

(2) General principle:

Consider X is polish space. $I: X \rightarrow \mathbb{R}^{20}$ is a proper rate func. i.e. $K_\alpha = \{I(x) \leq \alpha\}$ is cpt for \mathcal{H}^1 . and I is l.s.c.

Def: For (P_n) seq of laws on X . it satisfy large deviation principle (LDP) on X at speed n . w.r.t. I if:

$$i) \lim_{n \rightarrow \infty} \frac{1}{n} \log P_n(\omega) \leq -\inf_c I(c), \quad \forall c \subset X.$$

$$ii) \lim_{n \rightarrow \infty} \frac{1}{n} \log P_n(\omega) \geq -\inf_o I(o), \quad \forall o \subset X.$$

Rmk: i) Note for i.i.d seq case.

it satisfies LDP.

ii) Note $P_n(x) > 1 \Rightarrow \inf I(x)$ must

be 0. for (P_n) .

Thm: For (P_n) satisfies LDP at speed n .

w.r.t I on X . Then. $\forall \epsilon < \infty. \exists D^\epsilon$
 $\subset X$. st. $P_n(D^\epsilon) \geq 1 - e^{-n\epsilon}. \forall n$.

Pf: consider $A^\epsilon = \{I(x) < \epsilon + 2\}$. cpt

Suppose A_k is finite union of open balls with $r = 1/k$ cover A^ϵ .

By LDP: $\lim_{n \rightarrow \infty} \frac{1}{n} \log P_n(A_k^c) \stackrel{I \geq \ell+2}{\underset{\text{on } A_k^c}{=}} -(\ell+2)$

$\Rightarrow \exists n_0 = n_0(\ell), \forall n \geq n_0: P_n(A_k^c) \leq e^{-n(\ell+1)}$

WLOG. Set $n_0(k) \geq k, \forall k \geq 1$.

$\Rightarrow P_n(A_k^c) \leq e^{-\ell k} \cdot e^{-k}$.

On the other hand. For each $k \geq 1$.

$\exists (B_{k,i})_{i \in \text{m}(k)}$ seq of opt sets. St.

$P_i(A_{k,i}^c) \leq e^{-k} \cdot e^{-i\ell}, \forall i \in \text{m}(k)$.

Let $D^\ell = \bigcap_k (\bar{A}_k \cup \bigcup_{j=1}^{\text{m}(k)} B_{k,j})$

$P_n(D^\ell)^c \leq \sum_{k \geq 1} P_n(A_k^c \cap (\cap B_{k,j})) \leq e^{-\ell k} \sum_k e^{-k}$

Since $P_n(A_k^c \cap (\cap B_{k,j})) \leq \begin{cases} P_n(A_k^c), & n \geq n_0(k) \\ P_n(B_{k,n}), & n < n_0(k) \end{cases}$

Thus (Product of LDP)

If $(P_n), (Q_n)$ are two seq of laws

satisfy LDP w.r.t. $I(x)$ and $J(y)$ on X .

and Y . Then $R_n := P_n \times Q_n$ satisfies LDP

w.r.t $I(x, y) = I(x) + J(y)$ on $X \times Y$.

Pf: i) Consider $Z = (x, y) \in X \times Y$.

By I.S.C. of I . $\exists U_x$ nbd of x .

St. $J(x') \geq I(x) - \varepsilon, \forall x' \in U_x$.

By separable. $\exists U_x$ open. $U_x \subset \bar{U}_x \subset U_x$.

$\Rightarrow \lim_{n \rightarrow \infty} \frac{1}{n} \log P_n(U_x) \leq -I(x) + \varepsilon$.

Similarly, $\exists V_Y$ of y . Set $N_Z = N_X \times V_Y$.

$$\Rightarrow \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log R_n(N_Z) \leq -k(z) + \epsilon.$$

2) For $D \subset X \times Y$, opt evt. cover D
by finite nbd's N_Z (ns above).

$$R_n(\tilde{V}N_Z) \leq N \max_{1 \leq j \leq N} R_n(N_Z)$$

Note N will disappear. after taking
logarithm and divide n . set $n \rightarrow \infty$

$$\Rightarrow \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log R_n(\tilde{V}N_Z) \leq -\inf_D k(z) + \epsilon.$$

3) For general $C \subset X \times Y$.

Note $\exists A^C \subset X$, $B^C \subset Y$, st.

$$(P_n(A^C))^c \leq e^{-n\varepsilon}, \quad (P_n(B^C))^c \leq e^{-n\varepsilon}.$$

$$\text{Set } C^C = A^C \times B^C. \quad C = C \cap C^C + C \cap C^{C^C}.$$

$$\Rightarrow \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log R_n(C) \leq \max_C (-\inf k(z), -\varepsilon)$$

$$\leq \max_C \{-\inf k(z), -\varepsilon\} \xrightarrow[n \rightarrow \infty]{C} -\inf k(z),$$

4) For lower bound, since it's local.

find $N_X \times V_Y = N$, nbd of $z = (x, y)$.

By LDP of P_n on \mathcal{A}_n , we have:

$$\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log R_n(N) \geq -k(z) \quad \forall z \in \Omega.$$

Thm: (Contraction principle)

$F: X \rightarrow Y$. is contr. between polish spaces.

If $\{P_n\}$ satisfies LDP with $I(\cdot)$ on X .

Then. $\{Q_n = P_n \circ F'$ satisfies LDP on Y

w.r.t. $\bar{I}(y) = \inf_{x \in F^{-1}(y)} I(x)$.

Pf: Note F' retains open/close. easy to check.

Thm (Varadhan's Thm)

If $\{P_n\}$ satisfies LDP on X w.r.t. $I(x)$.

$F \in C_B^{(k)}(X)$. Then $a_n := \int_X e^{-nF(x)} dP_n(x)$.

satisfies $\lim_{n \rightarrow \infty} \frac{1}{n} \log a_n = \sup_x (F(x) - I(x))$.

Rank: It's some kind of Laplace transf.

Pf: procedure: nbd of point \Rightarrow opt set \Rightarrow general.

i) $\forall x \in X. \exists N_x$. nbd of x . s.t.

$$|F(x') - F(x)| \leq \varepsilon. \quad I(x') + \varepsilon \geq I(x). \quad \forall x' \in N_x.$$

$$\text{Let } a_n(A) = \int_A e^{-nF(x)} dP_n(x)$$

$$\text{Note } a_n(N_x) \leq e^{-n(F(x)+\varepsilon)} P_n(N_x)$$

$$\Rightarrow \overline{\lim_{n \rightarrow \infty}} \frac{1}{n} \log a_n(N_x) \leq (F(x) - I(x)) + 2\varepsilon.$$

ii) $\forall D$ opt in X . similar as before

$$\overline{\lim_{n \rightarrow \infty}} \frac{1}{n} \log a_n(D) \leq \sup_x (F(x) - I(x)) + 2\varepsilon.$$

3') Find $\limsup_{n \rightarrow \infty} \ln c_n(k^n)$. s.t. $(P_n c_n(k^n))^c \leq e^{-n\lambda}$

$$\Rightarrow \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log \ln c_n(k^n) \leq \|f\|_{\infty} - \lambda.$$

$$\text{So } \lim_{n \rightarrow \infty} \frac{1}{n} \log \ln c_n \leq \max \left\{ \|f\|_{\infty} - \lambda, \sup_{x \in X} (F(x) - I(x)) + 2\varepsilon \right\}$$

$$\xrightarrow[\varepsilon \rightarrow 0]{} \sup_{x \in X} (F(x) - I(x))$$

4') Fix lower bound. $\forall x \in X. \exists N_x$ nbd of x .

$$a_n \geq \ln c_n(N_x) \geq e^{n(F(x) - \lambda)} P_n(N_x)$$

$$\Rightarrow \liminf_{n \rightarrow \infty} \frac{1}{n} \log a_n \geq \sup_x (F(x) - I(x)) - \varepsilon.$$

Remk: Note bdd and u.s.c. are for upper bound. while l.s.c. is for lower.

Def: $Q_n(A) := \frac{\int_A e^{-nF(x)} \lambda P_n(dx)}{\int_X e^{-nF(x)} \lambda P_n(dx)}$. for $A \in \mathcal{B}_X$.

Thm. Fix c_{P_n} satisfies LDP with rate function $I(x)$ and $F \in C_{\mathcal{B}(X)}$. Then c_{Q_n} also satisfies LDP on X . w.r.t $J(x) := \sup_x (F - I)$

$$= (F(x) - I(x))$$

Pf: Note that $Q_n(x) = a_n(x) / n$.

argue as before. $\Rightarrow \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \log a_n(x) \leq -\inf_c J$.