VI Semester MIDTERM TEST PARALLEL COMPUTER ARCHITECTURE AND PROGRAMMING (CSE_3252)

Time Duration: 2 Hours Date: 19.03.2024 Max marks: 30 MARKS

Scheme and Solution

Question	Topic	Marks	BL	СО
No	·			
1	Which type of memory in OpenCL is shared among all work-items?	1	B1	CO3
	A) Local memory			
	B) Global memory			
	C) Constant memory			
	D) Private memory			
2	Which component of OpenCL is responsible for managing device	1	B1	CO3
	resources?			
	A) Compiler			
	B) Runtime			
	C) Scheduler			
	D) Kernel			
3	What does the term "host" refer to in the context of OpenCL?	1	B1	CO3
	A) The main processing unit of a GPU			
	B) The central processing unit (CPU) of a computer			
	C) A specialized type of kernel			
	D) A cloud computing service			
4	In OpenCL, what is a kernel?	1	B1	CO3
	A) A type of memory storage			
	B) An execution unit for parallel computing tasks			
	C) A graphical user interface element			
	D) A data structure for file management			
5	model defines how the concurrency model is mapped to	1	B1	CO3
ı	physical hardware			
	A) Platform model			
	B) Execution model			
	C) Memory model			
	D) Programming model			
6	Which generation of computer systems used Small-scale integrated	1	B1	CO1
	(SSI) and Medium-scale integrated (MSI) circuits as the basic building			
	blocks			
	A) Second generation			
	B) Third generation			
	C) First generation			
	D) Fourth Generation			
7	#include <mpi.h></mpi.h>	1	В3	CO2

```
#include <stdio.h>
       int main(int argc, char *argv[]) {
         int rank, size, value, sum;
         MPI Init(&argc, &argv);
         MPI Comm rank(MPI COMM WORLD, &rank);
         MPI Comm size(MPI COMM WORLD, &size);
         value = rank;
         MPI Reduce(&value,
                                &sum, 1, MPI INT, MPI SUM,
       MPI COMM WORLD);
         if (rank == 0) {
           printf("Sum of ranks: %d\n", sum);
         MPI_Finalize();
         return 0;
      }
      Assuming we run this MPI program with 4 processes, what will be the
      output?
      A) Sum of ranks: 6
      B) Sum of ranks: 10
      C) Sum of ranks: 4
      D) Sum of ranks: 0
                                                                             1
                                                                                   В3
                                                                                       CO<sub>2</sub>
8
       #include <mpi.h>
       #include <stdio.h>
       int main(int argc, char *argv[]) {
         int rank, size, data;
         MPI Init(&argc, &argv);
         MPI Comm rank(MPI COMM WORLD, &rank);
         MPI Comm size(MPI COMM WORLD, &size);
         if (rank == 0) {
           data = 100;
           MPI_Send(&data, 1, MPI_INT, 1, 0, MPI_COMM_WORLD);
         } else if (rank == 1) {
           MPI_Recv(&data, 1, MPI_INT, 0, 0, MPI_COMM_WORLD,
       MPI STATUS IGNORE);
           printf("Process 1 received data: %d\n", data);
         MPI_Finalize();
         return 0;
       }
       What will be the output of the program when executed with 2 MPI
       processes?
       A) Process 1 received data: 0
       B) Process 1 received data: 100
       C) Process 1 received data: 1
       D) Compilation Error
```

		1		1
9	Which of the following is a characteristic of GPU architecture?	1	B1	CO1
	A) High clock speeds and large cache sizes			
	B) Large numbers of cores optimized for parallel processing			
	C) Low power consumption and compact design			
	D) Support for complex branch prediction and out-of-order execution			
10	In which architecture all PEs receive the same instruction broadcast	1	B1	CO1
	from the control unit but operate on different data sets from distinct			
	data streams			
	A) Single instruction stream-single data stream (SISD)			
	B) Single instruction stream-multiple data stream (SIMD)			
	C) Multiple instruction stream-single date stream (MISD)			
	D) Multiple instruction stream-multiple data stream (MIMD)			
11 a.	Using collective communication routines, implement an efficient MPI	4	В3	CO2
	program which reads matrix A of size 4x4 and it produces a resultant	•		-
	matrix <i>RES</i> of size 4x4 such that every row elements of <i>A</i> are added			
	with a key value. The key value for the first row is the minimum			
	· · ·			
	element of the last row and the key value for remaining row is the			
	minimum element from the previous row of matrix A. Make use of 4			
	processes (including root) to perform this task.			
	Example: A RES			
	1234 3456			
	5678 6789			
	2435 79810			
	2 346 4568			
	*use of collective communication-MPI_Bcast() and MPI_Gather()-2 marks *correct logic in code – 2 marks			
	Solution:			
	#include <stdio.h></stdio.h>			
	#include <stdlib.h></stdlib.h>			
	#include <mpi.h></mpi.h>			
	#define ROWS 4			
	#define COLS 4			
	// Function to find the minimum element in a row			
	int min_row(int row[]) {			
	int min = row[0];			
	for (int i = 1; i < COLS; i++) {			
	if (row[i] < min) {			
	min = row[i];			
	}			
	}			
	return min;			
	recurrently	<u> </u>	<u> </u>	

```
}
int main(int argc, char** argv) {
 int rank, size;
  MPI_Init(&argc, &argv);
  MPI Comm rank(MPI COMM WORLD, &rank);
  MPI_Comm_size(MPI_COMM_WORLD, &size);
 // Check if the number of processes is not equal to 4
 if (size != 4) {
    if (rank == 0) {
      printf("Error: Number of processes must be 4\n");
    }
    MPI Finalize();
    return 1;
 }
 int A[ROWS][COLS];
 int key;
 // Root process reads the matrix A
  if (rank == 0) {
    printf("Enter elements of matrix A (4x4):\n");
    for (int i = 0; i < ROWS; i++) {
      for (int j = 0; j < COLS; j++) {
        scanf("%d", &A[i][j]);
      }
    }
 }
 // Broadcast matrix A from root process to all other processes
  MPI Bcast(A, ROWS * COLS, MPI INT, 0, MPI COMM WORLD);
 // Calculate key value based on rank
  if (rank == 0) {
    key = min row(A[ROWS - 1]); // Key for the first row
 } else {
    key = min row(A[rank - 1]); // Key for the remaining rows
  }
 // Broadcast key value from root process to all other processes
  MPI Bcast(&key, 1, MPI INT, 0, MPI COMM WORLD);
 // Add key value to each element of the row
 for (int i = 0; i < COLS; i++) {
    A[rank][i] += key;
```

```
}
          // Gather the resultant matrix RES to root process
          int RES[ROWS][COLS];
          MPI_Gather(A[rank], COLS, MPI_INT, RES, COLS, MPI_INT, 0,
        MPI COMM WORLD);
          // Print resultant matrix RES in root process
          if (rank == 0) {
             printf("Resultant Matrix (RES):\n");
            for (int i = 0; i < ROWS; i++) {
               for (int j = 0; j < COLS; j++) {
                 printf("%d ", RES[i][j]);
               printf("\n");
            }
          }
          MPI Finalize();
          return 0;
        }
11 b.
                                                                                             CO2
                                                                                        В4
         Differentiate between following collective communication routines
         with the help of example code
         i) MPI Allgather() and MPI Alltoall() -2 marks
         ii) MPI Reduce() and MPI Scan() -2 marks
         *Function prototype and sample code
        MPI Allgather() gathers data from all processes and scatters it to all processes,
        while MPI Alltoall() sends distinct data from each process to all other
        processes.
        MPI_Reduce() performs a reduction operation on data from all processes and
        sends the result to one process,
        while MPI_Scan() performs a parallel prefix operation and sends the partial
        result to each process.
        MPI Allgather(&send data,
                                    1, MPI INT,
                                                    recv_data,
                                                                     MPI INT,
        MPI COMM WORLD);
                                                   recv_data,
        MPI Alltoall(send data,
                                 1,
                                       MPI_INT,
                                                                     MPI INT,
        MPI COMM WORLD);
                                  &recv_data, 1, MPI_INT,
        MPI_Reduce(&send_data,
                                                                MPI_SUM,
        MPI_COMM_WORLD);
```

	MPI_Scan(&send_data, &recv_data, 1, MPI_INT, MPI_SUM, MPI_COMM_WORLD);			
11 c.	Demonstrate the process of creating kernel and setting the kernel arguments.	2	В3	CO3
	ClCreatekernel and clSetkernelArg function protype with working -2 marks			
	To create a kernel in OpenCL, you typically define a kernel function in a separate .cl file. Here's an example kernel function that performs element-wise addition of two arrays:			
	<pre>kernel void add_arrays(global const float* a,global const float* b,global float* result, const int size) { int i = get_global_id(0); if (i < size) { result[i] = a[i] + b[i]; } }</pre>			
	In this kernel,kernel is the kernel specifier, add_arrays is the name of the kernel function, and it takes input arrays a and b, as well as an output array result, along with the size of the arrays. Setting Kernel Arguments: After creating the kernel, you need to set kernel arguments before enqueuing the kernel for execution. Here's how you set kernel arguments using the OpenCL API in C: // Assume you have already created an OpenCL context, command queue, and program object			
	// Create kernel object cl_kernel clCreateKernel(
	cl_kernel kernel = clCreateKernel(program, "add_arrays", &err);			
	// Set kernel arguments cl_int clSetKernelArg(
	clSetKernelArg(kernel, 0, sizeof(cl_mem), (void*)&buffer_a); clSetKernelArg(kernel, 1, sizeof(cl_mem), (void*)&buffer_b); clSetKernelArg(kernel, 2, sizeof(cl_mem), (void*)&buffer_result); clSetKernelArg(kernel, 3, sizeof(int), (void*)&size);			

	command queue have been removed from the			
	This function blocks until all of the commands in a			
	o o			
	✓ cl_int clFlush (cl_command_queue cmdQueue);			
	command queue have completely executed.			
	This function blocks until all of the commands in a			
	✓ cl_int clFinish (cl_command_queue cmdQueue);			
	This function blocks until all of the commands in a			
	This function blocks until all of the commands in a			
	command queue have completely executed			
	command gueue have completely executed.			
	command queue have completely executed.			
	✓ cl_int clFlush (cl_command_queue cmdQueue);			
	✓ cl_int_clFlush (cl_command_queue cmdQueue);			
	✓ cl_int_clFlush (cl_command_queue cmdQueue);			
	v ci_iiit ciriusii (ci_coiiiiianu_queue cinaqueue),			
	si_me on tash (si_command_queue on aqueue)			
	This function blocks until all of the commands in a			
	command queue have been removed from the			
	command queue.			
	·			
	This means that the commands will defienetly be in-flight but will not			
	necessarily have completed.			
	necessarily have completed.			
	kernel void reneat string/ global const char* input string global char*			
	kernel void repeat_string(global const char* input_string,global char*			
	output_string, const int n) {			
	int global_id = get_global_id(0);			
	· - · · · · · · · · · · · · · · · · · ·			
	<pre>int input_length = strlen(input_string);</pre>			
	int output_index = global_id * input_length;			
	, = , = , = , .			
	// Repeat the input string N times			
	for (int i = 0; i < n; i++) {			
	for (int j = 0; j < input length; j++) {			
	output_string[output_index++] = input_string[j];			
	and a second to a			
	}			
	;			
	, f			
	}			
	Compare multi-core trajectory and many-core trajectory.	2	B4	CO1
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12 c.	compare mater core trajectory and many core trajectory.	-		

The multi-core trajectory:

- It seeks to maintain the execution speed of sequential programs while moving into multiple cores.
- *Multicore trajectory* began as two-core processors, with the number of cores doubling every generation of MP.
- ➤ Eg: Intel Corei7 MP which has 4 processor cores, each of which is an out-of-order, multiple instruction issue processor implementing the full x86 instruction set.
- This MP supports *hyperthreading* with two *hardware threads* per core and is designed to maximize the exec speed of sequential programs.

The many-core (many-thread) trajectory:

- It focuses more on the *execution throughput* of parallel applications
- The many-cores began as a *large number of much smaller cores*, and, once again, the number of cores doubles with each generation