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Performance Comparison

Solving MC Challenges

M4GB: An Efficient Gröbner Basis Algorithm

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Problem

 $\mathbb{F}[x_1,\ldots,x_n]$ - a polynomial ring over a field \mathbb{F} together with an admissible monomial ordering <.

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Problem

 $\mathbb{F}[x_1,\ldots,x_n]$ - a polynomial ring over a field \mathbb{F} together with an admissible monomial ordering <.

Problem (MQ-problem)

Let $n, m \in \mathbb{Z}_{>0}$. Given $f_1, \ldots, f_m \in \mathbb{F}[x_1, \ldots, x_n]$ with f_i be quadratic polynomials, find a $(a_1, \ldots, a_n) \in \mathbb{F}^n$ such that $f_i(a_1, \ldots, a_n) = 0$ for all $i = 1, \ldots, m$.

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Notations

$$f = -15x^2 + 8xy - 13z^2 - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

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Example

$$f = -15x^2 + 8xy - 13z^2 - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

• $LM(f) = x^2$ (the leading monomial of f)

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$$f = -15x^{2} + 8xy - 13z^{2} - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

- LM(f) = x^2 (the leading monomial of f)
- LC(f) = -15 (the leading coefficient of f)

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$$f = -15x^2 + 8xy - 13z^2 - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

- LM(f) = x^2 (the leading monomial of f)
- LC(f) = -15 (the leading coefficient of f)
- LT(f) = $-15x^2$ (the leading term of f)

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$$f = -15x^2 + 8xy - 13z^2 - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

- $LM(f) = x^2$ (the leading monomial of f)
- LC(f) = -15 (the leading coefficient of f)
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- Tail $(f) = 8xy 13z^2 4x + 11z$ (the tail of f)

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$$f = -15x^2 + 8xy - 13z^2 - 4x + 11z \in \mathbb{F}_{31}[x, y, z]$$

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Polynomial Reduction

Theorem

Let $G = (g_1, \ldots, g_t)$ be a nonempty ordered finite subset of $\mathbb{F}[x_1, \ldots, x_n]$. Then every polynomial $f \in \mathbb{F}[x_1, \ldots, x_n]$ can be written as

$$f = q_1g_1 + \ldots + q_tg_t + r,$$

where $q_1, \ldots, q_t, r \in \mathbb{F}[x_1, \ldots, x_n]$ and either r = 0 or none of terms of r is divisible by any of $\mathsf{LT}(g_1), \ldots, \mathsf{LT}(g_t)$.

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Polynomial Reduction

Theorem

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$$r \leftarrow \text{FullReduce}(f, G)$$

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Gröbner basis

Definition

Let $I \neq \{0\}$ be an ideal of $\mathbb{F}[x_1, \dots, x_n]$. A finite subset $G \subseteq I$ that generates I is a Gröbner basis of I if for all $f \in I$, there exists $g \in G$ such that $\mathsf{LT}(g) \mid \mathsf{LT}(f)$.

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S-polynomial

Definition

Let $f, g \in \mathbb{F}[x_1, \dots, x_n]$ be nonzero polynomials and let $x^{\gamma} = \mathsf{LCM}(\mathsf{LM}(f), \mathsf{LM}(g))$. The S-polynomial of f and g is defined as

$$\mathsf{Spoly}(f,g) = \frac{x^{\gamma}}{\mathsf{LT}(f)} \cdot f - \frac{x^{\gamma}}{\mathsf{LT}(g)} \cdot g.$$

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Buchberger's Algorithm

Input: A finite ordered subset $F \subseteq \mathbb{F}[x_1, \dots, x_n]$ **Result:** A Gröbner basis G such that $\langle G \rangle = \langle F \rangle$

1
$$P \leftarrow \{\{p,q\} : \forall p,q \in F \text{ and } p \neq q\}$$

$$2 G \leftarrow F$$

6

7

8

3 while
$$P \neq \{\}$$
 do

4
$$\{p,q\} \leftarrow \text{SELECT}(P)$$

5 $P \leftarrow P \setminus \{\{p,q\}\}\}$

$$P \leftarrow P \setminus \{\{p,q\}\}$$

$$r \leftarrow \text{FullReduce}(\mathsf{Spoly}(p,q),G)$$

if
$$r \neq 0$$
 then

$$P \leftarrow P \cup \{\{r,g\} : \forall g \in G\}$$

$$9 \quad \bigsqcup \quad G \leftarrow G \cup (r)$$

10 return G

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M4GB (Simplified)

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Solving MQ Challenges • Implemented using C++11

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- Implemented using C++11
- Comparison with existing implementations
 - 1 FGb C Interface Implementation by Jean Charles Faugere¹
 - 2 Magma v2.20-6
 - 3 OpenF4 v1.0.1 Open source implementation by Coladon, Vitse and Joux².

¹Available at http://www-polsys.lip6.fr/~jcf/FGb/C/index.html

²Available at https://github.com/nauotit/openf4

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- Implemented using C++11
- Comparison with existing implementations
 - 1 FGb C Interface Implementation by Jean Charles Faugere¹
 - 2 Magma v2.20-6
 - 3 OpenF4 v1.0.1 Open source implementation by Coladon, Vitse and $Joux^2$.
- Test cases
 - $\ensuremath{\text{1}}$ Dense polynomials with coefficients in $\ensuremath{\mathbb{F}}_{31}$
 - 2 m = 2n and m = n + 1.

¹Available at http://www-polsys.lip6.fr/~jcf/FGb/C/index.html

²Available at https://github.com/nauotit/openf4

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Benchmark for m = 2n

			Total CPU time (sec)				
n	m	OpenF4 FGb Magma (projected)		Magma (projected)	M4GB		
20	40	206	470	232.17	57		
21	42	472	1002	500.26	170		
22	44	1145	3118	1616.73	424		
23	46	2274	6849	3184.82	1060		
24	48	10293	64700	31167.61	2556		
25	50	-	151653	77678.58	5575		
26	52	-	360055	183628.74	15517		
27	54	-	767543	409451.87	46548		

		Memory (MB)				
20	40	4240	112	361.84	73	
21	42	6640	165	577.34	121	
22	44	14368	525	853.84	226	
23	46	26135	918	1324.16	395	
24	48	161945	1561	8872.94	663	
25	50	-	2765	19718.78	1471	
26	52	-	4607	25197	3328	
27	54	-	8180	39844.84	6799	

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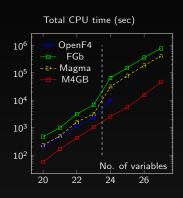
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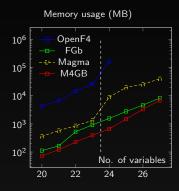
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Graph for m = 2n





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Benchmark for m = n + 1

		Total CPU time (sec)				
n	m	OpenF4	FGb	FGb Magma (projected)		
10	11	2.99	5	3.29	0.98	
11	12	8.73	21	11.172	2.6	
12	13	36.76	134	59.08	13.92	
13	14	172.49	642	286.4	58.18	
14	15	1258	5850	2810.75	393.19	
15	16	7225	36361	17265.5	2424	
Memory (MB)						
10	11	101	33	32.09	17	
11	12	341	50	64.12	16	
12	13	1463	112	113.59	31	
13	14	7622	323	281.53	74	
14	15	33460	1098	1104	250	
15	16	117396	4118	3320	837	

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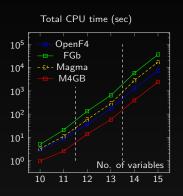
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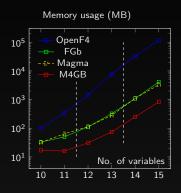
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Graph for m = n + 1





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- MQ-based public key and digital signature are candidates of post-quantum cryptography.
- Their security relies on the difficulty of finding a solution of an MQ problem.
- Need to understand its difficulty in practice

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Solving MQ Challenges

- Started on 1st April 2015
- Six different type of challenges

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Solving MQ Challenges

- Started on 1st April 2015
- Six different type of challenges
- Type I, II, and III are encryption-type parameter (m=2n) and coefficients in $\mathbb{F}_2, \mathbb{F}_{2^8}, \mathbb{F}_{31}$ respectively.

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- Started on 1st April 2015
- Six different type of challenges
- Type I, II, and III are encryption-type parameter (m = 2n) and coefficients in F₂, F₂₈, F₃₁ respectively.
- Type IV, V, and VI are signature-type parameter $(n \approx 1.5m)$ and coefficients in $\mathbb{F}_2, \mathbb{F}_{2^8}, \mathbb{F}_{31}$ respectively.

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- Started on 1st April 2015
- Six different type of challenges
- Type I, II, and III are encryption-type parameter (m = 2n) and coefficients in F₂, F₂₈, F₃₁ respectively.
- Type IV, V, and VI are signature-type parameter $(n \approx 1.5m)$ and coefficients in $\mathbb{F}_2, \mathbb{F}_{2^8}, \mathbb{F}_{31}$ respectively.
- Parameter Choice: Require at least one month for Magma 2.19-9 to solve using Four 6-cores Intel(R) Xeon(R) CPU E5-4617 @ 2.9GHz and 1TB of RAM.

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Fukuoka MQ Challenge

- Started on 1st April 2015
- Six different type of challenges
- Type I, II, and III are encryption-type parameter (m = 2n) and coefficients in F₂, F₂₈, F₃₁ respectively.
- Type IV, V, and VI are signature-type parameter $(n \approx 1.5m)$ and coefficients in $\mathbb{F}_2, \mathbb{F}_{2^8}, \mathbb{F}_{31}$ respectively.
- Parameter Choice: Require at least one month for Magma 2.19-9 to solve using Four 6-cores Intel(R) Xeon(R) CPU E5-4617 @ 2.9GHz and 1TB of RAM.

https://www.mqchallenge.org

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Solving Signature-type MQ Challenge

- Hybrid approach : trade-off between exhaustive search and computing Gröbner bases
- Idea :
 - **1** Select a random vector $(a_1, \ldots, a_{n-m}) \in \mathbb{F}_q^{n-m}$
 - 2 Construct a new system with n = m

$$\tilde{F} = \{f(x_1, \ldots, x_m, a_1, \ldots, a_{n-m}) : \forall f \in F\}$$

- 3 Select $k \in \{1, ..., m\}$ and construct q^k subsystems from \tilde{F} by substituting k variables with all elements of \mathbb{F}_a^k .
- 4 Each subsystem generated can be solved in parallel.

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Computational Resources

A) Desktop machine with Intel(R) Core(TM) i7-2600K CPU @ $3.40 \mathrm{GHz}$ and $16 \mathrm{GB}$ RAM

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Computational Resources

- A) Desktop machine with Intel(R) Core(TM) i7-2600K CPU @ 3.40GHz and 16GB RAM
- B) NUMA machine with two nodes of Intel(R) Xeon(R) CPU E5-2650 v3 @ 2.30GHz and 128GB RAM each.

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Туре	n/m	Machine Used	# Node	Duration

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Туре	n/m	Machine Used	# Node	Duration
V	24/16 25/17			
V	25/17			
V	27/18			

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Туре	n/m	Machine Used	# Node	Duration
V	24/16	А	1	pprox 9.3 hours
V	25/17			
V	27/18			

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Туре	n/m	Machine Used	# Node	Duration
V	24/16	А	1	pprox 9.3 hours
V	25/17	В	1	pprox 46.33 hours
V	27/18	В	2	pprox 10.9 days

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Туре	n/m	Machine Used	# Node	Duration
V	24/16	А	1	pprox 9.3 hours
V	25/17	В	1	pprox 46.33 hours
V	27/18	В	2	pprox 10.9 days
VI	24/16			
VI	25/17			
VI	27/18			
VI	28/19			

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Туре	n/m	Machine Used	# Node	Duration
V	24/16	А	1	pprox 9.3 hours
V	25/17	В	1	pprox 46.33 hours
V	27/18	В	2	pprox 10.9 days
VI	24/16	А	1	pprox 1.2 hours
VI	25/17			
VI	27/18			
VI	28/19			

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Туре	n/m	Machine Used	# Node	Duration
V	24/16	А	1	pprox 9.3 hours
V	25/17	В	1	pprox 46.33 hours
V	27/18	В	2	pprox 10.9 days
VI	24/16	А	1	pprox 1.2 hours
VI	25/17	В	1	pprox 9.87 hours
VI	27/18	В	1	pprox 31.48 hours
VI	28/19	В	2	pprox 7.61 days

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Solving MQ Challenges https://github.com/cr-marcstevens/m4gb

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Solving MQ Challenges Question ?