

Combinational Circuit Building Blocks

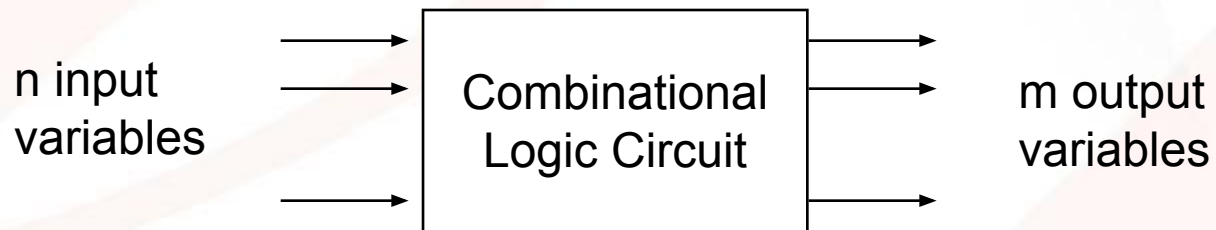
Chapter 6

Combinational Circuits

- Logic circuits for digital system
 - **Combinational circuits**
 - the outputs are a function of the current inputs
 - **Sequential circuits**
 - contain memory elements (Flip-Flops)
 - the outputs are a function of the current inputs and the state of the memory elements
 - the outputs also depend on past inputs

Combinational Logic Circuit

- A combinational circuits
 - 2^n possible combinations of input values

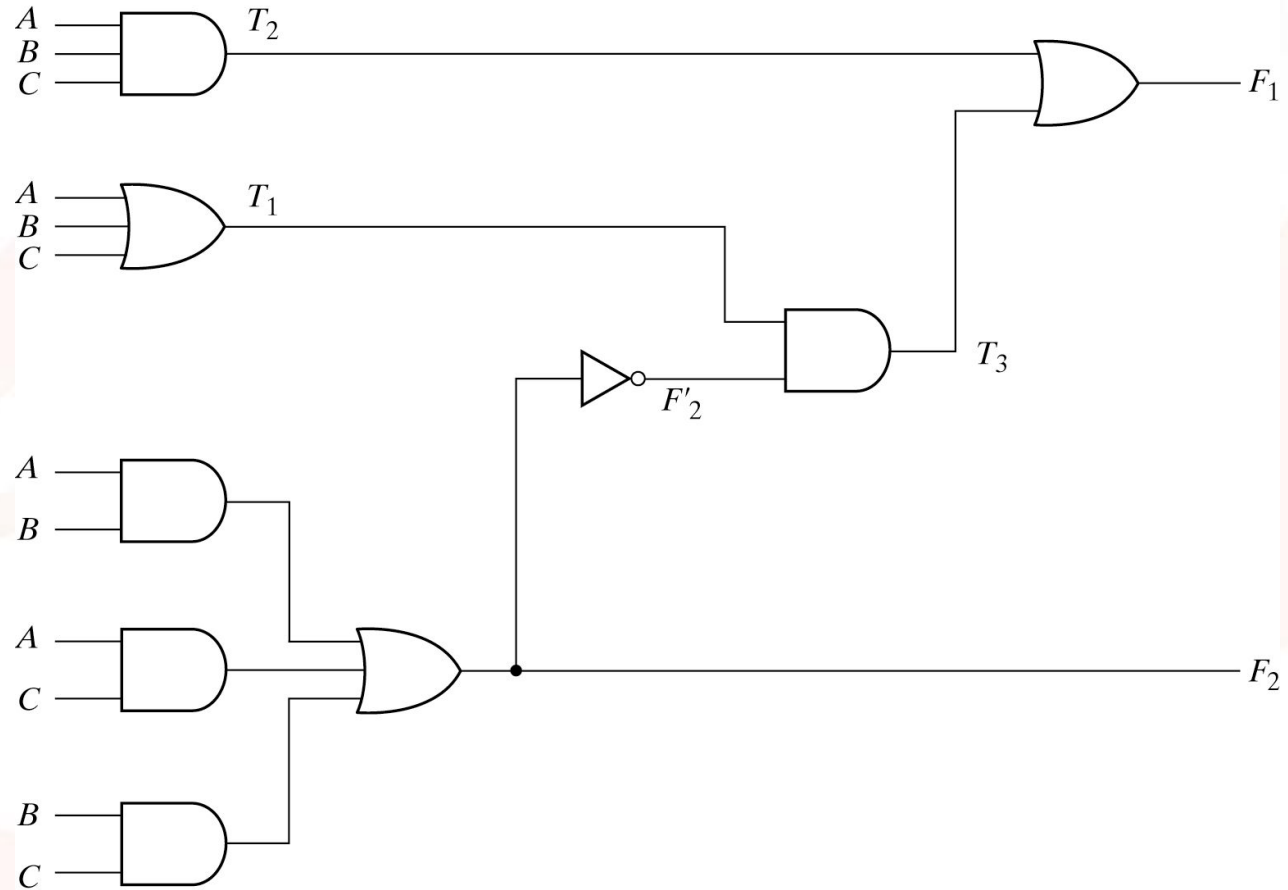


- Specific functions
 - Adders, subtractors, comparators, decoders, encoders, and multiplexers
 - Medium-scale integration (MSI) circuits or standard cells

Analysis Procedure

- A combinational circuit
 - make sure that it is combinational not sequential
 - No feedback path
 - derive its Boolean functions (truth table)
 - design verification
 - a verbal explanation of its function

Analysis Procedure (cont.)



Analysis Procedure (cont.)

- A straightforward procedure

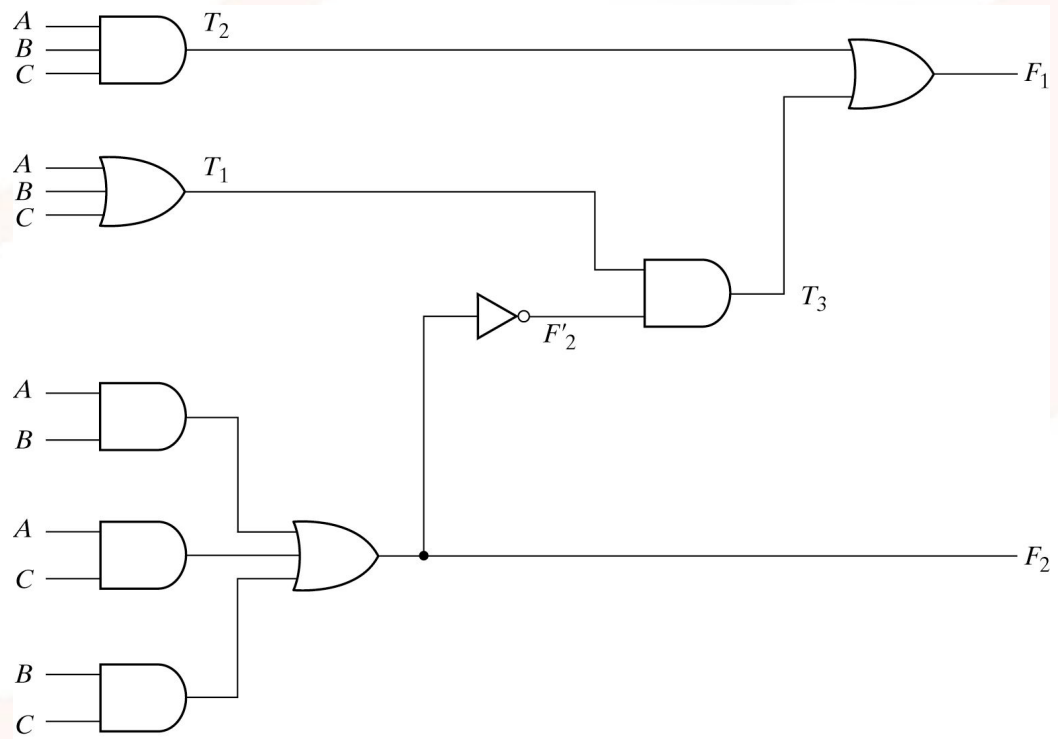
- $F_2 = AB + AC + BC$

- $T_1 = A + B + C$

- $T_2 = ABC$

- $T_3 = F_2' T_1$

- $F_1 = T_3 + T_2$



Analysis Procedure (cont.)

- $$\begin{aligned} F_1 &= T_3 + T_2 = F_2' T_1 + ABC \\ &= (AB + AC + BC)'(A + B + C) + ABC \\ &= (A' + B')(A' + C')(B' + C')(A + B + C) + ABC \\ &= (A' + B'C')(AB' + AC' + BC' + B'C) + ABC \\ &= A'BC' + A'B'C + AB'C' + ABC \end{aligned}$$

Analysis Procedure (end)

- The truth table

- $F_1 = A'BC' + A'B'C + AB'C' + ABC$

- $F_2 = AB + AC + BC$

A	B	C	F_2	F_2'	T_1	T_2	T_3	F_1
0	0	0	0	1	0	0	0	0
0	0	1	0	1	1	0	1	1
0	1	0	0	1	1	0	1	1
0	1	1	1	0	1	0	0	0
1	0	0	0	1	1	0	1	1
1	0	1	1	0	1	0	0	0
1	1	0	1	0	1	0	0	0
1	1	1	1	0	1	1	0	1

Design Procedure

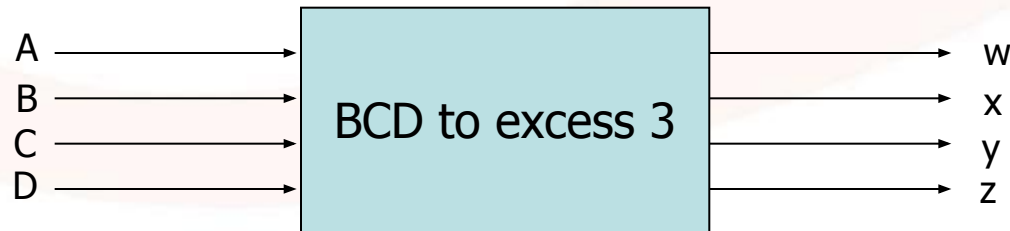
- The design procedure of combinational circuits
 - State the problem (system spec.)
 - determine the inputs and outputs
 - the input and output variables are assigned symbols
 - derive the truth table
 - derive the simplified Boolean functions
 - draw the logic diagram and verify the correctness

Combinational Circuit Design

- Functional description
 - Boolean function
 - HDL (Hardware description language)
 - Verilog HDL
 - VHDL
 - Schematic entry
- Logic minimization
 - number of gates
 - number of inputs to a gate (Fan-in)
 - propagation delay
 - number of interconnection
 - limitations of the driving capabilities (Fan-out)

Design Example: Code conversion

- BCD to excess-3 code
 - System Specification
 - excess-3 code is $\text{BCD} + 3$
 - Inputs and Outputs
 - Inputs = BCD (4 digits)
 - Outputs = excess-3 code (4 digits)
 - Symbols for input and output variables
 - Input variables = ABCD
 - Output variables = wxyz



Code conversion example

- The truth table

A	B	C	D	w	x	y	z
0	0	0	0	0	0	1	1
0	0	0	1	0	1	0	0
0	0	1	0	0	1	0	1
0	0	1	1	0	1	1	0
0	1	0	0	0	1	1	1
0	1	0	1	1	0	0	0
0	1	1	0	1	0	0	1
0	1	1	1	1	0	1	0
1	0	0	0	1	0	1	1
1	0	0	1	1	1	0	0
1	0	1	0	x	x	x	x
1	0	1	1	x	x	x	x
1	1	0	0	x	x	x	x
1	1	0	1	x	x	x	x
1	1	1	0	x	x	x	x
1	1	1	1	x	x	x	x

■ The maps

		CD		C	
AB		00	01	11	10
A	00	1			1
	01	1			1
	11	X	X	X	X
	10	1		X	X

D
 $z = D'$

		CD		C	
AB		00	01	11	10
A	00	1		1	
	01	1		1	
	11	X	X	X	X
	10	1		X	X

D
 $y = CD + C'D'$

		CD		C	
AB		00	01	11	10
A	00		1	1	1
	01	1			
	11	X	X	X	X
	10		1	X	X

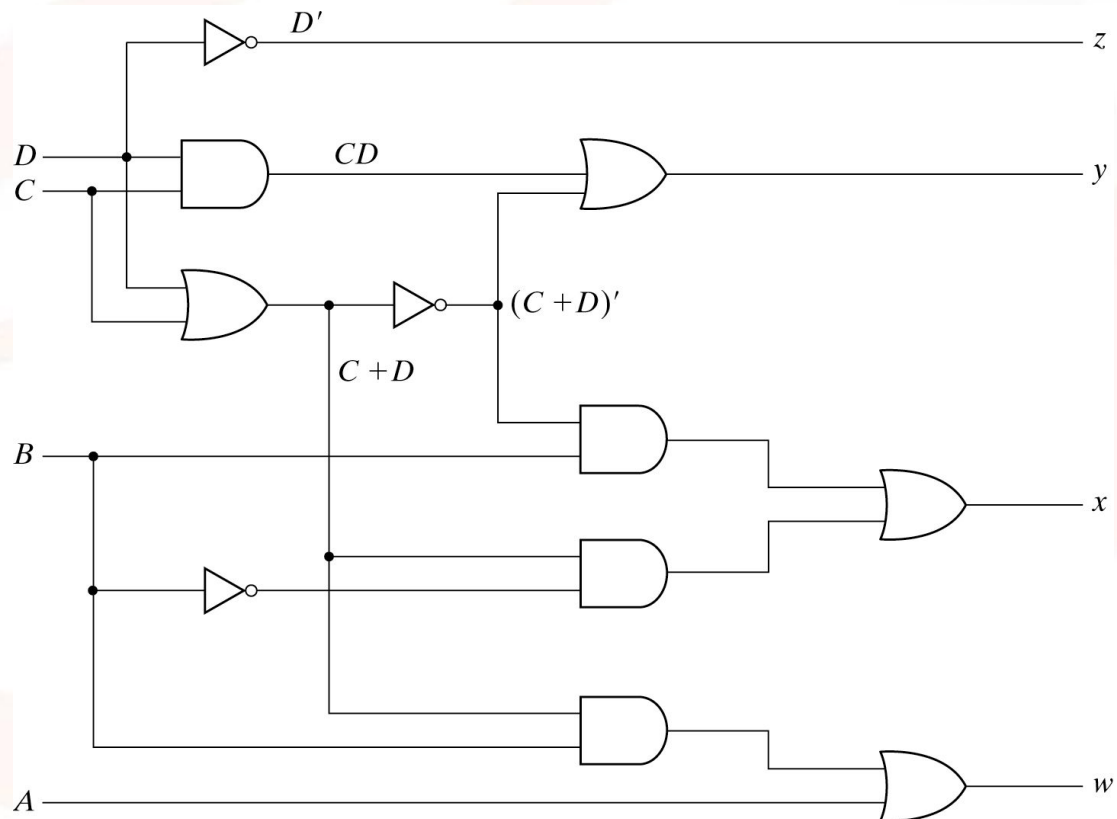
D
 $X = B'C + B'D + BC'D'$

		CD		C	
AB		00	01	11	10
A	00				
	01		1	1	1
	11	X	X	X	X
	10	1	1	X	X

D
 $w = A + BC + BD$

- The simplified functions
 - $z = D'$
 - $y = CD + C'D'$
 - $x = B'C + B'D + BC'D'$
 - $w = A + BC + BD$
- Another implementation
 - $z = D'$
 - $y = CD + C'D' = CD + (C+D)'$
 - $x = B'C + B'D + BC'D' = B'(C+D) + B(C+D)'$
 - $w = A + BC + BD$

- The logic diagram

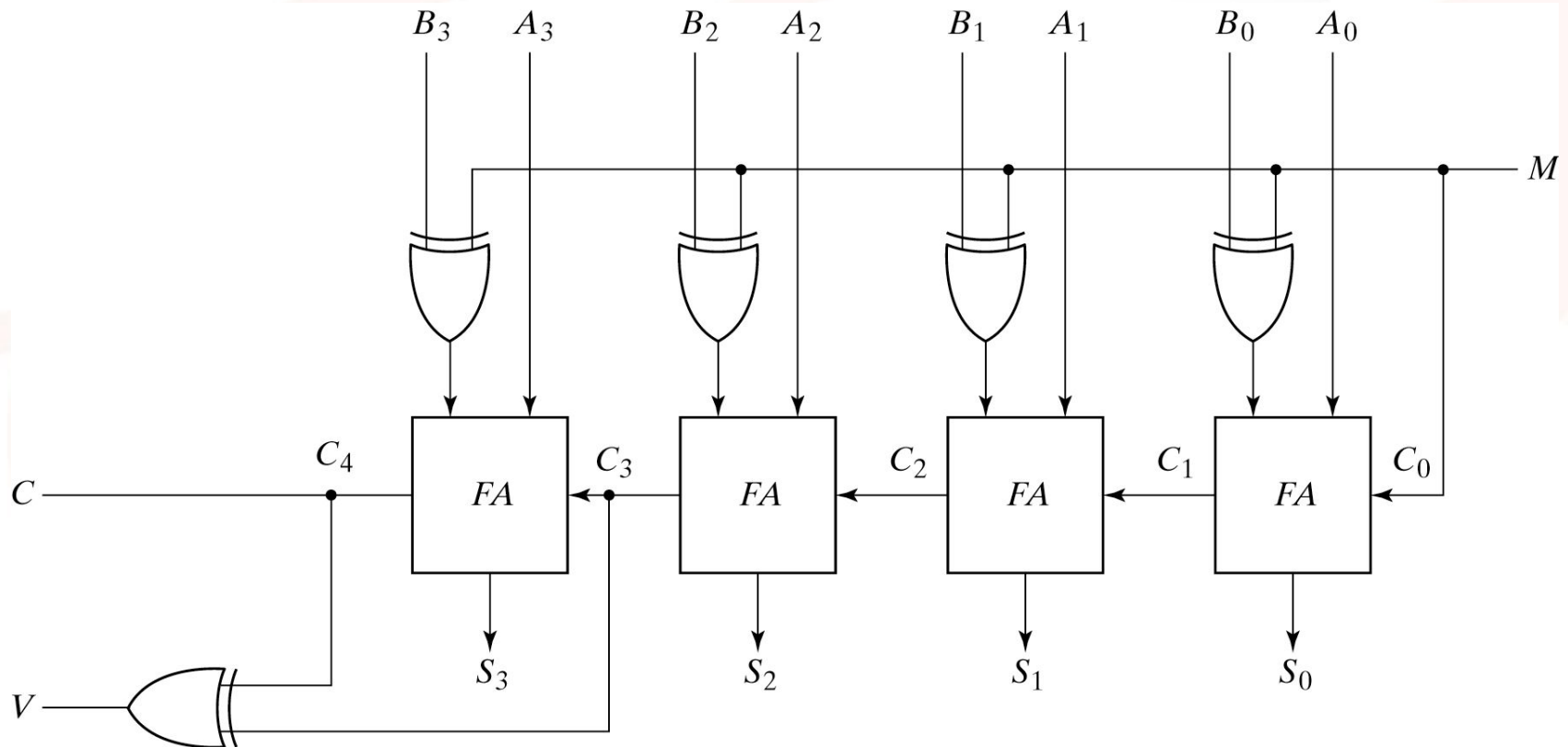


Arithmetic combinational circuits

- Adder/Subtractor
 - Full Adder
 - Half Adder
 - Carry Look Ahead (CLA) Adder
- Multiplier
- Divider

Binary Subtractor (Recall)

- $A - B = A + (2\text{'s complement of } B)$
- 4-bit Adder-subtractor



Conditions of Overflow (Recall)

- **Overflow**
 - The storage is limited
 - Add two positive numbers and obtain a negative number
 - Add two negative numbers and obtain a positive number
 - $V=0$, no overflow; $V=1$, overflow

Decimal Adder (BCD Adder)

- Add two BCD's
 - **9 inputs**: two BCD's and one carry-in
 - **5 outputs**: one BCD and one carry-out
- Design approaches
 - A truth table with 2^9 entries
 - use binary full Adders
 - the sum $\leq 9+9+1 = 19$
 - binary to BCD

Decimal Adder (BCD Adder)

■ The truth table

Derivation of a BCD Adder

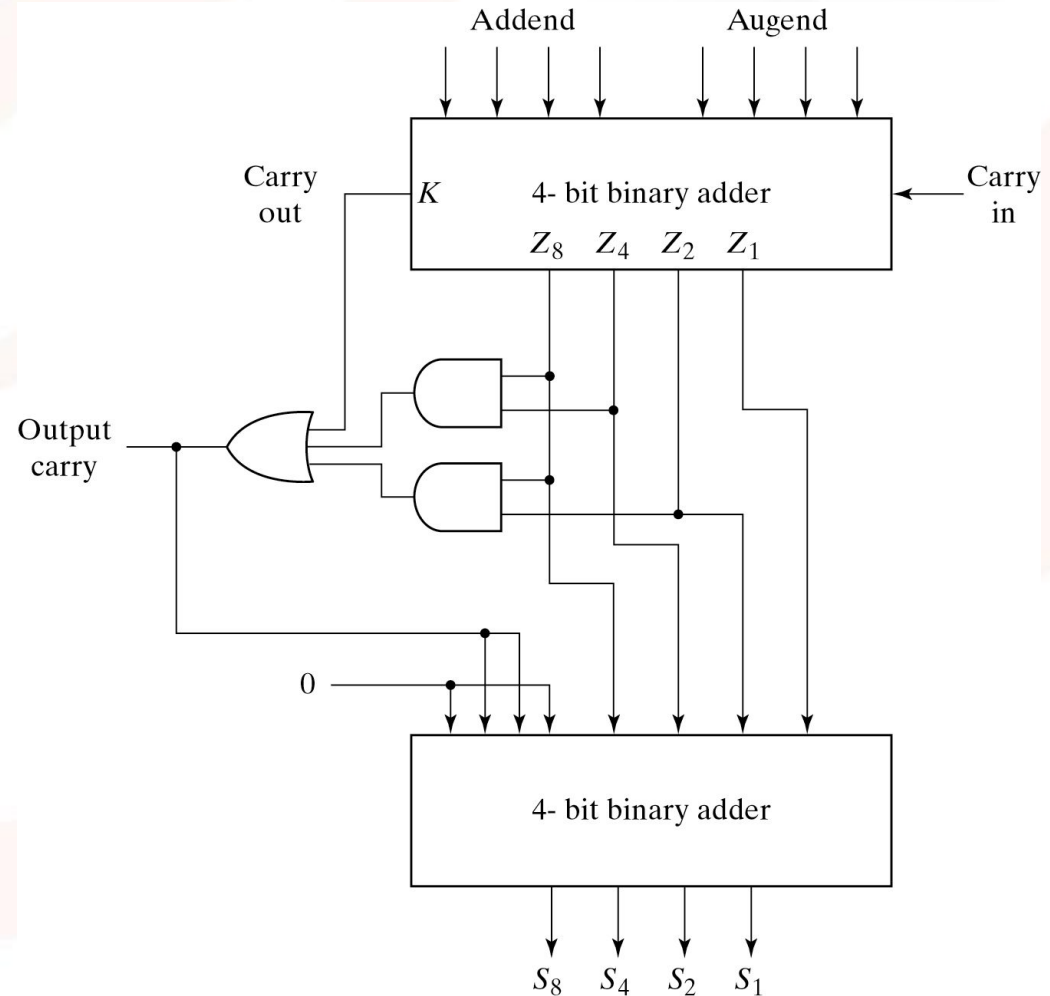
K	Binary Sum				BCD Sum					Decimal
	Z _B	Z ₄	Z ₂	Z ₁	C	S _B	S ₄	S ₂	S ₁	
0	0	0	0	0	0	0	0	0	0	0
0	0	0	0	1	0	0	0	0	1	1
0	0	0	1	0	0	0	0	1	0	2
0	0	0	1	1	0	0	0	1	1	3
0	0	1	0	0	0	0	1	0	0	4
0	0	1	0	1	0	0	1	0	1	5
0	0	1	1	0	0	0	1	1	0	6
0	0	1	1	1	0	0	1	1	1	7
0	1	0	0	0	0	1	0	0	0	8
0	1	0	0	1	0	1	0	0	1	9
0	1	0	1	0	1	0	0	0	0	10
0	1	0	1	1	1	0	0	0	1	11
0	1	1	0	0	1	0	0	1	0	12
0	1	1	0	1	1	0	0	1	1	13
0	1	1	1	0	1	0	1	0	0	14
0	1	1	1	1	1	0	1	0	1	15
1	0	0	0	0	1	0	1	1	0	16
1	0	0	0	1	1	0	1	1	1	17
1	0	0	1	0	1	1	0	0	0	18
1	0	0	1	1	1	1	0	0	1	19

Decimal Adder (BCD Adder)

- Modifications are needed if the sum > 9
 - $C = 1$
 - $K = 1$
 - $Z_8 Z_4 = 1$
 - $Z_8 Z_2 = 1$
 - modification: $-(10)_d$ or $+6$

Decimal Adder (BCD Adder)

■ Logic Circuit Diagram



Magnitude Comparator

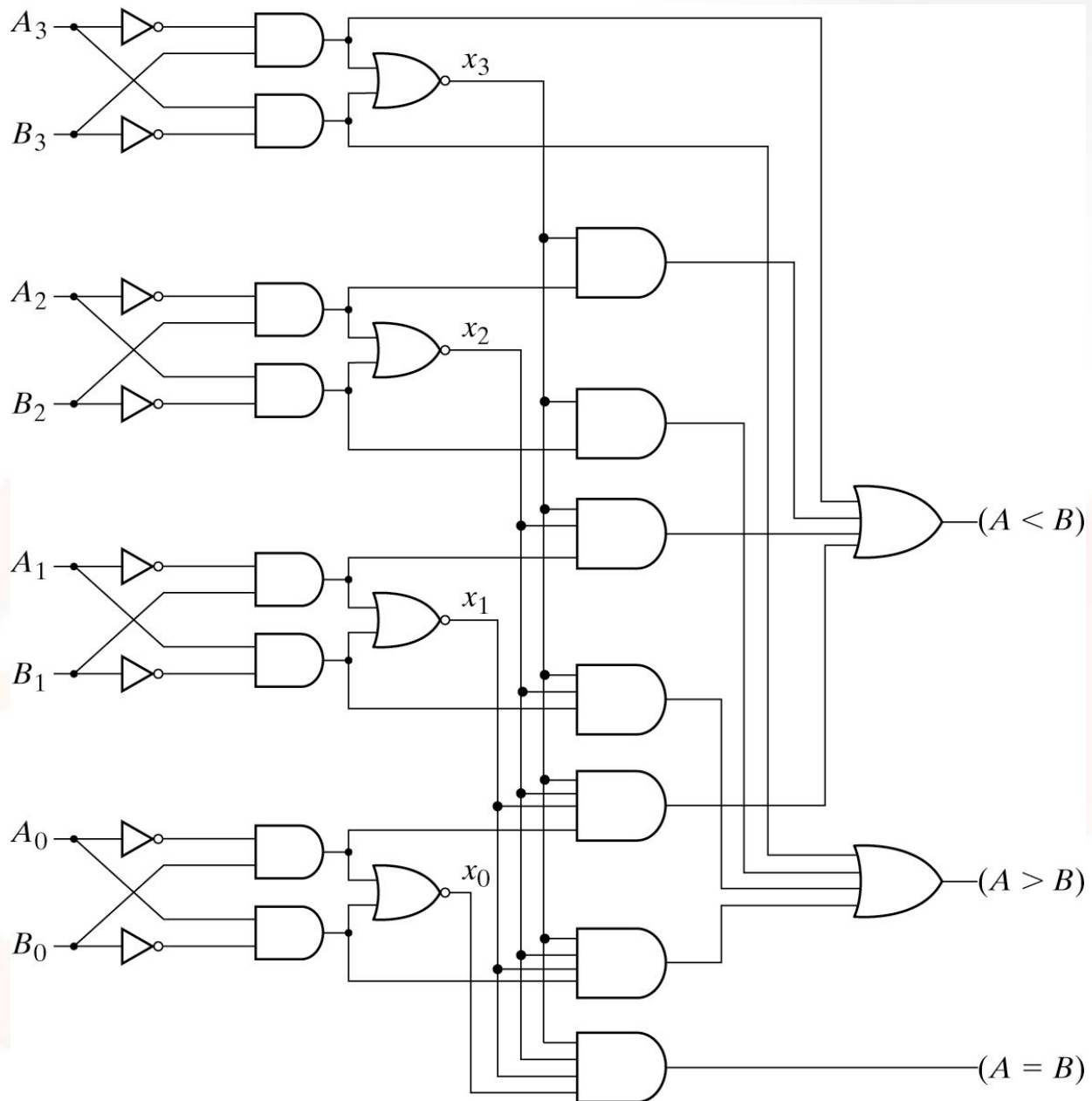
- The comparison of two numbers
 - outputs: $A > B$, $A = B$, $A < B$
- Design Approaches
 - the truth table
 - 2^{2n} entries - too cumbersome for large n
 - use inherent regularity of the problem
 - reduce design efforts
 - reduce human errors

■ Algorithm -> logic

- $A = A_3A_2A_1A_0$; $B = B_3B_2B_1B_0$
- $A=B$ if $A_3=B_3$, $A_2=B_2$, $A_1=B_1$ and $A_0=B_0$
 - equality: $x_i = A_iB_i + A_i'B_i'$
 - $(A=B) = x_3x_2x_1x_0$
- $(A>B) = A_3B_3' + x_3A_2B_2' + x_3x_2A_1B_1' + x_3x_2x_1A_0B_0'$
- $(A<B) = A_3'B_3 + x_3A_2'B_2 + x_3x_2A_1'B_1 + x_3x_2x_1A_0'B_0$

■ Implementation

- $x_i = (A_iB_i' + A_i'B_i)'$



Decoder

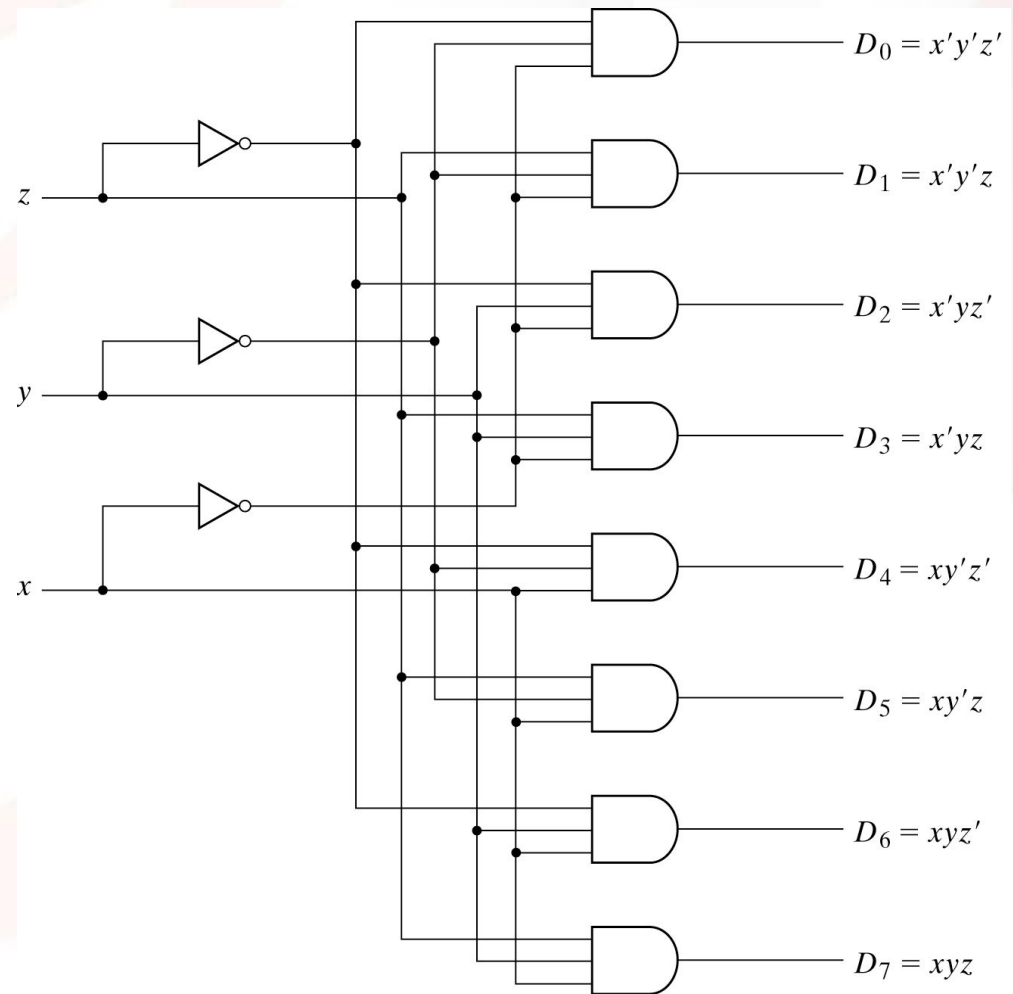
- A n-to-m decoder
 - a binary code of n bits = 2^n distinct information
 - n input variables; up to 2^n output lines
 - only one output is active at time

Truth Table of a 3-to-8-Line Decoder

Inputs			Outputs							
x	y	z	D_0	D_1	D_2	D_3	D_4	D_5	D_6	D_7
0	0	0	1	0	0	0	0	0	0	0
0	0	1	0	1	0	0	0	0	0	0
0	1	0	0	0	1	0	0	0	0	0
0	1	1	0	0	0	1	0	0	0	0
1	0	0	0	0	0	0	1	0	0	0
1	0	1	0	0	0	0	0	1	0	0
1	1	0	0	0	0	0	0	0	1	0
1	1	1	0	0	0	0	0	0	0	1

Decoder (cont.)

- An implementation of 3-to-8-line decoder



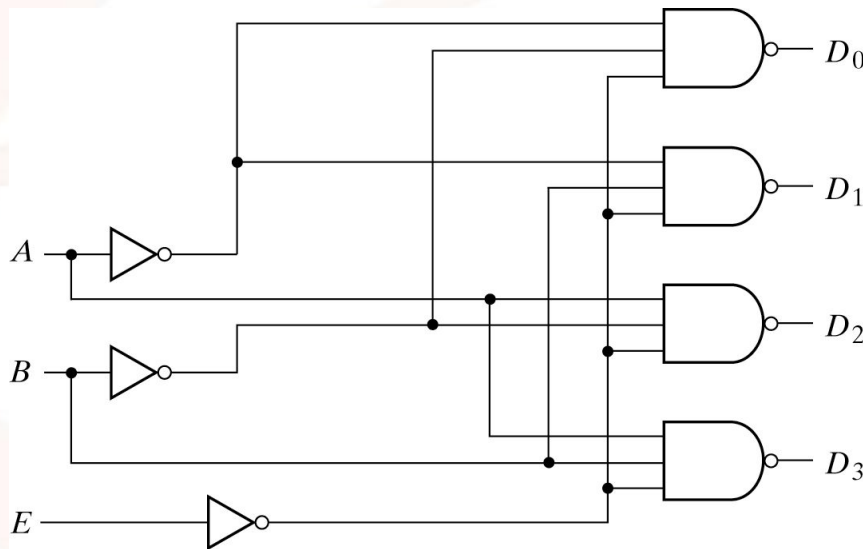
Decoder

- Combinational logic implementation
 - each output = a minterm
 - use a decoder and an external OR gate to implement any Boolean function of n input variables

Demultiplexers (DEMUX)

■ Demultiplexers

- a decoder with an enable input
- receive information on a single line and transmits it on one of 2^n possible output lines



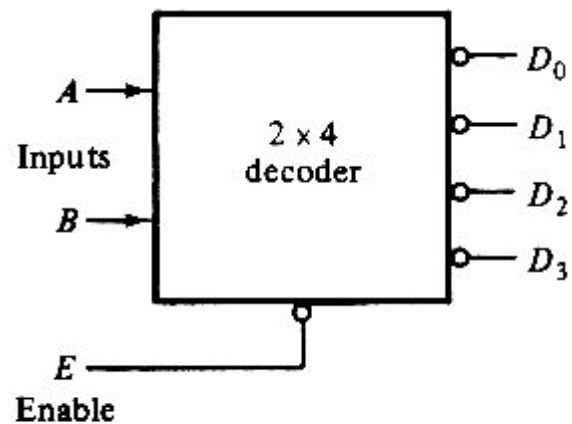
(a) Logic diagram

<i>E</i>	<i>A</i>	<i>B</i>	<i>D</i> ₀	<i>D</i> ₁	<i>D</i> ₂	<i>D</i> ₃
1	<i>X</i>	<i>X</i>	1	1	1	1
0	0	0	0	1	1	1
0	0	1	1	0	1	1
0	1	0	1	1	0	1
0	1	1	1	1	1	0

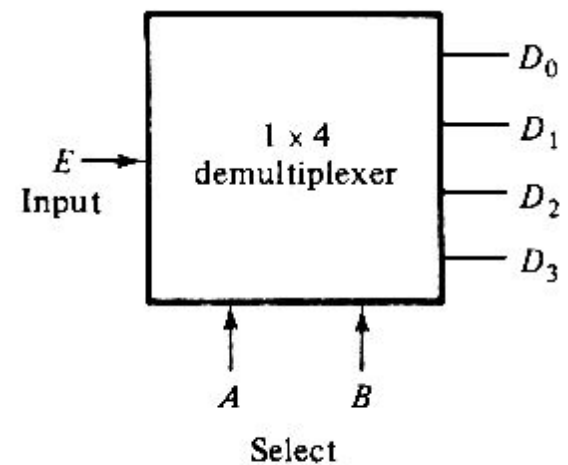
(b) Truth table

Demultiplexer (cont.)

- Decoder/demultiplexers



(a) Decoder with enable

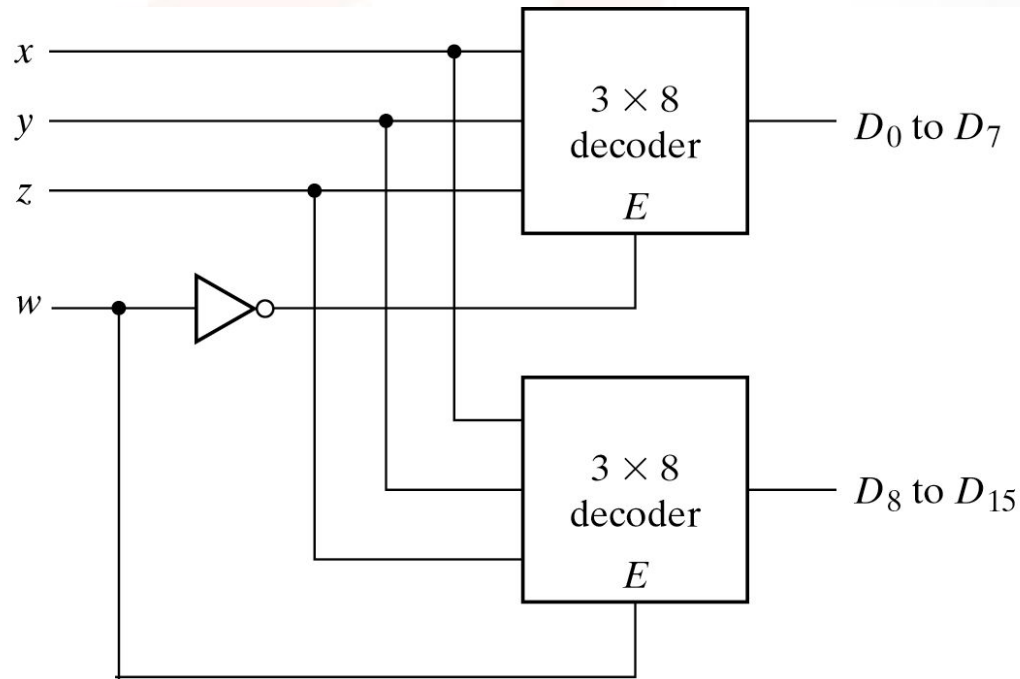


(b) Demultiplexer

Demultiplexer (cont.)

■ Expansion

- two 3-to-8 decoder: a 4-to-16 deocder

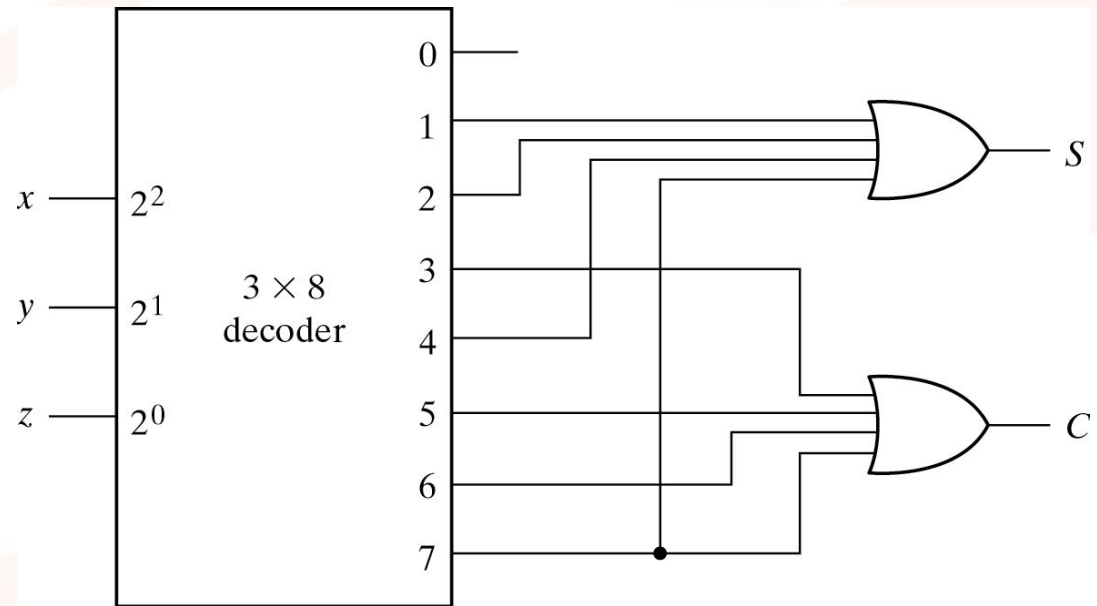


- a 5-to-32 decoder?

Combinational Logic Implementation using a decoder

- each output = a minterm
- use a decoder and an external OR gate to implement any Boolean function of n input variables
- A full-adder

- $S(x,y,z) = \Sigma(1,2,4,7)$
- $C(x,y,z) = \Sigma(3,5,6,7)$



Combinational Logic Implementation using a decoder

- two possible approaches using decoder
 - OR(minterms of F): k inputs
 - NOR(minterms of F'): $2^n - k$ inputs
- In general, it is not a practical implementation

Encoder

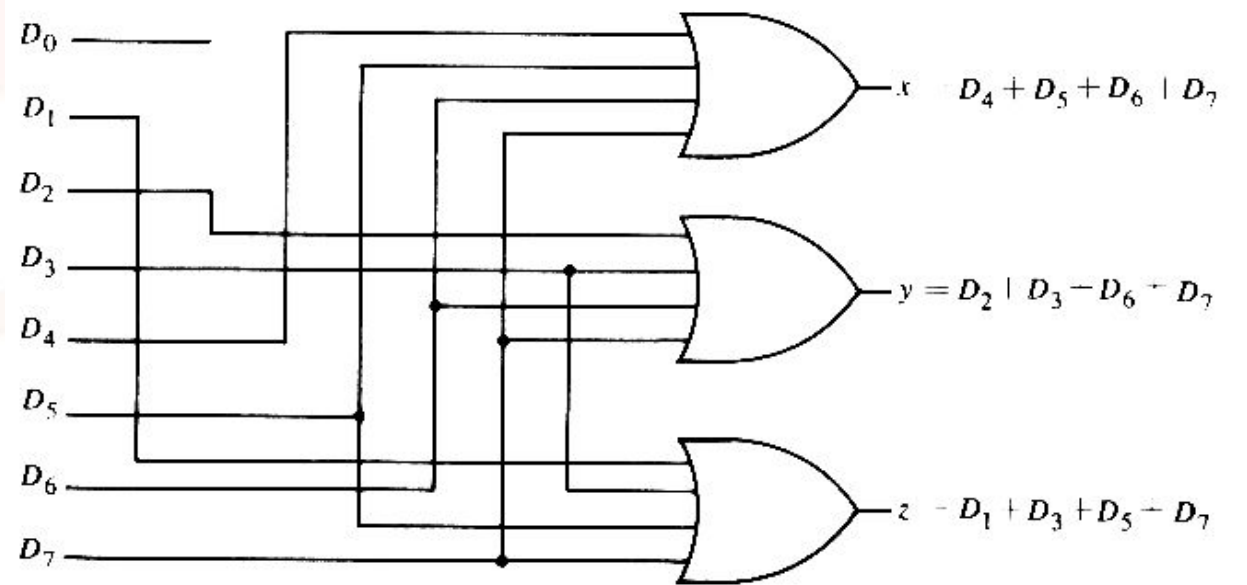
- The inverse function of a decoder

Truth Table of Octal-to-Binary Encoder

Inputs								Outputs		
D_0	D_1	D_2	D_3	D_4	D_5	D_6	D_7	x	y	z
1	0	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	0	1
0	0	1	0	0	0	0	0	0	1	0
0	0	0	1	0	0	0	0	0	1	1
0	0	0	0	1	0	0	0	1	0	0
0	0	0	0	0	1	0	0	1	0	1
0	0	0	0	0	0	1	0	1	1	0
0	0	0	0	0	0	0	1	1	1	1

Encoder

■ An implementation



■ limitations

- illegal input: e.g. $D_3 = D_6 = 1$
- the output = 111 ('13 and '16)

Priority Encoder

- resolve the ambiguity of illegal inputs
- only one of the input is encoded

Truth Table of a Priority Encoder

Inputs				Outputs		
D_0	D_1	D_2	D_3	x	y	V
0	0	0	0	X	X	0
1	0	0	0	0	0	1
X	1	0	0	0	1	1
X	X	1	0	1	0	1
X	X	X	1	1	1	1

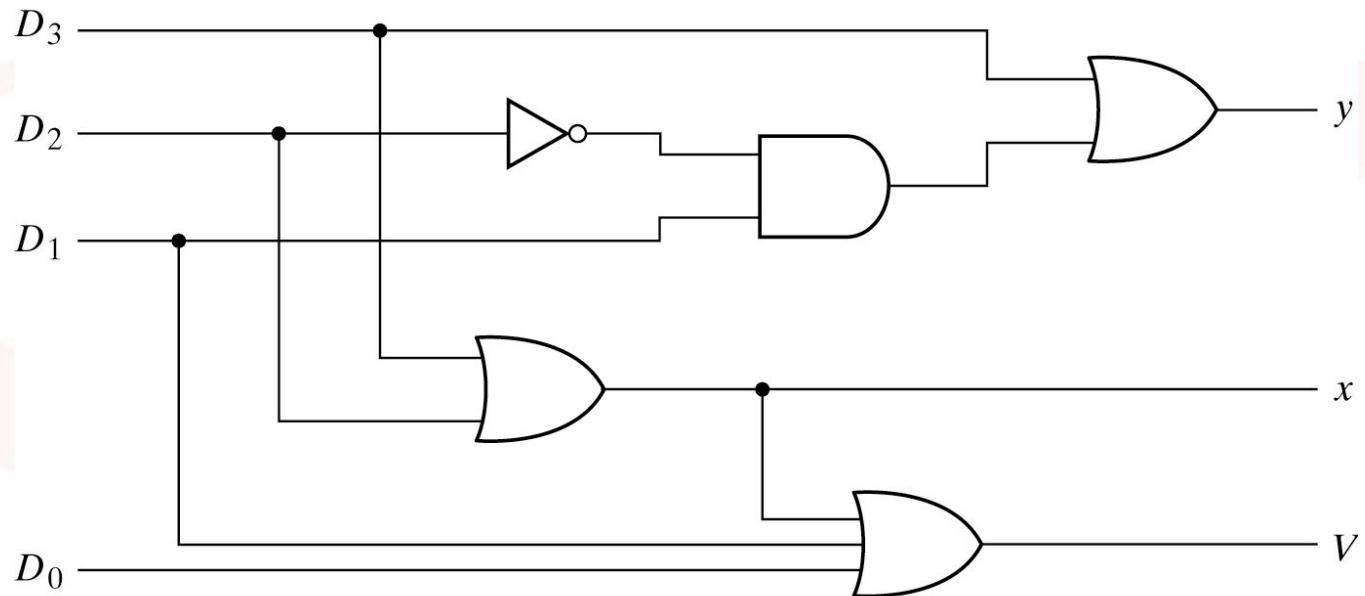
- D_3 has the highest priority
- D_0 has the lowest priority
- X: don't-care conditions
- V: valid output indicator

		D_2				
		00	01	11	10	
D_0	00	X	1	1	1	D_1
	01		1	1	1	
	11		1	1	1	
	10		1	1	1	
		D_3				

$$x = D_2 + D_3$$

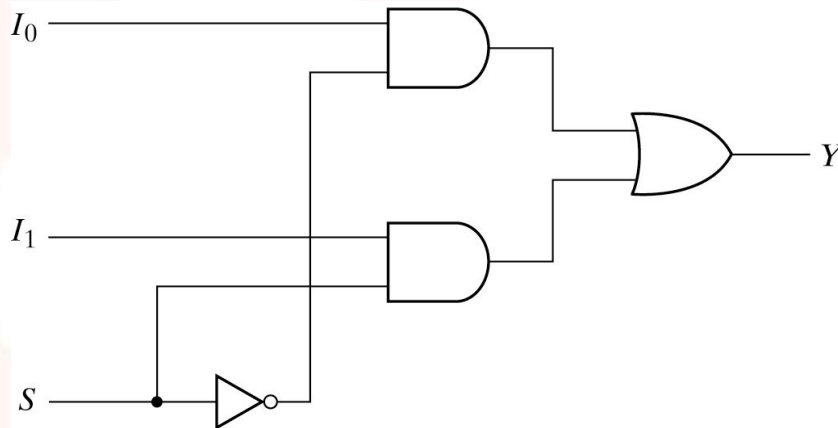
		D_2				
		00	01	11	10	
D_0	00	X	1	1	1	D_1
	01	1	1	1	1	
	11	1	1	1	1	
	10		1	1	1	
		D_3				

$$y = D_3 + D_1 D'_2$$

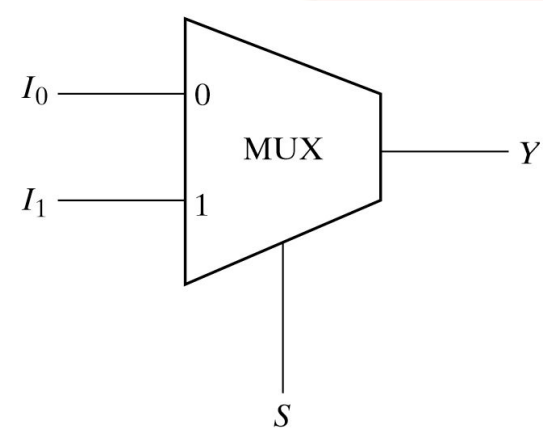


Multiplexer

- select binary information from one of many input lines and direct it to a single output line
- 2^n input lines, n selection lines and one output line
- e.g.: 2-to-1-line multiplexer



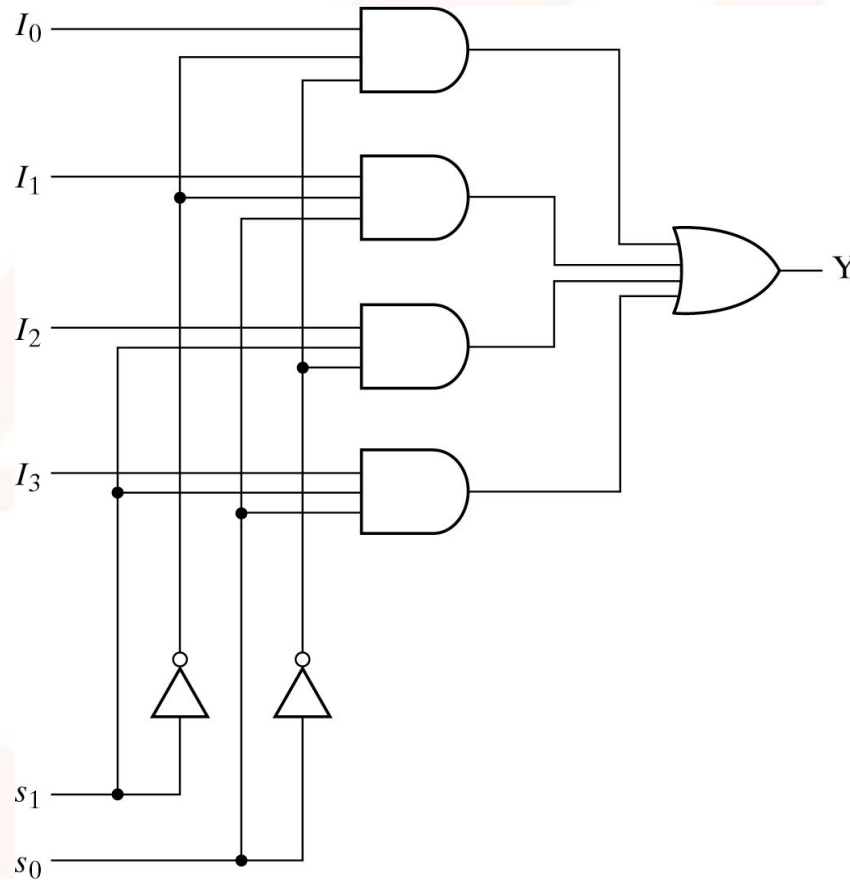
(a) Logic diagram



(b) Block diagram

Multiplexer

■ 4-to-1-line multiplexer



(a) Logic diagram

s_1	s_0	Y
0	0	I_0
0	1	I_1
1	0	I_2
1	1	I_3

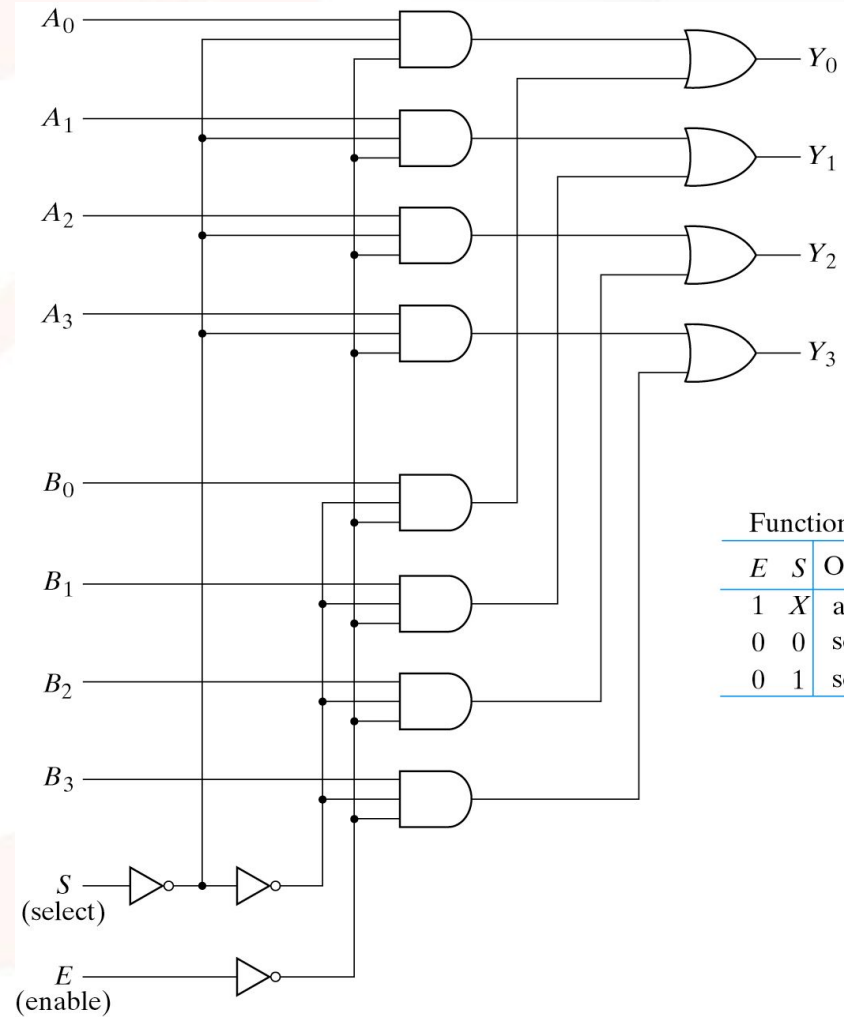
(b) Function table

Multiplexer

- Implementation note
 - n-to- 2^n decoder
 - add the 2^n input lines to each AND gate
 - OR(all AND gates)
 - an enable input (an option)

Multiplexer

Quadruple 2-to-1 MUX



Function table		
E	S	Output Y
1	X	all 0's
0	0	select A
0	1	select B

Boolean function implementation using a multiplexer

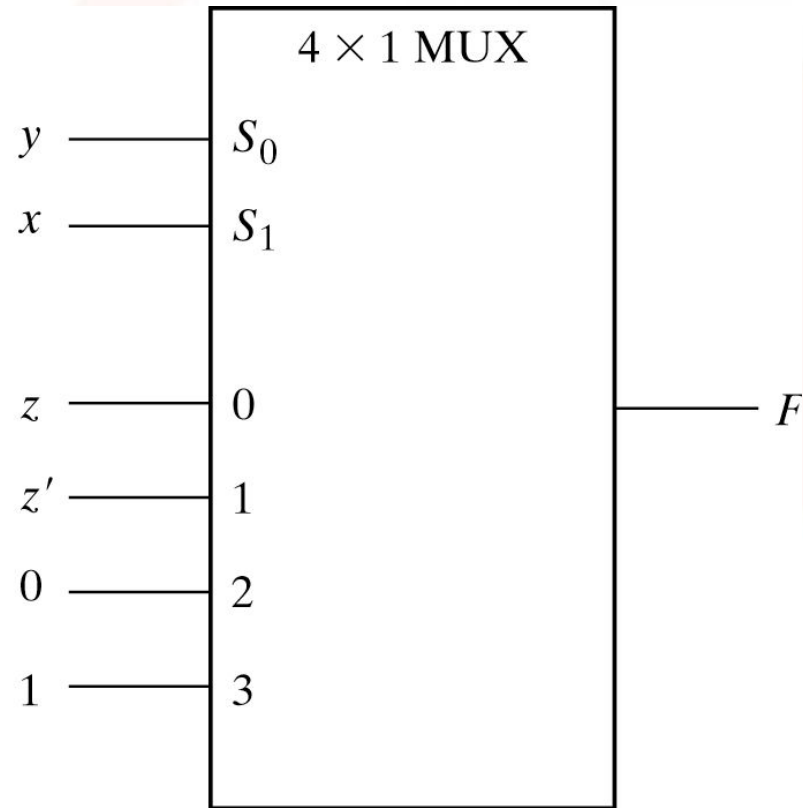
- MUX: a decoder + an OR gate
- 2^n -to-1 MUX can implement any Boolean function of n input variable
- a better solution: implement any Boolean function of $n+1$ input variable
 - n of these variables: the selection lines
 - the remaining variable: the inputs

Boolean function implementation using a multiplexer

- an example: $F(A,B,C)=\Sigma(1,2,6,7)$

x	y	z	F	
0	0	0	0	
0	0	1	1	$F = z$
0	1	0	1	
0	1	1	0	$F = z'$
1	0	0	0	
1	0	1	0	$F = 0$
1	1	0	1	
1	1	1	1	$F = 1$

(a) Truth table



(b) Multiplexer implementation

Boolean function implementation using a multiplexer

■ Procedure:

- assign an ordering sequence of the input variable
- the rightmost variable (D) will be used for the input lines
- assign the remaining $n-1$ variables to the selection lines w.r.t. their corresponding sequence
- construct the truth table
- consider a pair of consecutive minterms starting from m_0
- determine the input lines

Boolean function implementation using a multiplexer

<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>F</i>	
0	0	0	0	0	
0	0	0	1	1	$F = D$
0	0	1	0	0	
0	0	1	1	1	$F = D$
0	1	0	0	1	
0	1	0	1	0	$F = D'$
0	1	1	0	0	
0	1	1	1	0	$F = 0$
1	0	0	0	0	
1	0	0	1	0	$F = 0$
1	0	1	0	0	
1	0	1	1	1	$F = D$
1	1	0	0	1	
1	1	0	1	1	$F = 1$
1	1	1	0	1	
1	1	1	1	1	$F = 1$

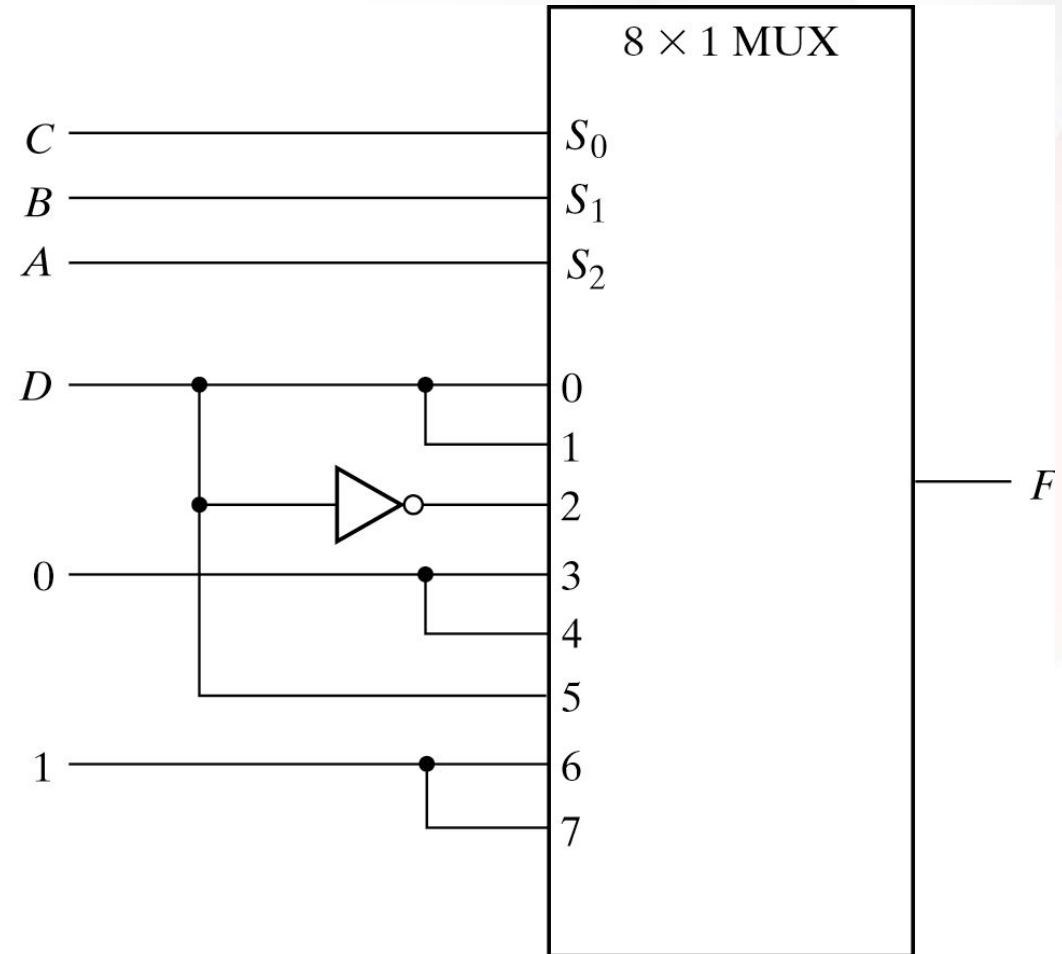
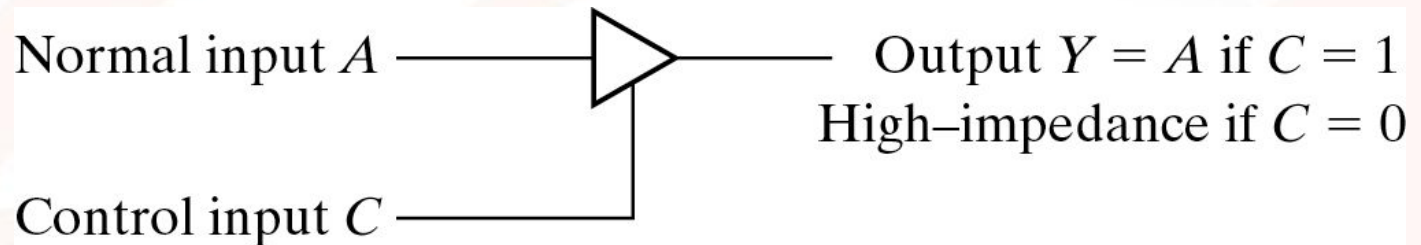


Fig. 4-28 Implementing a 4-Input Function with a Multiplexer

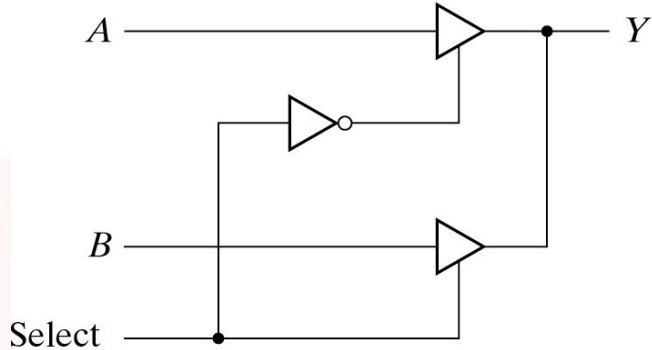
Three-state gates

- A multiplexer can be constructed with three-state gates
- Output state: 0, 1, and high-impedance (open circuit)

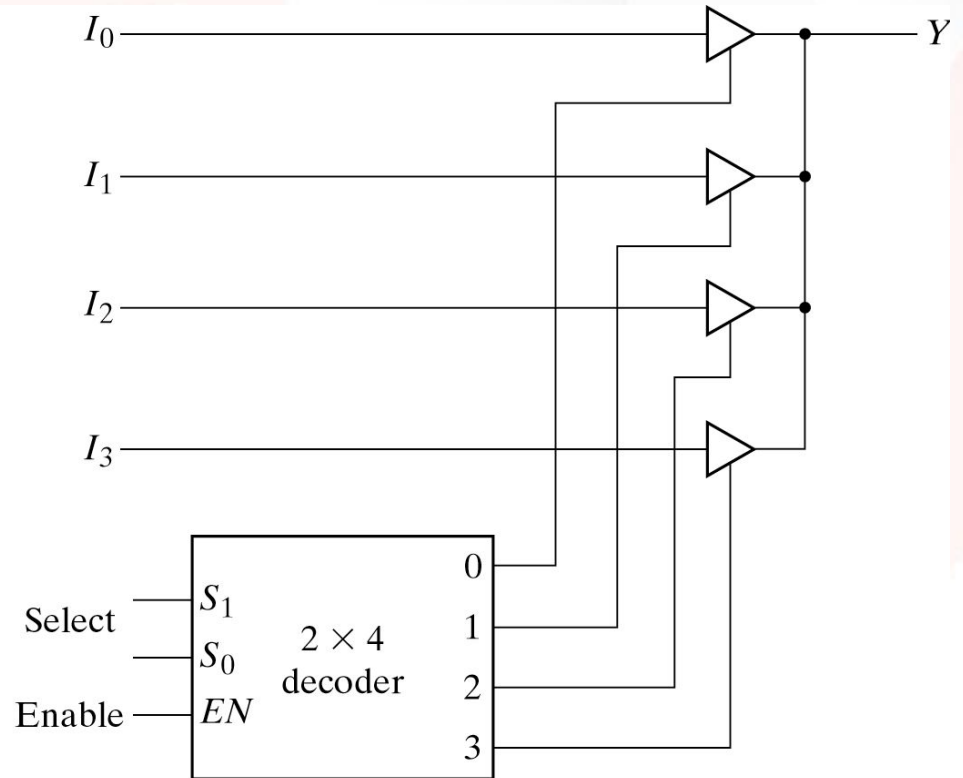


Tri state buffer

Multiplexer using three state gates



(a) 2-to-1- line mux



(b) 4 - to - 1 line mux