

## Assignment No. 4

To solve the problem of finding the maximum-length common subsequence of  $X$  and  $Y$ , where  $X$  and  $Y$  are two sequences, I constructed two 2-d arrays  $b$  and  $c$ . Array  $b$  is used to print the Longest common subsequence and Array  $c$  is used to calculate the length of the longest common subsequence.

As per the Theorem 15.1 on page 392 of the textbook - "Introduction to Algorithms", we are recursively solving the substructure of the 2 sequences to find common elements which will eventually build up the subsequence. The Theorem states that - Let  $X$  and  $Y$  be sequences and  $Z$  be any LCS of  $X$  and  $Y$ .

Then

1. If  $x_m = y_n$ , then  $z_k = x_m = y_n$  and  $Z_{k-1}$  is an LCS of  $X_{m-1}$  and  $Y_{n-1}$
2. If  $x_m \neq y_n$ , then  $z_k \neq x_m$  implies that  $Z$  is an LCS of  $X_{m-1}$  and  $Y$ .
3. If  $x_m \neq y_n$ , then  $z_k \neq y_n$  implies that  $Z$  is an LCS of  $X$  and  $Y_{n-1}$ .

Therefore, array  $c$  is basically used to solve these subproblems which are shown by the above theorem

Recursive Formulation

$$\begin{aligned} C[i,j] &= 0, && \text{if } i==0 \text{ or } j==0 \\ &= c[i-1,j-1]+1, && \text{if } i,j > 0 \text{ and } x_i = y_j \\ &= \max(c[i,j-1], c[i-1,j]), && \text{if } i,j > 0 \text{ and } x_i \neq y_j \end{aligned}$$