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Advanced Data Structures Lab  
Batch - 5  
Program - 5

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Insertion :-

insert (node, compare value)

locate a leaf to put value in it  
if leaf is a 2 node, make it a 3 node  
inserting the value appropriately  
if leaf is a 3 node, split the node

Split (node N)

if N is root  
make middle child into a 2 node  
make small & large key into 2 node  
Reduce children

else N has a parent P → move middle key to P  
make small and large children into a  
2 node  
Reducing children.

insert (a, r) {

if ( r consists of a single leaf labeled b )  
create a new ~~node~~ root r'  
create a new leaf v labeled a  
make l & v children of r'  
update L & M for r'

else

set f to search (a, r)  
create a new leaf l labeled a

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```
if f has 2 children
    insert l into proper position
    update L & H
else
    create a transitory node tree at f
    Addchild(f)
}
```

```
Addchild(v) {
    create new node v'
    move 2 right most children of v to v'
    if (v has no parent)
        make new root r'
        make v: left child & v': right child
        update L & H
    else
        let f be parent of v
        make v' child of f immediately to right of v
        if f now has 4 children
            Addchild(f)
        else
            update L and H.
}
```

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delete () {

let  $f$  be parent of node just deleted  
while ( $f$  is an illegal interior node) {

if ( $f$  has no parent)

make single child of  $f$ . new node  
delete  $f$ .

set  $of$  to the root

else

let  $g$  be parent of  $f$

if (one of  $f$ 's sibling is a 3 node)  
move one child from 3-node into  
 $f$ , update  $k_1$  &  $k_2$  everywhere

else

give  $f$ 's remaining child to one  $f$ 's  
sibling

delete  $f$

update  $k_1$  &  $k_2$  everywhere

set  $f$  to  $g$

}  
}

Autograph.