

- Breadth-First Search (BFS)

- Queue ADT

Ring buffer : Implementation technique for queues.

Queue $\langle T \rangle$

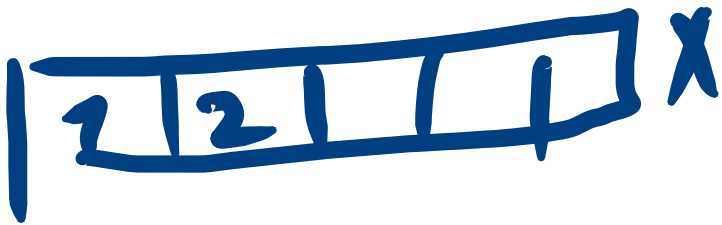
FIFO ~ First-in First-out

is.empty(q)

enqueue(q, e) \rightarrow adds elt to queue

dequeue(q) \rightarrow returns the earliest element in the queue and removes it.

Array Implementation of Queue



front back

$e(3)$

$e(1)$

$e(2)$

$d() \rightarrow 3$

$\leftarrow e() \text{ \& } d() \text{ are now } O(1) \text{ operations.}$



$e(3)$
 $e(4)$



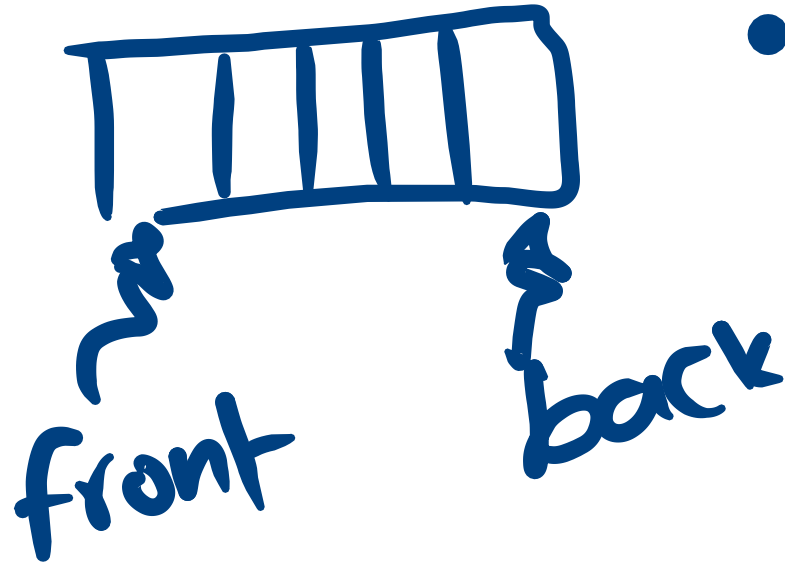
$d()$

$d()$

$d()$

front back

How to solve? Imagine first elt of array follows the last.
(Ring buffer)



• len of the queue.

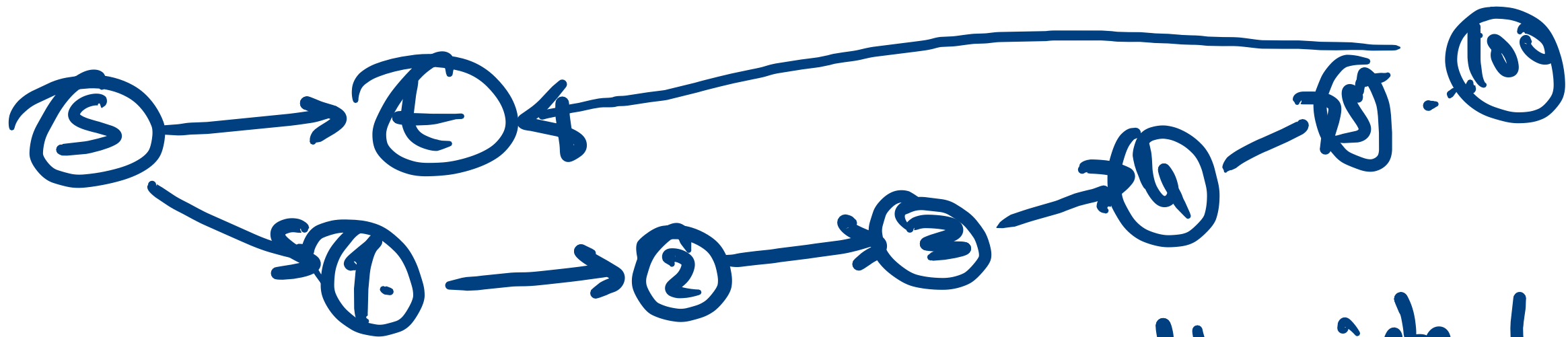
Why BFS?

8-Puzzle

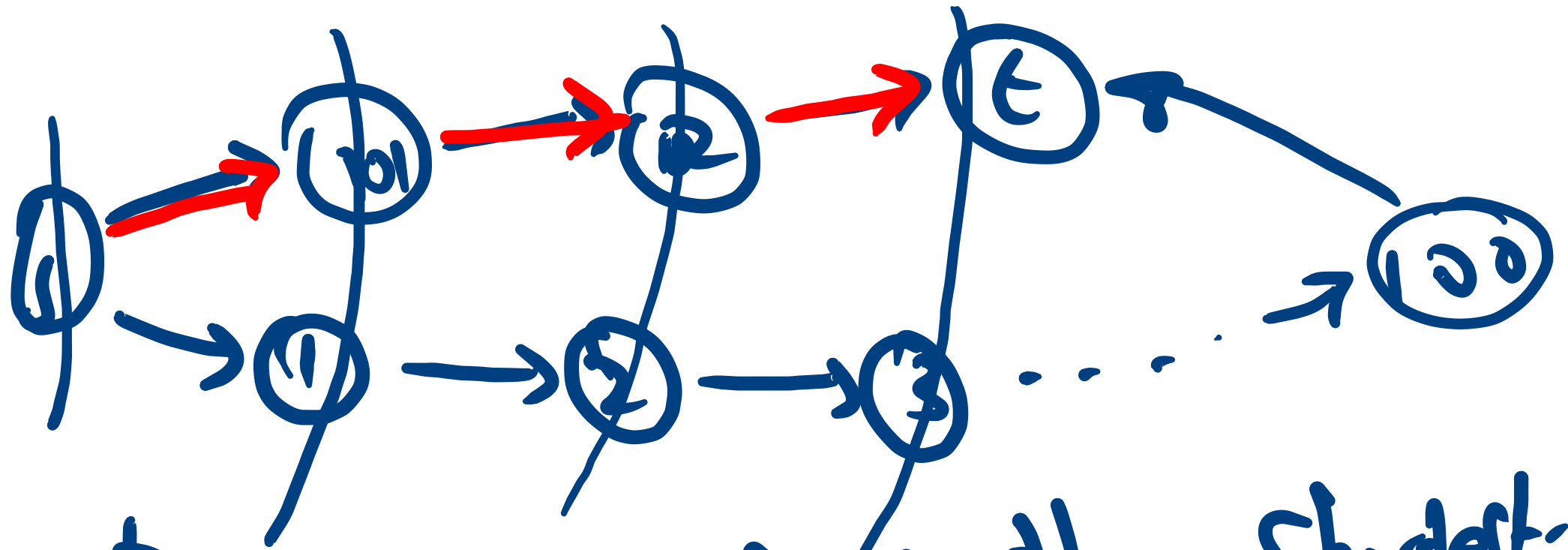
$G = (V = \text{Board state})$

$E = \leq 4 \text{ neighbors/board state}$

Why not a DFS to find a path to the solution?



DFS may find long paths instead of short paths.



BFS would find the Shortest-path
from s to t .

'BFS is obtained by switching
Stack in DFS with a queue.

bfs(G, source)

queue $\leftarrow \{ \text{source} \}$

visited[v] \leftarrow false $\forall v \in V(G)$

Parent[v] \leftarrow \perp $\forall v \in V(G)$

while queue is not empty

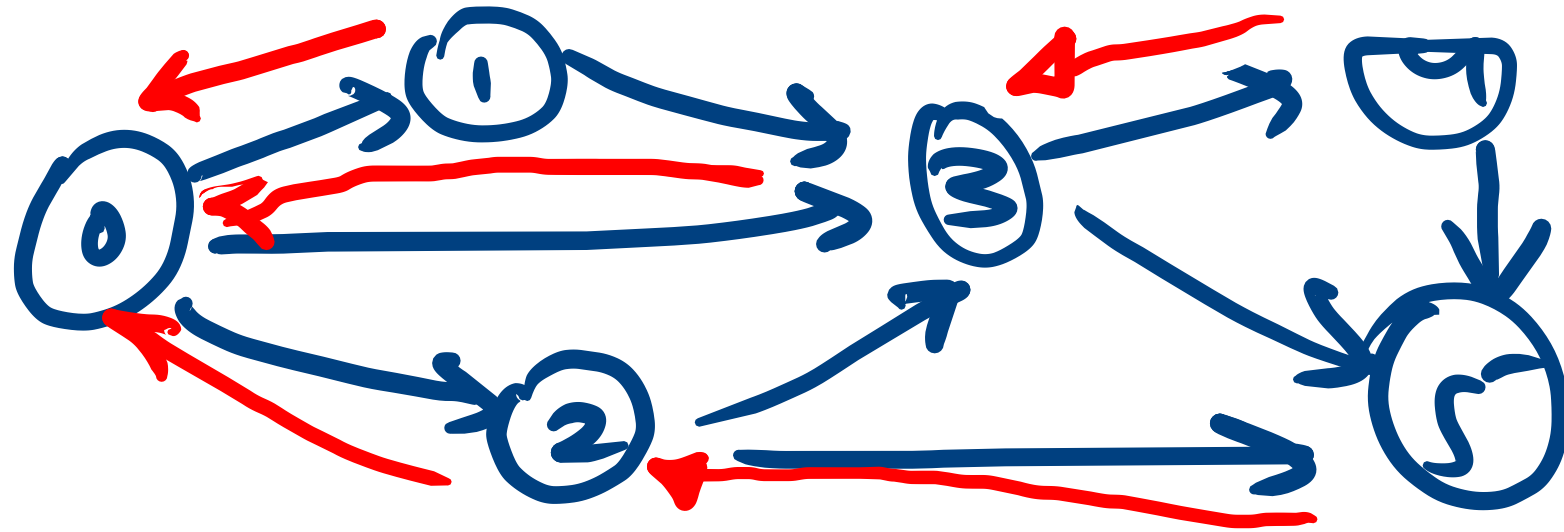
$u \leftarrow \text{dequeue}(q)$

 if visited[u] is false

 visited[u] \leftarrow true

 for each $u \rightarrow v$

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if visited[v] = false  
    enqueue(q, v)  
    Parent[v] ← u
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$q \leftarrow \{0\}$
 $P[v] \leftarrow v \quad v \in 0..5$

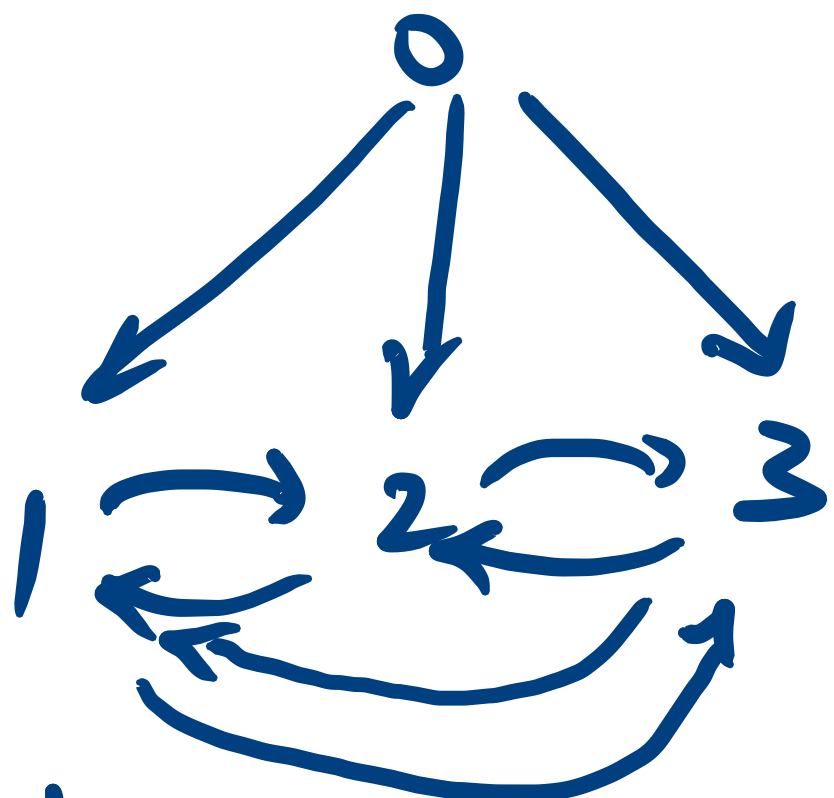
4) $u \leftarrow 3$
 $V[u] \leftarrow \text{true}$
 $P[u] \leftarrow 3$
 $e(u)$

1: $V[0] \leftarrow \text{true}$
 $e(1)$
 $e(2)$
 $e(3)$

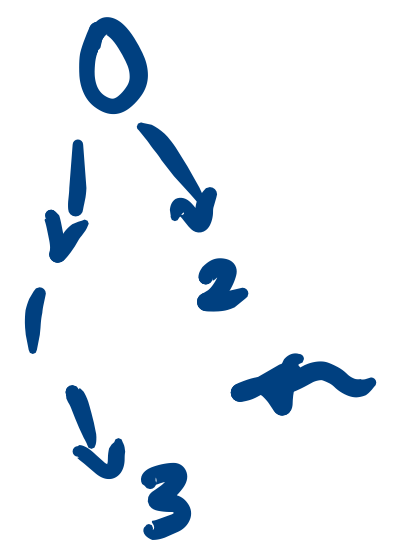
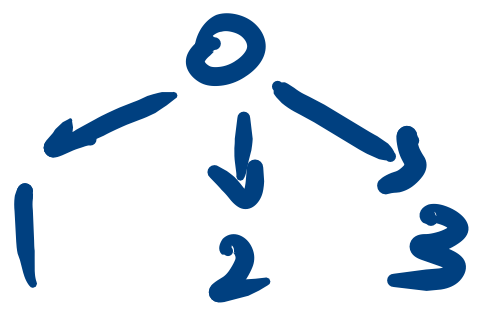
$q \leftarrow 1 \ 2 \ 3$
 $V[1..3] \leftarrow \text{true}$
 $P[1..3] \leftarrow 0$

2: $u \leftarrow 1$

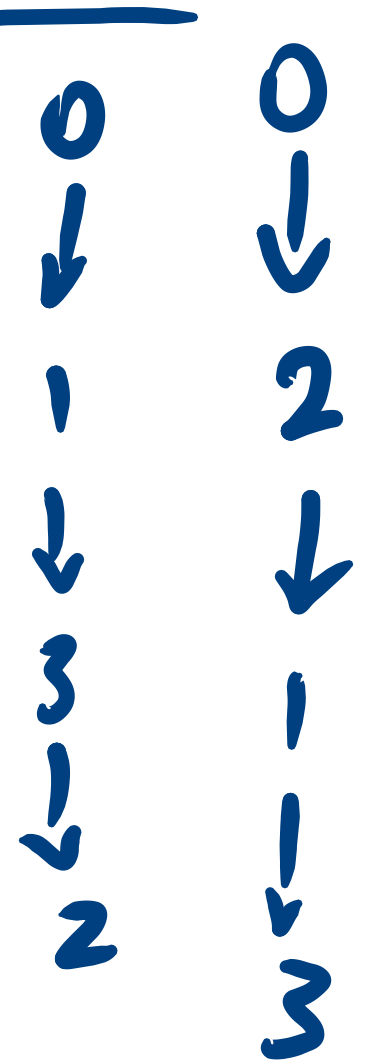
3: $u \leftarrow 2$
 $e(5)$
 $V[5] \leftarrow \text{true}$
 $P[5] \leftarrow 2$



BFS tree



DFS trees



Not DFS tree,
Not BFS tree.

DFS

- Explore From the latest visited vertex \Rightarrow Stack.

BFS.

Explore from the earliest (closest) visited vertex \Rightarrow queue.