$$

$$(\int \frac{1}{(1+|m|^2)})^{\frac{1}{2}} < \infty.$$

Applications:

① Consider  $(S_t)$  price of assets modeled as exponential growth.  $S_t = S_0 e^{X_t}$ .

Next, assume  $r=0$ .  $S_t$  is IP-mart.

$\Rightarrow E(S_T) = f(X_T + \log S_0)$  payoff  $F$  of  $S_T$   
where  $f = F \circ \exp$ . We want to compute:

$$\begin{aligned} E(F(S_T)) &= E(f(X_T + \log S_0)) \\ &= \frac{1}{2\pi} \int M_{X_T + \log S_0}(R-iu) \hat{f}(u+iR) du \\ &= \frac{1}{2\pi} \int S_0^{R-iu} M_{X_T}(R-iu) \hat{f}(u+iR) du. \end{aligned}$$

Next, we consider vanilla option, i.e. it  
can't be exercised earlier before expire  
time  $T$  and only depends on  $S_T$ .

e.g.) Call option

$f(x) = (e^x - k)_+$  is payoff func.

$$\Rightarrow \hat{f}(z) = k^{1+i^z} / (iz + i)$$

Note  $\{R \in \mathbb{C}, n\} \subset \{f_R \in L_1 \cap L^2\}$

And  $\bar{f}_R(x) = e^{-Rx} \langle e^{-Cx} - Re^{-Rx} + Rk \rangle I_{\{x > \ln k\}}$

in weak sense.  $\Rightarrow f_R \in L^2$  as well.

So  $\hat{f}_R \in L^1$  since  $f_R \in H^1(\mathbb{R})$ .

$$\Rightarrow (1, n) \subset \mathcal{I}$$

Krk: For put option  $f(x) = (k - e^x)_+$ .

We have  $(-\infty, 0) \subset \mathcal{I}$ .

ii) Fourier method is quite suitable  
for exp. Levy process. Since we  
know  $m_x$  explicitly.

Assume its triplet  $(b, c, \nu)$  and

$m_x < \infty$  on  $\{n \in \mathbb{C} : Re(n) \in [a, b]\}$ .

i.e.  $\mathcal{I} = [a, b]$ . So if consider  $f$  is  
call option above and  $[a, b] \cap (1, n)$   
 $\neq \emptyset$ . Then we can use the formula.

Rmk:  $[a, b] \supset [0, 1]$  in fact.

### ① Computation of Greek:

We want to find delta A for ①:

$$Af(x, s_0) = \frac{\partial}{\partial s_0} E(f(x_t + \log s_0))$$

Then: If "i)  $\lim_{t \rightarrow R^-} m_{x_t}(R-i, n) \in L'$ " and  $\hat{f}(c+iR)$  is bad" or "ii)  $\lim_{t \rightarrow R^-} \hat{f}(c+iR, t) \in L'$  and  $m_{x_t}(R-i, \cdot)$  is bad" holds. Then:

$$Af(x, s_0) = \frac{1}{2\pi} \int_{\mathbb{R}} S_0^{\frac{R-1-i}{2}} M_{x_t}(R-i, n) \frac{\hat{f}(c+iR)}{(R-i, n)^{-1}}$$

Pf: Apply DCT. It follows from:

$$\left| \frac{\partial}{\partial s_0} \right| \leq C(1+|n|) |m_{x_t}(R-i, n)| |\hat{f}(c+iR)|$$

Rmk: Condition i) implies  $m_{x_t}$  admits a continuity  $\in C'$ .

### ② Multidimension:

For  $f: \mathbb{R}^k - \mathbb{R}'$  profit and  $(x_t)$  k-lim

v.v.'s. Replace " $R \cdot x$ " by " $\langle R, x \rangle$ ",  $R \in \mathbb{R}^k$

Set  $R, R' \in \mathbb{R}^k$  as above.

Thm. If  $R \neq \emptyset$  or  $R' \neq \emptyset$ .  $R \subset R' \cup R$ . Then:

$$I(f, x) = (2\pi)^{-k} \int_{R^k} M_x(R-iu) \hat{f}(u+ik) du$$

Cor.  $E \in F(S_T)$

$$= (2\pi)^{-k} \int_{R^k} e^{(R-iu, \log s_0)} M_x(R-iu) \hat{f}(u+ik) du$$

Emb. It will suffer curse of dim. So

We mostly consider  $k \leq 3$ .

Lem.  $f_i : \mathbb{R}' \rightarrow \mathbb{R}'$ . Let  $f(x) = \sum_{i=1}^k f_i(x_i)$ . Then:

$$\mathcal{Z}' \supset \bigcirc_{i=1}^k \mathcal{Z}'_i.$$

C.G. C (all option)

$$f(x) = (e^{x_1} \wedge \cdots \wedge e^{x_k} - k)_+$$

$$\Rightarrow \hat{f}(z) = -k^{1+i \sum z_k} / (-1)^k (1 + i \sum z_k) \prod_{i=1}^k \pi(z_i)$$

$$\mathcal{Z}' = \{ k_i > 0, \forall i \in \mathbb{N}, \sum_{i=1}^k k_i > 1 \}.$$

$$\text{Put: For put option } (k - e^{x_1} \wedge \cdots \wedge e^{x_k})_+$$

$$\Rightarrow \hat{f}(z) = k^{1+i \sum z_k} / (1 + i \sum z_k) \prod_{i=1}^k \pi(z_i)$$

$$\text{and } \mathcal{Z}' = \{ k_i < 0, \forall i \in \mathbb{N} \}.$$

(3) Fast Fourier Transf.:

Set  $\sigma = \log s_0$ . Next. We want to implement

the formula into computer:

$$O_f(s) = \overline{E} e^{F(s)} = \frac{e^{ks}}{2\pi} \int_{k'} e^{-ins} M_{X_1} \circ R - ins \hat{f}(c_n), n.$$

$$= e^{ks/2\pi} \cdot \int_{k'} e^{-ins} \chi(c_n) dn.$$

$$\approx \frac{e^{ks}}{2\pi} \int_a^b e^{-ins} \chi(c_n) dn.$$

(i.e. fix on finite interval to approx.)

Set  $\eta = \frac{b-a}{N-1}$ .  $u_k = a + k\eta$ ,  $0 \leq k \leq N-1$ .

$$\int_a^b e^{-ins} \chi(c_n) \approx \eta \left[ \frac{e^{-ins}\chi(a)}{2} + \sum_{k=1}^{N-1} e^{-ins} \chi(u_k) + e^{-ibs}\chi(b)/2 \right] = \sum_0^{N-1} e^{-ius} \chi(u_k)$$

We also choose a uniform grid on

$$S_{\text{hommin}} \left[ -\frac{\lambda N}{2}, \frac{\lambda N}{2} - \lambda \right] : \text{let } \beta = -\frac{\lambda N}{2}.$$

$s_j = \beta + \lambda j$ ,  $0 \leq j \leq N-1$ . So: we want to

compute:  $\sum_{k=1}^{N-1} e^{-i(c_n+k\lambda)\lambda j} e^{i\beta u_k} \chi(u_k) \eta$ ,  $j \leq N-1$

We choose  $\lambda, \eta$ , so  $\lambda\eta = 2\pi/N$ . (Nyquist rel.)

Let  $\phi_j := \sum_{k=0}^{N-1} e^{-i\frac{2\pi}{N}kj} \chi_k$ ,  $\chi_k = e^{i\beta u_k} \chi(u_k)$ .

So:  $O_f(s_j) \approx e^{\beta k - \lambda j(c_k - i\eta)} \sum \phi_j / 2\pi$ .

Rank:  $\Phi = (\phi_0 \dots \phi_{N-1})$  is discrete Fourier

transf. of vector  $\ell = (\ell_0, \dots, \ell_{N-1})$

We want to find a fast algo. to compute  $\varphi$ .

Next, we introduce fast Fourier transf.

(FFT). We can reduce the computation cost  $\sim N^2$  above to cost  $\sim N \log N$ .

Let  $\omega_N := e^{-2\pi i/N}$  and define the  $N \times N$ -matrix  $T_N$  by

$$T_N := \begin{pmatrix} \omega_N^0 & \omega_N^0 & \omega_N^0 & \cdots & \omega_N^0 \\ \omega_N^0 & \omega_N^1 & \omega_N^2 & \cdots & \omega_N^{N-1} \\ \omega_N^0 & \omega_N^2 & \omega_N^4 & \cdots & \omega_N^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \omega_N^0 & \omega_N^{N-1} & \omega_N^{2(N-1)} & \cdots & \omega_N^{(N-1)(N-1)} \end{pmatrix}. \Rightarrow \varphi = \text{Tr } \ell$$

Lem.  $\ell \in \mathbb{C}^N$ .  $\varphi = \text{Tr } \ell$ . Let  $\ell' := (\ell_0, \ell_2, \dots, \ell_{2n-1})$ .

$\ell'' := (\ell_1, \ell_3, \dots, \ell_{2n})$ .  $\varphi' := \text{Tr } \ell'$

$\varphi'' := \text{Tr } \ell''$ . And

$D_n := \text{diag}(\omega_{2n}, \dots, \omega_{2n}^{n-1})$ . Then:

$\varphi' = c + \lambda$ .  $\varphi'' = c - \lambda$  where  $c = \text{Tr } \ell'$

and  $\lambda = D_n \text{Tr } \ell''$ .

Pf: Using  $W_n^{j\ell} = W_{2n}^{2j\ell}$

It forms a divide-and-conquer algo.:

**Algorithm 7.26 (FFT).** Assume that  $N = 2^J$ ,  $J \geq 1$ . Given  $\phi \in \mathbb{C}^N$ , apply the following recursive algorithm to compute its discrete Fourier transform  $\Phi = T_N \phi$ :

1. If  $N = 2$  go to 2, otherwise: split  $\phi$  into  $\phi'$  and  $\phi''$  as in Lemma 7.25, apply the FFT to compute  $c = T_{N/2} \phi'$ ,  $d = D_{N/2} T_{N/2} \phi''$  and return  $\Phi = (\Phi', \Phi'')$  given by  $\Phi' = c + d$  and  $\Phi'' = c - d$ .
2. If  $N = 2$  compute  $\Phi = T_2 \phi$  directly.

Let  $N_k = 2^k$ ,  $1 \leq k \leq J$ . implement  $J$  times.

Lem.  $N = 2^k$ .  $C$  is cost of floating point operations (addition, multipli. ...). Then:

$$\text{Computation work } W(N) \leq C + \frac{3}{2}(j_2 N + \frac{1}{2})N.$$

Pf: Note for FFT in  $N$ -dim.  $W_C$

need one vector addition and  
one subtraction in  $N/2$ -dim.

With one elementwise multiplicate  
of two vectors in dim  $N/2$ .

$$\Rightarrow W(N) \leq 2W(N/2) + \frac{3}{2}CN.$$

$$W(2) \leq 4C.$$

$$\text{Let } \tilde{W}(N) = W(N)/CN.$$

$$\begin{aligned} \text{So: } \tilde{W}(N) &\leq \tilde{W}(N/2) + \frac{3}{2} \\ &\leq (k-1)\frac{3}{2} + \tilde{W}(2) \\ &\leq \frac{3}{2}k + \frac{1}{2} \end{aligned}$$

$$\text{where } k = \log_2 N.$$

Rmk: i) we can also compute the inverse discrete Fourier transf.

ii) Variant of FFT exists i.e.  $N$  doesn't need to be  $2^k$ .

(4) Gsinc-Series expansion:

Note for even func.  $f$ . We have:

$$\hat{f}(z) = 2 \int_0^\infty f(x) \cos(zx) dx.$$

Assume  $\mathcal{Z}_T$  density of  $X_T$  decay very fast to 0. So: WLOG.  $\text{Supp}(\mathcal{Z}_T) \subset [0, 2]$ .

Rmk: Recall for locally cpt abelian group.

i. the dual group  $\widehat{G}$  consists of all characters of  $G$ . i.e. all anti. group homo. from  $G$  to  $\mathbb{T} \subset \mathbb{C}^*$ .

$$a) G = \mathbb{K}' \Rightarrow \widehat{G} \cong \mathbb{K}' \text{ i.e. } \chi_{ex} = e^{iux} \quad u \in \mathbb{K}'.$$

$$b) G = [-\pi, \pi] \Rightarrow \widehat{G} \cong \mathbb{Z}, \text{ i.e. } \chi_{ex} = e^{inx} \quad n \in \mathbb{Z}.$$

Let  $\mu$  denote Haar measure on  $G$ .

$$\hat{f}(x) := \int_G f(x) \overline{\chi_{ex}} \mu(dx) \in L^2(\widehat{G}) \text{ is}$$

Fourier transf. of  $f \in L^1([a, b]; \mathbb{C})$ .

a)  $h = ik' \Rightarrow n = \omega$ . So  $\hat{f} = \int_{[a, b]} e^{-inx} f(x) dx$

b)  $h = [-\pi, \pi] \Rightarrow n = \lfloor \frac{1}{2\pi} \int_{[-\pi, \pi]} f(x) dx \rfloor$ .  
 $= \frac{1}{2\pi} \int_{-\pi}^{\pi} f(x) e^{-inx} dx$ .  $\hat{f} = \frac{1}{2\pi} \int_{-\pi}^{\pi} f(x) e^{inx} dx$ .

Consider cosine represent. of  $g = z\tau$ :

$$g(\theta) = \sum_k A_k \cos(k\theta) + \frac{i}{2} A_0. A_k = \frac{1}{2\pi} \int_{-\pi}^{\pi} g(x) \cos(kx) dx$$

Rmk: Fix  $z$  supp. on  $[a, b]$ . By change vari.:

$$g(\theta) = \sum_k A_k \cos(kz \frac{\theta-a}{b-a}) + \frac{i}{2} A_0 \text{ with}$$

$$A_k = \frac{1}{b-a} \int_a^b z(x) \cos(kz \frac{x-a}{b-a}) dx.$$

If we know  $\varphi = \tilde{z}$  but not  $z$ . First

$$\text{Let } \int_a^b e^{inx} z(x) dx = q_i \approx \varphi.$$

$$\Rightarrow A_k = \frac{1}{b-a} \operatorname{Re} \left( q_i \cdot e^{-ikz \frac{\theta-a}{b-a}} \right) \stackrel{-ikz \frac{\theta-a}{b-a}}{\longrightarrow}$$

$$\approx \frac{1}{b-a} \operatorname{Re} \left( q_i \cdot e^{-ikz \frac{\theta-a}{b-a}} \right) \stackrel{-ikz \frac{\theta-a}{b-a}}{\longrightarrow} =: F_k$$

$$\text{So: } z(x) \approx \tilde{z}(x) = \sum_{k=0}^{N-1} F_k \cos(kz \frac{x-a}{b-a}) + \frac{1}{2} F_0.$$

Rmk: There're three different errors:

a) Truncated the integral on  $[a, b]$ .

b) Replace  $A_k$  by  $F_k$ .

c) Truncate infinite sum by finite

sum  $\sum_0^{n-1}$ .

So next we can consider the option

valuation of  $f: C(S, T) = e^{-rT} \int_{-\infty}^{+\infty} f(x) Z_T dx$

Replace  $Z_T$  by approx. above. We have:

$$C(S, T) \approx e^{-rT} \sum_0^{n-1} f(x) \rho_T \left( \frac{kx}{b-a} \right) e^{-ikx \frac{a}{b-a}} c_k.$$

$$\text{where } \rho_T = \tilde{Z}_T, c_k = \frac{2}{b-a} \int_a^b f(x) \cos(kx \frac{a}{b-a}) dx$$

Rmk: i) The error of approx. mostly depend on smoothness of density  $Z_T$ :

a)  $Z_T$  smooth on  $[a, b]$   $\Rightarrow$   $Z_T$  decay exponentially  $\sim e^{-\nu c(N-1)}$

b)  $Z_T \in C^1$  on  $[a, b] \Rightarrow$   $Z_T$   $\|/\|$  decay algebraically  $\sim (N-1)^{-\beta}$

ii) Choice of  $a, b$  will depend on cumulants  $c_n$  of dist.:

(7.41)

$$a = c_1 - L \sqrt{c_2 + \sqrt{c_4}}, \quad b = c_1 + L \sqrt{c_2 + \sqrt{c_4}}$$

with  $L = 10$ .

ii) For use the approx.' We need to compute  $C_k$  at first. In some case  $C_k$  is explicit.

e.g. (Call option)

$$f(x) = (K(e^x - 1))_+ \cdot x = \log(s\tau/k)$$

log-moneyness. Then:

$$C_k^{\text{call}} = \frac{2}{b-a} K (\chi_k(0, b) - \psi_k(0, b)),$$

with

$$\begin{aligned} \chi_k(c, d) := & \frac{1}{1 + \left(\frac{k\pi}{b-a}\right)^2} \left[ \cos\left(k\pi \frac{d-a}{b-a}\right) e^d - \cos\left(k\pi \frac{c-a}{b-a}\right) e^c + \right. \\ & \left. + \frac{k\pi}{b-a} \sin\left(k\pi \frac{d-a}{b-a}\right) e^d - \frac{k\pi}{b-a} \sin\left(k\pi \frac{c-a}{b-a}\right) e^c \right] \end{aligned}$$

and

$$\psi_k(c, d) := \begin{cases} \left( \sin\left(k\pi \frac{d-a}{b-a}\right) - \sin\left(k\pi \frac{c-a}{b-a}\right) \right) \frac{b-a}{k\pi}, & k \neq 0, \\ d - c, & k = 0. \end{cases}$$

For the put-option, we obtain

$$C_k^{\text{put}} = \frac{2}{b-a} K (\psi_k(a, 0) - \chi_k(a, 0)).$$

$k \neq 0$ : It's only valid for options written in log-moneyness.