

# Written Assignment 1

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Q1.

$$f(x) = 3x^2 + 5x + 3$$

$$g(x) = 2x^3 + x - 100$$

$$f(x) = O(g(x))$$

To show that  $\exists c, n_0$  such that

$$0 \leq 3n^2 + 5n + 3 \leq c(2n^3 + n - 100) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 5$

$$3(25) + 25 + 3 \leq 1(2(125) + 5 - 100)$$

$$103 \leq 155$$

*Q.E.D*

Therefore,  $0 \leq 3n^2 + 5n + 3 \leq 1(2n^3 + n - 100) \quad \forall n \geq 5$

Q2.

The running time of the program given is :  $\left(\frac{n}{2}\right)^2$

Consider  $n = 10$

The outer loop will run 5 times

The inner loop will run 5 times

In total the convoluted loop will run  $5 \times 5 = 25$  times

Even for other values of  $n$  we can verify that

The running time of the program will be  $\left(\frac{n}{2}\right)^2$

Therefore, the program will have its upper bound as  $O(n^2)$

Q3.

To prove that  $f(n) = \frac{1}{n} = O(1)$

Therefore,  $f(n) = \frac{1}{n}$  and  $g(n) = O(1)$

To show that  $\exists c, n_0$  such that

$$0 \leq \frac{1}{n} \leq c(1) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 1$

$$\frac{1}{1} \leq 1(1)$$

$$1 \leq 1$$

Q.E.D

Therefore,  $0 \leq \frac{1}{n} \leq 1(1) \quad \forall n \geq 1$

Q4.

Given  $f(n) = O(g(n))$  and  $f(n) = O(h(n))$

Example for  $g(n) = O(h(n))$

Let  $f(n) = x$  and  $g(n) = x^2$  and  $h(n) = x^3$

Here,  $f(n) = O(g(n))$  i.e.  $x = O(x^2)$  and

$f(n) = O(h(n))$  i.e.  $x = O(x^3)$

Also,  $x^2 \leq x^3 \quad \forall n \geq 0$

Therefore,  $x^2 = O(x^3)$  by definition :  $0 \leq x^2 \leq c(x^3) \quad \forall n \geq 1$  and  $c = 1$

Q.E.D

Example for  $g(n) \neq O(h(n))$

Let  $f(n) = x$  and  $g(n) = x^8$  and  $h(n) = x^3$

Here,  $f(n) = O(g(n))$  i.e.  $x = O(x^8)$  and

$f(n) = O(h(n))$  i.e.  $x = O(x^3)$

But,  $x^8 \geq x^3 \quad \forall n \geq 0$

Therefore,  $x^8 \neq O(x^3)$

Since,  $x^3 = O(x^8)$  by definition :  $0 \leq x^3 \leq c(x^8) \quad \forall n \geq 1$  and  $c = 1$

Therefore,  $g(n) \neq O(h(n))$

Since  $g(n) = \Omega(h(n))$

Q.E.D

Q5. i)

$$f(n) = 2^n$$

$$g(n) = n^2$$

$$f(n) = O(g(n))$$

To show that  $\exists c, n_0$  such that

$$0 \leq 2^n \leq c(n^2) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 2$

$$2^2 \leq 1(2^2)$$

$$4 \leq 4$$

*Q.E.D*

Therefore,  $0 \leq 2^n \leq 1(n^2) \quad \forall n \geq 2$

ii)

$$f(n) = 2^n$$

$$g(n) = 3^n$$

$$f(n) = O(g(n))$$

To show that  $\exists c, n_0$  such that

$$0 \leq 2^n \leq c(3^n) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 1$

$$2^1 \leq 1(3^1)$$

$$2 \leq 3$$

*Q.E.D*

Therefore,  $0 \leq 2^n \leq 1(3^n) \quad \forall n \geq 1$

iii)

$$f(n) = \log n$$

$$g(n) = \log^2 n$$

$$f(n) = O(g(n))$$

To show that  $\exists c, n_0$  such that

$$0 \leq \log n \leq c(\log^2 n) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 4$

$$\log 4 \leq 1(\log^2 4)$$

$$2 \leq 2^2$$

$$2 \leq 4$$

Q.E.D

Therefore,  $0 \leq \log n \leq 1(\log^2 n) \quad \forall n \geq 2$

iv)

$$f(n) = n^{\sqrt{n}}$$

$$g(n) = n^n$$

$$f(n) = O(g(n))$$

To show that  $\exists c, n_0$  such that

$$0 \leq n^{\sqrt{n}} \leq c(n^n) \quad \forall n \geq n_0$$

Set  $c = 1$  and  $n_0 = 1$

$$1^{\sqrt{1}} \leq 1(1^1)$$

$$1 \leq 1$$

Q.E.D

Therefore,  $0 \leq n^{\sqrt{n}} \leq c(n^n) \quad \forall n \geq 1$