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## Written Assignment 3

## 1. Binomial Coefficients

Input :-  $n, k$ Output :-  $C[n][k]$ 

$C$  is the array to store the computed Binomial Coefficients.

~~Now  $C[n][k] =$~~

$$\text{Now } C(n, k) = C(n-1, k-1) + C(n-1, k)$$

$$C(n, 0) = C(n, n) = 1 \leftarrow \text{Stopping Condition}$$

procedure Binomial Coef ( $n, k$ ):

for  $i$  from 0 to  $n$

for  $j$  from 0 to  $k$

if  $j == 0$  or  $j == i$ :

$$C[i][j] = 1$$

else:

$$C[i][j] = C[i-1][j-1] + C[i-1][j]$$

return  $C[n][k]$

¶ The Time Complexity of the above algorithm is  $O(nk)$ . Since inner loop runs  $k$  times and outer loop runs  $n$  times.

2. Algorithm to count the number of ways in which the road can be paved.

Input =  $n$  = length of road in meters

Output = number of ways in which the road can be paved.

$$\text{roadpaving}[n] = \text{roadpaving}[n-2] + \text{roadpaving}[n-3] + \text{roadpaving}[n-5]$$

Stopping Condition  $\rightarrow$   $\text{roadpaving}[0] = 1$   
&  $\text{roadpaving}[1] = 0$

[Sanity Check] & if  $n < 0$  then  $\text{roadpaving}[n] = 0$

def pavingRoad( $n$ ):

if  $n < 0$ : return 0

~~paving~~  $\text{roadpaving}[0] = 1$

$\text{roadpaving}[1] = 0$

for  $i$  from 2 to  $n$ :

sum = 0

if  $i$  is greater than or equal to 2:  
sum +=  $\text{roadpaving}[i-2]$

if  $i$  is greater than or equal to 3:  
sum +=  $\text{roadpaving}[i-3]$

if  $i$  is greater than or equal to 5:  
sum +=  $\text{roadpaving}[i-5]$

$\text{roadpaving}[i] = \text{sum}$

return  $\text{roadpaving}[n]$



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3. Number of ways of partitioning a set of  $n$  elements

Reference :- [en.wikipedia.org/wiki/Bell\\_triangle](https://en.wikipedia.org/wiki/Bell_triangle)  
 Array  $B$  Bell stores the bell triangle.  
 It is a 2d matrix [lower triangular matrix]

Input =  $n$  = number of elements in a set

Output = number of ways they can be partitioned

$$\text{Bell}[i][j] = \text{Bell}[i-1][j-1] + \text{Bell}[i][j-1]$$

procedure Partition ( $n$ ):

~~for~~ for  $i$  from 0 to  $n$

for  $j$  from 0 to  $i$

if ( $i$  and  $j$  are equal to zero)

or ( $i$  is equal to one and  $j$  is equal to zero))

then  $\text{Bell}[i][j] = 1$

// element  $(0,1)$  and  $(0,0)$  are set to 1

else if  $j$  is equal to 0

$\text{Bell}[i][0] = \text{Bell}[i-1][i-1]$

// if column = 0 then  $\text{Bell}[\text{row}][0] = \text{Bell}[\text{row}-1][\text{row}-1]$

// , i.e., set first element of the row to equal to

// the last element of previous row.



else :

$$\text{Bell}[i][j] = \text{Bell}[i-1][j-1] - \text{Bell}[i][j-1]$$

return Bell[n-1][n-1]

Note: Stopping condition here is that if the element is the first element of row 0 and row 1 then set it to 1.

4 Grid pathway problem.

Input =  $n, m$   $n$  = number of rows  
 $m$  = number of cols.

$\therefore n \times m$  = Size of grid.

Output = Number of shortest paths.

A = Array of size  $n \times m$

$$\begin{aligned} A[n][m] &= 1 \text{ if } [(n=1 \& m=0) \text{ or } (n=0 \& m=1)] \\ A[n][m] &= \begin{cases} 1, & \text{if } [(n=1 \& m=0) \text{ or } (n=0 \& m=1)] \\ \text{path}(n, m-1), & \text{if } n=0 \\ \text{path}(n-1, m), & \text{if } m=0 \\ \text{path}(n, m-1) + \text{path}(n-1, m), & \text{otherwise} \end{cases} \end{aligned}$$

Note:- Initially all elements in array A are equal to 1

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procedure path(n, m):
    if A[n][m] > -1 // if it stores a value
        return A[n][m]
    else if [(n is equal to 1 and m is equal to zero)
            or n is equal to 0 and m is equal to one]
        return 1
    else if (n is equal to zero)
        A[n][m] = path(n, m-1)
    else if (m is equal to zero)
        A[n][m] = path(n-1, m)
    else
        A[n][m] = path(n, m-1) + path(n-1, m)
    return A[n][m]

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### Explanation

1. Checks if  $A[n][m]$  is filled. If yes returns that value
2. Checks if element is  $A[0][1]$  or  $A[1][0]$  since path from  $A[0][0]$  to the elements above is equal to 1, This is the stopping condition
3. Checks if the element is in the first column if yes check the element below
4. Checks if the element is in the first row, if yes check the element to the left



5 Else return the addition of path  $(n, m-1)$   
and path  $(n-1, m)$