



GRENOBLE INP – ENSIMAG, UGA

M2 CySEC – ADVANCED CRYPTOGRAPHY

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## Elliptic curves lab session

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## 1 Part 1 : function over E(F)

We consider the following equation for the curve E :

$$Y^2 = X^3 + aX + b$$

And F denote the field  $\mathbb{Z}/p\mathbb{Z}$ . For the compilation and execution I used python3.8.

### 1.1 Question 1

The function to compute the inverse element over P is **inverse\_element**.

### 1.2 Question 2

The function to test if an element is on the curve is **is\_on\_curve**. We have defined the InfinityPoint by the point of coordinates (0,0).

### 1.3 Question 3

To compute the map  $P \rightarrow 2P$  we have to use the function **doubling**.

### 1.4 Question 4

To compute the addition over the elliptic curve we used **\_\_add\_\_**. For example if P and Q are two point the program instruction  $P + Q$  will compute the addition of the two points over E(F).

### 1.5 Question 5

To compute the multiplication over the elliptic curve we used **\_\_rmul\_\_**. For example if P is a point and n a scalar the program instruction  $n * P$  will compute the multiplication over E(F).

### 1.6 Quesiton 6

The small program that perform the Diffie-Hellman key exchange works as follow : With A and B the two party who want to exchange a secret :

Given two scalars a and b (choose randomly for A and B with a only known by A and b only known by B) and a public point G.

A send  $X = aG$  to B (over E(F))

B send  $Y = bG$  to A (over E(F))

A compute  $aY = \text{secret\_key}$

B compute  $aX = \text{secret\_key}$

The algorithm is compute by the function **diffie\_hellman**.

## 2 Part 2

For this section we take the following value :

$a = 0$

$b = 7$

$x = 0x79BE667EF9DCBBAC55A06295CE870B07029BFCDB2DCE28D959F2815B16F81798$

$y = 0x483ADA7726A3C4655DA4FBFC0E1108A8FD17B448A68554199C47D08FFB10D4B8$

$p = 0xFFEFFFFFC2F$

### 2.1 Quesiton 7

We compute the following point  $P = (x,y)$  and we see that is belong to the curve. Test with the assertion **assert G.is\_on\_curve() == True**.

### 2.2 Quesiton 8

With  $o = 0xFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFFEBAAEDCE6AF48A03BBFD25E8CD0364141$

We compute  $oP$  and we get the InfinityPoint.

### 2.3 Quesiton 9

We have used this website <https://www.dcode.fr/primality-test> to check if  $o$  is prime. And we get that  $o$  is prime.

### 2.4 Quesiton 10

With the last question we can conclude that the order of  $P$  is  $o$ . Because a point on the curve multiplied by the scalar  $o$  give the neutral point Infinity and  $o$  divide  $p-1$ .

### 2.5 Quesiton 11

For this question we have change the Diffie Helllman protocol to choose the secret keys of  $A$  and  $B$  randomly. The new function is **diffie\_hellman\_subgroup**.