Enhancing Ethereum PoA Clique Network with DAGbased BFT Consensus

Yongrae Jo and Chanik Park

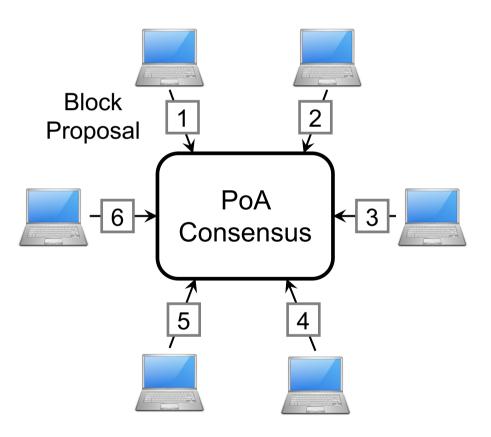
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Proof of Authority (PoA)

Authorities



- Consensus for permissioned blockchain
 - Only a few messages for block agreement
 - Efficient than proof-of-work (PoW)



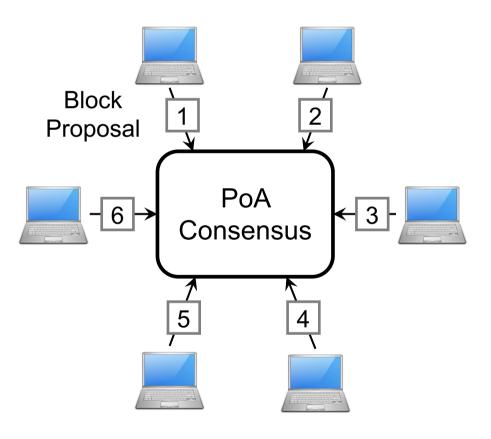






Proof of Authority (PoA)

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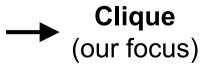


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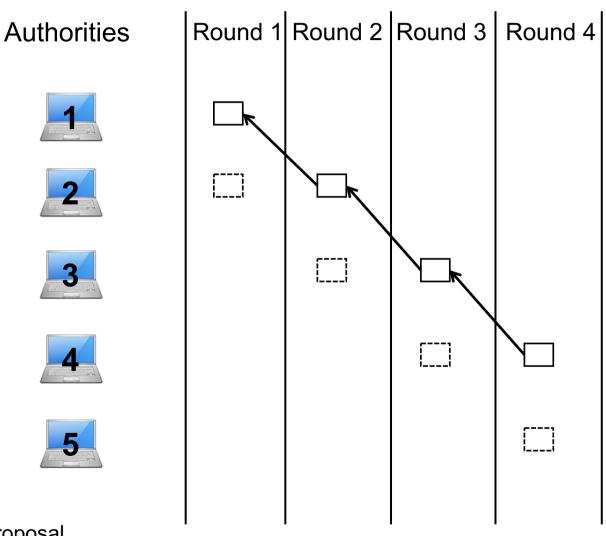








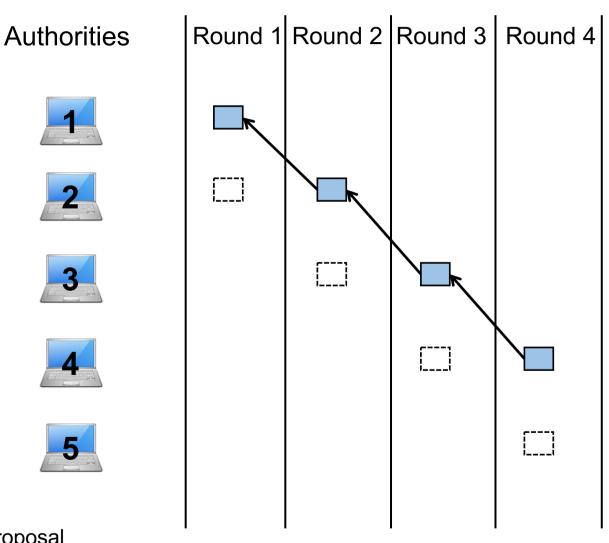




- Round-based
 - (rotating) leader & round timeout Δ
- Multiple block proposals in a round
 - by a leader
 - by non-leaders* with <u>random delays</u>

Block proposal

by leader



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- Single block acceptance in a round
 - Normal: a leader block

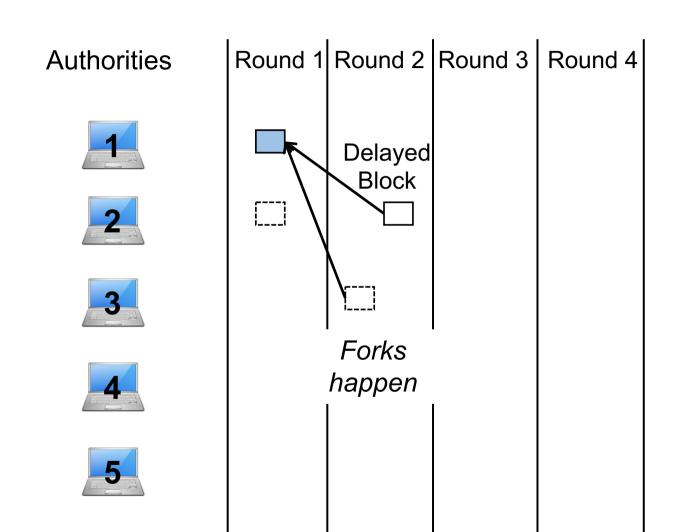
Block proposal

by leader

Accepted

by non-leader

*#non-leaders = up to $\left| \frac{N}{2} \right| - 2$



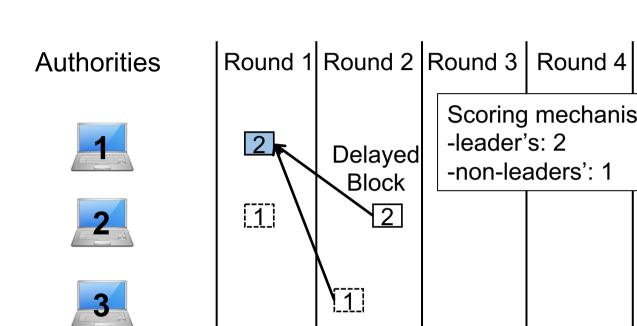
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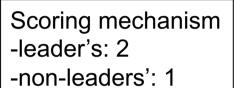
by leader



Accepted







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Block proposal

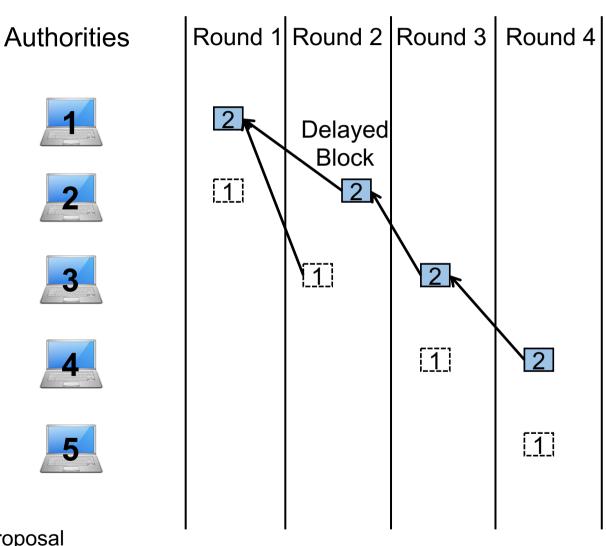
by leader



Accepted

Forks

happen



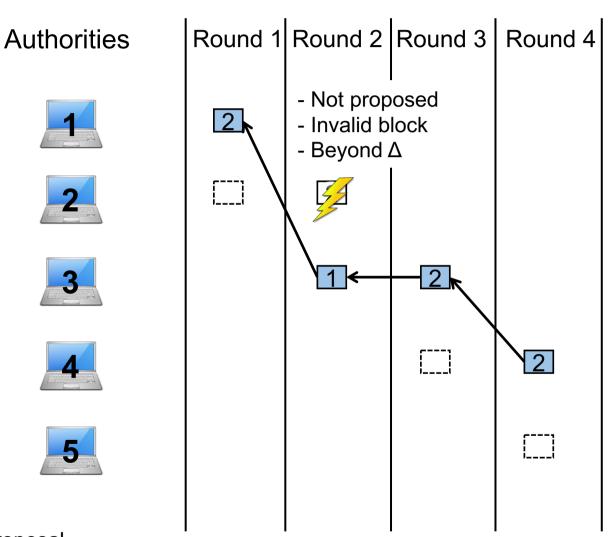
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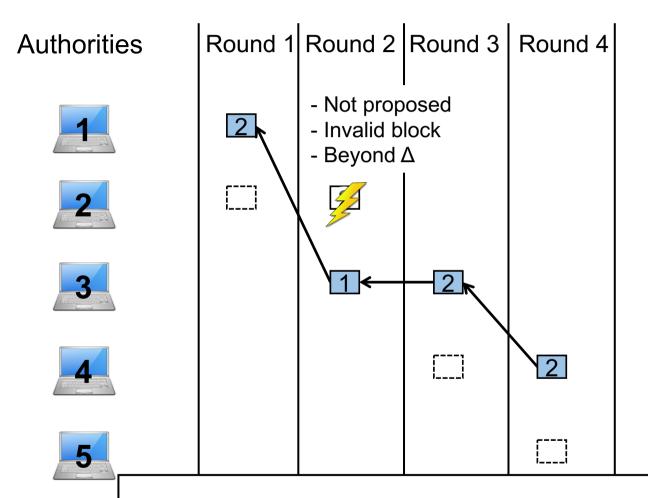
Block proposal

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Accepted

PoA Clique: Issues



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Multiple block proposals, but a single block acceptance

Block proposal

____ by leader

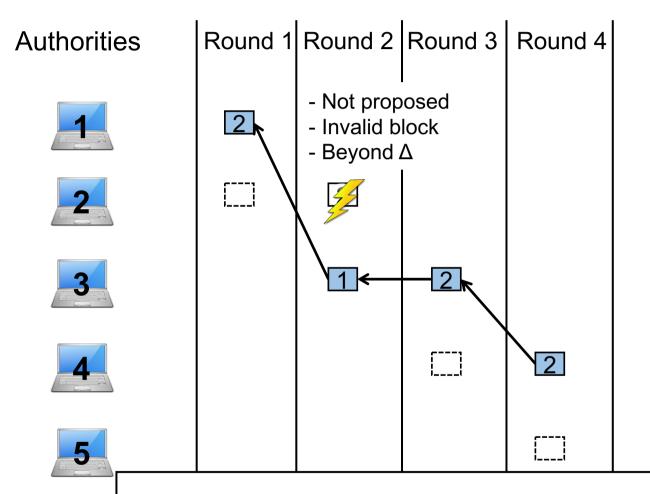


Accepted

by non-leader

*#non-leaders = up to $\left[\frac{N}{2}\right] - 2$

PoA Clique: Issues



- Round-based
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Multiple block proposals, but a single block acceptance

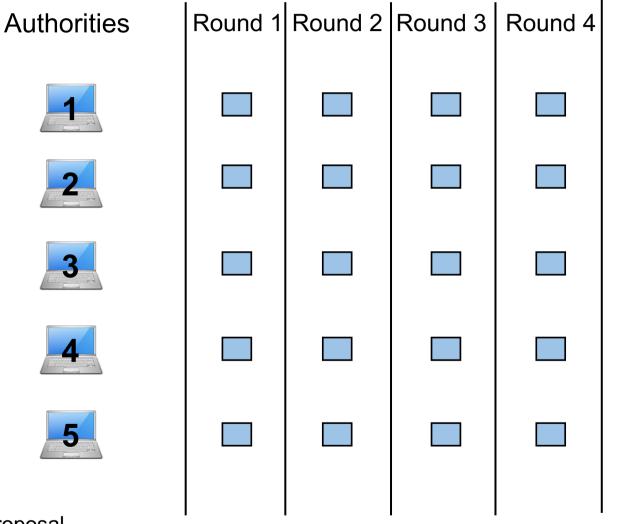
Block proposal

by leader

by non-leader

Limited throughput & Wasted resources

Our Idea



- Round-based
 - leader-less & round timeout Δ
- Multiple block proposals in a round
 - By all authorities

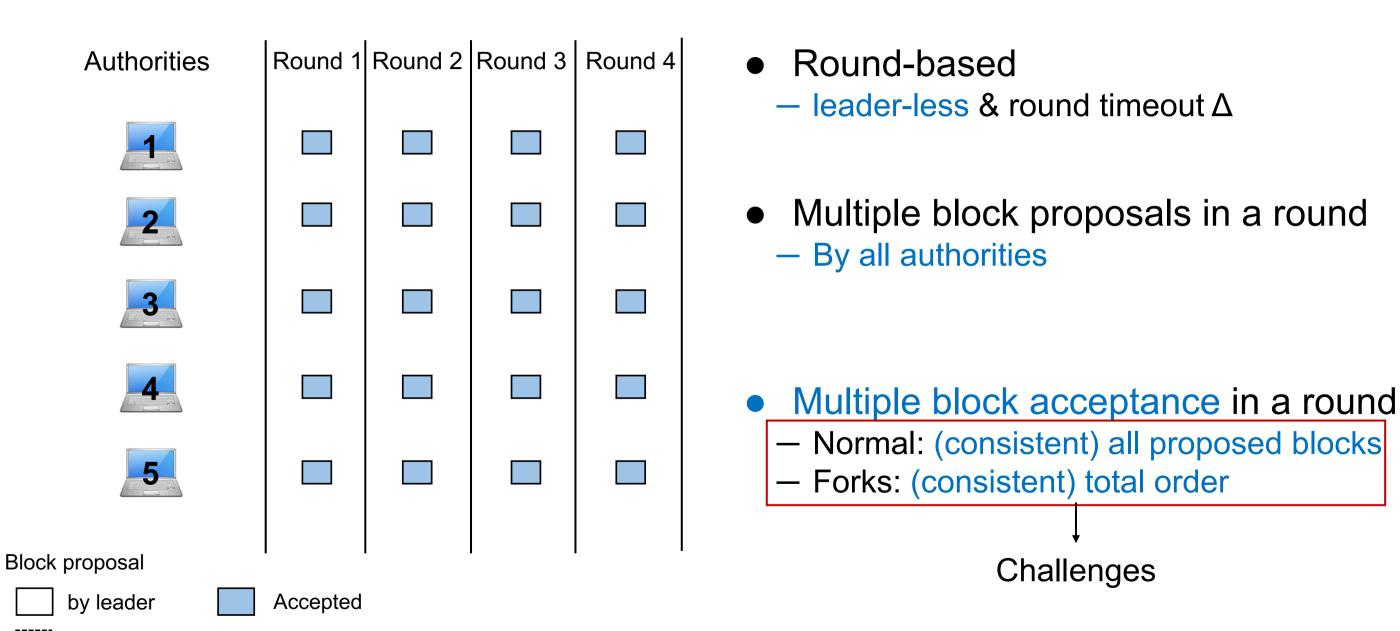
- Multiple block acceptance in a round
 - Normal: (consistent) all proposed blocks
 - Forks: (consistent) total order

Block proposal

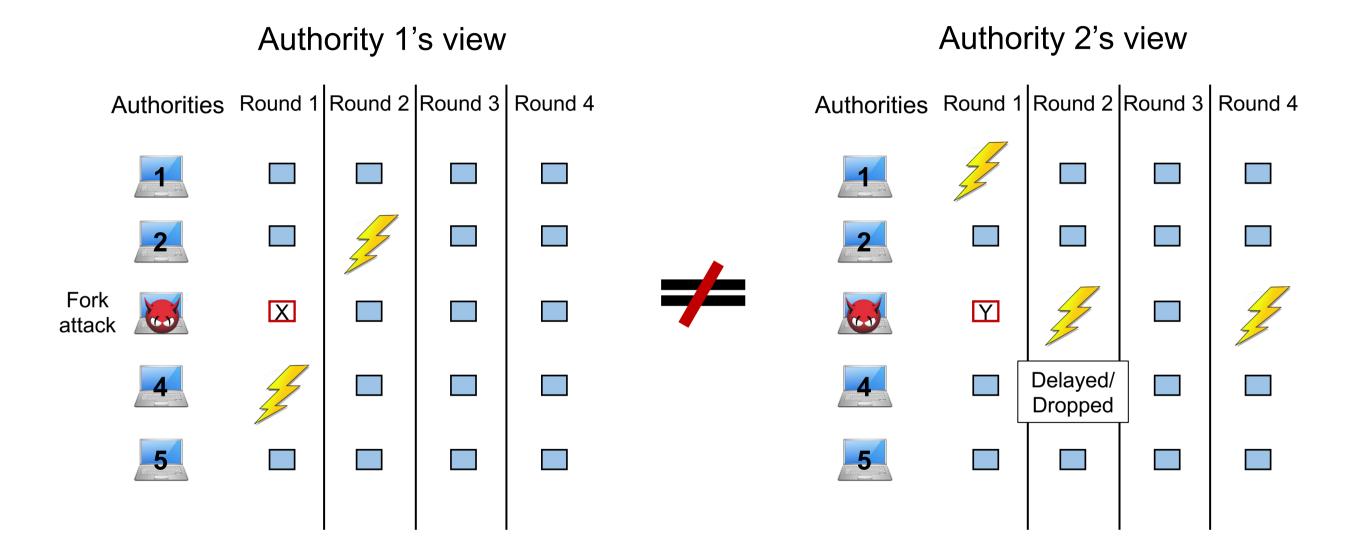
by leader

Accepted

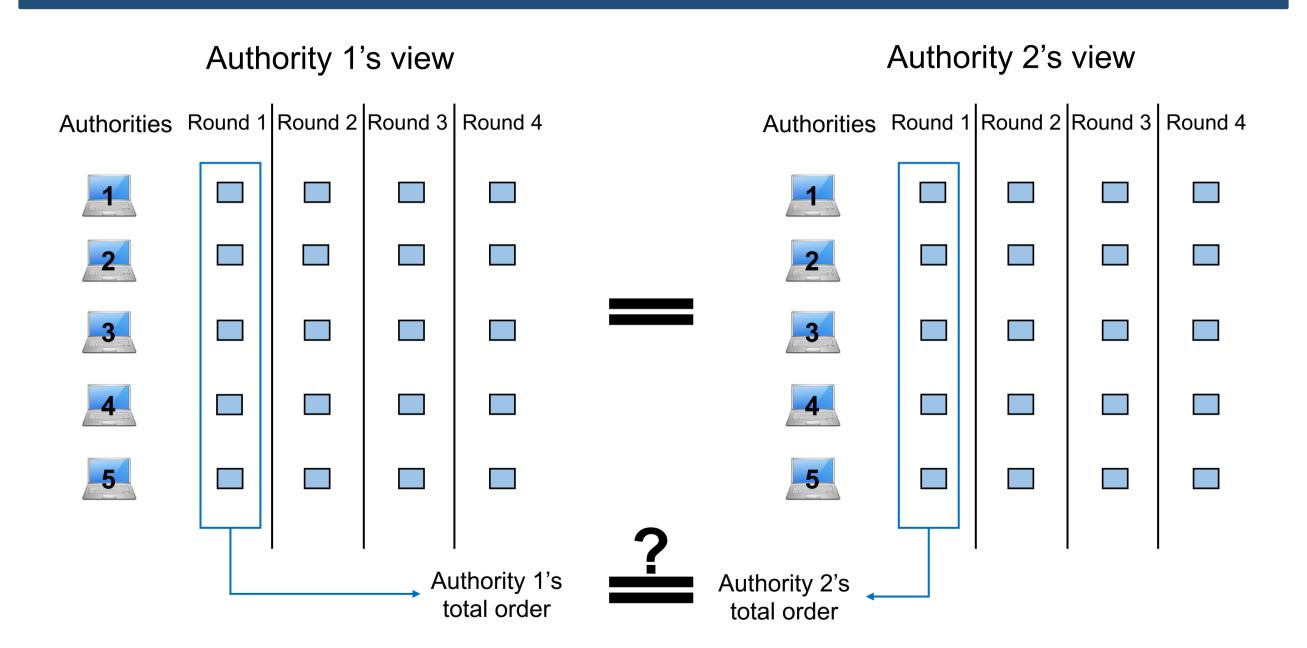
Our Idea



Challenge (1) Consistent view across nodes



Challenge (2) Consistent total ordering



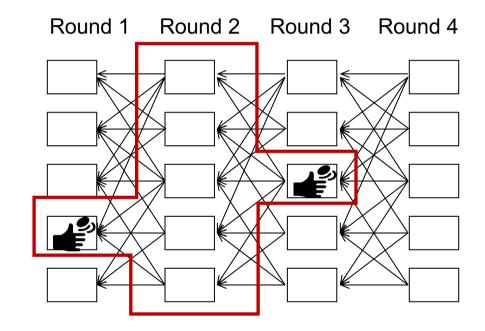
DAG-based BFT Consensus

DAG Mempool

Round 1 Round 2 Round 3 Round 4 2 4 5

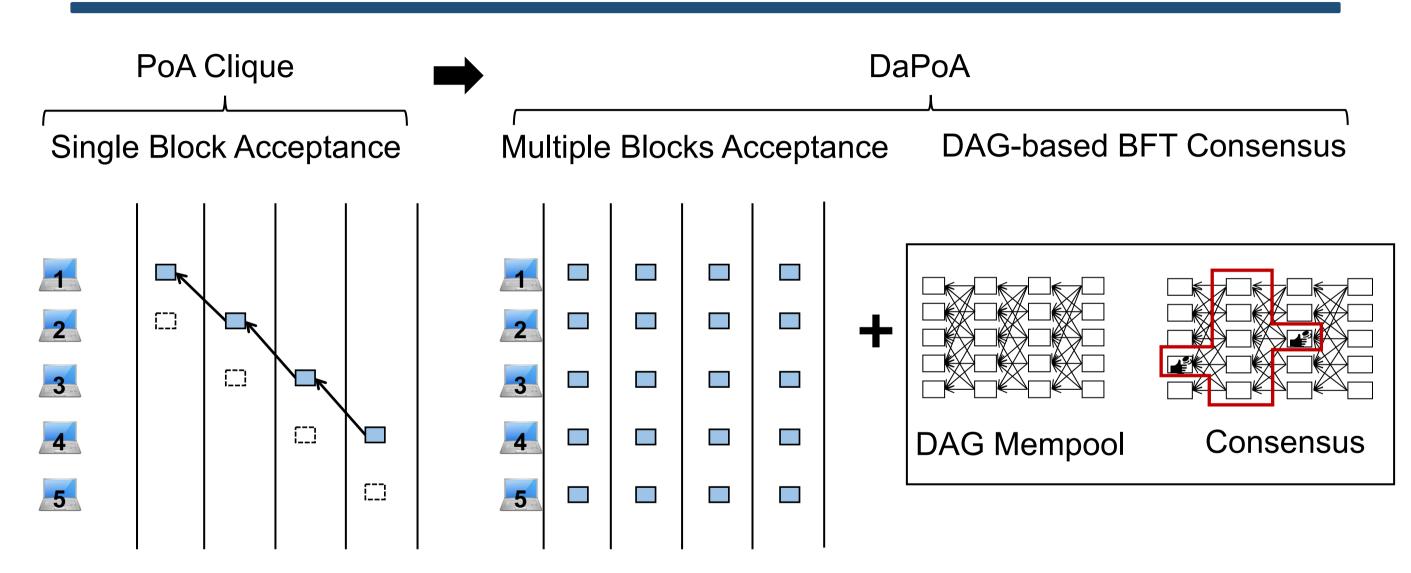
- Consistent blocks across nodes
- Reliable block distribution
- Parallel block proposals & Acceptance

Consensus



- Consistent total order
- Local & deterministic
- (shared) randomness

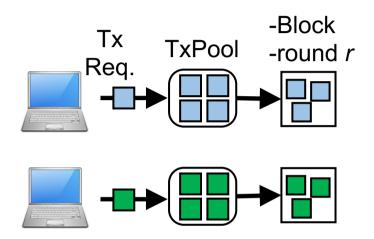
Contributions



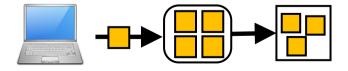
achieves 2.47x higher throughput, 5.76x lower latency than Clique

Assumptions

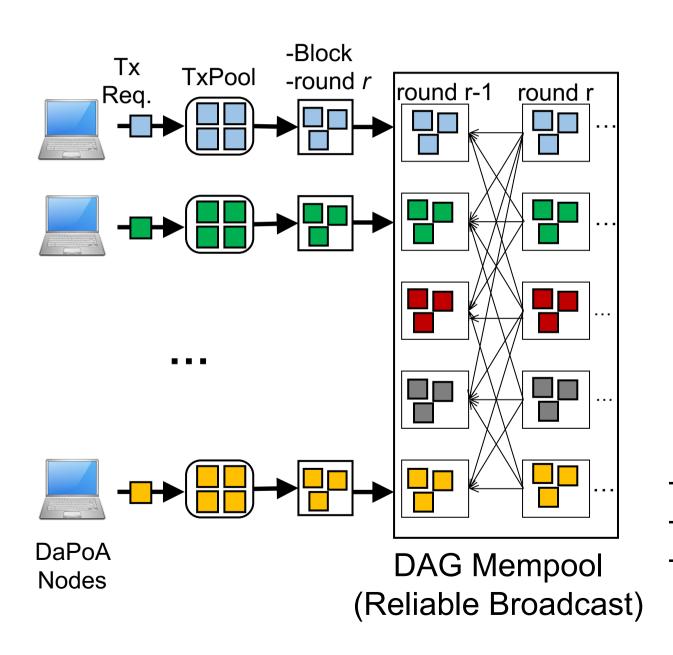
- A consortium blockchain
 - Known Identities
 - Members operates DaPoA node
 - Clients submit Txs to a trusted node
 - f Byzantine members out of 3f+1 members
- Partially asynchronous network
 - unknown time bound Δ
- Crypto cannot be subverted



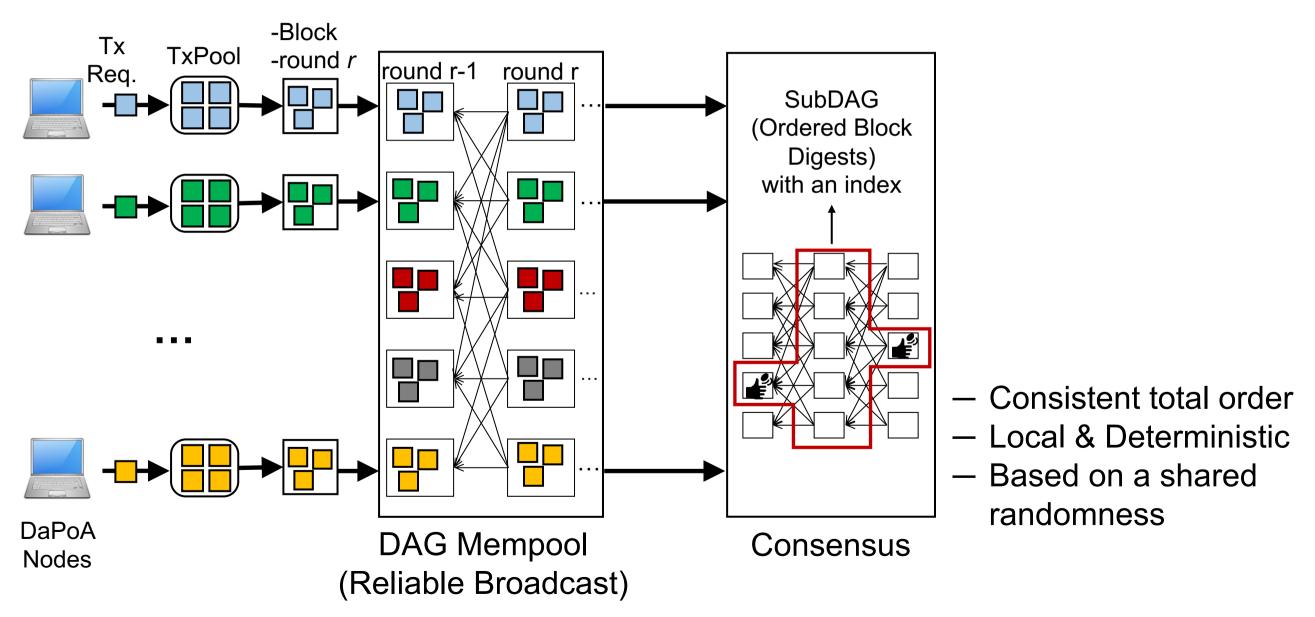
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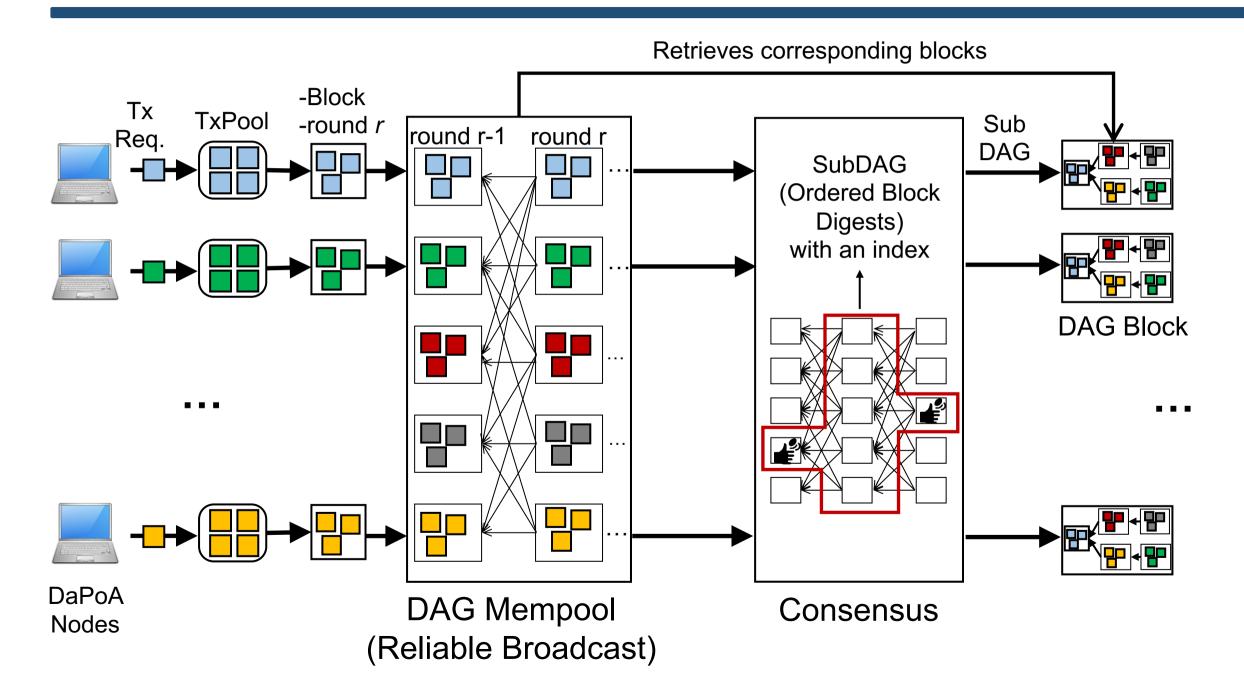


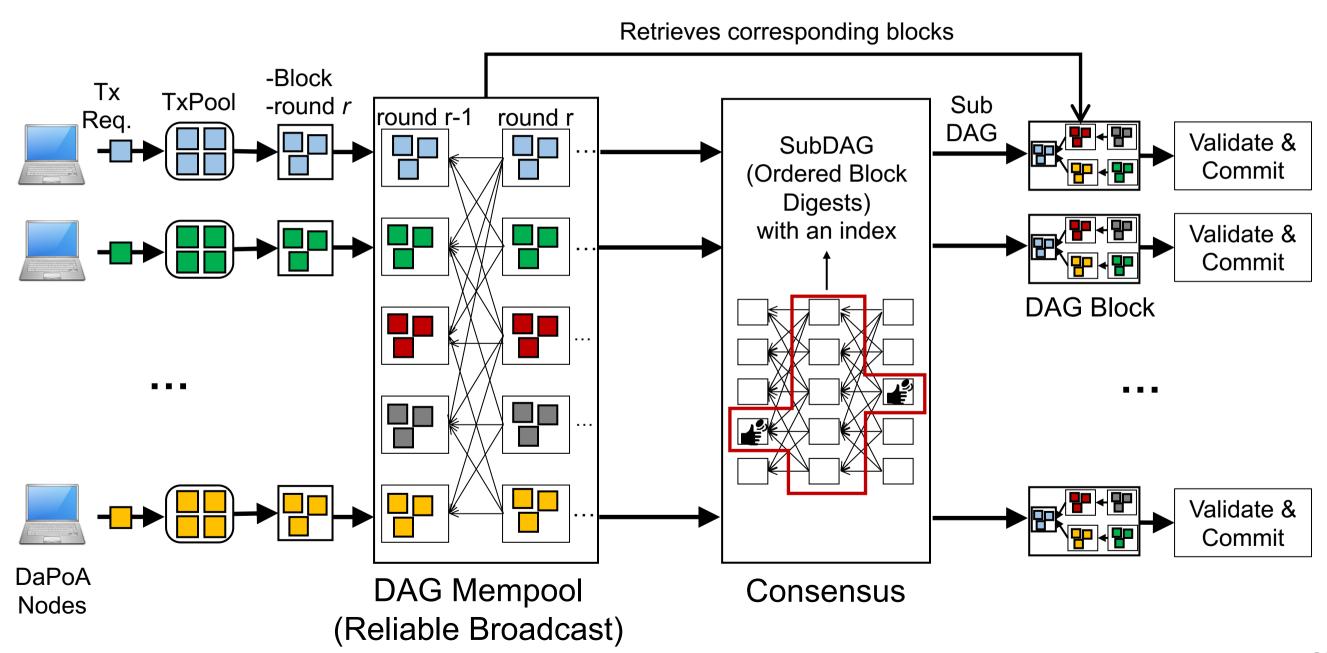
DaPoA Nodes



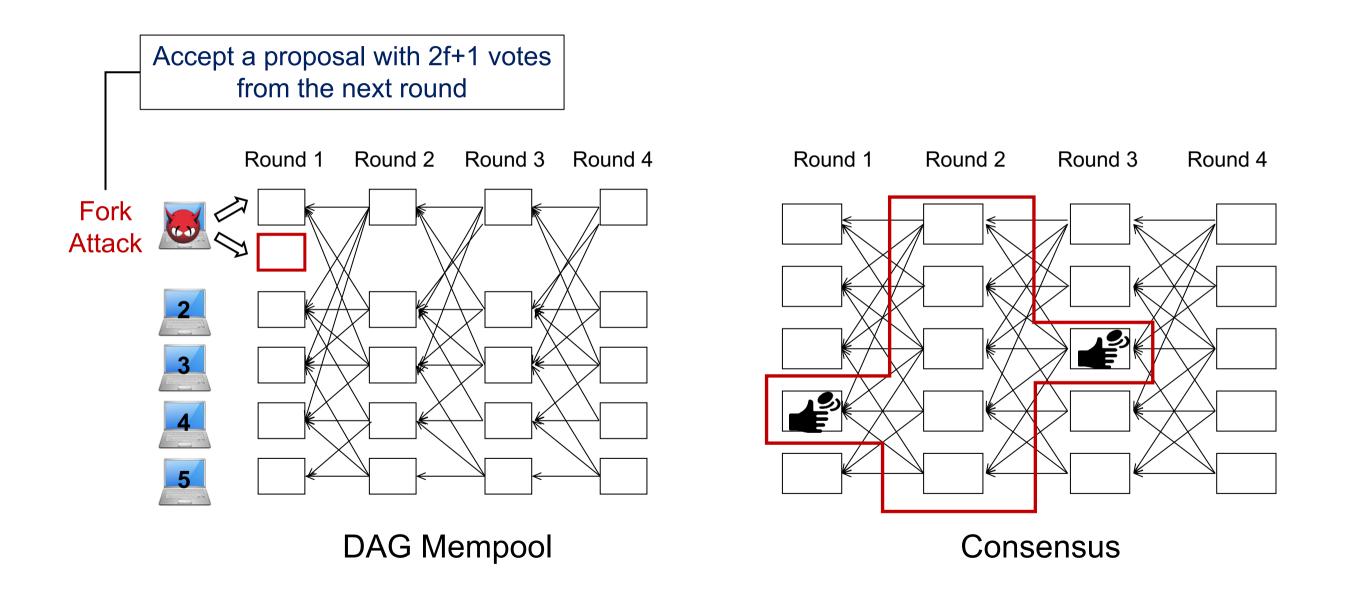
- Reliable block replication across nodes
- Consistent block views
- Multiple block proposals & acceptance



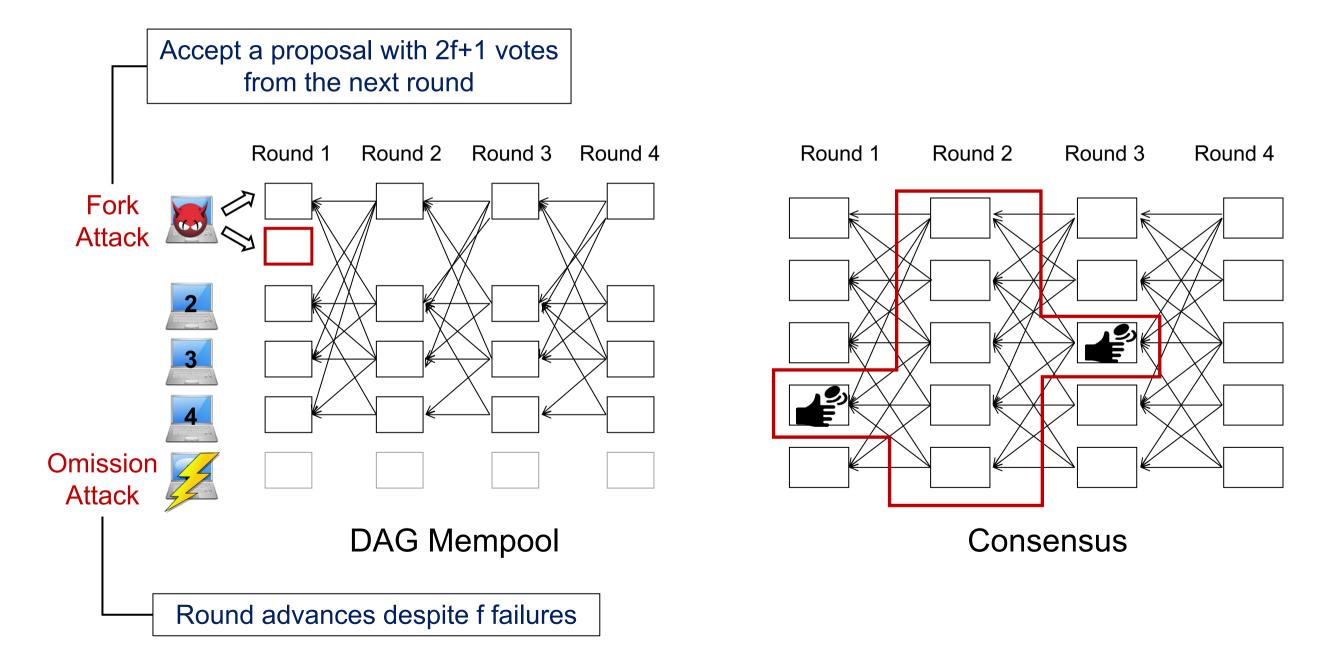




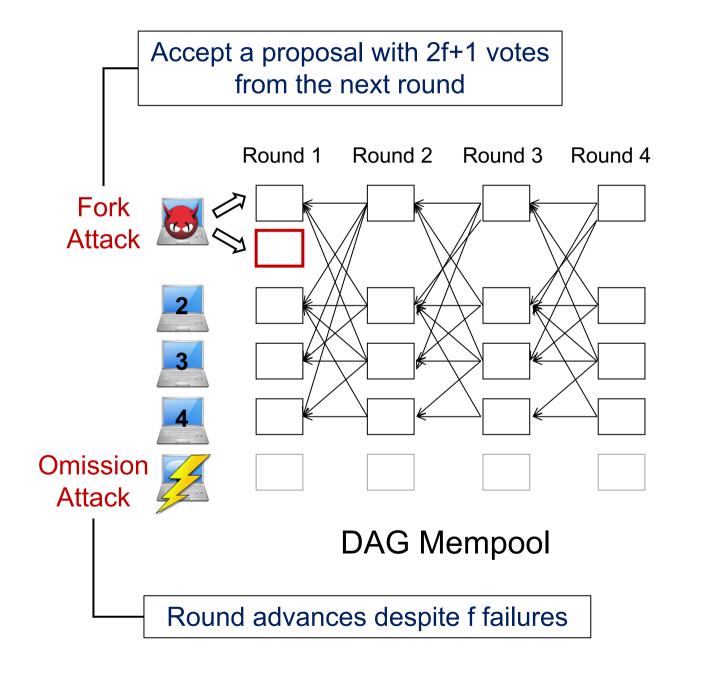
Security Analysis

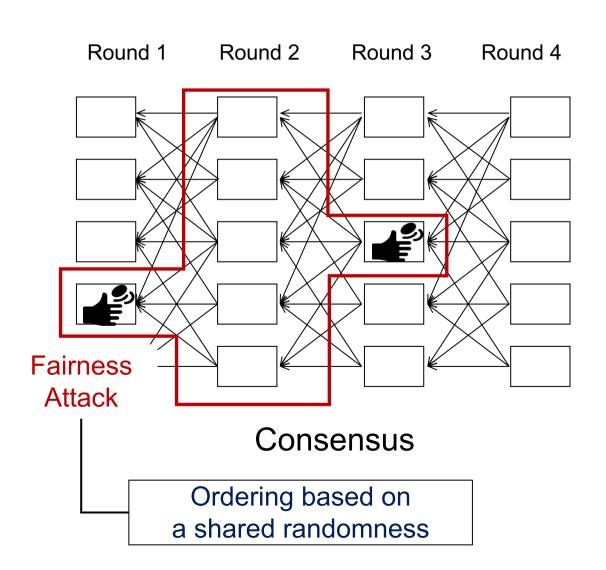


Security Analysis



Security Analysis



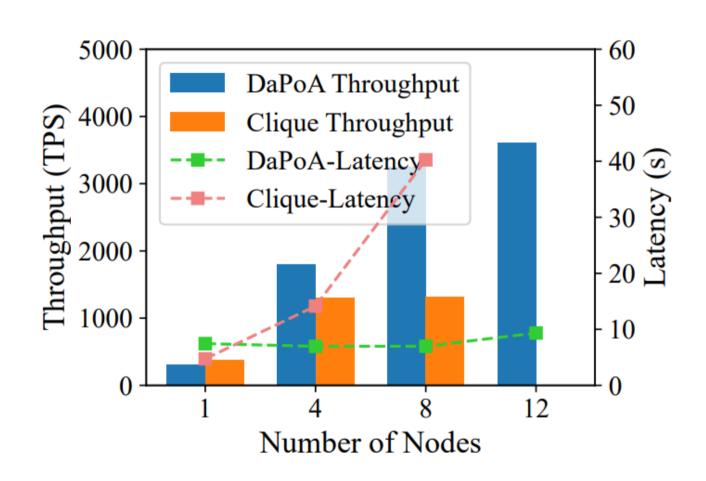


Evaluation: vs. Clique

- Implement DaPoA based on Ethereum Geth*
- Narwhal/Tusk** for DAG BFT
- Benchmark using Hyperledger Caliper
- Hardware: AMD Ryzen 3990X CPU, 256 GB RAM

Parameters

- —Simple payment contract
- -TX confirmation blocks: 2
- -Block Period: 1
- -Gaslimit: 10⁹
- —DAG BFT (Batch Size: 2MB, 1 worker, 1 primary)
- -Send rate: 700

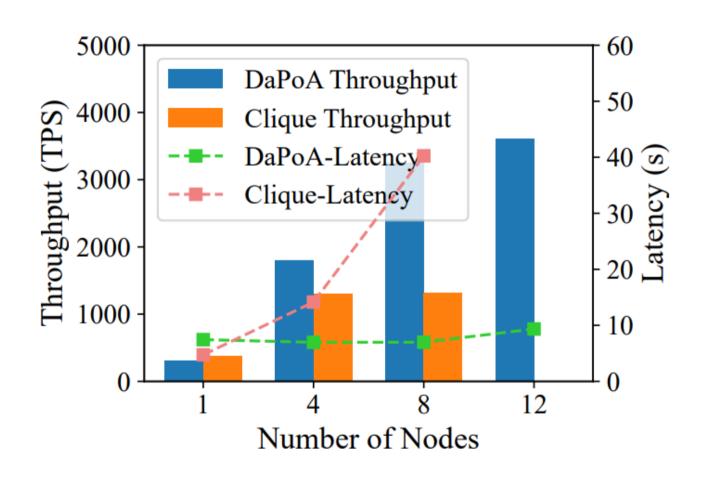


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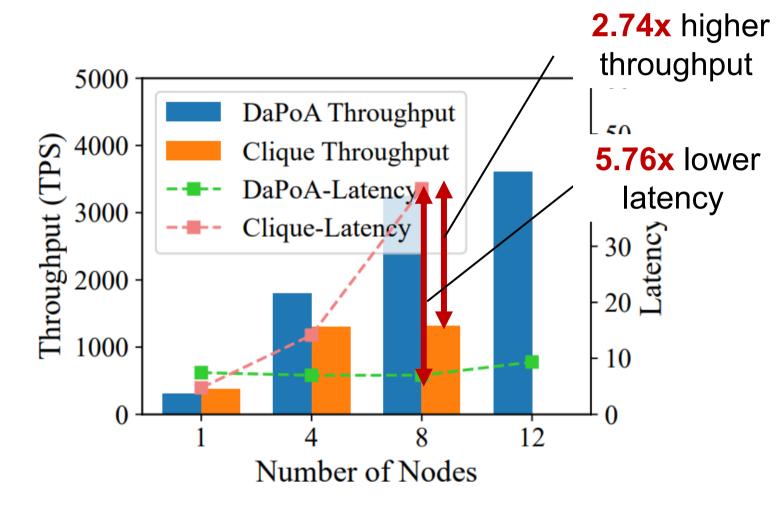


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Conclusion

DaPoA

- An effort to enhance Ethereum Clique for performance scalability
 - Multiple block acceptance
- Leveraging DAG-based BFT consensus
 - A consistent block view across nodes using DAG mempool.
 - A consistent block ordering using local consensus based on shared randomness