

B. Tech. I CSE (Sem-2)  
Data Structures  
CS102

# Queue: A Data Structure

\*Some slides are kept with logical or syntax error. So, if you are absent in theory class, please also refer the slides provided at the end of the presentation.

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# Queue Data Structure

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- Basic principles
- Operation of queue
- Queue using Array
- Applications of queue
- Queue using Linked List

# Queue Data Structure

- Data structure, somewhat similar to Stacks. Unlike stacks, a queue is open at both its ends. One end is always used to insert data (enqueue) and the other is used to remove data (dequeue).
- An ordered group of homogeneous items or elements.



- Elements are added at one end.
- Elements are removed from the other end.
- The element added first is also removed first
- FIFO: First In, First Out



# Queue Data Structure

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- Traffic Management Systems
  - In traffic control and toll booths, vehicles often form a queue as they wait for their turn
  - An example of queue real-time applications where vehicles are processed in the order they arrive
- Keyboards' keypress sequence
- Queue of network data packets to send
- Used as buffers on MP3 players and portable CD players, iPod playlist

# Queue Data Structure



## Queue

- As in stacks, a queue can also be implemented using Arrays, Linked-lists, Pointers and Structures.



# Queue Data Structure

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## Operations

enqueue

dequeue

create

isempty

size



QUEUE



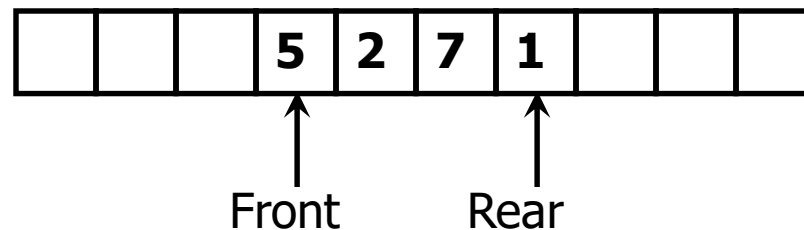
# Queue Data Structure

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- Like stacks, *queues* are lists.
- With a queue, however, insertion is done at one end, whereas deletion is performed at the other end.
- The basic operations on a queue are *Enqueue*, which inserts an element at the end of the list (called the rear), and *Dequeue*, which deletes (and returns) the element at the start of the list (known as the front).

# Array Implementation of Queues

- For each queue data structure, we keep an array, *Queue[]*, and the positions *Front* and *Rear*, which represent the ends of the queue.
- We also keep track of the number of elements that are actually in the queue, *Size*. All this information is part of one structure.
- The following figure shows a queue in some intermediate state. The cells that are blanks have undefined values in them:







# Algorithm Enqueue

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Algorithm Enqueue (Q, F, R, Y)

//Queue Q consisting of N elements

If(  $R \geq N$ ) /\* check for the overflow condition \*/

    Printf ("Overflow"); return;

Else

{

$R = R + 1$ ; /\*Increment rear pointer \*/

$Q[R] = Y$ ; /\* insert element \*/

    If ( $F == 0$ ) /\* set the front pointer \*/

$F = 1$ ;

    Return;

}



# Algorithm Dequeue

---

Algorithm DeQueue (Q, F, R)

If (F==0) /\* check for the underflow condition \*/

    Printf ("underflow");

    return 0; /\* 0 denotes the empty queue \*/

Else

{

Y = Q [F]; /\* delete element \*/

If (F == R)

F = R = 0;

Else

F = F + 1; /\* increment front pointer \*/

Return Y;

}

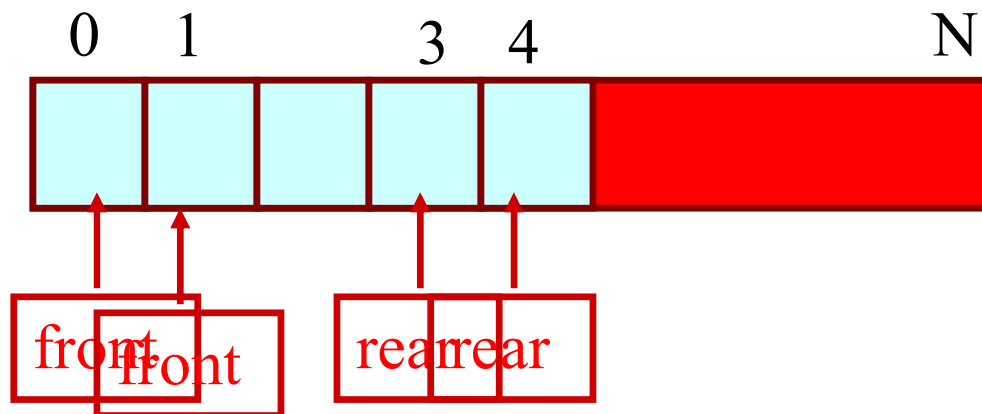
# Array Implementation of Queues

- The size of the queue depends on the number and order of enqueue and dequeue.
- It may be situation where memory is available but enqueue is not possible.

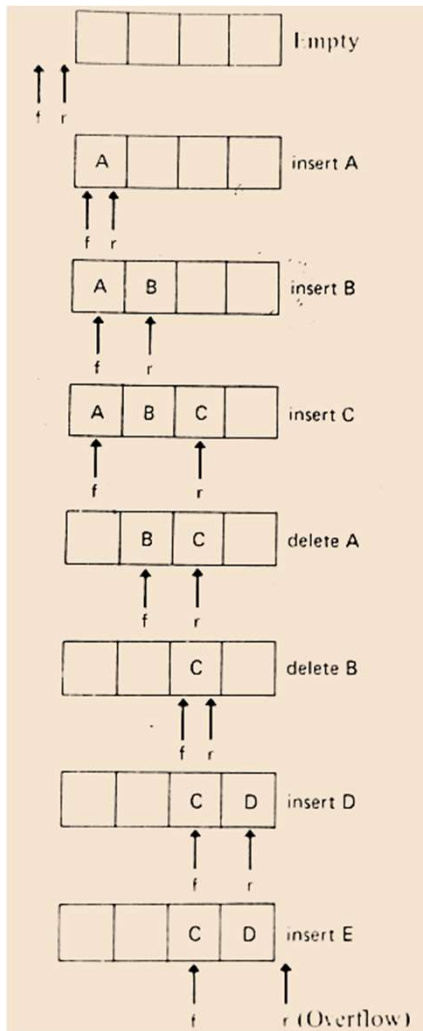
ENQUEUE

DEQUEUE

Overflow occurs on trying to insert , even though the first few locations are not being used  
Effective queuing storage area of array gets reduced.



# Array Implementation of Queues



- How many elements are available?
- How many MORE elements CAN BE INSERTED?
- Show the different situation(s)
  - When only 1 element is available?
  - When 2 elements are available?
  - Etc.



# Array Implementation of Queues

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- Very wasteful of storage
- Arbitrary large amount of memory would be required to accommodate the elements
- ONLY BE USED
  - When the queue is emptied at certain intervals



# Array Implementation of Queues

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- Performance
  - Let  $n$  be the number of elements in the queue
  - The space used is  $O(n)$
  - Each operation runs in time  $O(1)$
- Limitations
  - The maximum size of the queue must be defined *a priori*, and cannot be changed
  - Trying to enqueue a new element into a full queue causes an implementation-specific exception



# Array Implementation of Queues

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## Forwarding Elements Array-based Queue

- On element removal from array front location, remaining elements are shifted forward one position in the array
- Similar to what we did for an array-based stack
- The enqueue operation has running time
  - $O(n+k)$  with the  $k$  shifting



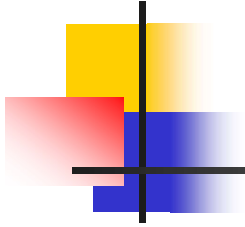
# Array Implementation of Queues

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## Growable Array-based Queue

- In an enqueue operation, when the array is full, instead of throwing an exception, we can replace the array with a larger one
- Similar to what we did for an array-based stack
- The enqueue operation has running time
  - $O(n)$  with the incremental strategy
  - $O(1)$  with the doubling strategy





# Applications of Queue



# Reading a String of Characters

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- A queue can retain characters in the order in which they are typed

```
Queue.createQueue()  
while (not end of line) {  
    Read a new character ch  
    Queue.enqueue(ch)  
} //end while
```

- Once the characters are in a queue, the system can process them as necessary



# Recognizing Palindromes

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- A palindrome
  - A string of characters that reads the same from left to right as it does from right to left
- To recognize a palindrome, a queue can be used in conjunction with a stack
  - A stack reverses the order of occurrences
  - A queue preserves the order of occurrences



# Recognizing Palindromes

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- A nonrecursive recognition algorithm for palindromes
  - As traverse the character string from left to right, insert each character into both a queue and a stack
  - Compare the characters at the front of the queue and the top of the stack

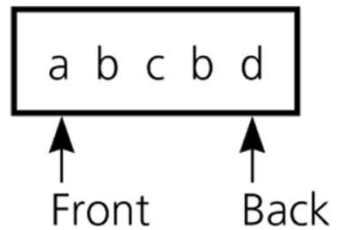


# Recognizing Palindromes

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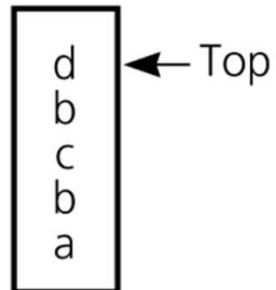
String: abcbd

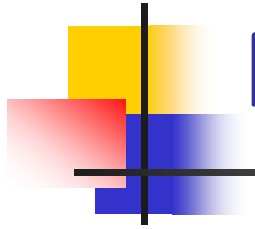
Queue:



Inserting a string into both  
a queue and a stack

Stack:





# Radix Sort

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- Radix is a synonym for base. base 10, base 2
- Multi pass sorting algorithm that only looks at individual digits during each pass
- Use queues as buckets to store elements
- Create an array of 10 queues
- Starting with the least significant digit place value in queue that matches digit
- Empty queues back into array
- Repeat, moving to next least significant digit



# Radix Sort

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- Original values in array

9, 113, 70, 86, 12, 93, 37, 40, 252, 7, 79, 12

- Look at ones place

9, 113, 70, 86, 12, 93, 37, 40, 252, 7, 79, 12

- Array of Queues (all empty initially):

0	5
1	6
2	7
3	8
4	9



# Radix Sort

---

- Original values in array

9, 113, 70, 86, 12, 93, 37, 40, 252, 7, 79, 12

- Look at ones place

9, 113, 70, 86, 12, 93, 37, 40, 252, 7, 79, 12

- Queues:

0    70, 40

5

1

6   86

2    12, 252, 12

7   37, 7

3    113, 93

8

4

9   9, 79





# Radix Sort

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- Empty queues in order from 0 to 9 back into array

70, 40, 12, 252, 12, 113, 93, 86, 37, 7, 9, 79

- Now look at 10's place

70, 40, 12, 252, 12, 113, 93, 86, 37, 7, 9, 79

- Queues:

0 7, 9

1 12, 12, 113

2

3 37

4 40

5 252

6

7 70, 79

8 86

9 93



# Radix Sort

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- Empty queues in order from 0 to 9 back into array

7, 9, 12, 12, 113, 37, 40, 252, 70, 79, 86, 93

- Now look at 100's place

7, 9, 12, 12, 113, 37, 40, 252, 70, 79, 86, 93

- Queues:

0    7, 9, 12, 12, 37, 40, 70, 79, 86, 93 5

1    113 6

2    252 7

3 8

4 9



# Radix Sort

---

- Empty queues in order from 0 to 9 back into array  
7, 9, 12, 12, 40, 70, 79, 86, 93, 113, 252



# Point of Sale (POS)

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- Maintains a list of customers waiting for service
  - Can add a new customer at the end
  - Can display name of next customer to serve
  - Can display the length of the line
  - Can determine how many people are ahead of a particular waiting customer



# Queue Variants

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- Circular Queue
- Doubly Ended Queue
- Priority Queue