# Type inference for PHP

## The value of annotations in a dynamic language

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# Abstract

Dynamic language are generally hard to statically analyse because of run-time dependencies. Without running the program there are many things unknown. Because dynamic languages are PHP are widely used, the need for decent analysis tool grows. This research examines the value of adding annotations to PHP code to improve the analysability. In the results we see that annotations improve the analysability of software code (this is a guess). Here I should state something about the correctness of the annotations. And end with a general conclusion.

# Preface

In this section I will thank everyone who has helped me. Maybe also introduce some anecdote on how this research came to be.

## Introduction

#### 1.1 PHP

PHP¹is a server-side programming language created by Rasmus Lerdorf in 1995. The original name 'Personal Home Page' changed to 'PHP: Hypertext Preprocessor' in 1998. PHP source files are executed using the PHP Interpreter. The language is dynamically typed, the types of variables are examined during run-time. In statically typed languages all variable types are known at compile time. PHP supports duck-typing, which allows variables to change types during execution.

Evolution The programming language PHP evolved after its creation in 1995. In the year 2000 Object-Oriented (OO) language structures were added to the langue with the release of PHP 4.0. The 5th version of PHP was released in 2004 and provided an improved OO structure. Namespace were added in PHP 5.3 in 2009, to be able to resolve class naming conflicts between library and create better readable class names. Namespaces are comparable to packages in Java. OPcache extension is added was added in PHP 5.5 and speeds up the performance of including files on run-time by storing precompiled script byte-code in shared memory. The most recent stable version is 5.6 and includes more internal performance optimisations and introduces a new debugger.

**Popularity** According to the Tiobe Index<sup>2</sup> of December 2014, PHP is the 6th most popular language of all programming languages. The language has been in the top 10 since it's introduction in the Tiobe Index in 2001. More than 80 precent of the websites have a php backend<sup>3</sup>. The majority of these websites use PHP version 5, rather than version 4 or older versions. It is therefor wise to focus on PHP version from 5 and discard the older unused versions.

**Analysability** Although the popularity for more than a decade, there is still a lack of good PHP code analysis tools. Tools can help to reveal security vulnerabilities or find vulnerabilities or bugs in source code. The tools can also provide code completions or do automatic transformations which can be used to execute refactoring patterns. Source code analysis can be performed statically or dynamically or a combination of the two. More information on the analysability of php can be found in section 2.1.

#### 1.2 Position

In this research we investigate how we can improve the analysability of PHP programs. We will show that the use of annotated source code can help to improve the analysability. The correctness of the annotations can also be examined by checking the implementation of the code.

These annotations can help to improve the analysis. The results can be used to find security issues, and if they are highly reliable we can even make compiler optimisations.

<sup>&</sup>lt;sup>1</sup>http://php.net

<sup>&</sup>lt;sup>2</sup>http://www.tiobe.com/index.php/content/paperinfo/tpci/index.html, December 2014

<sup>&</sup>lt;sup>3</sup>http://w3techs.com/technologies/details/pl-php/all/all, December 2014

As far as we know, there is no constraint based type inference research like this one performed for PHP. That makes this research unique. There have been similar analysis for other dynamic languages, like smalltalk, ruby and javascript, but none like this.

#### 1.3 Contribution

TODO Review this part when the result of the analysis are performed.

- Created an M3 for php.
- Constraint system.
- Show the value of annotations for analysis.

Some idea's are that this analysis can help IDE tools to perform transformations on the source code. (But the performance may not be sufficient.)

The creation of the M3 model can help to compare researchers compare PHP programs with other programming languages. For now only Java is implemented, but more can follow (unchecked statement).

#### 1.4 Plan

The rest of this thesis is as follows: chapter 2 contains background and related work. Background exists of important language constructs, information about annotations in PHP, introduction to  $M^3$  and type systems. The last section of this chapter shows similar research and their relation to this research. Chapter 3, research method, defines the research question and the research context. Chapter 4 describes the performed actions of this research. The analysis is presented in chapter 5, with the results in chapter 6. This thesis ends with the conclusions in chapter 8.

# Background and Related Work

Chapter 2 describes seven important language constructs which the reader needs to understand in order to understand the difficulties of analysing PHP in section 2.1. The second section, section 2.2, introduces the syntax and usage of annotations in PHP. Section 2.3 explains Rascal, the programming language used for the analysis.  $M^3$  is described in more details in section 2.4.  $M^3$  is a programming language independent meta model which holds various facts about programs. Section 2.5 introduces type systems and how type systems relate to this research. The last section of this chapter, section 2.7, describes the related work and how these researches are related to this thesis.

### 2.1 PHP Language Constructs

PHP has various language constructs which complicate statical analysis. This section presents language constructs and why these constructs are important for this research. Explanations of these constructs help you to understand the performed analysis. The discussed parts are scope, file includes, conditional classes and functions, dynamic features, late static binding, magic methods and dynamic class properties.

**Scope** In PHP, all classes and functions are globally accessible once they are declared. All classes and functions are implicitly public, inner classes are not allowed, and conditional functions (see paragraph about conditional classes and functions) will be available in the global scope. If a class or function is declared inside a namespace, their name is prefixed with the name of the namespace.

Variables have three scope levels: global-, function-, and method-scope. Under normal circumstances when a variable is declared inside a function or method, their scope is limited to this function or method. Variables declared outside function or methods are available in the global scope, but not in the method or function scope. There is an exception for some predefined global variables which are available everywhere. Examples are \$GLOBALS, \$\_POST, and \$\_GET. Variables inside a function or method can be aliased to a global variable by adding the keyword GLOBAL in front of the variable name. The variable are then linked to the global variable in the symbol table 1.

Closures (anonymous functions in PHP) have the same scoping rules as variables, but they can inherit variables from outside their scope by providing them in the use statement.

**Includes** PHP allows files to include other PHP-files during execution of the program. The content of these files will be loaded inline. This means that if you use an include in the middle of a file, the source code of this file will be inserted virtually at that place. In this research we will not perform an include analysis. Instead we will assume that all files in the project are included during execution.

According to the php coding standard<sup>2</sup>, function- and class-name classes should not appear when using namespaces and autoloading. When a class which is not already loaded in memory, the autoloading function will try to include a file and load the class. The structure of the autoloading is meant to include classes, interfaces, traits, and functions and should not have inline code executions which would lead to side-effects.

<sup>&</sup>lt;sup>1</sup>http://php.net/manual/en/language.variables.scope.php, July 2014

<sup>&</sup>lt;sup>2</sup>http://www.php-fig.org/psr/psr-0/, July 2014

Conditional classes and functions Once a file is included in the execution, all the classes and functions in the top scope are declared. All class and function declarations within condition statements or within a method or function scope are only declared when the code is executed.

An example of an conditional statement can be found in listing 2.1. If the class Foo or function bar to not exist before the statements is executed, then the class and function will not yet declared. When you try to use the class or function before the code is executed, the script will exit with an fatal error.

```
if (!class_exists("Foo"))
class Foo { /* ... */ }

if (!function_exists("bar"))
function bar() { /* ... */ }
```

Listing 2.1: Conditional class and function definitions

More examples of dynamic function and class loading is displayed in listing 2.2. If the first call is g() as you can see in line 8, the script will result in a fatal error because function g() will only be declared after function f() is executed. The class C will be declared once function g() is executed. As soon as the functions and classes are declared, they are available in the top scope, possibly prefixed with the name of the namespace they are declared within.

```
function f() {
    function g() {
3
       class C {}
4
    }
5
  }
6
7
  // Execution examples:
8
                       // will fail because 'g();' is not declared yet
  g();
9
  f(); g();
                                                    is declared when calling
                       // will work because 'g();'
10
  f(); new C();
                       // will fail because 'g();'
                                                    needs to be called first
  f(); g(); new C();
                       // will work because 'g();' is called and has declared
```

Listing 2.2: Conditional function declaration

**Dynamic features** PHP comes with dynamic features like: include dynamic variables, dynamic class instantiations, dynamic function calls, dynamic function creation, reflection, and eval. These features allow code to change during compile time. New functions and classes can be declared on the fly. Method calls, or even whole pieces of code, can be executed based on variable strings.

A previous study by Mark Hills[HKV13] has shown that most real applications make use of dynamic features. Dynamic features powerful, but can complicate the statical analysis. Additional analysis like constant propagation is needed to resolve most of these dynamic features. This is not in scope for this research.

Late static binding Late static binding<sup>3</sup> is implemented in PHP since version 5.3 by adding the keyword static to the language. It's usage is similar to the keyword self, which refers to the current class. The main difference is that self refers to the class where the code is located, while static refers to the actual instantiated class. The keyword self can be easily resolved while static can only be resolved on runtime. For the keyword self we will have to refer to the class it's declared in, plus all their descended classes.

Magic methods PHP allows calls and property access on methods and fields that don't exist on a class. Normally a call to a non-existing method or property would result in a fatal error, but with the use of magic methods you can specify the wanted behavior. Listing 2.3 shows an example of the \_\_call method. This method is triggered when a non-accessible or non-existing method is called. In this example the code will try to return the value of a private property based on the provided name. The

 $<sup>^3</sup> http://php.net/manual/en/language.oop5.late-static-bindings.php,\ July\ 2014$ 

full list of magic methods is \_\_construct, \_\_destruct, \_\_call, \_\_callStatic, \_\_get, \_\_set, \_\_isset, \_\_unset, \_\_wakeup, \_\_toString, \_\_invoke, \_\_set\_state, \_\_clone, and \_\_debugInfo.

```
class Car {
   private $maxSpeed = 210;
   function __get($name) { return $this->$name; }
}

var_dump((new Car)->maxSpeed); // 210
var_dump((new Car)->numberOfWheels); // NULL
```

Listing 2.3: Magic methods in PHP

Dynamic class properties Although it is a good practice to define your class properties, it is not required to do so in PHP. After instantiating a class it is possible to add properties to classes, even without the implementation of magic methods. In listing 2.4 you can see a code sample of adding a property after instantiation of a class. The access of the non-existing property nonExistingProperty will result in a warning, but code execution will continue and will just return NULL. The code on line 4 is where the property is written. The object \$c will have the nonExistingProperty publicly available now. But in a new class instantiation, like you can see on line 6, will not have the property there.

```
class C {}

c
```

Listing 2.4: Dynamic class property

#### 2.2 Annotations

Annotations are pieces of meta data, defined on class, method, function, or statement level. Despite the proposal for official support of annotations, PHP has still no native support them. But PHP has a getDocComment method in the ReflectionClass since version 5.1 in 2005. The getDocComment method returns the complete doc block of a certain element as a string. A doc block in php has the format /\*\*...\*/. Listing 2.5 shows an example of two doc blocks in PHP. The first doc block is defined above the class and contains information about the class. The second doc block is related to the method getSomething. The block contains a short description of the method, provides type hints for the parameter and the return type, and provides information which possible exceptions can be thrown by the method.

```
namespace Thesis;
2
3
4
   * Class Example
5
   * @package Thesis
6
   */
7
  class Example
8
  {
9
10
        * This is a description of the method getSomething
11
        * @param SomeTypeHint $someObject
12
13
        * @return string
        * @throws NoNameException
14
15
```

<sup>&</sup>lt;sup>4</sup>https://wiki.php.net/rfc/annotations-in-docblock

<sup>&</sup>lt;sup>5</sup>http://php.net/manual/en/reflectionclass.getdoccomment.php

```
public function getSomething(SomeTypeHint $someObject)
{
    if (null === $someObject->getName()) {
        throw new NoNameException();
    }
}
return $someObject->getName();
}
```

Listing 2.5: Examples of PHP DocBlocks

Annotations are mainly used for type hinting, documentation, and code execution. Software analysis tools and IDE's can use the type hints to aid understanding code and in finding bugs and security issues. Available tools can generate documentation based on the doc blocks. Programs like Symfony2, ZEND Framework, and Doctrine ORM use annotations for controller routing, templating information, ORM mappings, filters, and validation configuration.

This research focusses on the first type of annotations which can help developers and IDE's to better understand how code behaves within a program. For example a programmer can see what kind of input and output is expected for a method. Doc blocks with annotations can be placed on top of classes, methods, functions, and variables.

A standard on using annotations is not in the PHP Standard Recommendations (PSR) yet, but there is a proposal<sup>6</sup>. For this research we will only focus on the @param, @return, @var, and @inheritDoc annotations. The annotations @return and @param are only useful for functions, class methods, and closures. Type hints are described with @var and can be used on all structures, but mainly occur on variables and class fields.

There is no official standard for the use of annotations, but most projects follow the phpDocumentor<sup>7</sup> syntax. For this research the following annotations are considered:

$$\operatorname{@return} = \Big\{ \operatorname{@return} \, type, \quad \text{unconditionally read } \operatorname{@return} \, type. \\ \tag{2.1}$$

$$\operatorname{@var} type \$var, \quad \text{if `@var} type \$var' \text{ occurs at least once.}$$

$$\operatorname{@var} \$var type, \quad \text{else if `@var} \$var type' \text{ occurs at least once.}$$

$$\operatorname{@var} type, \quad \text{otherwise try to match `@var} type'.$$

$$(2.3)$$

$$type = \begin{cases} type|type, & \text{if '|' in } type. \\ type, & \text{otherwise} \end{cases}$$
 (2.4)

#### 2.3 Rascal

Rascal [KSV09] is a meta programming language developed by Centrum Wiskunde & Informatica (CWI). Rascal is designed to analyse, transform and visualise source code. The language is build on top of Java and implements various concepts of existing programming languages. In this research, Rascal is the main programming language. The only part that does not use Rascal is the parser of PHP code. Although this could be easily implemented in Rascal, there was an existing library written in PHP available.

 $<sup>^6\</sup>mathrm{https://github.com/php-fig/fig-standards/pull/169/files,\ July\ 2014}$ 

 $<sup>^7 \</sup>mathrm{http://www.phpdoc.org/}$ 

### **2.4** $M^3$

The  $M^3$ -model[Izm+13] is a model which holds various information of source code and is implemented in Rascal. This model is created to gain insights in the quality of open-source projects. For our research we use the  $M^3$ -model to store facts about the program in a structured way, so we can easily use it at a later stage.

The core element of the  $M^3$ -model contains containment, declarations, documentation, modifiers, names, types, uses, messages. The declarations relation contains class, method, variable-information with their logical name and their real location. The type of the relation are locations and represent the logical name of the declaration and will be used in the rest of the  $M^3$ . The containment relation has information on what declarations are contained in each other. For example a package can contain a class; a class can contain fields and methods or an inner class; a method can contain variables. The documentation relation contains all comments from the source code and its source location. The modifiers relation has information on the modifiers of declarations. Modifiers are abstract, final, public, protected, or private. The names relation contains a simplified name of the full declarations. The types relation has information about the type of the source code elements. The uses relation describes what references use an object. For instance when a field of a class is used in some expression, the uses relation links the field in the expression to the declaration of the field in the class. And lastly, messages contains errors, warnings, and info statements.

### 2.5 Type system

Type systems define how a set of rules are applied to types in their context. The system validates the type usage with **type checking**. The process of resolving types is called **type inference**. On one hand the system needs to determine the type of variables, on the other hand it will check the type of the variables.

Type checking Type checking is a mechanism which validates and/or enforces the constraints of a type in their specific context. There is a difference between static type checking and dynamic type checking. Static type checking is a process of checking the types based on the source code. The static type checker will ensure that a program is type safe before executing the program, which means that there will occur no type errors during runtime. Dynamic type checking performs the type checking during runtime. This means that the program needs to run to gain feedback on the usage of types. PHP is a dynamically typed language, which means that there are no types checked before actually running the program. There are languages, like JAVA, which have use a combination of static and dynamic type checking. The advantage of static type checking is that type errors can be caught early in the development process and allows for compiler optimisations. Dynamic type checking systems have more flexibility and allow dynamic loading of new code.

Type inference Type inference is the process of resolving the type of a variable or expression. This process is required to be able perform type checking. Being able to infer the type before running the program enables you to optimise code execution by applying compiler optimisations. These optimisation allow performance improvements or reduce memory usage. In dynamic languages like PHP it can be difficult to resolve the type of a variable or expression without running the program. In statically typed languages, type inference happens at compile time. In this paragraph we explain four types of type inference algorithms.

The Hindley-Milner (HM) type system was found in 1969 by Roger J. Hinldey[Hin69] and almost 10 years later rediscovered by Robin Milner[Mil78]. The first implementation, **Algorithm W**[DM82], was four years later created by pHd student Luis Damas in the language ML. The HM type system deduces the types of the variables based on their usage. Haskell uses a version of the HM type system.

kth-order control-flow analysis!!need citation!! (k-CFA) and cartesian product algorithm (CPA) are two other ways of type inference. In k-CFA.

The analysis we perform is flow insensitive, which means that we do not analyse the flows of the program. We do take the context of the programs into account, making the analysis context sensitive.

In this research we use **constraint based type inference**. Every typeable language construct will have

a set of possible types. At the start of the analysis every variable and expression will have all possible types, we call this Universe. Then we will read all the constraints and will try to solve them. Every time a constraint is solved, the number of possible types for one variable will decrease. This process will be repeated until we have solved all the constraints and can no longer limit the number of types for a variable. This is explained in more details in chapter 4.3.

### 2.6 Constraint solvers

todo: Give answers to the following questions: What is a constraint solver? What types of constraint solvers are there? Which one did I use and why?

#### 2.7 Related work

Describe these:

- 'The HipHop Compiler for PHP' [Zha+12] (not much information available on their type inference, only source code)
- 'Phantm: PHP analyzer for type mismatch' [KSK10a; KSK10b] (investigate this in more details, their focus is on finding type errors)
- PHPLint <sup>8</sup>(uses a different kind of annotations, not the Java like PhpDocs)
- 'Soft typing and analyses on PHP programs' [], code implementations: https://github.com/henkerik/typing and https://github.com/marcelosousa/soft-typing-PHP5 (created for php4, code for php5, should check this out, might be able to compare results with this)
- 'Design and Implementation of an Ahead-of-Time Compiler for PHP'[Big10] (to check in detail)

Also describe their differences with my research.

<sup>&</sup>lt;sup>8</sup>http://www.icosaedro.it/phplint, July 2014

## Research Method

In this research we will try to answer research questions. Section 3.1 describes the research questions. In section 3.2 we will explain more about types, subtypes and their relationship. Section 3.3 explains in which context the research is executed.

### 3.1 Research question

The main research question we want to answer is:

Will the use of annotations<sup>1</sup> help to do better static source code analysis?

Subquestions:

- What is the accuracy (recall and precision) of our analysis?
- Are annotations used in programs reliable?

The subquestions are here to support the main question. In the first question we want to measure the precision and recall of this research. With precision we want to measure if the results we have are reliable. With recall we want to have the number of results we were able to reveal.

In the second subquestion we want to know if the annotations conform to the documentation. We will measure this by doing the analysis without using annotations, and see if the results fit the annotations.

### 3.2 Types

PHP is dynamically typed and allow coercions and duck-typing. Coercion make it possible for variables to hold different types Duck-typing checks the object instead of the class.

```
module lang::php::m3::TypeSymbol
```

<sup>&</sup>lt;sup>1</sup>The annotations are limited to: <code>@param</code>, <code>@return</code>, <code>@var</code>, and <code>@inheritDoc</code>.

#### 3.2.1 PHP types

PHP has a similar class inheritance structure and interface implementation as Java. The main difference is that in PHP all class are public and that inner classes are not allowed in PHP.

The basis types in PHP are integers, floats (similar to doubles and reals), booleans, strings, arrays, resources and null. When variables are initialised without a values, they are null. The recourse type is a special one which is not important for this research.

#### 3.2.2 Subtypes

Explain something about subtypes here. For now, only this figure 3.1



Figure 3.1: Subtype hierarchy

The subtype relation of class inheritance is a reflexive transitive closure relation. A class extension of class A on class C will define class A as a subtype of class C in our analysis, as you can see in figure 3.2. If a class does not extend another class, it will implicitly extend the stdClass class. You can see that this happens with class D in the example. The stdClass is represented as the type object() in our analysis.

The basic PHP types also contain a subtype relation. Integers are subtypes of floats.



Figure 3.2: Relation of subtypes among classes

#### 3.3 Research context

In order to let our research take place, we need to make sure that some environment variables are constant.

**Program Correctness** In order to be able to execute this research we will need to assume that the program is correct and works as intended. We will assume that the system contains no bugs. This is needed to be able to say something about the programs we analyse.

**File includes** In this research we will assume that all file are included during runtime. When a PHP system is constructed of classes with namespaces, the files will be logically loaded using PHP's autoloader. Because most recent systems use namespaces, we will assume that all files are included. For legacy systems, this can influence the results of this research.

**Register globals** Register globals allows variables to be magically be created from GET and POST values. Since it is discouraged to use this setting, we will assume that all software products have this setting disabled.

**PHP Warnings** For this research we will ignore all warnings. Warnings do not alter the behaviour of the program. In a normal production environment, the warnings are not shown.

**Sensitivity** Our analysis is flow-, control-, and context-insensitive. Flow-insensitive means that we do not look at messages between objects, and only look in the body of a method. Control-insensitive means that we ignore all control structures. Examples of control structures are if, else, and switch. Context-insensitive means that we do not look at the order of which code is executed.

## Research

The research is executed in 3 main steps. The first step creates an  $M^3$  model for a PHP program and contains various facts about the program. The creation of an  $M^3$  for PHP is explained in more details in section 4.1. Once the  $M^3$  model is constructed, the second step is to extract constraints from the program using the  $M^3$  model. How and which constraints are extracted is described in section 4.2. In the third step, in section 4.3, the constraints are solved and resolve the types of variables used in the program. The final section 4.4 of this chapter explains how annotations can be included in the process to gain more precise results.

### 4.1 $M^3$ for PHP

As explained in section 2.4,  $M^3$  is a language independent meta model which holds facts about programs. The model can be extended with language specific elements and will be used to query the system for facts about the system. An overview how an  $M^3$  for PHP is build is shown in figure 4.1. Independent  $M^3$  models are build each PHP file in the program. All these individual  $M^3$ 's are in the end combined to collect all the facts about a program. This results in one  $M^3$  for the whole program.  $M^3$ 's are first created for each file is because it is not defined what the dependencies of a individual file are. There is no main file, and all files can load each other. In this research we assume that all files in a program are loaded when needed using autoloaders or manual includes.



Figure 4.1:  $M^3$  Creation

#### 4.1.1 The algorithm

In order to get a better understanding of the creation of an  $M^3$  for PHP, the algorithm is provided in Algorithm 1. The input of the algorithm is all the PHP files of a program, which are all files ending on .php. The output is an  $M^3$  with facts about the provided program. The algorithm starts with initialising an empty  $M^3$  collection in line 1 which with be filled with the result of each individual files. In the big loop on line 2 to 13 we iterate over all the PHP files of the program. The first thing that needs to be done is to create an Abstract Syntax Tree (AST) of the program using an external PHP parser which returns an Rascal AST in parseUsingPhpParserAndReturnRascalAST on line 3. From this AST we create an

<sup>&</sup>lt;sup>1</sup>https://github.com/ruudvanderweijde/PHP-Parser

empty  $M^3$  in createEmptyM3 on line 4. In order to be able to refer to any source code element, we need to have the declarations of the elements. These elements are extracted in addDeclarations on line 5 can be namespace, class, interface, trait, method, function, variable, field, or constant. For all functions we need to add scope information in case functions are declared inside another function or method. In line 6 we add the scope information to all nodes to the AST by visiting the AST. This is needed in order to extract the right facts about the logical containment which is done in line 7. The next three lines (8-10) extract facts about class inheritance and the interface implementations, modifiers, PHP doc blocks, and annotations. Once all the basic information is collected, calculateUsesFlowInsensitive on line 11 tries to find the declarations of the used objects, methods, functions and variables. This is flow and context insensitive and only tries to resolve non-dynamic language constructs. The last step of the iterator is to add the constructed  $M^3$  to the  $M^3$  collection. Finally when an  $M^3$  is constructed for all files, all facts of the individual  $M^3$ 's are merged into one  $M^3$  model in line 14.

```
Algorithm 1: PHP program to M^3
  Input: PHP files of a program
  Output: M^3 for the PHP program
\mathbf{1} m3Collection = [];
2 forall the file \in program do
      ast = parseUsingPhpParserAndReturnRascalAST(file);
3
      m3 = createEmptyM3(ast);
4
      m3 = addDeclarations(m3, ast);
5
      ast = addScopeInformation(m3, ast);
6
      m3 = addContainment(m3, ast);
      m3 = addExtendsAndImplements(m3, ast);
8
      m3 = addModifiers(m3, ast);
10
      m3 = addRawDocBlocksAndAnnotations(m3, ast);
      m3 = calculateUsesFlowInsensitive(m3, ast);
11
      m3Collection += m3;
12
13 end
14 return projectM3 = composePhpM3(m3Collection);
```

#### 4.1.2 Core elements

The  $M^3$  model has the following core elements: Declarations, Containment, Modifiers, Extends, and Uses. All these core elements are explained in more detail in the paragraphs below.

Declarations defines the declarations of namespaces, classes, interfaces, traits, methods, functions, and variables and holds the relation between the logical name (which is used to refer to the declaration) and the actual file location (which is the physical place in the file system. For example the logical name of a class can be |php+class://SomeNameSpace/ClassX| while the actual location might be |file:///project/SomeNameSpace/ClassX.php|.

Containment holds information about what elements logically contain other elements. For example, a property or method is contained in a class and a class is contained in a package. When a function is declared in another function, they are both logically contained in the global namespace (the highest level) because all functions are declared as first class citizens in PHP.

Modifiers element contains information about the modifiers of classes, fields, and methods. Classes can only be abstract, fields can be public, private, or protected, and methods can be all of them. Abstract methods can only be declared in abstract classes. Classes are implicitly public.

Extends contains information about what classes and interface extend other classes of interfaces. Please note that we do not hold information about which class implements which interface, because that information is contained in the implements relation. Interface extensions work just like class extensions.

Uses relation holds information about the usages of certain elements. It is the relation between the usage and the declaration. For example when you instantiate a class, in that case you 'use' that specific class.

#### 4.1.3 PHP specific elements

Because every programming language differs in syntax and semantics, the  $M^3$  model is extensible to provide language specific modeling elements. The following php specific items are added: Implements, TraitUses, Parameters, Constructors, Aliases, and Annotations. These PHP specific modeling elements are described in more detail in the paragraphs below.

Implements holds information on what class implements which interfaces. One class can implement multiple interfaces. Because the information is a relation between the class and the interface, we can easily add multiple interfaces to one class in the model.

TraitUses knows which traits a class uses. Traits are a collection of reusable functions, are defined on class level, and one class can have multiple trait usages. All the methods of a trait are on runtime imported in the class.

Parameters keeps track of the parameters of a method or function. This information is stored to make it easier to resolve the call to a method or function.

Constructors lists all the constructors for classes. This information is needed because it is not always clear what constructor is used, due to legacy PHP4 way of using class constructors. In PHP4 the constructor was defined as a method with the same name as the class. Since PHP5 the language is provided with a magic method \_\_construct(), which results in two ways to have constructors, but only one constructor will be called. The PHP4 constructors will be removed in PHP7.

Aliases has a relation between aliases and the actual implementation. For instance the function class\_alias defines a new name for the same class. This relation is also used to keep track of references.

Annotations contain a relation between a declaration and the annotations that are known. Annotations are defined in the raw doc blocks and are parsed using regular expressions (regex). Regex was in this case easier then parsing because there is no official grammar or standard defined. For this research we only use <code>@param</code>, <code>@var</code>, <code>@returns</code>.

#### **Rascal 1** $M^3$ definitions in Rascal

#### 4.2 Fact extraction

Besides the just discussed items of the  $M^3$  model, we need to add more facts to the  $M^3$  model. We extract fact about classes, class-constants, class-fields, methods, functions, and parameters. We store these facts in a relation in Rascal. The left side of the relation contains the class, function, or method.

Other facts that will be used:

#### 4.2.1 Type extraction

In order to define the subtype relations in class extensions, we will need to declare all existing class types. We can do this in rascal like is done in the example below:

```
visit (system) {
    case c:class(_, _, _, _, _): types += class(c@decl);
}
     Once all types are defined, we can add the subtype relation. We will need to have the subtype of
int() and float() and the class extensions. You can see that in the code below:
public rel[TypeSymbol, TypeSymbol] getSubTypes(M3 m3, System system)
{
    rel[TypeSymbol, TypeSymbol] subtypes
        // add int() as subtype of float()
        = { <\int(), float()> }
        // use the extends relation from M3
        + { <class(c), class(e)> | <c,e> <- m3@extends }
        // add subtype of object for all classes which do not extends a class
        + \{ \langle class(c@decl), object() \rangle \mid 1 \langle -system, /c:class(n,_,noName(),_,_) \langle -system[l] \}
};
    // compute reflexive transitive closure and return the result
    return subtypes*;
}
```

#### 4.2.2 Constraint extraction

The constraints used like defined in the book [PS94]. (based on these rules, we can add constraints to the source code)

Introduction is needed here... for now I will just list the types that I have found. Maybe this needs to be moved to a different chapter.

This is a list of items which are not supported (yet):

- References (in PHP they are symbol table aliases)
  - on expression assignments :: a = &b
  - on functions :: function &f() {...}
  - on parameters :: function  $f(\&\$a) \{...\}$
- Variable structures:
  - Variable variables :: \$\$a;
  - Variable class instantiation :: new \$a;
  - Variable method or function calls :: a();
- List assign :: list(\$a,\$b) = array("one","two"); (we can assume that the rhs is of type array, when the program is correct)
- Method or function parameters (including type hints)
- Class structures, method calls
- Class Constants
- The global statement (should be resolved by the usage relation from M3)
- Casts of expressions
- Parameters
- Predefined variables (\$this, self, parent, static)
- Eval (will not be supported)
- Closures (not used much in production code)
- Traits (not used much in production code)
- Callable (introduced in 5.4 as typehint, not used much in production code)
- Foreach(a as ... (=> ...)) => a is an array or an object;
- return; => return type is null (is added to the situation when there are no return statements)
- add predefined globals (and their type: \$[GLOBALS, \_SERVER, \_GET, \_POST, \_REQUEST, \_COOKIE, \_ENV, \_SESSION, php\_errormsg] (all in global scope))
- add magic constants: \_\_[DIR, FILE, LINE, NAMESPACE, FUNCTION, CLASS, METHOD]\_\_
- $\bullet \ \, \text{predefined constants: } \ \, \text{TRUE(b), } \ \, \text{FALSE(b), } \ \, \text{NAN(f), } \ \, \text{INF(f), } \ \, \text{NULL(n), } \ \, \text{STDIN(r), } \ \, \text{STDOUT(r), } \\ \, \text{STDERR(r)}$
- define("name", value) mixed with constants (?out of scope?)
- keywords: self, parent, static in a class (is included in method and property calls )
- ADD CONSTANTS! RECORD THE TYPE OF THE DEFINED CONSTANTS AND TRY TO READ THEIR TYPE.

#### Legend

```
Equal to (expression)
  \equiv
           Equal to (type)
                                                        A class
 <:
           (lhs) is subTypeOf (rhs)
                                                        A class constant
 E_k
           An expression
                                                        A class property
[E_k]
           Type of some expression
                                                        A class method
                                           \rightarrow m
           A function
                                             [m]
                                                        (Return) type of a method call
 [f]
           (Return) type of a function
                                                        The n'th actual argument
                                           (A_n)
           Static property fetch
                                           (P_n)
                                                        The n'th formal parameter
:: c
           Static method call
                                              th
                                                        Type hint
:: m
                                                        Default value
           Static property fetch
:: p
Mfs
           Modifiers
                                               \Gamma
                                                        The program
```

Table 4.1: Constraint legend

A list of predefined items can be found here:

- Constants
- Variables

todo: functionstodo: classes

#### Expressions

Normal assignment

```
E \equiv (E_1 = E_2)[E_2] <: [E_1],[E_1] = [E]
```

Listing 4.1: Normal assignment

Ternary

$$E \equiv (E_1 ? E_2 : E_3)$$

$$[E] <: [E_2] \lor [E] <: [E_3]$$

$$E \equiv (E_1 ? : E_3)$$

$$[E] <: [E_1] \lor [E] <: [E_3]$$

```
$\frac{1}{2} \text{$\frac{\pmatrix}{\pmatrix}} \text{$\frac{\pmatrix}
```

Listing 4.2: Ternary

Assignments with operators (1) always resulting in ints

$$E_1 \& = E_2$$
 $E_1 \mid = E_2$ 
 $E_1 \hat{} = E_2$ 
 $E_1 <<= E_2$ 
 $E_1 >>= E_2$ 
 $E_1 \% = E_2$ 
 $E_1 \mid = integer()$ 

Listing 4.3: Assignments with operators (1)

Assignments with operators (2) string concat (. =)

$$\frac{E_1 := E_2}{[E_1] = string(),}$$

$$if([E_2] <: object()) => hasMethod([E_2], "\_tostring")$$

```
1  $a .= $b;  /* $a = string() */
2  // An error occurs when $b is of type object() and
3  // __toString is not defined or does not return a string
```

Listing 4.4: Assignments with operators (2)

Assignments with operators (3) resulting in int where rhs is no array

$$E_1 /= E_2$$

$$E_1 -= E_2$$

$$[E_1] = integer(),$$

$$[E_2] \neq array(any())$$

```
$\ \begin{aligned}
\begin
```

Listing 4.5: Assignments with operators (3)

Assignments with operators (4) resulting in int or float

$$\frac{E_1 *= E_2}{E_1 += E_2} \frac{E_1 = E_2}{[E_1] <: float()}$$

Listing 4.6: Assignments with operators (4)

Unary operators

$$E \equiv (+E_1) \lor (-E_1)$$

$$[E] <: float(),$$

$$[E_1] \neq array(\backslash any())$$

$$E \equiv (!E_1)$$

$$[E] = boolean()$$

```
[E_1] = float() \lor [E_1] = integer() \lor [E_1] = string(),
                                      [E] = integer() \lor [E] = string()
                                         E \equiv (E_1 + +) \vee (E_1 - -)
                           if([E_1] <: array(\any()) => [E] <: array(\any()),
                                 if([E_1] = boolean()) => [E] = boolean(),
                                   if([E_1] = float()) => [E] = float(),
                                 if([E_1] = integer()) => [E] = integer(),
                            if([E_1] = null()) => [E] = integer() \vee [E] = null,
                                  if([E_1] <: object()) => [E] <: object(),
                               if([E_1] = resource()) => [E] = resource(),
                 if([E_1] = string()) = > [E] = integer() \lor [E] = float() \lor [E] = string()
                                        E \equiv (+ + E_1) \vee (- - E_1)
                           if([E_1] <: array(\any()) => [E] <: array(\any()),
                                 if([E_1] = boolean()) => [E] = boolean(),
                                   if([E_1] = float()) => [E] = float(),
                                 if([E_1] = integer()) => [E] = integer(),
                                     if([E_1] = null()) \Longrightarrow [E] = null,
                                  if([E_1] <: object()) => [E] <: object(),
                               if([E_1] = resource()) => [E] = resource(),
                 if([E_1] = string()) = > [E] = integer() \lor [E] = float() \lor [E] = string()
+$a // positive
-$a // negation
!$a // not
~$a // bitwise not
$a++ // post increase
$a-- // post decrease
++$a // pre increase
```

 $E \equiv (\sim E_1)$ 

Listing 4.7: Unary operators

#### Binary operators

--\$a // pre decrease

$$\frac{E \equiv (E_1 + E_2)}{[E] <: array(\_) \lor [E] <: float(),}$$

$$if([E_1] <: array(\_) \land [E_2] <: array(\_)) => [E] <: array(\_),$$

$$if([E_1]! <: array(\_) \lor [E_2]! <: array(\_)) => [E] <: float()$$

$$\frac{E \equiv (E_1 - E_2) \lor (E_1 * E_2) \lor (E_1 / E_2)}{[E] <: float(),}$$

$$[E_1] \neq array(\_),$$

$$[E_2] \neq array(\_)$$

$$\frac{E \equiv (E_1 \% E_2) \lor (E_1 << E_2) \lor (E_1 >> E_2)}{[E] = integer()}$$

```
E \equiv (E_1 \& E_2) \lor (E_1 \mid E_2) \lor (E_1 \hat{E}_2)
E[] = string() \lor [E] = integer(),
if([E_1] = string() \land [E_2] = string()) => [E] = string(),
if([E_1] \neq string() \lor [E_2] \neq string()) => [E] = integer()
```

```
1 $a + $b // addition
2 $a - $b // subtraction
3 $a * $b // mulitiplication
4 $a / $b // division
5 $a % $b // modulus
6 $a & $b // bitwise And
7 $a | $b // bitwise Or
8 $a ^ $b // bitwise Shift left
10 $a >> $b // bitwise shift right
```

Listing 4.8: Binary operators

#### Comparison operators

```
E \equiv (E_1 == E_2)
E \equiv (E_1 == E_2)
E \equiv (E_1 != E_2)
E \equiv (E_1 <= E_2)
E \equiv (E_1 <= E_2)
E \equiv (E_1 < E_2)
E \equiv (E_1 >= E_2)
```

```
$a == $b
         /* boolean() */
$a === $b /* boolean() */
$a != $b
          /* boolean() */
$a <> $b
          /* boolean() */
$a !== $b /* boolean() */
$a < $b
          /* boolean() */
$a > $b
          /* boolean() */
         /* boolean() */
$a <= $b
         /* boolean() */
$a >= $b
```

Listing 4.9: Comparison operators

#### Logical operators

$$E \equiv (E_1 \text{ and } E_2)$$

$$E \equiv (E_1 \text{ or } E_2)$$

$$E \equiv (E_1 \text{ xor } E_2)$$

$$E \equiv (E_1 \&\& E_2)$$

$$E \equiv (E_1 || E_2)$$

$$E \equiv boolean()$$

Listing 4.10: Logical operators

#### Array

Array value fetch

```
\begin{split} \underline{E &\equiv E_1[E_2]} \\ \hline [E_1] &\neq object(), \\ if([E_1] &= string()) => [E] = string(), \\ if([E_1] &= array(\{types\})) => [E] <: \{types\}, \\ if([E_1] &\neq string() \land [E_1] \neq array(\_)) => [E] = null() \end{split}
```

Listing 4.11: Array value fetch

Array declaration

```
E', where E' is an array declaration |E'| <: array(any())
```

```
1 array(/*...*/); // typeOf() = array();
2 // Rascal: array(_) => array(\any())
```

Listing 4.12: Array declaration

#### **Scalars**

Scalars

$$E, E ext{ is a string}$$
 $E = string()$ 
 $E, E ext{ is a float}$ 
 $E = float()$ 
 $E, E ext{ is a integer}$ 
 $E = integer()$ 

```
1 "Str" // string()
2 'abc' // string()
3 100 // integer()
4 1.4 // float()
```

Listing 4.13: Scalars

Encapsulated strings

$$\frac{E, E \text{ is an encapsed string*}}{[E] = string()}$$

\* When a string contains expression(/variables), it is processed as encapsed.

```
"$var"
```

Listing 4.14: Encapsulated strings

#### Casts

Casts

Note: PHP Warnings are ignored

$$E \equiv (array)E_1$$

$$E \equiv (bool)E_1 \lor (boolean)E_1$$

$$E \equiv (bool)E_1 \lor (double)E_1 \lor (real)E_1$$

$$E \equiv (float)E_1 \lor (double)E_1 \lor (real)E_1$$

$$E \equiv (int)E_1 \lor (integer)E_1$$

$$E \equiv (int)E_1 \lor (integer)E_1$$

$$E \equiv (object)E_1$$

$$E \equiv (object)E_1$$

$$E \equiv (string)E_1$$

$$E \equiv (string)E_1$$

$$E \equiv string(),$$

$$if([E_1] \lessdot object()) => hasMethod([E_1], `\__tostring')$$

$$E \equiv (unset)E_1$$

$$E \equiv (unset)E_1$$

$$E \equiv (unset)E_1$$

$$E \equiv (unset)E_1$$

```
1 (array)$a // <: array(\any())
2 (bool)$a // boolean()
3 (float)$a // float()
4 (int)$a // integer()
5 (object)$a // object()
6 (string)$a // string(), when $a == object() the object needs to have __toString()
7 (unset)$a // null()</pre>
```

Listing 4.15: Casts

#### Clone

Clone

$$\frac{E \equiv clone(E_1)}{[E] <: object(), [E_1] <: object()}$$

```
1 clone($a) // typeOf($a) = object, typeOf(clone($a)) = object
```

Listing 4.16: Clone

#### Class

Class instantiation (1) matching the class name

$$E \equiv new C_1()$$
$$[E] = class(C.decl)$$

```
1 new C;
```

Listing 4.17: Class instantiation (1)

Class instantiation (2) of an expression

$$\frac{E \equiv new \ E_1}{[E] <: object(),}$$
$$[E_1] <: object() \lor [E_1] = string()$$

```
1  $c = "C";
new $c;
```

Listing 4.18: Class instantiation (2)

Special keywords self parent parent static

$$E \equiv \$this \in C$$

$$[E] <: object()$$

$$[E] = class(C) \lor [E] :> class(C)$$

$$\frac{E \equiv self \in C}{[E] <: object()}$$

$$[self] = class(C)$$

$$\frac{E \equiv parent \in C}{[E] <: object()}$$

$$[E] :> class(C)$$

$$\frac{E \equiv static \in C}{[E] <: object()}$$

$$([E] <: class(C) \lor [E] :> class(C))$$

```
1
// $this can only be used within a class
2 $this // in class C -> class(C)
3 self // in class C -> class(C)
4 parent // in class C -> parentOf(class(C))
5 static // in class C -> class(C) or parentOf(class(C))
```

Listing 4.19: Special keywords

#### Class property fetch

\* Possible add fact that the field E is declared in class C, when it is on the left side of an assignment.

$$\$this \to E_1 \subseteq C_1$$
 
$$\overline{[E_1] = C_1.\text{hasProperty}(E_1.\text{name, static} \notin \text{Mfs}) \lor}$$
 
$$[E_1] = C_1.\text{parent.hasProperty}(E_1.\text{name, public|protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor$$
 
$$[E_1] = C_1[\text{.parent}].\text{hasMethod}("\_\_\text{get"})$$
 
$$\frac{self :: E_1 \subseteq C_1}{\overline{[E_1] = C_1.\text{hasProperty}(E_1.\text{name, static} \in \text{Mfs})}}$$
 
$$\underline{parent :: E_1 \subseteq m}$$
 
$$\overline{[E_1] = C.\text{parent.hasProperty}(E_1.\text{name, static} \in \text{Mfs})}$$
 
$$\underline{E_1 \to E_2 \subseteq C_1}^*$$
 
$$\overline{[E_1] = C_1.\text{hasProperty}(E_2.\text{name, static} \notin \text{Mfs}) \lor}$$
 
$$[E_1] = C_1.\text{parent.hasProperty}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor}$$
 
$$[E_1] = C.\text{hasProperty}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor}$$

\*The same goes for static property fetches, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E_1 \to E_2 \not\subseteq C \subseteq \Gamma^*$$
 
$$[E_1] = \text{C.hasProperty}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \notin \text{Mfs})$$

\*Property fetch outside a class scope, also for static properties.

```
$\frac{1}{2} \text{$\frac{1}{2} \text{$\frac{1} \text{$\frac{1} \text{$\frac{1} \text{$\frac{1} \text{$\frac{1} \text{$\fr
```

Listing 4.20: Class property fetch

Class property fetch variable

$$E \equiv E_1 \to E_2, E_2 \text{ is an expression}$$
$$[E_1] <: \text{ object()}$$

```
1  $b = "b";
2  $a->$b
```

Listing 4.21: Class property fetch variable

Class method call

$$E \equiv \$this \to E_1 \subseteq C_1$$
 
$$[\$this] <: \text{ object}(), [\$this] = \text{class}(C) \lor [E] :> \text{class}(C),$$
 
$$[E_1] \text{ isMethod}(), [E_1] \text{ hasName}(E_1.name \lor "\__call"),$$
 
$$[E_1] = C_1.\text{hasMethod}(E_1.\text{name}, \text{static} \notin \text{Mfs}) \lor$$
 
$$[E] <: [E_1]$$

$$E \equiv self :: E_1 \subseteq C_1$$
 
$$E_1] = C_1.\text{hasMethod}(E_1.\text{name, static} \in \text{Mfs}) \vee$$
 
$$[E_1] = C_1.\text{parent.hasMethod}(E_1.\text{name, public|protected} \in \text{Mfs} \land \text{static} \in \text{Mfs}) \vee$$
 
$$[E_1] = C_1.\text{hasMethod}("\_\_\text{callStatic"})$$
 
$$E \equiv parent :: E_1 \subseteq C_1$$
 
$$E_1] = C_1.\text{parent.hasMethod}(E_1.\text{name, public|protected} \in \text{Mfs}) \vee$$
 
$$[E_1] = C_1.\text{parent.hasMethod}("\_\_\text{callStatic"})$$
 
$$E \equiv E_1 \rightarrow E_2 \subseteq C_1^*$$
 
$$E_1] = C_1.\text{hasMethod}(E_2.\text{name, static} \notin \text{Mfs}) \vee$$
 
$$E_1] = C_1.\text{parent.hasMethod}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \vee$$
 
$$[E_1] = C_1.\text{parent.hasMethod}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \vee$$
 
$$[E_1] = C.\text{hasMethod}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \vee$$

\*The same goes for static method calls, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E \equiv E_1 \to E_2 \not\subseteq C \subseteq \Gamma^*$$
 
$$[E_1] = \text{C.hasMethod}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \notin \text{Mfs})$$

\*method call outside a class scope, also for static methods. This stuff is old!!!!!!

$$\begin{array}{c} \$this \rightarrow E_1 \subseteq C_1 \\ \hline [E_1] = C_1. \text{hasMethod}(E_1. \text{name, static} \notin \text{Mfs}) \lor \\ [E_1] = C_1. \text{parent.hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C_1[. \text{parent}]. \text{hasMethod}("\_\_\text{call"}) \\ \hline \frac{self :: E_1 \subseteq C_1}{[E_1] = C_1. \text{hasMethod}(E_1. \text{name, static} \in \text{Mfs}) \lor } \\ [E_1] = C_1. \text{hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \in \text{Mfs}) \lor \\ \hline [E_1] = C_1. \text{hasMethod}("\_\_\text{callStatic"}) \\ \hline \frac{parent :: E_1 \subseteq C_1}{[E_1] = C_1. \text{parent.hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs}) \lor } \\ \hline [E_1] = C_1. \text{parent.hasMethod}("\_\_\text{callStatic"}) \\ \hline \frac{E_1 \rightarrow E_2 \subseteq C_1 *}{[E_1] = C_1. \text{hasMethod}(E_2. \text{name, static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C_1. \text{parent.hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs}) \lor } \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{ static} \notin \text{Mfs})$$

\*The same goes for static method calls, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E_1 \to E_2 \not\subseteq C \subseteq \Gamma^*$$
 
$$[E_1] = \text{C.hasMethod}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \not\in \text{Mfs})$$

\*method call outside a class scope, also for static methods.

```
1 $this->methodCall();
2 self::methodCall();
3 parent::methodCall();
4 \sa->methodCall();
5  $a::methodCall();
```

Listing 4.22: Class method call

Class method call variable

$$\frac{E \equiv E_1 \to E_2(), E_2 \text{ is an expression}}{[E_1] <: \text{ object()}}$$

```
1 | $a->$methodCall()
```

Listing 4.23: Class method call variable

Class constants (needs to be reviewed)

$$\frac{\text{self}::c_1 \subseteq \Gamma}{[\text{self}::c_1] = C_1.\text{hasConstant}(E_2.\text{name}) \vee}$$
 
$$[\text{self}::c_1] = C_1.\text{parent.hasConstant}(E_2.\text{name}, \text{public}|\text{protected} \in \text{Mfs})$$
 
$$\text{parent}::c_1 \subseteq \Gamma$$
 
$$[\text{self}::c_1] = C_1.\text{parent.hasConstant}(E_2.\text{name}, \text{public}|\text{protected} \in \text{Mfs})$$
 
$$\frac{E_1::c_1 \subseteq \Gamma}{[E_1] = \text{object}()}$$

```
self::CONST
parent::CONST
SOMECLASS::CONST
```

Listing 4.24: Class constants (needs to be reviewed)

#### **Parameters**

Parameters in class instantiation

\*These parameters are just examples for what happens if they have typeHints (th), default values(v) or none \*The constructor can be found in the m3 model (@constructors(loc classDecl, loc constructorMethodDecl))

```
new \ C_1 \ (A_1, \ A_2, \dots, \ A_k) \subseteq \Gamma  class \ C \ (th_1 \ P_1, \ P_2 = v, \dots, P_k) \subseteq \Gamma  sa \rightarrow m() \ (A_1, \ A_2, \dots, \ A_k) \subseteq \Gamma public function  m() \ (th_1 \ P_1, \ P_2 = v, \dots, P_k) \subseteq \Gamma  function_1 \ (A_1, \ A_2, \dots, \ A_k) \subseteq \Gamma function  (th_1 \ P_1, \ P_2 = v, \dots, P_k) \subseteq \Gamma  [P_1] <: [A_1], \ [A_1] <: [th_1], \ [P_1] <: [th_1], \ has Required Param(P_1), \ has Required Param(P_k)
```

```
1 new C($foo);
```

Listing 4.25: Parameters in class instantiation

#### Scope

Type of a certain variable within some scope this applies to global- class- function- and method- scope

$$E, E', E'', E''' \dots etc \subseteq f \qquad E \text{ is a variable}$$
$$[E] = [E] \vee [E'] \vee [E''] \vee [E'''] \dots etc$$

```
function f() {
    $a = 1;
    $a = "true";
}
// typeOf($a) is typeOf($a1, $a2, ..., $an);
```

Listing 4.26: Type of a certain variable within some scope

Return type of function or method (1) having no return statements or return;

$$\frac{\text{return} \not\subseteq f \vee \text{return}; \subseteq f}{[f] = null()}$$

```
function f() {} // no return = null()
function f() { return; } // return; = null()
```

Listing 4.27: Return type of function or method (1)

Return type of function or method (2) every exit path ends with a return statement

```
\frac{(\text{return } E_1) \vee (\text{return } E_2) \vee \cdots \vee (\text{return } E_k) \subseteq f}{[f] <: [E_1] \vee [E_2] \vee \cdots \vee [E_k]}
```

```
function f() {
   if (rand(0,1))
   return $a;
else
   return $b;
}
// returns typeOf($a) or typeOf($b)
```

Listing 4.28: Return type of function or method (2)

Return type of function or method (3) possible no return value

```
(\text{return } E_1) \vee (\text{return } E_2) \vee \cdots \vee (\text{return } E_k) \vee (\neg \text{ return}) \subseteq f[f] <: [E_1] \vee [E_2] \vee \cdots \vee [E_k] \vee null()
```

```
function f() {
   if (rand(0,1))
   return $a;
   else if (rand(0,1))
    return $b;
}
// returns typeOf($a) or typeOf($b) or null()
```

Listing 4.29: Return type of function or method (3)

#### Function calls

Function call

$$\frac{f() \subseteq \Gamma}{[f()] <: \text{return of } [f]}$$

```
function f() {}
f();
```

Listing 4.30: Function call

Function call variable

$$\begin{split} \underline{E \equiv E_1() \subseteq \Gamma} \\ \underline{[E] = any(),} \\ [E_1] <: object() \vee [E_1] = string(), \\ \text{i} f([E_1] <: object()) => hasMethod("\__invoke") \end{split}$$

```
function f() {}

f
```

Listing 4.31: Function call variable

How to resolve expressions:

- Find all expressions which are defined above and annotate them with @type.
- Annotate the rest of the expressions with @type = any(); (should only be for relevant expressions)

### 4.3 Constraint solving

When we have all the constraints from the source code as facts, we will solve the constraints until we can no longer solve any constraints. The result will be a list of possible types for each class, method, fields, functions, variable and expression. The first step is to initialise all type-able objects. In this initial phase the types of each object is of type any, except for the literal types which can be resolved already. When we solve the constraints step by step, we will be able to limit the number of possible types for a certain object. We do this by taking the intersection of the constraint result and the possible types we have. This way there should be less and less possible types for each variable.

When we take the intersection of none overlapping types, we have a type error. There is no possible type for this object, resulting in an empty set of possible types.

The algorithm In algorithm 2 we show the algorithm to solve the constraints. The algorithm input is a set of constraints and the output is a map of solutions. The set of constraints are the constraints defined in constraint extracting section, and the map of results is the location of all type-able objects with their possible solutions. The init function on line 1 adds all type-able objects to the set of possible solutions. The initial possible types of the objects will be the collection of all possible types, called Universe(). Only the literal types can be resolved already, like strings, numbers and boolean constants

Line starts a loop until a fixed point is reached. In Rascal this method is called Solve() and will loop until there are no more changes in the given data. The methods of line 3-4 may alter the values of constraints or solutions.

The first function of the loop on line 3, deiveMore, visits all constraints and checks if there can be new constraints extracted based on the latest constraints and solutions. Example of these new constraints are conditional constraints, which could only be added if more information of the conditional was available.

Function propagate Solutions on line 4 propagates resolved estimates. Here we take the intersection of the known solutions with the given constraints.

#### **Algorithm 2:** Constraint solving algorithm

#### 4.4 Annotations

After we gathers all the facts from the source code, we will add additional information which we read from the annotations. For this we use regex to match @return type and @param type var for methods and functions. We read @var type for variables and class attributes. In our first analysis we do not include the facts we gathered from the annotations. In the second analysis we do include the facts. This way we can compare the end results.

In order to gain some knowledge about the reliability of the annotations, we compare the result our or initial analysis with the provided provided annotation information. Here the implementation should comply to the used annotations.

# Analysis

In order to validate the performed research we have tested them on the most popular packages of Packagist<sup>1</sup>, which are listed in table 5.1. The statistics are generated using phploc<sup>2</sup>. All packages have between 2 and 6 million downloads. Packagist is a repository for Composer<sup>3</sup> projects. Composer is a dependency manager for PHP projects. All external plugins for PHP projects can be managed via a composer json file. You only need know the name and a version of the external package in the repository of your project.

	Product	F	iles	C	bjects		Lines of code					
Vendor Project*		Version	$D^1$	$F^2$	$C_3$	$\mathrm{I}^4$	$\mathrm{T}^5$	Total ↑	Logical		$\mathrm{Global}^6$	
doctrine	lexer	v1.0	2	7	3	0	0	733	128	(17.46%)	13	(10.16%)
phpunit	php-timer	1.0.5	5	11	5	0	0	740	117	(15.81%)	17	(14.53%)
phpunit	php-text-template	1.2.2	5	11	5	0	0	768	125	(16.28%)	15	(12.00%)
doctrine	inflector	v1.0	2	7	3	0	0	853	130	(15.24%)	13	(10.00%)
psr-fig	log	1.0.0	3	15	8	2	2	1039	155	(14.92%)	22	(14.19%)
phpunit	php-file-iterator	1.3.4	5	13	7	0	0	1071	176	(16.43%)	15	(8.52%)
symfony	filesystem	v2.5.3	3	11	5	2	0	1 090	193	(17.71%)	19	(9.84%)
symfony	yaml	v2.5.3	3	16	11	1	0	2270	509	(22.42%)	28	(5.50%)
phpunit	php-token-stream	1.2.2	6	13	169	0	0	2360	377	(15.97%)	15	(3.98%)
doctrine	collections	v1.2	3	18	11	3	0	2504	394	(15.73%)	33	(8.38%)
symfony	process	v2.5.3	3	19	14	1	0	3198	604	(18.89%)	37	(6.13%)
symfony	finder	v2.5.3	8	43	36	3	0	4976	909	(18.27%)	80	(8.80%)
symfony	dom-crawler	v2.5.3	12	63	53	6	0	7825	1296	(16.56%)	157	(12.11%)
symfony	translation	v2.5.3	21	121	97	20	0	12345	2299	(18.62%)	257	(11.18%)
symfony	console	v2.5.3	17	84	66	13	2	13546	2556	(18.87%)	246	(9.62%)
symfony	http-foundation	v2.5.3	16	90	76	10	0	14179	2262	(15.95%)	154	(6.81%)
twig	twig	v1.16.0	18	172	148	19	0	14689	2630	(17.90%)	15	(0.57%)
symfony	event-dispatcher	v2.5.3	27	170	133	31	3	20230	3629	(17.94%)	418	(11.52%)
swiftmailer	swiftmailer	v5.2.1	37	238	170	52	0	28965	4645	(16.04%)	144	(3.10%)
phpunit	php-code-coverage	2.0.1	62	259	381	24	0	50371	6579	(13.06%)	87	(1.32%)
phpunit	phpunit	4.2.2	65	270	388	26	0	51516	6764	(13.13%)	129	(1.91%)
phpunit	phpunit-mock-objects	2.2.0	66	271	393	27	0	51735	6 801	(13.15%)	132	(1.94%)
doctrine	annotations	v1.2.0	69	306	423	28	0	57325	7718	(13.46%)	188	(2.44%)
doctrine	common	v2.4.2	76	337	440	45	0	62406	8 3 2 6	(13.34%)	298	(3.58%)
symfony	http-kernel	v2.5.3	96	565	471	90	3	79294	14169	(17.87%)	1449	(10.23%)
doctrine	cache	1.3.0	152	687	729	102	2	103024	16667	(16.18%)	1355	(8.13%
doctrine	dbal	v2.4.2	121	557	628	63	0	104 630	15234	(14.56%)	1 033	(6.78%)
guzzle	guzzle	v3.9.2	150	832	828	141	7	117699	19772	(16.80%)	1787	(9.04%)
doctrine	orm	v2.4.4	175	1007	875	119	2	158530	27932	(17.62%)	2866	(10.26%
monolog	monolog	1.10.0	350	1911	1904	135	2	288507	31415	(10.89%)	4221	(13.44%)
werkspot	old-Website	07-2014	928	6225	4907	224	0	1054686	167978	(15.93%)	22693	(13.51%)

<sup>\*</sup>This is a list of the 30 most popular packages of packagist ordered by total lines of code, in July 2014.  $^1$  = Directories,  $^2$  = Files,  $^3$  = Classes,  $^4$  = Interfaces,  $^5$  = Traits,  $^6$  = Not in class or function

Table 5.1: List of analysed projects.

To collect the source code for each project, we have executed the following steps:

1. git clone the github repo.

<sup>&</sup>lt;sup>1</sup>https://packagist.org/explore/popular, July 2014

<sup>&</sup>lt;sup>2</sup>https://github.com/sebastianbergmann/phploc, July 2014

 $<sup>^3</sup>$ https://getcomposer.org/, July 2014

- 2. Run composer install, and the source code including dependencies will be downloaded in the /vendor folder.
- 3. Remove the autoload.php and composer folder, as we don't need them.
- 4. Remove the test folders by removing all folders matching Tests or tests.

#### To measure the coverage:

- 1. git clone the github repo.
- 2. Run composer install, and the source code including dependencies will be downloaded in the /vendor folder.
- 3. Run the unit tests and use xdebug to resolve the types.
- 4. Compare the results.

# Results

The results below are based on our analysis on PhpStorm. The first table contains information about the source folder of the test projects. The second table includes vendor folders, which makes is harder to pin point the problems, but does give a better indication of the impact of using annotations. In either cases the vendor folders are used in the analysis to resolve types. The type resolution is based on the type inference implementation of PhpStorm.

#### 6.1 Results

	Unresolved types											
					Ţ	Jnique types		Multiple types				
Product	Total	w/o doc	with doc	$\Delta$	w/o doc	with doc	$\Delta$	w/o doc	with doc	$\Delta$		
php-timer	21	12	7	(-83%)	8	13	(77%)	1	1	(0%)		
php-code-coverage	1444	583	396	(-64%)	809	991	(37%)	52	57	(18%)		
phpunit-mock-objects	800	243	117	(-104%)	539	642	(32%)	18	41	(112%)		
Twig	3720	1 346	1 081	(-39%)	2 286	2458	(14%)	88	181	(103%)		
doctrine2	11452	5 3 2 0	2843	(-93%)	5 929	8 057	(53%)	203	552	(126%)		
php-text-template	41	13	8	(-77%)	27	32	(31%)	1	1	(0%)		
lexer	65	27	8	(-141%)	37	50	(52%)	1	7	(171%)		
php-token-stream	398	123	90	(-54%)	270	300	(20%)	5	8	(75%)		
inflector	39	25	6	(-152%)	10	29	(131%)	4	4	(0%)		
php-file-iterator	96	37	23	(-76%)	57	59	(7%)	2	14	(171%)		
dbal	7770	2993	1 848	(-77%)	4 658	5516	(31%)	119	406	(141%)		
monolog	2018	689	390	(-87%)	1 304	1536	(30%)	25	92	(146%)		
common	1583	736	316	(-114%)	826	1151	(56%)	21	116	(164%)		
collections	518	240	174	(-55%)	270	312	(27%)	8	32	(150%)		
cache	524	282	226	(-40%)	235	272	(27%)	7	26	(146%)		
phpunit	5010	2115	1065	(-99%)	2 798	3796	(53%)	97	149	(70%)		
annotations	709	286	171	(-80%)	416	520	(40%)	7	18	(122%)		
guzzle3	0	0	0	(0%)	0	0	(0%)	0	0	(0%)		
log	185	68	8	(-176%)	117	177	(68%)	0	0	(0%)		

Table 6.1: Results of type usage, source folder only

-		Ur	resolved typ	es			Resolve	d types		
					U	nique types		Multiple types		
Product	Total	w/o doc	with doc	$\Delta$	w/o doc	with doc	$\Delta$	$\rm w/o~doc$	with doc	$\Delta$
php-timer	13751	4671	2 998	(-72%)	8 615	10157	(30%)	465	596	(44%)
php-code-coverage	12717	4174	2621	(-74%)	8 103	9536	(30%)	440	560	(43%)
phpunit-mock-objects	13559	4633	2984	(-71%)	8 468	9989	(30%)	458	586	(44%)
Twig	5518	1856	1570	(-31%)	3 533	3717	(10%)	129	231	(88%)
doctrine2	63148	22830	15408	(-65%)	39 425	45307	(26%)	893	2433	(127%)
php-text-template	162	65	39	(-80%)	91	108	(31%)	6	15	(120%)
lexer	174	75	35	(-107%)	93	118	(42%)	6	21	(143%)
php-token-stream	723	228	174	(-47%)	485	527	(16%)	10	22	(109%)
inflector	177	83	43	(-96%)	85	116	(53%)	9	18	(100%)
php-file-iterator	217	89	54	(-79%)	121	135	(21%)	7	28	(150%)
dbal	31828	10901	6988	(-72%)	20 202	23397	(27%)	725	1443	(100%)
common	17997	6280	3986	(-73%)	11 169	13153	(30%)	548	858	(72%)
collections	866	343	255	(-51%)	510	565	(19%)	13	46	(143%)
cache	36379	11970	7974	(-67%)	23 615	26830	(24%)	794	1575	$(99\%)^{\circ}$
phpunit	13201	4452	2841	(-72%)	8 298	9783	(30%)	451	577	(44%)
annotations	16341	5546	3641	(-69%)	10 234	11978	(29%)	561	722	(45%)
log	294	116	35	(-140%)	173	245	(59%)	5	14	(129%)

Table 6.2: Results of type usage, including vendor folder

### 6.2 Validation of the results

Say something about:

- Soundness (what we measured, is it correct?)
- Completeness (how much did we measure?)
- Accuracy (how precise are the results?)

### 6.3 Annotations

The results of the analysis when adding the annotations to the analysis. Compare the results with the results of the analysis without the annotation information.

# Case Study

Explain how the case study is performed.

This chapter will show the case study, but I just need to place this information somewhere. A list of the 40 most popular packages from packages.

- ullet create composer file
- composer install
- mkdir phploc
- $\bullet$  list all packages: find ./vendor/\* -maxdepth 1 -mindepth 1 -type d -exec ls -d ""
- $\bullet$  prefix the list with "phploc" and postfix with " > phploc/<file>.phploc"

## Conclusion

Summary of the whole work, with conclusions. T.B.A.

#### 8.1 Conclusion

#### 8.2 Future work

These items will not be covered by the analysis (maybe add this to threats/future work)

- Analysis is flow insensitive
- Closure
- References
- Variable constructs (variable -variable, -method/function calls, -class instantiation, eval) :: todo: explain WHY not.
- Yields

Explain something about combining this analysis to other analysis (like dead code elimination, constant folding/propagation resolve, alias analysis, array analysis) to gain more precise results.

Something about performance optimisations... Explain what is already done to boost the performance and what still can be done.

Use a bigger corpus to gains better results of the analysis by doing analysis on more programs.

### 8.3 Threats to validity

When relying on annotations for analysis, the annotations need to be correct and sound. The hard part is that the annotations are not examined on runtime, and are therefor not easily validated.

# Glossary

#### AST

An abstract representation of the structure of the source code. .

#### PSR

PHP Standard Recommendation (PSR) is project which provides rules for commonalities between PHP projects. Autoloading, coding style guide, logging, and HTTP Message interface are the first few accepted standards. The PHPDoc standard describes how and what to use in doc blocks and is currently in draft phase. See http://www.php-fig.org/ for more information. .

#### Rascal

Rascal is a meta-programming language developed by SWAT (Software analyse and transformation) team at CWI in the Netherlands. See http://www.rascal-mpl.org/ for more information.

#### reflexive transitive closure

A relation is transitive if  $\langle a, b \rangle \in R$  then  $\langle b, a \rangle \in R$ .

A relation is reflexive if  $\langle a, b \rangle \in R$  and  $\langle b, c \rangle \in R$  then  $\langle a, c \rangle \in R$ .

A reflexive transitive closure can be established by creating direct paths for all indirect paths and adding self references, until a fixed point is reached.

#### stdClass

A predefined class in the PHP library. The class is the root of the class hierarchy. It is comparable to the Object class in Java.

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