

Higher-ranked Exception Types

Ruud Koot

February 27, 2015

1 The λ^{\cup} -calculus

Types

$$\begin{array}{lll} \tau \in \mathbf{Ty} & ::= & \mathcal{P} \quad \text{(base type)} \\ & | & \tau_1 \rightarrow \tau_2 \quad \text{(function type)} \end{array}$$

Terms

$$\begin{array}{lll} t \in \mathbf{Tm} & ::= & x, y, \dots \quad \text{(variable)} \\ & | & \lambda x : \tau. t \quad \text{(abstraction)} \\ & | & t_1 t_2 \quad \text{(application)} \\ & | & \emptyset \quad \text{(empty)} \\ & | & \{c\} \quad \text{(singleton)} \\ & | & t_1 \cup t_2 \quad \text{(union)} \end{array}$$

Values Values v are terms of the form

$$\lambda x_1 : \tau_1. \dots \lambda x_i : \tau_i. \{c_1\} \cup (\dots \cup (\{c_j\} \cup (x_1 v_{11} \dots v_{1m} \cup (\dots \cup x_k v_{k1} \dots v_{kn}))))$$

Environments

$$\Gamma \in \mathbf{Env} ::= \cdot \quad | \quad \Gamma, x : \tau$$

1.1 Typing relation

$$\frac{}{\Gamma, x : \tau \vdash x : \tau} [\mathbf{T-VAR}] \quad \frac{\Gamma, x : \tau_1 \vdash t : \tau_2}{\Gamma \vdash \lambda x : \tau_1. t : \tau_1 \rightarrow \tau_2} [\mathbf{T-ABS}] \quad \frac{\Gamma \vdash t_1 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash t_2 : \tau_1}{\Gamma \vdash t_1 t_2 : \tau_2} [\mathbf{T-APP}]$$

$$\frac{}{\Gamma \vdash \emptyset : \mathcal{P}} [\mathbf{T-EMPTY}] \quad \frac{}{\Gamma \vdash \{c\} : \mathcal{P}} [\mathbf{T-CON}] \quad \frac{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : \tau}{\Gamma \vdash t_1 \cup t_2 : \tau} [\mathbf{T-UNION}]$$

1.2 Semantics

1.3 Reduction relation

Definition 1. Let \prec be a strict total order on $\mathbf{Con} \cup \mathbf{Var}$, with $c \prec x$ for all $c \in \mathbf{Con}$ and $x \in \mathbf{Var}$.

$$\begin{aligned}
& (\lambda x : \tau. t_1) \ t_2 \longrightarrow [t_2/x] \ t_1 && (\beta\text{-reduction}) \\
& (t_1 \cup t_2) \ t_3 \longrightarrow t_1 \ t_3 \cup t_2 \ t_3 \\
& (\lambda x : \tau. t_1) \cup (\lambda x : \tau. t_2) \longrightarrow \lambda x : \tau. (t_1 \cup t_2) && (\text{congruences}) \\
& x \ t_1 \cdots t_n \cup x' \ t'_1 \cdots t'_n \longrightarrow x \ (t_1 \cup t'_1) \cdots (t_n \cup t'_n) \\
& (t_1 \cup t_2) \cup t_3 \longrightarrow t_1 \cup (t_2 \cup t_3) && (\text{associativity}) \\
& \emptyset \cup t \longrightarrow t \\
& t \cup \emptyset \longrightarrow t && (\text{unit}) \\
& x \cup x \longrightarrow x \\
& x \cup (x \cup t) \longrightarrow x \cup t && (\text{idempotence}) \\
& \{c\} \cup \{c\} \longrightarrow \{c\} \\
& \{c\} \cup (\{c\} \cup t) \longrightarrow \{c\} \cup t \\
& x \ t_1 \cdots t_n \cup \{c\} \longrightarrow \{c\} \cup x \ t_1 \cdots t_n && (1) \\
& x \ t_1 \cdots t_n \cup (\{c\} \cup t) \longrightarrow \{c\} \cup (x \ t_1 \cdots t_n \cup t) && (2) \\
& x \ t_1 \cdots t_n \cup x' \ t'_1 \cdots t'_n \longrightarrow x' \ t'_1 \cdots t'_n \cup x \ t_1 \cdots t_n && \text{if } x' \prec x \quad (3) \\
& x \ t_1 \cdots t_n \cup (x' \ t'_1 \cdots t'_n \cup t) \longrightarrow x' \ t'_1 \cdots t'_n \cup (x \ t_1 \cdots t_n \cup t) && \text{if } x' \prec x \quad (4) \\
& \{c\} \cup \{c'\} \longrightarrow \{c'\} \cup \{c\} && \text{if } c' \prec c \quad (5) \\
& \{c\} \cup (\{c'\} \cup t) \longrightarrow \{c'\} \cup (\{c\} \cup t) && \text{if } c' \prec c \quad (6)
\end{aligned}$$

Conjecture 1. The reduction relation \longrightarrow preserves meaning.

Conjecture 2. The reduction relation \longrightarrow is strongly normalizing.

Conjecture 3. The reduction relation \longrightarrow is locally confluent.

Corollary 1. The reduction relation \longrightarrow is confluent.

Proof. Follows from SN, LC and Newman's Lemma. \square

Corollary 2. The λ^\cup -calculus has unique normal forms.

Corollary 3. Equality of λ^\cup -terms can be decided by normalization.

2 Completion

$$\begin{aligned}
\kappa \in \mathbf{Kind} & ::= \mathbf{E} && (\text{exception}) \\
& \mid \kappa_1 \Rightarrow \kappa_2 && (\text{exception operator})
\end{aligned}$$

$\varphi \in \mathbf{Exn}$	$::=$	e	(exception variables)
		$\lambda e : \kappa. \varphi$	(exception abstraction)
$\widehat{\tau} \in \mathbf{ExnTy}$	$::=$	$\forall e :: \kappa. \widehat{\tau}$	(exception quantification)
		\mathbf{bool}	(boolean type)
		$[\widehat{\tau}(\varphi)]$	(list type)
		$\widehat{\tau}_1(\varphi_1) \rightarrow \widehat{\tau}_2(\varphi_2)$	(function type)

The completion procedure as a set of inference rules:

$$\begin{array}{c}
\frac{}{\overline{e_i} :: \kappa_i \vdash \mathbf{bool} : \widehat{\mathbf{bool}} \ \& \ e \ \overline{e_i} \triangleright e :: \kappa_i \Rightarrow_{\mathbf{E}}} \text{[C-BOOL]} \\
\\
\frac{\overline{e_i} :: \kappa_i \vdash \tau : \widehat{\tau} \ \& \ \varphi \triangleright \overline{e_j} :: \kappa_j}{\overline{e_i} :: \kappa_i \vdash [\tau] : [\widehat{\tau}(\varphi)] \ \& \ e \ \overline{e_i} \triangleright e :: \kappa_i \Rightarrow_{\mathbf{E}}, \overline{e_j} :: \kappa_j} \text{[C-LIST]} \\
\\
\frac{\vdash \tau_1 : \widehat{\tau}_1 \ \& \ \varphi_1 \triangleright \overline{e_j} :: \kappa_j \quad \overline{e_i} :: \kappa_i, \overline{e_j} :: \kappa_j \vdash \tau_2 : \widehat{\tau}_2 \ \& \ \varphi_2 \triangleright \overline{e_j} :: \kappa_j}{\overline{e_i} :: \kappa_i \vdash \tau_1 \rightarrow \tau_2 : \forall \overline{e_j} :: \kappa_j. (\widehat{\tau}_1(\varphi_1) \rightarrow \widehat{\tau}_2(\varphi_2)) \ \& \ e \ \overline{e_i} \triangleright e :: \kappa_i \Rightarrow_{\mathbf{E}}, \overline{e_k} :: \kappa_k} \text{[C-ARR]}
\end{array}$$

Figure 1: Type completion ($\Gamma \vdash \tau : \widehat{\tau} \ \& \ \varphi \triangleright \Gamma'$)

The completion procedure as an algorithm:

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complete :: Env × Ty → ExnTy × Exn × Env
complete  $\overline{e_i} :: \kappa_i \ \mathbf{bool} =$ 
  let  $e$  be fresh
  in  $\langle \widehat{\mathbf{bool}}; e \ \overline{e_i}; e :: \kappa_i \Rightarrow_{\mathbf{E}} \rangle$ 

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3 Type system

3.1 Terms

$t \in \mathbf{Tm}$	$::=$	x	(term variable)
		c_τ	(term constant)
		$\lambda x : \tau. t$	(term abstraction)
		$t_1 t_2$	(term application)
		$t_1 \oplus t_2$	(operator)
		if t_1 then t_2 else t_3	(conditional)
		\downarrow_τ^ℓ	(exception constant)
		t_1 seq t_2	(forcing)
		fix t	(anonymous fixpoint)
		\square_τ	(nil constructor)
		$t_1 :: t_2$	(cons constructor)
		case t_1 of $\{\square \mapsto t_2; x_1 :: x_2 \mapsto t_3\}$	(list eliminator)

3.2 Underlying type system

$$\begin{array}{c}
\frac{}{\Gamma, x : \tau \vdash x : \tau} [\text{T-VAR}] \quad \frac{}{\Gamma \vdash c_\tau : \tau} [\text{T-CON}] \quad \frac{}{\Gamma \vdash \downarrow_\tau^\ell : \tau} [\text{T-CRASH}] \\
\\
\frac{\Gamma, x : \tau_1 \vdash t : \tau_2}{\Gamma \vdash \lambda x : \tau_1. t : \tau_1 \rightarrow \tau_2} [\text{T-ABS}] \quad \frac{\Gamma \vdash t_1 : \tau_2 \rightarrow \tau \quad \Gamma \vdash t_2 : \tau_2}{\Gamma \vdash t_1 t_2 : \tau} [\text{T-APP}] \\
\\
\frac{\Gamma \vdash t : \tau \rightarrow \tau}{\Gamma \vdash \mathbf{fix} \, t : \tau} [\text{T-FIX}] \\
\\
\frac{\Gamma \vdash t_1 : \mathbf{int} \quad \Gamma \vdash t_2 : \mathbf{int}}{\Gamma \vdash t_1 \oplus t_2 : \mathbf{bool}} [\text{T-OP}] \quad \frac{\Gamma \vdash t_1 : \tau_1 \quad \Gamma \vdash t_2 : \tau_2}{\Gamma \vdash t_1 \mathbf{seq} \, t_2 : \tau_2} [\text{T-SEQ}] \\
\\
\frac{\Gamma \vdash t_1 : \mathbf{bool} \quad \Gamma \vdash t_2 : \tau \quad \Gamma \vdash t_3 : \tau}{\Gamma \vdash \mathbf{if} \, t_1 \mathbf{then} \, t_2 \mathbf{else} \, t_3 : \tau} [\text{T-IF}] \\
\\
\frac{}{\Gamma \vdash \square_\tau : [\tau]} [\text{T-NIL}] \quad \frac{\Gamma \vdash t_1 : \tau \quad \Gamma \vdash t_2 : [\tau]}{\Gamma \vdash t_1 :: t_2 : [\tau]} [\text{T-CONS}] \\
\\
\frac{\Gamma \vdash t_1 : [\tau_1] \quad \Gamma \vdash t_2 : \tau \quad \Gamma, x_1 : \tau_1, x_2 : [\tau_1] \vdash t_3 : \tau}{\Gamma \vdash \mathbf{case} \, t_1 \mathbf{of} \, \{\square \mapsto t_2; x_1 :: x_2 \mapsto t_3\} : \tau} [\text{T-CASE}]
\end{array}$$

Figure 2: Underlying type system ($\Gamma \vdash t : \tau$)

3.3 Declarative exception type system

- In T-Abs and T-AnnAbs, should the term-level term-abstraction also have an explicit effect annotation?
- In T-AnnAbs, might need a side condition stating that e is not free in Δ .
- In T-App, note the double occurrence of φ when typing t_1 . Is subeffecting sufficient here? Also note that we do *not* expect an exception variable in the left-hand side annotation of the function space constructor.
- In T-AnnApp, note the substitution. We will need a substitution lemma for annotations.
- In T-Fix, the might be some universal quantifiers in our way. Do annotation applications in t take care of this, already? Perhaps we do need to change **fix** t into a binding construct to resolve this? Also, there is some implicit subeffecting going on between the annotations and effect.
- In T-Case, note the use of explicit subeffecting. Can this be done using implicit subeffecting?
- For T-Sub, should we introduce a term-level coercion, as in Dussart–Henglein–Mossin? We now do shape-conformant subtyping, is subeffecting sufficient?
- Do we need additional kinding judgements in some of the rules? Can we merge the kinding judgement with the subtyping and/or -effecting judgement? Kind-preserving substitutions.

3.4 Type elaboration system

- For T-Fix: how would a binding fixpoint construct work?

3.5 Type inference algorithm

$$\begin{aligned}
\mathcal{R} : \mathbf{Env} \times \mathbf{KiEnv} \times \mathbf{Tm} &\rightarrow \mathbf{ExnTy} \times \mathbf{Exn} \\
\mathcal{R} \Gamma \Delta x &= \Gamma_x \\
\mathcal{R} \Gamma \Delta c_\tau &= \langle \perp_\tau; \emptyset \rangle \\
\mathcal{R} \Gamma \Delta \not\downarrow_\tau^\ell &= \langle \perp_\tau; \{\ell\} \rangle \\
\mathcal{R} \Gamma \Delta \lambda x : \tau. t = \mathbf{let} \langle \widehat{\tau}_1; e_1; \overline{e_i : \kappa_i} \rangle = \mathcal{C} \oslash \tau & \\
&\quad \langle \widehat{\tau}_2; \varphi_2 \rangle = \mathcal{R} \Gamma, \widehat{\tau}_1 \ \& \ e_1 \ \Delta, \overline{e_i : \kappa_i} \ t & \\
&\quad \mathbf{in} \langle \forall \overline{e_i : \kappa_i}. \widehat{\tau}_1 \langle e_1 \rangle \rightarrow \widehat{\tau}_2 \langle \varphi_2 \rangle; \emptyset \rangle & \\
\mathcal{R} \Gamma \Delta (t_1 \ t_2) = \mathbf{let} \langle \widehat{\tau}_1; \varphi_1 \rangle &= \mathcal{R} \Gamma \Delta t_1 & \\
&\quad \langle \widehat{\tau}_2; \varphi_2 \rangle = \mathcal{R} \Gamma \Delta t_2 & \\
&\quad \langle \widehat{\tau}'_2 \langle e'_2 \rangle \rightarrow \widehat{\tau}' \langle \varphi' \rangle; \overline{e_i : \kappa_i} \rangle = \mathcal{I} \ \widehat{\tau}_1 & \\
&\quad \theta = [e'_2 \mapsto \varphi_2] \circ \mathcal{M} \oslash \widehat{\tau}_2 \ \widehat{\tau}'_2 & \\
&\quad \mathbf{in} \langle [\theta \widehat{\tau}']_\Delta; [\theta \varphi' \cup \varphi_1]_\Delta \rangle &
\end{aligned}$$

- Note that we do not construct the elaborated term, as it is not useful other than for metatheoretic purposes.
- Simplification does not exactly match the prototype (update the latter).
- Lemma: The algorithm maintains the invariant that exception types and exceptions are in normal form.

$$\begin{array}{c}
\overline{\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash x : \hat{\tau} \& \varphi} \text{ [T-VAR]} \\
\\
\overline{\Gamma; \Delta \vdash c_\tau : \perp_\tau \& \emptyset} \text{ [T-CON]} \quad \overline{\Gamma; \Delta \vdash \not\downarrow_\tau^\ell : \perp_\tau \& \{\ell\}} \text{ [T-CRASH]} \\
\\
\frac{\Gamma, x : \hat{\tau}_1 \& \varphi_1; \Delta \vdash t : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash \lambda x : \hat{\tau}_1 \& \varphi_1. t : \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \emptyset} \text{ [T-ABS]} \\
\\
\frac{\Gamma; \Delta, e : \kappa \vdash t : \hat{\tau} \& \varphi}{\Gamma; \Delta \vdash \Lambda e : \kappa. t : \forall e : \kappa. \hat{\tau} \& \varphi} \text{ [T-ANNAbs]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \hat{\tau}_2 \langle \varphi_2 \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi \quad \Gamma; \Delta \vdash t_2 : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash t_1 t_2 : \hat{\tau} \& \varphi} \text{ [T-APP]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \forall e : \kappa. \hat{\tau} \& \varphi \quad \Delta \vdash \varphi_2 : \kappa}{\Gamma; \Delta \vdash t_1 \langle \varphi_2 \rangle : [\varphi_2/e] \hat{\tau} \& \varphi} \text{ [T-ANNApP]} \\
\\
\frac{\Gamma; \Delta \vdash t : \hat{\tau} \langle \varphi \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi}{\Gamma; \Delta \vdash \mathbf{fix} \, t : \hat{\tau} \& \varphi} \text{ [T-FIX]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \mathbf{i\hat{n}t} \& \varphi \quad \Gamma; \Delta \vdash t_2 : \mathbf{i\hat{n}t} \& \varphi}{\Gamma; \Delta \vdash t_1 \oplus t_2 : \mathbf{b\hat{o}ol} \& \varphi} \text{ [T-OP]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \hat{\tau}_1 \& \varphi \quad \Gamma; \Delta \vdash t_2 : \hat{\tau}_2 \& \varphi}{\Gamma; \Delta \vdash t_1 \mathbf{seq} \, t_2 : \hat{\tau}_2 \& \varphi} \text{ [T-SEQ]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \mathbf{b\hat{o}ol} \& \varphi \quad \Gamma; \Delta \vdash t_2 : \hat{\tau} \& \varphi \quad \Gamma; \Delta \vdash t_3 : \hat{\tau} \& \varphi}{\Gamma; \Delta \vdash \mathbf{if} \, t_1 \mathbf{then} \, t_2 \mathbf{else} \, t_3 : \hat{\tau} \& \varphi} \text{ [T-IF]} \\
\\
\overline{\Gamma; \Delta \vdash []_\tau : [\perp_\tau \langle \emptyset \rangle] \& \emptyset} \text{ [T-NIL]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : \hat{\tau} \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 : [\hat{\tau} \langle \varphi_1 \rangle] \& \varphi_2}{\Gamma; \Delta \vdash t_1 :: t_2 : [\hat{\tau} \langle \varphi_1 \rangle] \& \varphi_2} \text{ [T-CONS]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 : [\hat{\tau}_1 \langle \varphi_1 \rangle] \& \varphi' \quad \Delta \vdash \varphi' \leq \varphi \quad \Gamma; \Delta \vdash t_2 : \hat{\tau} \& \varphi \quad \Gamma, x_1 : \hat{\tau}_1 \& \varphi_1, x_2 : [\hat{\tau}_1 \langle \varphi_1 \rangle] \& \varphi'; \Delta \vdash t_3 : \hat{\tau} \& \varphi}{\Gamma; \Delta \vdash \mathbf{case} \, t_1 \mathbf{of} \, \{ [] \mapsto t_2; x_1 :: x_2 \mapsto t_3 \} : \hat{\tau} \& \varphi} \text{ [T-CASE]} \\
\\
\frac{\Gamma; \Delta \vdash t : \hat{\tau} \& \varphi \quad \Delta \vdash \hat{\tau} \leq \hat{\tau}' \quad \Delta \vdash \varphi \leq \varphi'}{\Gamma; \Delta \vdash t : \hat{\tau}' \& \varphi'} \text{ [T-SUB]}
\end{array}$$

Figure 3: Declarative type system ($\Gamma; \Delta \vdash t : \hat{\tau} \& \varphi$)

$$\begin{array}{c}
\overline{\Gamma, x : \hat{\tau} \& \varphi; \Delta \vdash x \rightsquigarrow x : \hat{\tau} \& \varphi} \text{ [T-VAR]} \\
\\
\overline{\Gamma; \Delta \vdash c_\tau \rightsquigarrow c_\tau : \tau \& \emptyset} \text{ [T-CON]} \quad \overline{\Gamma; \Delta \vdash \not\downarrow_\tau^\ell \rightsquigarrow \not\downarrow_\tau^\ell : \perp_\tau \& \{\ell\}} \text{ [T-CRASH]} \\
\\
\frac{\Delta, \overline{e_i} : \overline{\kappa_i} \vdash \hat{\tau}_1 \triangleright \tau_1 \quad \Delta, \overline{e_i} : \overline{\kappa_i} \vdash \varphi_1 : \mathbf{E} \quad \Gamma, x : \hat{\tau}_1 \& \varphi_1; \Delta, \overline{e_i} : \overline{\kappa_i} \vdash t \rightsquigarrow t' : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash \lambda x : \tau_1. t \rightsquigarrow \lambda x : \hat{\tau}_1 \& \varphi_1. t' : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau}_2 \langle \varphi_2 \rangle \& \emptyset} \text{ [T-ABS]} \\
\\
\frac{\Delta \vdash \hat{\tau}_2 \leq [\overline{\varphi_i}/\overline{e_i}] \hat{\tau} \quad \Delta \vdash \varphi_2 \leq [\overline{\varphi_i}/\overline{e_i}] \varphi \quad \overline{\Delta} \vdash \varphi_i : \overline{\kappa_i} \quad \Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau}_1 \langle \varphi_1 \rangle \rightarrow \hat{\tau} \langle \varphi \rangle \& \varphi' \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash t_1 t_2 \rightsquigarrow t'_1 \langle \overline{\varphi_i} \rangle t'_2 : [\overline{\varphi_i}/\overline{e_i}] \hat{\tau} \& [\overline{\varphi_i}/\overline{e_i}] \varphi \cup \varphi'} \text{ [T-APP]} \\
\\
\frac{\Gamma; \Delta \vdash t \rightsquigarrow t' : \forall \overline{e_i} : \overline{\kappa_i}. \hat{\tau} \langle \varphi \rangle \rightarrow \hat{\tau}' \langle \varphi' \rangle \& \varphi'' \quad \Delta \vdash [\overline{\varphi_i}/\overline{e_i}] \hat{\tau}' \leq [\overline{\varphi_i}/\overline{e_i}] \hat{\tau} \quad \Delta \vdash [\overline{\varphi_i}/\overline{e_i}] \varphi' \leq [\overline{\varphi_i}/\overline{e_i}] \varphi \quad \overline{\Delta} \vdash \varphi_i : \overline{\kappa_i}}{\Gamma; \Delta \vdash \mathbf{fix} \, t \rightsquigarrow \mathbf{fix} \, t' : [\overline{\varphi_i}/\overline{e_i}] \hat{\tau} \& [\overline{\varphi_i}/\overline{e_i}] \varphi \cup \varphi''} \text{ [T-FIX]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : \mathbf{\hat{int}} \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : \mathbf{\hat{int}} \& \varphi_2}{\Gamma; \Delta \vdash t_1 \oplus t_2 \rightsquigarrow t'_1 \oplus t'_2 : \mathbf{bool} \& \varphi_1 \cup \varphi_2} \text{ [T-OP]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : \hat{\tau}_1 \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : \hat{\tau}_2 \& \varphi_2}{\Gamma; \Delta \vdash t_1 \mathbf{seq} \, t_2 \rightsquigarrow t'_1 \mathbf{seq} \, t'_2 : \hat{\tau}_2 \& \varphi_1 \cup \varphi_2} \text{ [T-SEQ]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : \mathbf{bool} \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : \hat{\tau}_2 \& \varphi_2 \quad \Gamma; \Delta \vdash t_3 \rightsquigarrow t'_3 : \hat{\tau}_3 \& \varphi_3}{\Gamma; \Delta \vdash \mathbf{if} \, t_1 \mathbf{then} \, t_2 \mathbf{else} \, t_3 \rightsquigarrow \mathbf{if} \, t'_1 \mathbf{then} \, t'_2 \mathbf{else} \, t'_3 : \hat{\tau}_2 \sqcup \hat{\tau}_3 \& \varphi_1 \cup \varphi_2 \cup \varphi_3} \text{ [T-IF]} \\
\\
\overline{\Gamma; \Delta \vdash []_\tau \rightsquigarrow []_\tau : \perp_\tau \& \emptyset} \text{ [T-NIL]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : \hat{\tau}_1 \& \varphi_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : [\hat{\tau}'_1 \langle \varphi'_1 \rangle] \& \varphi_2}{\Gamma; \Delta \vdash t_1 :: t_2 \rightsquigarrow t'_1 :: t'_2 : [\hat{\tau}_1 \sqcup \hat{\tau}'_1 \langle \varphi_1 \cup \varphi'_1 \rangle] \& \varphi_2} \text{ [T-CONS]} \\
\\
\frac{\Gamma; \Delta \vdash t_1 \rightsquigarrow t'_1 : [\tau_1 \langle \varphi_1 \rangle] \& \varphi'_1 \quad \Gamma; \Delta \vdash t_2 \rightsquigarrow t'_2 : \hat{\tau}_2 \& \varphi_2 \quad \Gamma, x_1 : \hat{\tau}_1 \& \varphi_1, x_2 : [\tau_1 \langle \varphi_1 \rangle] \& \varphi'_1; \Delta \vdash t_3 \rightsquigarrow t'_3 : \hat{\tau}_3 \& \varphi_3}{\Gamma; \Delta \vdash \mathbf{case} \, t_1 \mathbf{of} \, \{ [] \mapsto t_2; x_1 :: x_2 \mapsto t_3 \} \rightsquigarrow \mathbf{case} \, t'_1 \mathbf{of} \, \{ [] \mapsto t'_2; x_1 :: x_2 \mapsto t'_3 \} : \hat{\tau}_2 \sqcup \hat{\tau}_3 \& \varphi'_1 \cup \varphi_2 \cup \varphi_3} \text{ [T-CASE]}
\end{array}$$

Figure 4: Syntax-directed type elaboration system ($\Gamma; \Delta \vdash t \rightsquigarrow t' : \hat{\tau} \& \varphi$)