Higher-ranked Exception Types

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1 The $\lambda^{\cup}$-calculus

Types

$$\begin{array}{rcll} \tau \in \mathbf{Ty} & ::= & \mathcal{P} & \text{(base type)} \\ & | & \tau_1 \to \tau_2 & \text{(function type)} \end{array}$$

Terms

$$\begin{array}{rcll} t \in \mathbf{Tm} & ::= & x, y, ... & \text{(variable)} \\ & | & \lambda x : \tau . t & \text{(abstraction)} \\ & | & t_1\ t_2 & \text{(application)} \\ & | & \emptyset & \text{(empty)} \\ & | & \{c\} & \text{(singleton)} \\ & | & t_1 \cup t_2 & \text{(union)} \end{array}$$

Values Values v are terms of the form

$$\lambda x_1 : \tau_1 . \cdots \lambda x_i : \tau_i . \{c_1\} \cup (\cdots \cup (\{c_j\} \cup (x_1\ v_{11} \cdots v_{1m} \cup (\cdots \cup x_k\ v_{k1} \cdots v_{kn}))))$$

Environments

$$\Gamma \in \mathbf{Env} ::= \cdot \quad | \quad \Gamma, x : \tau$$

1.1 Typing relation

$$\frac{}{\Gamma, x : \tau \vdash x : \tau}\ [\text{T-Var}] \quad \frac{\Gamma, x : \tau_1 \vdash t : \tau_2}{\Gamma \vdash \lambda x : \tau_1 . t : \tau_1 \to \tau_2}\ [\text{T-Abs}] \quad \frac{\Gamma \vdash t_1 : \tau_1 \to \tau_2 \quad \Gamma \vdash t_2 : \tau_1}{\Gamma \vdash t_1\ t_2 : \tau_2}\ [\text{T-App}]$$

$$\frac{}{\Gamma \vdash \emptyset : \mathcal{P}}\ [\text{T-Empty}] \quad \frac{}{\Gamma \vdash \{c\} : \mathcal{P}}\ [\text{T-Con}] \quad \frac{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : \tau}{\Gamma \vdash t_1 \cup t_2 : \tau}\ [\text{T-Union}]$$

1.2 Semantics

1.3 Reduction relation

Definition 1. *Let $\prec$ be a strict total order on* $\mathbf{Con} \cup \mathbf{Var}$*, with $c \prec x$ for all $c \in \mathbf{Con}$ and $x \in \mathbf{Var}$.*

$$
\begin{aligned}
(\lambda x : \tau . t_1)\ \ t_2 &\longrightarrow [t_2/x]\, t_1 && (\beta\text{-reduction}) \\
(t_1 \cup t_2)\ \ t_3 &\longrightarrow t_1\ t_3 \cup t_2\ t_3 \\
(\lambda x : \tau . t_1) \cup (\lambda x : \tau . t_2) &\longrightarrow \lambda x : \tau .\, (t_1 \cup t_2) && (\text{congruences}) \\
x\ t_1 \cdots t_n \cup x'\ t'_1 \cdots t'_n &\longrightarrow x\ \ (t_1 \cup t'_1) \cdots (t_n \cup t'_n) \\
(t_1 \cup t_2) \cup t_3 &\longrightarrow t_1 \cup (t_2 \cup t_3) && (\text{associativity}) \\
\varnothing \cup t &\longrightarrow t && \\
t \cup \varnothing &\longrightarrow t && (\text{unit}) \\
x \cup x &\longrightarrow x \\
x \cup (x \cup t) &\longrightarrow x \cup t \\
\{c\} \cup \{c\} &\longrightarrow \{c\} && (\text{idempotence}) \\
\{c\} \cup (\{c\} \cup t) &\longrightarrow \{c\} \cup t \\
x\ t_1 \cdots t_n \cup \{c\} &\longrightarrow \{c\} \cup x\ t_1 \cdots t_n && (1) \\
x\ t_1 \cdots t_n \cup (\{c\} \cup t) &\longrightarrow \{c\} \cup (x\ t_1 \cdots t_n \cup t) && (2) \\
x\ t_1 \cdots t_n \cup x'\ t'_1 \cdots t'_n &\longrightarrow x'\ t'_1 \cdots t'_n \cup x\ t_1 \cdots t_n \quad \text{if } x' \prec x && (3) \\
x\ t_1 \cdots t_n \cup \left(x'\ t'_1 \cdots t'_n \cup t\right) &\longrightarrow x'\ t'_1 \cdots t'_n \cup (x\ t_1 \cdots t_n \cup t) \quad \text{if } x' \prec x && (4) \\
\{c\} \cup \{c'\} &\longrightarrow \{c'\} \cup \{c\} \quad \text{if } c' \prec c && (5) \\
\{c\} \cup \left(\{c'\} \cup t\right) &\longrightarrow \{c'\} \cup (\{c\} \cup t) \quad \text{if } c' \prec c && (6)
\end{aligned}
$$

Conjecture 1. *The reduction relation $\longrightarrow$ preserves meaning.*

Conjecture 2. *The reduction relation $\longrightarrow$ is strongly normalizing.*

Conjecture 3. *The reduction relation $\longrightarrow$ is locally confluent.*

Corollary 1. *The reduction relation $\longrightarrow$ is confluent.*

Proof. Follows from SN, LC and Newman's Lemma. □

Corollary 2. *The $\lambda^{\cup}$-calculus has unique normal forms.*

Corollary 3. *Equality of $\lambda^{\cup}$-terms can be decided by normalization.*

2 Completion

$$
\begin{aligned}
\kappa \in \mathbf{Kind} \quad &::= \quad \textsc{e} && (\text{exception}) \\
&\quad\ \mid \quad \kappa_1 \Rightarrow \kappa_2 && (\text{exception operator})
\end{aligned}
$$

$$\begin{array}{llll} \varphi \in \mathbf{Exn} & ::= & e & \text{(exception variables)} \\ & \mid & \lambda e : \kappa.\varphi & \text{(exception abstraction)} \end{array}$$

$$\begin{array}{llll} \widehat{\tau} \in \mathbf{ExnTy} & ::= & \forall e :: \kappa.\widehat{\tau} & \text{(exception quantification)} \\ & \mid & \widehat{\mathrm{bool}} & \text{(boolean type)} \\ & \mid & [\widehat{\tau}\langle\varphi\rangle] & \text{(list type)} \\ & \mid & \widehat{\tau}_1\langle\varphi_1\rangle \to \widehat{\tau}_2\langle\varphi_2\rangle & \text{(function type)} \end{array}$$

The completion procedure as a set of inference rules:

$$\frac{}{\overline{e_i :: \kappa_i} \vdash \mathbf{bool} : \widehat{\mathrm{bool}} \;\&\; e\;\overline{e_i} \triangleright e :: \overline{\kappa_i \Rightarrow} \mathrm{E}}\;[\text{C-Bool}]$$

$$\frac{\overline{e_i :: \kappa_i} \vdash \tau : \widehat{\tau} \;\&\; \varphi \triangleright \overline{e_j :: \kappa_j}}{\overline{e_i :: \kappa_i} \vdash [\tau] : [\widehat{\tau}\langle\varphi\rangle] \;\&\; e\;\overline{e_i} \triangleright e :: \overline{\kappa_i \Rightarrow} \mathrm{E}, \overline{e_j :: \kappa_j}}\;[\text{C-List}]$$

$$\frac{\vdash \tau_1 : \widehat{\tau}_1 \;\&\; \varphi_1 \triangleright \overline{e_j :: \kappa_j} \qquad \overline{e_i :: \kappa_i}, \overline{e_j :: \kappa_j} \vdash \tau_2 : \widehat{\tau}_2 \;\&\; \varphi_2 \triangleright \overline{e_j :: \kappa_j}}{\overline{e_i :: \kappa_i} \vdash \tau_1 \to \tau_2 : \forall \overline{e_j :: \kappa_j}.\,(\widehat{\tau}_1\langle\varphi_1\rangle \to \widehat{\tau}_2\langle\varphi_2\rangle) \;\&\; e\;\overline{e_i} \triangleright e :: \overline{\kappa_j \Rightarrow} \mathrm{E}, \overline{e_k :: \kappa_k}}\;[\text{C-Arr}]$$

Figure 1: Type completion ($\Gamma \vdash \tau : \widehat{\tau} \;\&\; \varphi \triangleright \Gamma'$)

The completion procedure as an algorithm:

$complete :: \mathbf{Env} \times \mathbf{Ty} \to \mathbf{ExnTy} \times \mathbf{Exn} \times \mathbf{Env}$
$complete\; \overline{e_i :: \kappa_i}\; \mathbf{bool} =$
  let *e be fresh*
  in $\langle \widehat{\mathrm{bool}}; e\;\overline{e_i}; e :: \overline{\kappa_i \Rightarrow}\;\mathrm{E} \rangle$

3 Type system

3.1 Terms

$$
\begin{array}{llll}
t \in \mathbf{Tm} & ::= & x & \text{(term variable)} \\
& | & c_\tau & \text{(term constant)} \\
& | & \lambda x : \tau . t & \text{(term abstraction)} \\
& | & t_1\ t_2 & \text{(term application)} \\
& | & t_1 \oplus t_2 & \text{(operator)} \\
& | & \mathbf{if}\ t_1\ \mathbf{then}\ t_2\ \mathbf{else}\ t_3 & \text{(conditional)} \\
& | & \lightning_\tau^\ell & \text{(exception constant)} \\
& | & t_1\ \mathbf{seq}\ t_2 & \text{(forcing)} \\
& | & \mathbf{fix}\ t & \text{(anonymous fixpoint)} \\
& | & []_\tau & \text{(nil constructor)} \\
& | & t_1 :: t_2 & \text{(cons constructor)} \\
& | & \mathbf{case}\ t_1\ \mathbf{of}\ \{[] \mapsto t_2; x_1 :: x_2 \mapsto t_3\} & \text{(list eliminator)}
\end{array}
$$

3.2 Underlying type system

$$
\frac{}{\Gamma, x : \tau \vdash x : \tau}\ [\text{T-Var}] \qquad \frac{}{\Gamma \vdash c_\tau : \tau}\ [\text{T-Con}] \qquad \frac{}{\Gamma \vdash \lightning_\tau^\ell : \tau}\ [\text{T-Crash}]
$$

$$
\frac{\Gamma, x : \tau_1 \vdash t : \tau_2}{\Gamma \vdash \lambda x : \tau_1 . t : \tau_1 \to \tau_2}\ [\text{T-Abs}] \qquad \frac{\Gamma \vdash t_1 : \tau_2 \to \tau \quad \Gamma \vdash t_2 : \tau_2}{\Gamma \vdash t_1\ t_2 : \tau}\ [\text{T-App}]
$$

$$
\frac{\Gamma \vdash t : \tau \to \tau}{\Gamma \vdash \mathbf{fix}\ t : \tau}\ [\text{T-Fix}]
$$

$$
\frac{\Gamma \vdash t_1 : \mathbf{int} \quad \Gamma \vdash t_2 : \mathbf{int}}{\Gamma \vdash t_1 \oplus t_2 : \mathbf{bool}}\ [\text{T-Op}] \qquad \frac{\Gamma \vdash t_1 : \tau_1 \quad \Gamma \vdash t_2 : \tau_2}{\Gamma \vdash t_1\ \mathbf{seq}\ t_2 : \tau_2}\ [\text{T-Seq}]
$$

$$
\frac{\Gamma \vdash t_1 : \mathbf{bool} \quad \Gamma \vdash t_2 : \tau \quad \Gamma \vdash t_3 : \tau}{\Gamma \vdash \mathbf{if}\ t_1\ \mathbf{then}\ t_2\ \mathbf{else}\ t_3 : \tau}\ [\text{T-If}]
$$

$$
\frac{}{\Gamma \vdash []_\tau : [\tau]}\ [\text{T-Nil}] \qquad \frac{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : [\tau]}{\Gamma \vdash t_1 :: t_2 : [\tau]}\ [\text{T-Cons}]
$$

$$
\frac{\Gamma \vdash t_1 : [\tau_1] \quad \Gamma \vdash t_2 : \tau \quad \Gamma, x_1 : \tau_1, x_2 : [\tau_1] \vdash t_3 : \tau}{\Gamma \vdash \mathbf{case}\ t_1\ \mathbf{of}\ \{[] \mapsto t_2; x_1 :: x_2 \mapsto t_3\} : \tau}\ [\text{T-Case}]
$$

Figure 2: Underlying type system ($\Gamma \vdash t : \tau$)

3.3 Declarative exception type system

- In T-Abs and T-AnnAbs, should the term-level term-abstraction also have an explicit effect annotation?
- In T-AnnAbs, might need a side condition stating that e is not free in Δ.
- In T-App, note the double occurence of φ when typing t_1. Is subeffecting sufficient here? Also note that we do *not* expect an exception variable in the left-hand side annotation of the function space constructor.
- In T-AnnApp, note the substitution. We will need a substitution lemma for annotations.
- In T-Fix, the might be some universal quantifiers in our way. Do annotation applications in t take care of this, already? Perhaps we do need to change **fix** t into a binding construct to resolve this? Also, there is some implicit subeffecting going on between the annotations and effect.
- In T-Case, note the use of explicit subeffecting. Can this be done using implicit subeffecting?
- For T-Sub, should we introduce a term-level coercion, as in Dussart–Henglein–Mossin? We now do shape-conformant subtyping, is subeffecting sufficient?
- Do we need additional kinding judgements in some of the rules? Can we merge the kinding judgement with the subtyping and/or -effecting judgement? Kind-preserving substitutions.

3.4 Type elaboration system

- For T-Fix: how would a binding fixpoint construct work?

3.5 Type inference algorithm

- In R-App and R-Fix: check that the fresh variables generated by $\mathcal{I}$ are subsituted away by the substitution θ created by $\mathcal{M}$.
- In R-Fix we could get rid of the auxillary underlying type function if the fixpoint construct was replaced with a binding variant with an explicit type annotation.
- Note that we do not construct the elaborated term, as it is not useful other than for metatheoretic purposes.
- Simplification does not exactly match the prototype (update the latter).
- Lemma: The algorithm maintains the invariant that exception types and exceptions are in normal form.

$$\frac{}{\Gamma, x : \widehat{\tau}\ \&\ \varphi; \Delta \vdash x : \widehat{\tau}\ \&\ \varphi}\ [\text{T-Var}]$$

$$\frac{}{\Gamma; \Delta \vdash c_\tau : \bot_\tau\ \&\ \emptyset}\ [\text{T-Con}] \qquad \frac{}{\Gamma; \Delta \vdash \lightning_\tau^\ell : \bot_\tau\ \&\ \{\ell\}}\ [\text{T-Crash}]$$

$$\frac{\Gamma, x : \widehat{\tau}_1\ \&\ \varphi_1; \Delta \vdash t : \widehat{\tau}_2\ \&\ \varphi_2}{\Gamma; \Delta \vdash \lambda x : \widehat{\tau}_1\ \&\ \varphi_1.t : \widehat{\tau}_1\langle\varphi_1\rangle \to \widehat{\tau}_2\langle\varphi_2\rangle\ \&\ \emptyset}\ [\text{T-Abs}]$$

$$\frac{\Gamma; \Delta, e : \kappa \vdash t : \widehat{\tau}\ \&\ \varphi}{\Gamma; \Delta \vdash \Lambda e : \kappa.t : \forall e : \kappa.\widehat{\tau}\ \&\ \varphi}\ [\text{T-AnnAbs}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \widehat{\tau}_2\langle\varphi_2\rangle \to \widehat{\tau}\langle\varphi\rangle\ \&\ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}_2\ \&\ \varphi_2}{\Gamma; \Delta \vdash t_1\ t_2 : \widehat{\tau}\ \&\ \varphi}\ [\text{T-App}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \forall e : \kappa.\widehat{\tau}\ \&\ \varphi \quad \Delta \vdash \varphi_2 : \kappa}{\Gamma; \Delta \vdash t_1\ \langle\varphi_2\rangle : [\varphi_2/e]\,\widehat{\tau}\ \&\ \varphi}\ [\text{T-AnnApp}]$$

$$\frac{\Gamma; \Delta \vdash t : \widehat{\tau}\langle\varphi\rangle \to \widehat{\tau}\langle\varphi\rangle\ \&\ \varphi}{\Gamma; \Delta \vdash \mathbf{fix}\ t : \widehat{\tau}\ \&\ \varphi}\ [\text{T-Fix}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \widehat{\mathbf{int}}\ \&\ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\mathbf{int}}\ \&\ \varphi}{\Gamma; \Delta \vdash t_1 \oplus t_2 : \widehat{\mathbf{bool}}\ \&\ \varphi}\ [\text{T-Op}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \widehat{\tau}_1\ \&\ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}_2\ \&\ \varphi}{\Gamma; \Delta \vdash t_1\ \mathbf{seq}\ t_2 : \widehat{\tau}_2\ \&\ \varphi}\ [\text{T-Seq}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \widehat{\mathbf{bool}}\ \&\ \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}\ \&\ \varphi \quad \Gamma; \Delta \vdash t_3 : \widehat{\tau}\ \&\ \varphi}{\Gamma; \Delta \vdash \mathbf{if}\ t_1\ \mathbf{then}\ t_2\ \mathbf{else}\ t_3 : \widehat{\tau}\ \&\ \varphi}\ [\text{T-If}]$$

$$\frac{}{\Gamma; \Delta \vdash []_\tau : [\bot_\tau\langle\emptyset\rangle]\ \&\ \emptyset}\ [\text{T-Nil}]$$

$$\frac{\Gamma; \Delta \vdash t_1 : \widehat{\tau}\ \&\ \varphi_1 \quad \Gamma; \Delta \vdash t_2 : [\widehat{\tau}\langle\varphi_1\rangle]\ \&\ \varphi_2}{\Gamma; \Delta \vdash t_1 :: t_2 : [\widehat{\tau}\langle\varphi_1\rangle]\ \&\ \varphi_2}\ [\text{T-Cons}]$$

$$\frac{\begin{array}{c}\Gamma; \Delta \vdash t_1 : [\widehat{\tau}_1\langle\varphi_1\rangle]\ \&\ \varphi' \quad \Delta \vdash \varphi' \leqslant \varphi \quad \Gamma; \Delta \vdash t_2 : \widehat{\tau}\ \&\ \varphi \\ \Gamma, x_1 : \widehat{\tau}_1\ \&\ \varphi_1, x_2 : [\widehat{\tau}_1\langle\varphi_1\rangle]\ \&\ \varphi'; \Delta \vdash t_3 : \widehat{\tau}\ \&\ \varphi\end{array}}{\Gamma; \Delta \vdash \mathbf{case}\ t_1\ \mathbf{of}\ \{[] \mapsto t_2; x_1 :: x_2 \mapsto t_3\} : \widehat{\tau}\ \&\ \varphi}\ [\text{T-Case}]$$

$$\frac{\Gamma; \Delta \vdash t : \widehat{\tau}\ \&\ \varphi \quad \Delta \vdash \widehat{\tau} \leqslant \widehat{\tau}' \quad \Delta \vdash \varphi \leqslant \varphi'}{\Gamma; \Delta \vdash t : \widehat{\tau}'\ \&\ \varphi'}\ [\text{T-Sub}]$$

Figure 3: Declarative type system ($\Gamma; \Delta \vdash t : \widehat{\tau}\ \&\ \varphi$)

$$\frac{}{\Gamma, x : \widehat{\tau} \mathbin{\&} \varphi; \Delta \vdash x \rightsquigarrow x : \widehat{\tau} \mathbin{\&} \varphi}\;[\textsc{T-Var}]$$

$$\frac{}{\Gamma; \Delta \vdash c_\tau \rightsquigarrow c_\tau : \tau \mathbin{\&} \emptyset}\;[\textsc{T-Con}] \qquad \frac{}{\Gamma; \Delta \vdash \lightning^{\ell}_{\tau} \rightsquigarrow \lightning^{\ell}_{\tau} : \bot_\tau \mathbin{\&} \{\ell\}}\;[\textsc{T-Crash}]$$

$$\frac{\begin{array}{c} \Delta, \overline{e_i : \kappa_i} \vdash \widehat{\tau}_1 \triangleright \tau_1 \quad \Delta, \overline{e_i : \kappa_i} \vdash \varphi_1 : \textsc{e} \\ \Gamma, x : \widehat{\tau}_1 \mathbin{\&} \varphi_1; \Delta, \overline{e_i : \kappa_i} \vdash t \rightsquigarrow t' : \widehat{\tau}_2 \mathbin{\&} \varphi_2 \end{array}}{\Gamma; \Delta \vdash \lambda x : \tau_1 . t \rightsquigarrow \lambda x : \widehat{\tau}_1 \mathbin{\&} \varphi_1 . t' : \forall \overline{e_i : \kappa_i} . \widehat{\tau}_1 \langle \varphi_1 \rangle \to \widehat{\tau}_2 \langle \varphi_2 \rangle \mathbin{\&} \emptyset}\;[\textsc{T-Abs}]$$

$$\frac{\begin{array}{c} \Delta \vdash \widehat{\tau}_2 \leqslant [\overline{\varphi_i / e_i}]\, \widehat{\tau} \quad \Delta \vdash \varphi_2 \leqslant [\overline{\varphi_i / e_i}]\, \varphi \quad \overline{\Delta \vdash \varphi_i : \kappa_i} \\ \Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : \forall \overline{e_i : \kappa_i} . \widehat{\tau}_1 \langle \varphi_1 \rangle \to \widehat{\tau} \langle \varphi \rangle \mathbin{\&} \varphi' \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \widehat{\tau}_2 \mathbin{\&} \varphi_2 \end{array}}{\Gamma; \Delta \vdash t_1\ t_2 \rightsquigarrow t_1'\ \langle \overline{\varphi_i} \rangle\ t_2' : [\overline{\varphi_i / e_i}]\, \widehat{\tau} \mathbin{\&} [\overline{\varphi_i / e_i}]\, \varphi \cup \varphi'}\;[\textsc{T-App}]$$

$$\frac{\begin{array}{c} \Gamma; \Delta \vdash t \rightsquigarrow t' : \forall \overline{e_i : \kappa_i} . \widehat{\tau} \langle \varphi \rangle \to \widehat{\tau}' \langle \varphi' \rangle \mathbin{\&} \varphi'' \\ \Delta \vdash [\overline{\varphi_i / e_i}]\, \widehat{\tau}' \leqslant [\overline{\varphi_i / e_i}]\, \widehat{\tau} \quad \Delta \vdash [\overline{\varphi_i / e_i}]\, \varphi' \leqslant [\overline{\varphi_i / e_i}]\, \varphi \quad \overline{\Delta \vdash \varphi_i : \kappa_i} \end{array}}{\Gamma; \Delta \vdash \mathbf{fix}\ t \rightsquigarrow \mathbf{fix}\ t' : [\overline{\varphi_i / e_i}]\, \widehat{\tau} \mathbin{\&} [\overline{\varphi_i / e_i}]\, \varphi \cup \varphi''}\;[\textsc{T-Fix}]$$

$$\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : \widehat{\mathbf{int}} \mathbin{\&} \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \widehat{\mathbf{int}} \mathbin{\&} \varphi_2}{\Gamma; \Delta \vdash t_1 \oplus t_2 \rightsquigarrow t_1' \oplus t_2' : \widehat{\mathbf{bool}} \mathbin{\&} \varphi_1 \cup \varphi_2}\;[\textsc{T-Op}]$$

$$\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : \widehat{\tau}_1 \mathbin{\&} \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \widehat{\tau}_2 \mathbin{\&} \varphi_2}{\Gamma; \Delta \vdash t_1\ \mathbf{seq}\ t_2 \rightsquigarrow t_1'\ \mathbf{seq}\ t_2' : \widehat{\tau}_2 \mathbin{\&} \varphi_1 \cup \varphi_2}\;[\textsc{T-Seq}]$$

$$\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : \widehat{\mathbf{bool}} \mathbin{\&} \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \widehat{\tau}_2 \mathbin{\&} \varphi_2 \quad \Gamma; \Delta \vdash t_3 \rightsquigarrow t_3' : \widehat{\tau}_3 \mathbin{\&} \varphi_3}{\Gamma; \Delta \vdash \mathbf{if}\ t_1\ \mathbf{then}\ t_2\ \mathbf{else}\ t_3 \rightsquigarrow \mathbf{if}\ t_1'\ \mathbf{then}\ t_2'\ \mathbf{else}\ t_3' : \widehat{\tau}_2 \sqcup \widehat{\tau}_3 \mathbin{\&} \varphi_1 \cup \varphi_2 \cup \varphi_3}\;[\textsc{T-If}]$$

$$\frac{}{\Gamma; \Delta \vdash [\,]_\tau \rightsquigarrow [\,]_\tau : \bot_\tau \mathbin{\&} \emptyset}\;[\textsc{T-Nil}]$$

$$\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : \widehat{\tau}_1 \mathbin{\&} \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \left[\widehat{\tau}_1' \langle \varphi_1' \rangle\right] \mathbin{\&} \varphi_2}{\Gamma; \Delta \vdash t_1 :: t_2 \rightsquigarrow t_1' :: t_2' : \left[\widehat{\tau}_1 \sqcup \widehat{\tau}_1' \langle \varphi_1 \cup \varphi_1' \rangle\right] \mathbin{\&} \varphi_2}\;[\textsc{T-Cons}]$$

$$\frac{\begin{array}{c} \Gamma; \Delta \vdash t_1 \rightsquigarrow t_1' : [\tau_1 \langle \varphi_1 \rangle] \mathbin{\&} \varphi_1' \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t_2' : \widehat{\tau}_2 \mathbin{\&} \varphi_2 \\ \Gamma, x_1 : \widehat{\tau}_1 \mathbin{\&} \varphi_1, x_2 : [\tau_1 \langle \varphi_1 \rangle] \mathbin{\&} \varphi_1'; \Delta \vdash t_3 \rightsquigarrow t_3' : \widehat{\tau}_3 \mathbin{\&} \varphi_3 \end{array}}{\begin{array}{l} \Gamma; \Delta \vdash \mathbf{case}\ t_1\ \mathbf{of}\ \{[\,] \mapsto t_2; x_1 :: x_2 \mapsto t_3\} \rightsquigarrow \\ \qquad\qquad \mathbf{case}\ t_1'\ \mathbf{of}\ \{[\,] \mapsto t_2'; x_1 :: x_2 \mapsto t_3'\} : \widehat{\tau}_2 \sqcup \widehat{\tau}_3 \mathbin{\&} \varphi_1' \cup \varphi_2 \cup \varphi_3 \end{array}}\;[\textsc{T-Case}]$$

Figure 4: Syntax-directed type elaboration system ($\Gamma; \Delta \vdash t \rightsquigarrow t' : \widehat{\tau} \mathbin{\&} \varphi$)

$$
\begin{array}{lll}
\multicolumn{3}{l}{\mathcal{R} : \mathbf{TyEnv} \times \mathbf{KiEnv} \times \mathbf{Tm} \to \mathbf{ExnTy} \times \mathbf{Exn}} \\
\mathcal{R}\ \Gamma\ \Delta\ x & = & \Gamma_x \\
\mathcal{R}\ \Gamma\ \Delta\ c_\tau & = & \langle \bot_\tau ; \emptyset \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ \lightning^{\ell}_{\tau} & = & \langle \bot_\tau ; \{\ell\} \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ (\lambda x : \tau . t) & = & \mathbf{let}\ \langle \widehat{\tau}_1 ; e_1 ; \overline{e_i : \kappa_i} \rangle = \mathcal{C}\ \emptyset\ \tau \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2 ; \varphi_2 \rangle \qquad\quad = \mathcal{R}\ (\Gamma, x : \widehat{\tau}_1\ \&\ e_1)\ (\Delta, \overline{e_i : \kappa_i})\ t \\
& & \mathbf{in}\ \ \langle \forall \overline{e_i : \kappa_i} . \widehat{\tau}_1 \langle e_1 \rangle \to \widehat{\tau}_2 \langle \varphi_2 \rangle ; \emptyset \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ (t_1\ t_2) & = & \mathbf{let}\ \langle \widehat{\tau}_1 ; \varphi_1 \rangle \qquad\qquad\qquad\quad = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2 ; \varphi_2 \rangle \qquad\qquad\qquad\quad = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2' \langle e_2' \rangle \to \widehat{\tau}' \langle \varphi' \rangle ; \overline{e_i : \kappa_i} \rangle = \mathcal{I}\ \widehat{\tau}_1 \\
& & \phantom{\mathbf{let}}\ \theta \qquad\qquad\qquad\qquad\qquad\ \ = [e_2' \mapsto \varphi_2] \circ \mathcal{M}\ \emptyset\ \widehat{\tau}_2\ \widehat{\tau}_2' \\
& & \mathbf{in}\ \ \langle \llbracket \theta \widehat{\tau}' \rrbracket_\Delta ; \llbracket \theta \varphi' \cup \varphi_1 \rrbracket_\Delta \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ (\mathbf{fix}\ t) & = & \mathbf{let}\ \langle \widehat{\tau} ; \varphi \rangle \qquad\qquad\qquad\qquad\quad = \mathcal{R}\ \Gamma\ \Delta\ t \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}' \langle e' \rangle \to \widehat{\tau}'' \langle \varphi'' \rangle ; \overline{e_i : \kappa_i} \rangle = \mathcal{I}\ \widehat{\tau} \\
& & \mathbf{in}\ \ \langle \widehat{\tau}_0 ; \varphi_0 ; i \rangle \leftarrow \langle \bot_{\lfloor \widehat{\tau}' \rfloor} ; \emptyset ; 0 \rangle \\
& & \phantom{\mathbf{in}}\ \ \mathbf{do}\ \theta \qquad\qquad\quad \leftarrow [e' \mapsto \varphi_i] \circ \mathcal{M}\ \emptyset\ \widehat{\tau}_i\ \widehat{\tau}' \\
& & \phantom{\mathbf{in}\ \ \mathbf{do}}\ \langle \widehat{\tau}_{i+1} ; \varphi_{i+1} ; i \rangle \leftarrow \langle \llbracket \theta \widehat{\tau}'' \rrbracket_\Delta ; \llbracket \theta \varphi'' \rrbracket_\Delta ; i + 1 \rangle \\
& & \phantom{\mathbf{in}}\ \ \mathbf{until}\ \langle \widehat{\tau}_i ; \varphi_i \rangle \equiv \langle \widehat{\tau}_{i-1} ; \varphi_{i-1} \rangle \\
& & \phantom{\mathbf{in}}\ \ \mathbf{return}\ \langle \widehat{\tau}_i ; \llbracket \varphi \cup \varphi_i \rrbracket_\Delta \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ (t_1 \oplus t_2) & = & \mathbf{let}\ \langle \widehat{\mathbf{int}} ; \varphi_1 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\mathbf{int}} ; \varphi_2 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \mathbf{in}\ \ \langle \widehat{\mathbf{bool}} ; \llbracket \varphi_1 \cup \varphi_2 \rrbracket_\Delta \rangle \\
\multicolumn{3}{l}{\mathcal{R}\ \Gamma\ \Delta\ (t_1\ \mathbf{seq}\ t_2)} \\
& = & \mathbf{let}\ \langle \widehat{\tau}_1 ; \varphi_1 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2 ; \varphi_2 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \mathbf{in}\ \ \langle \widehat{\tau}_2 ; \llbracket \varphi_1 \cup \varphi_2 \rrbracket_\Delta \rangle \\
\multicolumn{3}{l}{\mathcal{R}\ \Gamma\ \Delta\ (\mathbf{if}\ t_1\ \mathbf{then}\ t_2\ \mathbf{else}\ t_3)} \\
& = & \mathbf{let}\ \langle \widehat{\mathbf{bool}} ; \varphi_1 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2 ; \varphi_2 \rangle \quad\ \ = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_3 ; \varphi_3 \rangle \quad\ \ = \mathcal{R}\ \Gamma\ \Delta\ t_3 \\
& & \mathbf{in}\ \ \langle \llbracket \widehat{\tau}_2 \sqcup \widehat{\tau}_3 \rrbracket_\Delta ; \llbracket \varphi_1 \cup \varphi_2 \cup \varphi_3 \rrbracket_\Delta \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ [\,]_\tau & = & \langle [\bot_\tau \langle \emptyset \rangle] ; \emptyset \rangle \\
\mathcal{R}\ \Gamma\ \Delta\ (t_1 :: t_2) & = & \mathbf{let}\ \langle \widehat{\tau}_1 ; \varphi_1 \rangle \qquad\quad\ = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle [\widehat{\tau}_2 \langle \varphi_2' \rangle] ; \varphi_2 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \mathbf{in}\ \ \langle \llbracket [(\widehat{\tau}_1 \sqcup \widehat{\tau}_2) \langle (\varphi_1 \cup \varphi_2') \rangle] \rrbracket_\Delta ; \varphi_2 \rangle \\
\multicolumn{3}{l}{\mathcal{R}\ \Gamma\ \Delta\ (\mathbf{case}\ t_1\ \mathbf{of}\ \{[\,] \mapsto t_2 ; x_1 :: x_2 \mapsto t_3\})} \\
& = & \mathbf{let}\ \langle [\widehat{\tau}_1 \langle \varphi_1' \rangle] ; \varphi_1 \rangle \qquad\ = \mathcal{R}\ \Gamma\ \Delta\ t_1 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_2 ; \varphi_2 \rangle = \mathcal{R}\ (\Gamma, x_1 : \widehat{\tau}_1\ \&\ \varphi_1', x_2 : [\widehat{\tau}_1 \langle \varphi_1' \rangle]\ \&\ \varphi_1)\ \Delta\ t_2 \\
& & \phantom{\mathbf{let}}\ \langle \widehat{\tau}_3 ; \varphi_3 \rangle = \mathcal{R}\ \Gamma\ \Delta\ t_2 \\
& & \mathbf{in}\ \ \langle \llbracket \widehat{\tau}_2 \sqcup \widehat{\tau}_3 \rrbracket_\Delta ; \llbracket \varphi_1 \cup \varphi_2 \cup \varphi_3 \rrbracket_\Delta \rangle
\end{array}
$$

Figure 5: Type inference algorithm