

Higher-ranked Exception Types

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Abstract

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Categories and Subject Descriptors **TO DO.**CR-number [subcategory]: third-level

General Terms **TO DO.**term1, term2

Keywords **TO DO.**keyword1, keyword2

1. Introduction

An often heard selling point of non-strict functional languages is that they provide strong and expressive type systems that make side-effects explicit. This supposedly makes software more reliable by lessening the mental burden of programmers. Many object-oriented programmers are quite surprised, then, that when they make the transition to a functional language, that they lose a feature their type system formerly did provide: tracking of uncaught exceptions.

There is a good excuse why this feature is missing from the type systems of contemporary non-strict functional languages: in a strict first-order language it is sufficient to annotate each function

with a single set of uncaught exceptions the function may throw, in a non-strict higher-order language the situation becomes significantly more complicated. Let us first consider the two aspects “higher-order” and “non-strict” in isolation:

Higher-order functions The set of exceptions that may be raised by a higher-order function are not given by a fixed set of exceptions, but depends on the set of exceptions that may be raised by the function that is passed as its functional argument.

Higher-order functions will thus end up being *exception polymorphic*.

TO DO.concrete example

Non-strict evaluation In non-strictly evaluated languages, exception are not a form of control flow, but a kind of value. Typically the set of values of each type are extended with an *exceptional value* \perp (more commonly denoted \bot , but we shall not do so for reasons of ambiguity), or family of exceptional values \perp^ℓ .

This means we do not only need to annotate all functions with a set of exceptions, but each expression.

TO DO.concrete example

Take as an example the *map* function:

$$\begin{aligned} \text{map} &:: (a \rightarrow b) \rightarrow [a] \rightarrow [b] \\ \text{map} &= \lambda f \rightarrow \lambda xs \rightarrow \text{case } xs \text{ of} \\ &\quad [] \rightarrow [] \\ &\quad (y : ys) \rightarrow f y : \text{map } f \text{ } ys \end{aligned}$$

1.1 Contributions

- A *type system* that precisely tracks the uncaught exceptions using higher-ranked types.
- An *inference algorithm* that automatically infers such higher-ranked exception types.

TO DO.

- **TO DO.**Untracked exceptions can break information flow security.

2. The λ^U -calculus

Types

$$\begin{array}{lll} \tau \in \mathbf{Ty} & ::= & \mathcal{P} \quad \text{(base type)} \\ & | & \tau_1 \rightarrow \tau_2 \quad \text{(function type)} \end{array}$$

Terms

$t \in \mathbf{Tm}$	$::=$	x, y, \dots	(variable)
		$\lambda x : \tau. t$	(abstraction)
		$t_1 t_2$	(application)
		\emptyset	(empty)
		$\{c\}$	(singleton)
		$t_1 \cup t_2$	(union)

Values Values v are terms of the form

$$\lambda x_1 : \tau_1. \dots \lambda x_i : \tau_i. \{c_1\} \cup (\dots \cup (\{c_j\} \cup (x_1 v_{11} \dots v_{1m} \cup (\dots \cup x_k v_{k1} \dots v_{kn}))))$$

Environments

$$\Gamma \in \mathbf{Env} ::= \cdot \mid \Gamma, x : \tau$$

2.1 Typing relation

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} [\text{T-VAR}] \quad \frac{\Gamma, x : \tau_1 \vdash t : \tau_2}{\Gamma \vdash \lambda x : \tau_1. t : \tau_1 \rightarrow \tau_2} [\text{T-ABS}]$$

$$\frac{}{\Gamma \vdash \emptyset : \mathcal{P}} [\text{T-EMPTY}] \quad \frac{}{\Gamma \vdash \{c\} : \mathcal{P}} [\text{T-CON}] \quad \frac{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : \tau}{\Gamma \vdash t_1 \cup t_2 : \tau} [\text{T-UNION}]$$

2.2 Semantics

2.3 Reduction relation

- **To do.** Do not match the rules in the prototype (those are sensitive to the order in which they are tried).
- **To do.** In the second rule only one term is applied; contrast this with the other rules involving applications.
- **To do.** Should make use of the fact the everything is fully applied (and η -expanded/-long?): all atoms are of the form $k \bar{t}_i$, where k is c or x and the number of arguments fixed by the arity of k . Then try to factor out the commutativity rules by taking “sets” of these atoms. That might simplify stuff a whole lot...
- **To do.** Can we restrict the typing rule T-Union to only allow sets and not functions on both sides? This would remove the 2nd and 3rd rewrite rules and make the system a more traditional higher-order rewrite system: it’s “just” higher-order pattern E-unification (decidable), boolean rings are easy to integrate, and higher-ranked dimension types becomes higher-order E-unification (semi-decidable). Open question: how to represent e.g. $U(e_2(e_1), e_1) = [e_2 \mapsto \lambda e_1. e_1]$ without abstractions? (Reinterpret e_1 as $f(e_1)$ with $f = id$?)

Definition 1. Let \prec be a strict total order on $\mathbf{Con} \cup \mathbf{Var}$, with $c \prec x$ for all $c \in \mathbf{Con}$ and $x \in \mathbf{Var}$.

$$\begin{aligned} & (\lambda x : \tau. t_1) t_2 \longrightarrow t_1[t_2/x] && (\beta\text{-reduction}) \\ & (t_1 \cup t_2) t_3 \longrightarrow t_1 t_3 \cup t_2 t_3 \\ & (\lambda x : \tau. t_1) \cup (\lambda x : \tau. t_2) \longrightarrow \lambda x : \tau. (t_1 \cup t_2) \\ & x t_1 \dots t_n \cup x t'_1 \dots t'_n \longrightarrow x (t_1 \cup t'_1) \dots (t_n \cup t'_n) && (\text{congruences}) \\ & (t_1 \cup t_2) \cup t_3 \longrightarrow t_1 \cup (t_2 \cup t_3) && (\text{associativity}) \\ & \emptyset \cup t \longrightarrow t \\ & t \cup \emptyset \longrightarrow t && (\text{unit}) \\ & x \cup x \longrightarrow x \\ & x \cup (x \cup t) \longrightarrow x \cup t && (\text{idempotence}) \\ & \{c\} \cup \{c\} \longrightarrow \{c\} \\ & \{c\} \cup (\{c\} \cup t) \longrightarrow \{c\} \cup t \\ & x t_1 \dots t_n \cup \{c\} \longrightarrow \{c\} \cup x t_1 \dots t_n && (1) \\ & x t_1 \dots t_n \cup (\{c\} \cup t) \longrightarrow \{c\} \cup (x t_1 \dots t_n \cup t) && (2) \\ & x t_1 \dots t_n \cup x t'_1 \dots t'_n \longrightarrow x t'_1 \dots t'_n \cup x t_1 \dots t_n && \text{if } x' \prec x \quad (3) \\ & x t_1 \dots t_n \cup (x t'_1 \dots t'_n \cup t) \longrightarrow x t'_1 \dots t'_n \cup (x t_1 \dots t_n \cup t) && \text{if } x' \prec x \quad (4) \\ & \{c\} \cup \{c'\} \longrightarrow \{c'\} \cup \{c\} && \text{if } c' \prec c \quad (5) \\ & \{c\} \cup (\{c'\} \cup t) \longrightarrow \{c'\} \cup (\{c\} \cup t) && \text{if } c' \prec c \quad (6) \end{aligned}$$

Conjecture 1. The reduction relation \longrightarrow preserves meaning.

Conjecture 2. The reduction relation \longrightarrow is strongly normalizing.

Conjecture 3. The reduction relation \longrightarrow is locally confluent.

Corollary 1. The reduction relation \longrightarrow is confluent.

Proof. Follows from SN, LC and Newman’s Lemma. \square

Corollary 2. The λ^\cup -calculus has unique normal forms.

Corollary 3. Equality of λ^\cup -terms can be decided by normalization.

3. Completion

$\kappa \in \mathbf{Kind}$	$::=$	EXN	(exception)
		$\kappa_1 \Rightarrow \kappa_2$	(exception operator)
$\varphi \in \mathbf{Exn}$	$::=$	e	(exception variables)
		$\lambda e : \kappa. \varphi$	(exception abstraction)
$\hat{\tau} \in \mathbf{ExnTy}$	$::=$	$\forall e :: \kappa. \hat{\tau}$	(exception quantification)
		$\text{b}\hat{\text{ool}}$	(boolean type)
		$[\hat{\tau}(\varphi)]$	(list type)
		$\hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle$	(function type)

The completion procedure as a set of inference rules:

The completion procedure as an algorithm:

$C :: \mathbf{Env} \times \mathbf{Ty} \rightarrow \mathbf{ExnTy} \times \mathbf{Exn} \times \mathbf{Env}$

$C \frac{e_i :: \kappa_i}{\text{bool}} =$

let e be fresh

in $\langle \text{b}\hat{\text{ool}}; e \bar{e}_i; e :: \kappa_i \Rightarrow \text{EXN} \rangle$

$$\begin{array}{c}
\overline{e_i :: \kappa_i \vdash \mathbf{bool} : \widehat{\mathbf{bool}} \ \& \ e \ \bar{e}_i \triangleright e :: \kappa_i \Rightarrow \text{EXN}} \text{ [C-BOOL]} \\
\\
\overline{e_i :: \kappa_i \vdash \tau : \widehat{\tau} \ \& \ \varphi \triangleright \bar{e}_j :: \kappa_j} \text{ [C-LIST]} \\
\overline{e_i :: \kappa_i \vdash [\tau] : [\widehat{\tau}(\varphi)] \ \& \ e \ \bar{e}_i \triangleright e :: \kappa_i \Rightarrow \text{EXN}, \bar{e}_j :: \kappa_j} \\
\\
\overline{\vdash \tau_1 : \widehat{\tau}_1 \ \& \ \varphi_1 \triangleright \bar{e}_j :: \kappa_j \quad \bar{e}_i :: \kappa_i, \bar{e}_j :: \kappa_j \vdash \tau_2 : \widehat{\tau}_2 \ \& \ \varphi_2 \triangleright \bar{e}_j :: \kappa_j} \text{ [C-ARR]} \\
\overline{e_i :: \kappa_i \vdash \tau_1 \rightarrow \tau_2 : \forall \bar{e}_j :: \kappa_j. (\widehat{\tau}_1(\varphi_1) \rightarrow \widehat{\tau}_2(\varphi_2)) \ \& \ e \ \bar{e}_i \triangleright e :: \kappa_j \Rightarrow \text{EXN}, \bar{e}_j :: \kappa_j} \text{ [T-ABS]}
\end{array}$$

Figure 1. Type completion ($\Gamma \vdash \tau : \widehat{\tau} \ \& \ \varphi \triangleright \Gamma'$)

4. Type system

4.1 Terms

$t \in \mathbf{Tm}$	$::=$	x	(term variable)
		c_τ	(term constant)
		$\lambda x : \tau. t$	(term abstraction)
		$t_1 t_2$	(term application)
		$t_1 \oplus t_2$	(operator)
		if t_1 then t_2 else t_3	(conditional)
		$\not\downarrow_\tau^\ell$	(exception constant)
		$t_1 \mathbf{seq} t_2$	(forcing)
		fix t	(anonymous fixpoint)
		$\llbracket _ \rrbracket_\tau$	(nil constructor)
		$t_1 :: t_2$	(cons constructor)
		case t_1 of $\{\llbracket _ \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\}$	(list eliminator)

4.2 Underlying type system

$$\begin{array}{c}
\overline{\Gamma, x : \tau \vdash x : \tau} \text{ [T-VAR]} \quad \overline{\Gamma \vdash c_\tau : \tau} \text{ [T-CON]} \quad \overline{\Gamma \vdash \not\downarrow_\tau^\ell : \tau} \text{ [T-CRASH]} \\
\\
\overline{\Gamma, x : \tau_1 \vdash t : \tau_2} \text{ [T-ABS]} \quad \overline{\Gamma \vdash t_1 : \tau_2 \rightarrow \tau \quad \Gamma \vdash t_2 : \tau_2} \text{ [T-APP]} \\
\\
\overline{\Gamma \vdash t : \tau \rightarrow \tau} \text{ [T-FIX]} \\
\\
\overline{\Gamma \vdash t_1 : \mathbf{int} \quad \Gamma \vdash t_2 : \mathbf{int}} \text{ [T-OP]} \quad \overline{\Gamma \vdash t_1 : \tau_1 \quad \Gamma \vdash t_2 : \tau_2} \text{ [T-SEQ]} \\
\\
\overline{\Gamma \vdash t_1 : \mathbf{bool} \quad \Gamma \vdash t_2 : \tau \quad \Gamma \vdash t_3 : \tau} \text{ [T-IF]} \\
\\
\overline{\Gamma \vdash \llbracket _ \rrbracket_\tau : [\tau]} \text{ [T-NIL]} \quad \overline{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : [\tau]} \text{ [T-CONS]} \\
\\
\overline{\Gamma \vdash t_1 : [\tau_1] \quad \Gamma \vdash t_2 : \tau \quad \Gamma, x_1 : \tau_1, x_2 : [\tau_1] \vdash t_3 : \tau} \text{ [T-CASE]}
\end{array}$$

4.3 Declarative exception type system

$$\begin{array}{c}
\overline{\Gamma, x : \widehat{\tau} \ \& \ \varphi; \Delta \vdash x : \widehat{\tau} \ \& \ \varphi} \text{ [T-VAR]} \\
\\
\overline{\Gamma; \Delta \vdash c_\tau : \perp_\tau \ \& \ \emptyset} \text{ [T-CON]} \quad \overline{\Gamma; \Delta \vdash \not\downarrow_\tau^\ell : \perp_\tau \ \& \ \{\ell\}} \text{ [T-CRASH]} \\
\\
\overline{\Gamma, x : \widehat{\tau}_1 \ \& \ \varphi_1; \Delta \vdash t : \widehat{\tau}_2 \ \& \ \varphi_2} \text{ [T-ABS]} \\
\\
\overline{\Gamma; \Delta, e : \kappa \vdash t : \widehat{\tau} \ \& \ \varphi \quad e \notin \text{fv}(\varphi)} \text{ [T-ANNAbs]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \widehat{\tau}_1(\varphi_1) \rightarrow \widehat{\tau}(\varphi) \ \& \ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}_2 \ \& \ \varphi_2} \text{ [T-APP]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \forall e : \kappa. \widehat{\tau} \ \& \ \varphi \quad \Delta \vdash \varphi_2 : \kappa} \text{ [T-ANNApP]} \\
\\
\overline{\Gamma; \Delta \vdash t : \widehat{\tau}(\varphi') \rightarrow \widehat{\tau}(\varphi') \ \& \ \varphi'' \quad \Delta \vdash \varphi' \leq \varphi \quad \Delta \vdash \varphi'' \leq \varphi} \text{ [T-FIX]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \mathbf{int} \ \& \ \varphi \quad \Gamma; \Delta \vdash t_2 : \mathbf{int} \ \& \ \varphi} \text{ [T-OP]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \widehat{\tau}_1 \ \& \ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}_2 \ \& \ \varphi} \text{ [T-SEQ]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \mathbf{bool} \ \& \ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau} \ \& \ \varphi \quad \Gamma; \Delta \vdash t_3 : \widehat{\tau} \ \& \ \varphi} \text{ [T-IF]} \\
\\
\overline{\Gamma; \Delta \vdash \llbracket _ \rrbracket_\tau : [\perp_\tau(\emptyset)] \ \& \ \emptyset} \text{ [T-NIL]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : \widehat{\tau} \ \& \ \varphi_1 \quad \Gamma; \Delta \vdash t_2 : [\widehat{\tau}(\varphi_1)] \ \& \ \varphi_2} \text{ [T-CONS]} \\
\\
\overline{\Gamma; \Delta \vdash t_1 : [\widehat{\tau}_1(\varphi_1)] \ \& \ \varphi' \quad \Delta \vdash \varphi' \leq \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau} \ \& \ \varphi} \text{ [T-CASE]} \\
\\
\overline{\Gamma; \Delta \vdash t : \widehat{\tau}' \ \& \ \varphi' \quad \Delta \vdash \widehat{\tau}' \leq \widehat{\tau} \quad \Delta \vdash \varphi' \leq \varphi} \text{ [T-SUB]}
\end{array}$$

In T-Abs and T-AnnAbs, should the term-level term-abstraction also have an explicit effect annotation?

In T-AnnAbs, might need a side condition stating that e is not free in Δ .

– In T-App, note the double occurrence of φ when typing t_1 . Is subeffecting sufficient here? Also note that we do *not* expect an exception variable in the left-hand side annotation of the function space constructor.

– In T-AnnApp, note the substitution. We will need a substitution lemma for annotations.

– In T-Fix, there might be some universal quantifiers in our way. Do annotation applications in t take care of this, already? Perhaps we do need to change **fix** t into a binding construct to resolve this? Also, there is some implicit subeffecting going on between the annotations and effect.

– In T-Case, note the use of explicit subeffecting. Can this be done using implicit subeffecting?

- For T-Sub, should we introduce a term-level coercion, as in Dussart–Henglein–Mossin? We now do shape-conformant subtyping, is subeffecting sufficient?
- Do we need additional kinding judgements in some of the rules? Can we merge the kinding judgement with the subtyping and/or -effecting judgement? Kind-preserving substitutions.

4.4 Type elaboration system

- In T-APP and T-FIX, note that there are substitutions in the premises of the rules. Are these inductive? (Probably, as these premises are not “recursive” ones.)

$$\frac{}{\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash x \hookrightarrow x : \hat{\tau} \& \varphi} \text{[T-VAR]}$$

$$\frac{}{\Gamma; \Delta \vdash c_\tau \hookrightarrow c_\tau : \tau \& \emptyset} \text{[T-CON]} \quad \frac{}{\Gamma; \Delta \vdash \underline{\ell}_\tau \hookrightarrow \underline{\ell}_\tau : \perp_\tau \& \{\ell\}} \text{[T-CRASH]}$$

$$\frac{\Delta, \overline{e_i} : \overline{\kappa_i} \vdash \hat{\tau}_1 \triangleright \tau_1 \quad \Delta, \overline{e_i} : \overline{\kappa_i} \vdash \varphi_1 : \text{EXN} \quad \Gamma, x : \hat{\tau}_1 \& \varphi_1; \Delta, \overline{e_i} : \overline{\kappa_i} \vdash t \hookrightarrow t' : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash \lambda x : \tau_1. t \hookrightarrow \Lambda \overline{e_i} : \overline{\kappa_i}. \lambda x : \hat{\tau}_1 \& \varphi_1. t' : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \emptyset} \text{[T-ABS]}$$

$$\frac{\Delta \vdash \hat{\tau}_2 \leq \hat{\tau}[\overline{\varphi_i}/\overline{e_i}] \quad \Delta \vdash \varphi_2 \leq \varphi[\overline{\varphi_i}/\overline{e_i}] \quad \overline{\Delta \vdash \varphi_i : \kappa_i}}{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi' \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : \hat{\tau}_2 \& \varphi_2} \text{[T-APP]}$$

$$\frac{\Gamma; \Delta \vdash t \hookrightarrow t' : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau} \langle \varphi \rangle \rightarrow \hat{\tau}' \langle \varphi' \rangle \& \varphi'' \quad \Delta \vdash \hat{\tau}'[\overline{\varphi_i}/\overline{e_i}] \leq \hat{\tau}[\overline{\varphi_i}/\overline{e_i}] \quad \Delta \vdash \varphi'[\overline{\varphi_i}/\overline{e_i}] \leq \varphi[\overline{\varphi_i}/\overline{e_i}] \quad \overline{\Delta \vdash \varphi_i : \kappa_i}}{\Gamma; \Delta \vdash \text{fix } t \hookrightarrow \text{fix } t' : \hat{\tau}[\overline{\varphi_i}/\overline{e_i}] \& \varphi[\overline{\varphi_i}/\overline{e_i}] \cup \varphi''} \text{[T-FIX]}$$

$$\frac{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : \mathbf{int} \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : \mathbf{int} \& \varphi_2}{\Gamma; \Delta \vdash t_1 \oplus t_2 \hookrightarrow t'_1 \oplus t'_2 : \mathbf{bool} \& \varphi_1 \cup \varphi_2} \text{[T-OP]}$$

$$\frac{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : \hat{\tau}_1 \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash t_1 \text{ seq } t_2 \hookrightarrow t'_1 \text{ seq } t'_2 : \hat{\tau}_1 \& \varphi_1 \cup \varphi_2} \text{[T-SEQ]}$$

$$\frac{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : \mathbf{bool} \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : \hat{\tau}_2 \& \varphi_2 \quad \Gamma; \Delta \vdash t_3 \hookrightarrow t'_3 : \hat{\tau}_3 \& \varphi_3}{\Gamma; \Delta \vdash \text{if } t_1 \text{ then } t_2 \text{ else } t_3 \hookrightarrow \text{if } t'_1 \text{ then } t'_2 \text{ else } t'_3 : \hat{\tau}_2 \sqcup \hat{\tau}_3 \& \varphi_1 \cup \varphi_2 \cup \varphi_3} \text{[T-IF]}$$

$$\frac{}{\Gamma; \Delta \vdash \square_\tau \hookrightarrow \square_\tau : \perp_\tau \& \emptyset} \text{[T-NIL]}$$

$$\frac{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : \hat{\tau}_1 \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : [\hat{\tau}'_1 \langle \varphi'_1 \rangle] \& \varphi_2}{\Gamma; \Delta \vdash t_1 :: t_2 \hookrightarrow t'_1 :: t'_2 : [\hat{\tau}_1 \sqcup \hat{\tau}'_1 \langle \varphi_1 \cup \varphi'_1 \rangle] \& \varphi_2} \text{[T-CONS]}$$

$$\frac{\Gamma; \Delta \vdash t_1 \hookrightarrow t'_1 : [\tau_1 \langle \varphi_1 \rangle] \& \varphi'_1 \quad \Gamma; \Delta \vdash t_2 \hookrightarrow t'_2 : \hat{\tau}_2 \& \varphi_2 \quad \Gamma, x_1 : \hat{\tau}_1 \& \varphi_1, x_2 : [\tau_1 \langle \varphi_1 \rangle] \& \varphi'_1; \Delta \vdash t_3 \hookrightarrow t'_3 : \hat{\tau}_3 \& \varphi_3}{\Gamma; \Delta \vdash \text{case } t_1 \text{ of } \{\square \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \hookrightarrow \text{case } t'_1 \text{ of } \{\square \mapsto t'_2; x_1 :: x_2 \mapsto t'_3\} : \hat{\tau}_2 \sqcup \hat{\tau}_3 \& \varphi'_1 \cup \varphi_2 \cup \varphi_3} \text{[T-CASE]}$$

- For T-Fix: how would a binding fixpoint construct work?

4.5 Type inference algorithm

$$\mathcal{R} : \mathbf{TyEnv} \times \mathbf{KiEnv} \times \mathbf{Tm} \rightarrow \mathbf{ExnTy} \times \mathbf{Exn}$$

$$\mathcal{R} \Gamma \Delta x = \Gamma_x$$

$$\mathcal{R} \Gamma \Delta c_\tau = \langle \perp_\tau; \emptyset \rangle$$

$$\mathcal{R} \Gamma \Delta \underline{\ell}_\tau = \langle \perp_\tau; \{\ell\} \rangle$$

$$\mathcal{R} \Gamma \Delta (\lambda x : \tau. t) = \text{let } \langle \hat{\tau}_1; e_1; \overline{e_i} : \overline{\kappa_i} \rangle = \mathcal{C} \emptyset \tau$$

$$\text{in } \langle \hat{\tau}_2; \varphi_2 \rangle = \mathcal{R} (\Gamma, x : \hat{\tau}_1 \& e_1) (\Delta, \overline{e_i} : \overline{\kappa_i}) \quad \langle \hat{\tau}_2 \langle \varphi_2 \rangle \rangle = \mathcal{R} (\Gamma, x : \hat{\tau}_1 \& e_1) (\Delta, \overline{e_i} : \overline{\kappa_i}. \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle; \emptyset)$$

$$\mathcal{R} \Gamma \Delta (t_1 t_2) = \text{let } \langle \hat{\tau}_1; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1 \quad \langle \hat{\tau}_2; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2 \quad \langle \hat{\tau}'_2 \langle \varphi'_2 \rangle \rightarrow \hat{\tau}' \langle \varphi' \rangle; \overline{e_i} : \overline{\kappa_i} \rangle = \mathcal{I} \hat{\tau}_1 \quad \theta = [e'_2 \mapsto \varphi_2] \circ \mathcal{M} \emptyset \hat{\tau}_2 \hat{\tau}'$$

$$\text{in } \langle \ll \theta \hat{\tau}' \rrangle_\Delta; \ll \theta \varphi' \cup \varphi_1 \rrangle_\Delta \rangle$$

$$\mathcal{R} \Gamma \Delta (\text{fix } t) = \text{let } \langle \hat{\tau}; \varphi \rangle = \mathcal{R} \Gamma \Delta t$$

$$\langle \hat{\tau}' \langle \varphi' \rangle \rightarrow \hat{\tau}'' \langle \varphi'' \rangle; \overline{e_i} : \overline{\kappa_i} \rangle = \mathcal{I} \hat{\tau}$$

$$\text{in } \langle \hat{\tau}_0; \varphi_0; i \rangle \leftarrow \langle \perp_{[\hat{\tau}']} ; \emptyset; 0 \rangle$$

$$\text{do } \theta \leftarrow [e' \mapsto \varphi_i] \circ \mathcal{M} \emptyset \hat{\tau}_i \hat{\tau}'$$

$$\langle \hat{\tau}_{i+1}; \varphi_{i+1}; i \rangle \leftarrow \langle \ll \theta \hat{\tau}' \rrangle_\Delta; \ll \theta \varphi'' \rrangle_\Delta; i + 1 \rangle$$

$$\text{until } \langle \hat{\tau}_i; \varphi_i \rangle \equiv \langle \hat{\tau}_{i-1}; \varphi_{i-1} \rangle$$

$$\text{return } \langle \hat{\tau}_i; \ll \varphi \cup \varphi_i \rrangle_\Delta \rangle$$

$$\mathcal{R} \Gamma \Delta (t_1 \oplus t_2) = \text{let } \langle \mathbf{int}; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1$$

$$\langle \mathbf{int}; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2$$

$$\text{in } \langle \mathbf{bool}; \ll \varphi_1 \cup \varphi_2 \rrangle_\Delta \rangle$$

$$\mathcal{R} \Gamma \Delta (t_1 \text{ seq } t_2) = \text{let } \langle \hat{\tau}_1; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1$$

$$\langle \hat{\tau}_2; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2$$

$$\text{in } \langle \hat{\tau}_2; \ll \varphi_1 \cup \varphi_2 \rrangle_\Delta \rangle$$

$$\mathcal{R} \Gamma \Delta (\text{if } t_1 \text{ then } t_2 \text{ else } t_3) = \text{let } \langle \mathbf{bool}; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1$$

$$\langle \hat{\tau}_2; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2$$

$$\langle \hat{\tau}_3; \varphi_3 \rangle = \mathcal{R} \Gamma \Delta t_3$$

$$\text{in } \langle \ll \hat{\tau}_2 \sqcup \hat{\tau}_3 \rrangle_\Delta; \ll \varphi_1 \cup \varphi_2 \cup \varphi_3 \rrangle_\Delta \rangle$$

$$\mathcal{R} \Gamma \Delta \square_\tau = \langle \langle \perp_\tau \rangle; \emptyset \rangle$$

$$\mathcal{R} \Gamma \Delta (t_1 :: t_2) = \text{let } \langle \hat{\tau}_1; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1$$

$$\langle [\hat{\tau}_2 \langle \varphi'_2 \rangle]; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2$$

$$\text{in } \langle \ll [(\hat{\tau}_1 \sqcup \hat{\tau}_2) \langle \varphi_1 \cup \varphi'_2 \rangle] \rrangle_\Delta; \varphi_2 \rangle$$

$$\mathcal{R} \Gamma \Delta (\text{case } t_1 \text{ of } \{\square \mapsto t_2; x_1 :: x_2 \mapsto t_3\}) = \text{let } \langle [\hat{\tau}_1 \langle \varphi'_1 \rangle]; \varphi_1 \rangle = \mathcal{R} \Gamma \Delta t_1$$

$$\langle \hat{\tau}_2; \varphi_2 \rangle = \mathcal{R} (\Gamma, x_1 : \hat{\tau}_1 \& \varphi'_1, x_2 : [\hat{\tau}_1 \langle \varphi'_1 \rangle] \& \varphi_1)$$

$$\langle \hat{\tau}_3; \varphi_3 \rangle = \mathcal{R} \Gamma \Delta t_3$$

$$\text{in } \langle \ll \hat{\tau}_2 \sqcup \hat{\tau}_3 \rrangle_\Delta; \ll \varphi_1 \cup \varphi_2 \cup \varphi_3 \rrangle_\Delta \rangle$$

- In R-APP and R-FIX: check that the fresh variables generated by \mathcal{I} are substituted away by the substitution θ created by \mathcal{M} . Also, we don't need those variables in the algorithm if we don't generate the elaborated term.

- In R-FIX we could get rid of the auxillary underlying type function if the fixpoint construct was replaced with a binding variant with an explicit type annotation.

- For R-FIX, make sure the way we handle fixpoints of exceptional value in a manner that is sound w.r.t. to the operational semantics we are going to give to this.

- Note that we do not construct the elaborated term, as it is not useful other than for metatheoretic purposes.

- Lemma: The algorithm maintains the invariant that exception types and exceptions are in normal form.

4.6 Subtyping

- **S-REFL** an admissible/derivable rule, or should we drop **S-BOOL** and **S-INT**?

$$\frac{}{\Delta \vdash \widehat{\tau} \leq \widehat{\tau}} [\text{S-REFL}] \quad \frac{\Delta \vdash \widehat{\tau}_1 \leq \widehat{\tau}_2 \quad \Delta \vdash \widehat{\tau}_2 \leq \widehat{\tau}_3}{\Delta \vdash \widehat{\tau}_1 \leq \widehat{\tau}_3} [\text{S-TRANS}]$$

$$\frac{}{\Delta \vdash \widehat{\text{bool}} \leq \widehat{\text{bool}}} [\text{S-BOOL}] \quad \frac{}{\Delta \vdash \widehat{\text{int}} \leq \widehat{\text{int}}} [\text{S-INT}]$$

$$\frac{\Delta \vdash \widehat{\tau}_1 \leq \widehat{\tau}_1 \quad \Delta \vdash \varphi'_1 \leq \varphi_1 \quad \Delta \vdash \widehat{\tau}_2 \leq \widehat{\tau}'_2 \quad \Delta \vdash \varphi_2 \leq \varphi'_2}{\Delta \vdash \widehat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \widehat{\tau}'_2 \langle \varphi'_1 \rangle \leq \widehat{\tau}'_1 \langle \varphi'_1 \rangle \rightarrow \widehat{\tau}'_2 \langle \varphi'_2 \rangle} [\text{S-ARR}]$$

$$\frac{\Delta \vdash \widehat{\tau} \leq \widehat{\tau}' \quad \Delta \vdash \varphi \leq \varphi'}{\Delta \vdash [\widehat{\tau} \langle \varphi \rangle] \leq [\widehat{\tau}' \langle \varphi' \rangle]} [\text{S-LIST}] \quad \frac{\Delta, e : \kappa \vdash \widehat{\tau}_1 \leq \widehat{\tau}_2}{\Delta \vdash \forall e : \kappa. \widehat{\tau}_1 \leq \forall e : \kappa. \widehat{\tau}_2} [\text{S-FORALL}]$$

– Possibly useful lemma: $\widehat{\tau}_1 = \widehat{\tau}_2 \iff \widehat{\tau}_1 \leq \widehat{\tau}_2 \wedge \widehat{\tau}_2 \leq \widehat{\tau}_1$.

5. Operational semantics

5.1 Evaluation

- The reduction relation is non-deterministic.
- We do not have a Haskell-style imprecise exception semantics (e.g. E-IF).
- We either need to omit the type annotations on \downarrow^ℓ_τ , or add them to **if then else** and **case of** $\{\llbracket \cdot \rrbracket; :: \mapsto\}$.
- We do not have a rule E-ANNAPPEXN. Check that the canonical forms lemma gives us that terms of universally quantified type cannot be exceptional values.

6. Interesting observations

- Exception types are not invariant under η -reduction.

7. Metatheory

7.1 Declarative type system

Lemma 1 (Canonical forms).

1. If \widehat{v} is a possibly exceptional value of type $\widehat{\text{bool}}$, then \widehat{v} is either **true**, **false**, or \downarrow^ℓ .
2. If \widehat{v} is a possibly exceptional value of type $\widehat{\text{int}}$, then \widehat{v} is either some integer n , or an exceptional value \downarrow^ℓ .
3. If \widehat{v} is a possibly exceptional value of type $[\widehat{\tau} \langle \varphi \rangle]$, then \widehat{v} is either $\llbracket \cdot \rrbracket$, $t :: t'$, or \downarrow^ℓ .
4. If \widehat{v} is a possibly exceptional value of type $\widehat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \widehat{\tau}_2 \langle \varphi_2 \rangle$, then \widehat{v} is either $\lambda x : \widehat{\tau}_1 \& \varphi_1. t'$ or \downarrow^ℓ .
5. If \widehat{v} is a possibly exceptional value of type $\forall e : \kappa. \widehat{\tau}$, then \widehat{v} is $\lambda e : \kappa. t$.

Proof. For each part, inspect all forms of \widehat{v} and discard the unwanted cases by inversion of the typing relation. Note that \perp_τ cannot give us a type of the form $\forall e : \kappa. \widehat{\tau}$. \square

TO DO.: Say something about T-SUB?

Theorem 1 (Progress). If $\Gamma; \Delta \vdash t : \widehat{\tau} \& \varphi$ with t a closed term, then t is either a possibly exceptional value \widehat{v} or there is a closed term t' such that $t \rightarrow t'$.

Proof. By induction on the typing derivation $\Gamma; \Delta \vdash t : \widehat{\tau} \& \varphi$.

The case T-VAR can be discarded, as a variable is not a closed term. The cases T-CON, T-CRASH, T-ABS, T-ANNABS, T-NIL and T-CONS are immediate as they are values.

Case T-APP: We can immediately apply the induction hypothesis to $\Gamma; \Delta \vdash t_1 : \widehat{\tau}_2 \langle \varphi_2 \rangle \rightarrow \widehat{\tau} \langle \varphi \rangle \& \varphi$, giving us either a t'_1 such that $t_1 \rightarrow t'_1$ or that $t_1 = \widehat{v}$. In the former case we can make progress using E-APP. In the latter case the canonical forms lemma

$$\frac{t_1 \rightarrow t'_1}{t_1 t_2 \rightarrow t'_1 t_2} [\text{E-APP}] \quad \frac{}{(\lambda x : \widehat{\tau} \& \varphi. t) t_2 \rightarrow t_1[t_2/x]} [\text{E-APPABS}]$$

$$\frac{t \rightarrow t'}{t \langle \varphi \rangle \rightarrow t' \langle \varphi \rangle} [\text{E-ANNAPP}] \quad \frac{}{(\lambda e : \kappa. t) \langle \varphi \rangle \rightarrow t[\varphi/e]} [\text{E-ANNABSABS}]$$

$$\frac{t \rightarrow t'}{\text{fix } t \rightarrow \text{fix } t'} [\text{E-FIX}] \quad \frac{}{\text{fix } (\lambda x : \widehat{\tau} \& \varphi. t) \rightarrow t[\text{fix } (\lambda x : \widehat{\tau} \& \varphi. t)/x]} [\text{E-FIXABS}]$$

$$\frac{}{t \downarrow^\ell \rightarrow t \downarrow^\ell} [\text{E-APPEN}] \quad \frac{}{\text{fix } \downarrow^\ell \rightarrow \downarrow^\ell} [\text{E-FIXEXN}]$$

$$\frac{t_1 \rightarrow t'_1}{t_1 \oplus t_2 \rightarrow t'_1 \oplus t_2} [\text{E-OP1}] \quad \frac{t_2 \rightarrow t'_2}{t_1 \oplus t_2 \rightarrow t_1 \oplus t'_2} [\text{E-OP2}]$$

$$\frac{}{v_1 \oplus v_2 \rightarrow \llbracket v_1 \oplus v_2 \rrbracket} [\text{E-OP}]$$

$$\frac{}{\downarrow^\ell \oplus t_2 \rightarrow \downarrow^\ell} [\text{E-OPEXN1}] \quad \frac{}{t_1 \oplus \downarrow^\ell \rightarrow \downarrow^\ell} [\text{E-OPEXN2}]$$

$$\frac{t_1 \rightarrow t'_1}{t_1 \text{ seq } t_2 \rightarrow t'_1 \text{ seq } t_2} [\text{E-SEQ1}] \quad \frac{}{v_1 \text{ seq } t_2 \rightarrow t_2} [\text{E-SEQ2}]$$

$$\frac{}{\downarrow^\ell \text{ seq } t_2 \rightarrow \downarrow^\ell} [\text{E-SEQEXN}]$$

$$\frac{t_1 \rightarrow t'_1}{\text{if } t_1 \text{ then } t_2 \text{ else } t_3 \rightarrow \text{if } t'_1 \text{ then } t_2 \text{ else } t_3} [\text{E-IF}]$$

$$\frac{}{\text{if true then } t_2 \text{ else } t_3 \rightarrow t_2} [\text{E-IFTRUE}]$$

$$\frac{}{\text{if false then } t_2 \text{ else } t_3 \rightarrow t_3} [\text{E-IFFALSE}]$$

$$\frac{}{\text{if } \downarrow^\ell \text{ then } t_2 \text{ else } t_3 \rightarrow \downarrow^\ell} [\text{E-IFEXN}]$$

$$\frac{t_1 \rightarrow t'_1}{\text{case } t_1 \text{ of } \{\llbracket \cdot \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \rightarrow \text{case } t'_1 \text{ of } \{\llbracket \cdot \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\}} [\text{E-CASE1}]$$

$$\frac{}{\text{case } \llbracket \cdot \rrbracket \text{ of } \{\llbracket \cdot \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \rightarrow t_2} [\text{E-CASENIL}]$$

$$\frac{}{\text{case } t_1 :: t'_1 \text{ of } \{\llbracket \cdot \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \rightarrow t_3[t_1/t'_1/x_1; x_2]} [\text{E-CASENIL}]$$

$$\frac{}{\text{case } \downarrow^\ell \text{ of } \{\llbracket \cdot \rrbracket \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \rightarrow \downarrow^\ell} [\text{E-CASEEXN}]$$

Figure 2. Operational semantics ($t_1 \rightarrow t_2$)

tells us that either $t_1 = \lambda x : \widehat{\tau}_2 \& \varphi_2. t'_1$ or $t_1 = \downarrow^\ell$, in which case we can make progress using E-APPABS or E-APPEN, respectively.

The remaining cases follow by analogous reasoning. \square

Lemma 2 (Annotation substitution).

1. If $\Delta, e : \kappa' \vdash \varphi : \kappa$ and $\Delta \vdash \varphi' : \kappa'$ then $\Delta \vdash \varphi[\varphi'/e] : \kappa$.
2. If $\Delta, e : \kappa' \vdash \varphi_1 \leq \varphi_2$ and $\Delta \vdash \varphi' : \kappa'$ then $\Delta \vdash \varphi_1[\varphi'/e] \leq \varphi_2[\varphi'/e]$.
3. If $\Delta, e : \kappa' \vdash \widehat{\tau}_1 \leq \widehat{\tau}_2$ and $\Delta \vdash \varphi' : \kappa'$ then $\Delta \vdash \widehat{\tau}_1[\varphi'/e] \leq \widehat{\tau}_2[\varphi'/e]$.
4. If $\Gamma; \Delta, e : \kappa' \vdash t : \widehat{\tau} \& \varphi$ and $\Delta \vdash \varphi' : \kappa'$ then $\Gamma; \Delta \vdash t[\varphi'/e] : \widehat{\tau}[\varphi'/e] \& \varphi$.

TO DO.: In part 4, either we need the assumption $e \notin \text{fv}(\varphi)$ (which seems to be satisfied everywhere we want to apply this lemma), or we also need to apply the substitution to φ (is this

$$\begin{array}{ll}
e[\varphi/e] \equiv \varphi & \\
e'[\varphi/e] \equiv e' & \text{if } e \neq e' \\
\{\ell\}[\varphi/e] \equiv \{\ell\} & \\
\emptyset[\varphi/e] \equiv \emptyset & \\
(\lambda e' : \kappa. \varphi') [\varphi/e] \equiv \lambda e' : \kappa. \varphi'[\varphi/e] & \text{if } e \neq e' \text{ and } e' \notin \text{fv}(\varphi) \\
(e_1 e_2) [\varphi/e] \equiv (e_1[\varphi/e]) (e_2[\varphi/e]) & \\
(e_1 \cup e_2) [\varphi/e] \equiv e_1[\varphi/e] \cup e_2[\varphi/e] &
\end{array}$$

Figure 3. Annotation substitution

$$\begin{array}{ll}
x[t/x] \equiv t & \\
x'[t/x] \equiv x' & \text{if } x \neq x' \\
c_\tau[t/x] \equiv c_\tau & \\
(\lambda x' : \hat{\tau}. t') [t/x] \equiv \lambda x' : \hat{\tau}. t'[t/x] & \text{if } x \neq x' \text{ and } x' \notin \text{fv}(t) \\
\dots &
\end{array}$$

Figure 4. Term substitution

expected or not in a type-and-effect system)? T-FIX seems to be to only rule where an exception variable can flow from $\hat{\tau}$ to φ

Proof. 1. By induction on the derivation of $\Delta, e : \kappa' \vdash \varphi : \kappa$. The cases A-VAR, A-ABS and A-APP are analogous to the respective cases in the proof of term substitution below. In the case A-CON one can strengthen the assumption $\Delta, e : \kappa' \vdash \{\ell\} : \text{EXN}$ to $\Delta \vdash \{\ell\} : \text{EXN}$ as $e \notin \text{fv}(\{\ell\})$, the result is then immediate; similarly for A-EMPTY. The case A-UNION goes analogous to A-APP.

2. **TO DO.**

3. **TO DO.**

4. By induction on the derivation of $\Gamma; \Delta, e : \kappa' \vdash t : \hat{\tau} \& \varphi$. Most cases can be discarded by a straightforward application of the induction hypothesis; we show only the interesting case.

Case T-ANNAPP: **TO DO.**

TO DO.

□

Lemma 3 (Term substitution). *If $\Gamma, x : \hat{\tau}' \& \varphi'; \Delta \vdash t : \hat{\tau} \& \varphi$ and $\Gamma; \Delta \vdash t' : \hat{\tau}' \& \varphi'$ then $\Gamma; \Delta \vdash t[t'/x] : \hat{\tau} \& \varphi$.*

Proof. By induction on the derivation of $\Gamma, x : \hat{\tau}' \& \varphi'; \Delta \vdash t : \hat{\tau} \& \varphi$.

Case T-VAR: We either have $t = x$ or $t = x'$ with $x \neq x'$. In the first case we need to show that $\Gamma; \Delta \vdash x[t'/x] : \hat{\tau} \& \varphi$, which by definition of substitution is equal to $\Gamma; \Delta \vdash x : \hat{\tau} \& \varphi$, but this is one of our assumptions. In the second case we need to show that $\Gamma, x' : \hat{\tau} \& \varphi; \Delta \vdash x'[t'/x] : \hat{\tau} \& \varphi$, which by definition of substitution is equal to $\Gamma, x' : \hat{\tau} \& \varphi; \Delta \vdash x' : \hat{\tau} \& \varphi$. This follows immediately from T-VAR.

Case T-ABS: Our assumptions are

$$\Gamma, x : \hat{\tau}' \& \varphi', y : \hat{\tau}_1 \& \varphi_1; \Delta \vdash t : \hat{\tau}_2 \& \varphi_2 \quad (7)$$

$$\Gamma; \Delta \vdash t' : \hat{\tau}' \& \varphi'. \quad (8)$$

By the Barendregt convention we may assume that $y \neq x$ and $y \notin \text{fv}(t')$. We need to show that $\Gamma; \Delta \vdash (\lambda y : \hat{\tau}_1 \& \varphi_1. t)[t'/x] : \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \emptyset$, which by definition of substitution is equal to

$$\Gamma; \Delta \vdash \lambda y : \hat{\tau}_1 \& \varphi_1. t[t'/x] : \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \emptyset. \quad (9)$$

We weaken (8) to $\Gamma, y : \hat{\tau}_1 \& \varphi_1; \Delta \vdash t' : \hat{\tau}' \& \varphi'$ and apply the induction hypothesis on this and (7) to obtain

$$\Gamma, y : \hat{\tau}_1 \& \varphi_1; \Delta \vdash t[t'/x] : \hat{\tau}_2 \& \varphi_2. \quad (10)$$

The desired result (9) can be constructed from (10) using T-ABS.

Case T-ANNABS: Our assumptions are $\Gamma, x : \hat{\tau}' \& \varphi'; \Delta, e : \kappa \vdash t : \hat{\tau} \& \varphi$ and $\Gamma; \Delta \vdash t' : \hat{\tau}' \& \varphi'$. By the Barendregt convention we may assume that $e \notin \text{fv}(t')$. We need to show that $\Gamma; \Delta \vdash (\lambda e : \kappa. t)[t'/x] : \hat{\tau} \& \varphi$, which is equal to $\Gamma; \Delta \vdash \lambda e : \kappa. t[t'/x] : \hat{\tau} \& \varphi$ by definition of substitution. By applying the induction hypothesis we obtain $\Gamma; \Delta, e : \kappa \vdash t[t'/x] : \hat{\tau} \& \varphi$. The desired result can be constructed using T-ANNABS.

Case T-APP: Our assumptions are

$$\Gamma, x : \hat{\tau}' \& \varphi'; \Delta \vdash t_1 : \hat{\tau}_2 \langle \varphi_2 \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi \quad (11)$$

$$\Gamma, x : \hat{\tau}' \& \varphi'; \Delta \vdash t_2 : \hat{\tau}_2 \& \varphi_2. \quad (12)$$

We need to show that $\Gamma; \Delta \vdash (t_1 t_2)[t'/x] : \hat{\tau} \& \varphi$, which by definition of substitution is equal to

$$\Gamma; \Delta \vdash (t_1[t'/x]) (t_2[t'/x]) : \hat{\tau} \& \varphi. \quad (13)$$

By applying the induction hypothesis to (11) respectively (12) we obtain

$$\Gamma; \Delta \vdash t_1[t'/x] : \hat{\tau}_2 \langle \varphi_2 \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi \quad (14)$$

$$\Gamma; \Delta \vdash t_2[t'/x] : \hat{\tau}_2 \& \varphi_2. \quad (15)$$

The desired result (13) can be constructed by applying T-APP to (14) and (15).

All other cases are either immediate or analogous to the case of T-APP. □

Lemma 4 (Inversion).

1. If $\Gamma; \Delta \vdash \lambda x : \hat{\tau} \& \varphi. t : \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \varphi_3$, then
 - $\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash t : \hat{\tau}' \& \varphi'$,
 - $\Delta \vdash \hat{\tau}_1 \leq \hat{\tau}$ and $\Delta \vdash \varphi_1 \leq \varphi$,
 - $\Delta \vdash \hat{\tau}' \leq \hat{\tau}_2$ and $\Delta \vdash \varphi' \leq \varphi_2$.
2. If $\Gamma; \Delta \vdash \lambda e : \kappa. t : \forall e : \kappa. \hat{\tau} \& \varphi$, then
 - $\Gamma; \Delta, e : \kappa \vdash t : \hat{\tau}' \& \varphi'$,
 - $\Delta, e : \kappa \vdash \hat{\tau}' \leq \hat{\tau}$,
 - $\Delta \vdash \varphi' \leq \varphi$.
 - **TO DO.** $e \notin \text{fv}(\varphi)$ and/or $e \notin \text{fv}(\varphi')$.

Proof. 1. By induction on the typing derivation.

Case T-ABS: We have $\hat{\tau} = \hat{\tau}_1$, $\varphi = \varphi_1$ and take $\hat{\tau}' = \hat{\tau}_2$, $\varphi' = \varphi_2$, the result then follows immediately from the assumption $\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash t : \hat{\tau}_2 \& \varphi_2$ and reflexivity of the subtyping and subeffecting relations.

Case T-SUB: We are given the additional assumptions

$$\Gamma; \Delta \vdash \lambda x : \hat{\tau} \& \varphi. t : \hat{\tau}_1' \langle \varphi_1' \rangle \rightarrow \hat{\tau}_2' \langle \varphi_2' \rangle \& \varphi_3', \quad (16)$$

$$\Delta \vdash \hat{\tau}_1' \langle \varphi_1' \rangle \rightarrow \hat{\tau}_2' \langle \varphi_2' \rangle \leq \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle, \quad (17)$$

$$\Delta \vdash \varphi_3' \leq \varphi_3. \quad (18)$$

Applying the induction hypothesis to (16) gives us

$$\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash t : \hat{\tau}_2'' \& \varphi_2'', \quad (19)$$

$$\Delta \vdash \hat{\tau}_1' \leq \hat{\tau}, \quad \Delta \vdash \varphi_1' \leq \varphi, \quad (20)$$

$$\Delta \vdash \hat{\tau}_2'' \leq \hat{\tau}_2', \quad \Delta \vdash \varphi_2'' \leq \varphi_2'. \quad (21)$$

Inversion of the subtyping relation on (17) gives us

$$\Delta \vdash \hat{\tau}_1' \leq \hat{\tau}, \quad \Delta \vdash \varphi_1' \leq \varphi, \quad (22)$$

$$\Delta \vdash \hat{\tau}_2'' \leq \hat{\tau}_2', \quad \Delta \vdash \varphi_2'' \leq \varphi_2'. \quad (23)$$

The result follows from (19) and combining (22) with (20) and (21) with (23) using the transitivity of the subtyping and subeffecting relations.

2. By induction on the typing derivation.

Case T-ANNABS: We need to show that $\Gamma; \Delta, e : \kappa \vdash t : \hat{\tau} \& \varphi$, which is one of our assumptions, and that $\Delta, e : \kappa \vdash \hat{\tau} \leq \hat{\tau}$ and $\Delta \vdash \varphi \leq \varphi$; this follows from the reflexivity of the subtyping, respectively subeffecting, relation (noting that $e \notin \text{fv}(\varphi)$).

Case T-SUB: Similar to the case T-SUB in part 1. \square

Theorem 2 (Preservation). *If $\Gamma; \Delta \vdash t : \hat{\tau} \& \varphi$ and $t \longrightarrow t'$, then $\Gamma; \Delta \vdash t' : \hat{\tau} \& \varphi$.*

Proof. By induction on the typing derivation $\Gamma; \Delta \vdash t : \hat{\tau} \& \varphi$.

The cases for T-VAR, T-CON, T-CRASH, T-ABS, T-ANNABS, T-NIL, and T-CONS can be discarded immediately, as they have no applicable evaluation rules.

To DO. \square

7.2 Syntax-directed type elaboration

7.3 Type inference algorithm

Theorem 3 (Syntactic soundness). *If $\mathcal{R} \Gamma \Delta t = \langle \hat{\tau}; \varphi \rangle$, then $\Gamma; \Delta \vdash t : \hat{\tau} \& \varphi$.*

Proof. By induction on the term t .

To DO. \square

Theorem 4 (Termination). *$\mathcal{R} \Gamma \Delta t$ terminates.*

Proof. By induction on the term t .

To DO. \square