

Type inference for PHP

The value of annotations in a dynamic language

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Abstract

Dynamic language are generally hard to statically analyse because of run-time dependencies. Without running the program there are many things unknown. Because dynamic languages are PHP are widely used, the need for decent analysis tool grows. This research examines the value of adding annotations to PHP code to improve the analysability. In the results we see that annotations improve the analysability of software code (this is a guess). Here I should state something about the correctness of the annotations. And end with a general conclusion.

Preface

In this section I will thank everyone who has helped me. Maybe also introduce some anecdote on how this research came to be.

Chapter 1

Introduction

1.1 PHP

PHP¹ is a server-side scripting language created by Rasmus Lerdorf in 1995. The original name ‘Personal Home Page’ changed to ‘PHP: Hypertext Preprocessor’ in 1998. PHP source files are executed using the PHP Interpreter. The language is dynamically typed, which means that the types of variables are examined during run-time. In statically typed languages all variable types are known at compile time. PHP supports duck-typing, which allows variables to change types during execution.

Evolution The programming language PHP evolved after its creation in 1995. In the year 2000 Object-Oriented (OO) language structures were added to the language with the release of PHP 4.0. The 5th version of PHP was released in 2004 and provided an improved OO structure. Namespaces were added in PHP 5.3 in 2009, to be able to resolve class naming conflicts between library and create better readable class names. Namespaces are comparable to packages in Java. OPcache extension is added was added in PHP 5.5 and speeds up the performance of including files on run-time by storing precompiled script bytecode in shared memory. The most recent stable version is 5.6 and includes more internal performance optimisations and introduces a new debugger.

Popularity According to the Tiobe Index² of december 2014, PHP is the 6th most popular programming language. The language has been in the top 10 since its introduction in the Tiobe Index in 2001. More than 80 percent of the websites have a php backend³. The majority of these websites use PHP version 5, rather than version 4 or older versions. It is therefore wise to focus on PHP version from 5 and discard the older unused versions.

Analysability Although the popularity for more than a decade, there is still a lack of good PHP code analysis tools. Tools can help to reveal security vulnerabilities or find vulnerabilities or bugs in source code. The tools can also provide code completions or do automatic transformations which can be used to execute refactoring patterns. Source code analysis can be performed statically or dynamically or a combination of the two. More information on the analysability of php can be found in section ??.

1.2 Position

In this research we want to find out how we can improve the analysability of programs written in PHP. We will show that the use of annotated source code can help to improve the analysability. The correctness of the annotations can also be examined by checking the implementation of the code. These annotations can help to improve the analysis. The results can be used to find security issues, and if they are highly reliable we can even make compiler optimisations.

¹<http://php.net>

²<http://www.tiobe.com/index.php/content/paperinfo/tpci/index.html>, December 2014

³<http://w3techs.com/technologies/details/pl-php/all/all>, December 2014

As far as we know, there is no constraint based type inference research like this one performed for PHP. That makes this research unique. There have been similar analysis for other dynamic languages, like smalltalk, ruby and javascript, but none like this.

1.3 Contribution

TODO Review this part when the result of the analysis are performed.

- Created an M3 for php.
- Constraint system.
- Show the value of annotations for analysis.

Some idea's are that this analysis can help IDE tools to perform transformations on the source code. (But the performance may not be sufficient.)

The creation of the M3 model can help to compare researchers compare PHP programs with other programming languages. For now only Java is implemented, but more can follow (unchecked statement).

1.4 Plan

The rest of this thesis is as follows: chapter 2 contains background and related work. Here we explain important language constructs and explain similar research. In the next chapter 3 the research method is explained, which contains the steps taken in this the research. Chapter 4 describes the actual research. The analysis is presented in chapter 5, with the results in chapter 6. This thesis ends with the conclusions in chapter 8.

Chapter 2

Background and related work

Chapter 2 describes seven important language constructs in section 2.1 which you need to understand in order to understand the difficulties of analysing PHP. The second section, 2.2, gives an introduction to the syntax and usage of annotation in PHP. Section 2.3 explains the used programming language *Rascal*. M^3 is a generic model designed to analyse programming languages and is described in more details in section 2.4. The final section, section 2.6 of this chapter describes the related work and how these researches are related to this thesis.

2.1 PHP Language Constructs

This section presents important language constructs for this research. Explanations of these constructs should help to understand the performed analysis.

Scoping In PHP, all classes and functions are globally accessible once they are declared. All classes and functions are implicitly public, inner classes are not allowed, and conditional functions (see paragraph about conditional classes and functions) will be available in the global scope. If a class or function is declared inside a namespace, their name will be prefixed with the namespace name. Closures (anonymous functions in PHP) have the same scoping rules as variables, but they can inherit variables from outside their scope by providing them in the use statement. In this research we will not support closures because they are fairly new and not much used in practice.

For variables there are three scopes: global-, function-, and method-scope. There is an exception for some predefined global variables which are available everywhere. Examples are `$GLOBALS`, `$_POST`, and `$_GET`. Variables inside a function or method can be aliased to a global variable by adding the keyword `GLOBAL` in front of the variable name. The variable will then be linked to the global variable in the symbol table¹.

Includes In PHP it is possible to include other PHP-files during execution of the program. These files will be loaded inline. This means that if you use an include in the middle of a file, the source code of this file will be inserted virtually at that place. In this research we will not perform an include analysis. Instead we will assume that all files in the project are included during execution.

According to the coding standard that is used in the php community², function- and class-name classes should not appear when using namespaces and autoloading. When a class which is not loaded in memory is instantiated, the autoloading will try to include a file and load the class. The structure of the autoloading is meant to include classes, interfaces, traits and functions and should not have inline code executions which would lead to side-effects.

Conditional classes and functions Once a file is included in the execution, all the classes and functions in the top scope are declared. All class and function declarations within condition statements or within a method or function scope are only declared when the code is executed.

¹<http://php.net/manual/en/language.variables.scope.php>, July 2014

²<http://www.php-fig.org/psr/psr-0/>, July 2014

An example of an conditional statement can be found in listing 2.1. If the class `Foo` or function `bar` to not exist before the statements is executed, then the class and function will not yet declared. When you try to use the class or function, the script will die with an fatal error (if the class or function was not defined before).

```
1 if (!class_exists("Foo"))
2     class Foo { /* ... */ }
3
4 if (!function_exists("bar"))
5     function bar() { /* ... */ }
```

Listing 2.1: Conditional class and function definitions

Listing 2.2 shows when functions and classes will be available. If the first call is `g()` as you can see in line 7, the script will result in a fatal error. When function `f` is executed, function `g` will be declared, but not yet class `C`. The class `C` will be declared once function `g` is executed. Once the functions and classes are declared, they are available in the top scope, possibly prefixed with the name of the namespace.

```
1 function f() {
2     function g() {
3         class C {}
4     }
5 }
6
7 g(); f(); // will fail because 'g();' is not declared yet
8 f(); g(); // will work because 'g();' is declared when calling 'f();'
9 f(); new C(); // will fail because 'g();' needs to be called first
10 f(); g(); new C(); // will work because 'g();' is called and has declared 'f();'
```

Listing 2.2: Conditional function declaration

Dynamic features PHP includes some dynamic features like: include dynamic variables, dynamic class instantiations, dynamic function calls, dynamic function creation, reflection, and `eval`. A previous study by Mark Hills[HKV13] has shown that the dynamic features are not used too much. (please double check this!!) Our focus will not be on trying to analyse these features, because we would need constant propagation. The downside of these dynamic features is that it will probably lower our precision. (please check this, and move this to another section)

Late static binding Late static binding³ is implemented in PHP since version 5.3 by adding the keyword `static` to the language. It is similar to the keyword `self`, but it does not refer to the class it is declared in. The main difference is that `self` can be resolved statically, because it refers to the class it is declared in. `static` can only be resolved on runtime and represents the exact class that is instantiated.

Magic methods In PHP it is allowed to call methods or use properties that do not exists. Normally this would result in a fatal error, but not with the use of magic methods. One of the magic methods is the constructor method `__call`. This method is called when a non-accessible or non-existing method is called.

Dynamic class properties Although it is a good practice to define your class properties, it is not required. On runtime it is possible to add properties to classes, even without the implementation of magic methods. In listing 2.3 you can see a code sample of adding a property after instantiation of a class. The access of the non-existing property `nonExistingProperty` will result in a warning, but code execution will continue and will just return `NULL`. The code on line 5 is where the property is written. The object `$c` will have the `nonExistingProperty` publicly available now. But in a new class instantiation, like you can see on line 6, will not have the property there.

³<http://php.net/manual/en/language.oop5.late-static-bindings.php>, July 2014

```

1 class C {}
2 $c = new C();
3 var_dump($c->nonExistingProperty); // NULL
4 $p = $c->nonExistingProperty = "property now exists";
5 var_dump($p); // string(19) "property now exists"
6 $d = new C;
7 var_dump($d->nonExistingProperty); // NULL

```

Listing 2.3: Dynamic class property

2.2 Annotations

PHP has no native support for annotations. But PHP has a `getDocComment`⁴ method in the `ReflectionClass`.

This `getDocComment` method returns the complete doc block of a certain element as a string. A doc block in php has the format `/**...*/`. Listing 2.4 shows an example of a doc block in PHP.

```

1 namespace Thesis;
2
3 /**
4  * Class Example
5  * @package Thesis
6  */
7 class Example
8 {
9     /**
10     * This is a description of the method getSomething
11     *
12     * @param SomeTypeHint $someTypeHint
13     * @return bool
14     * @throws CustomException
15     */
16     public function getSomething(SomeTypeHint $someTypeHint)
17     {
18         if ($someTypeHint->getType() == 'something') {
19             throw new CustomException(sprintf('%s is not supported', $someTypeHint
20                 ->getType()));
21         }
22         return true;
23     }
24 }

```

Listing 2.4: PHP DocBlock

In PHP Annotations are not part of the official language. They are however widely used. For instance in ZEND, Symfony and Doctrine you can write business logic rules in the form of annotations. These annotations will be parsed and used in real code.

Other annotations are placed on top of classes, methods, functions, and variables. These annotations will help the developers to better understand what the code does. For example you can see what kind of input and output is expected for a method. IDE's will also use this information to better analyse the source code.

Writing annotations is not yet in the PSR standards for PHP, but there is a proposal⁵. For this research we will only focus on the `@param`, `@return`, `@var`, and `@inheritDoc` annotations. The annotations `@return` and `@param` are only useful for functions and class methods. Type hints are described with `@var` and can be used on all structures, but mainly occur on variables and class fields.

⁴<http://php.net/manual/en/reflectionclass.getdoccomment.php>

⁵<https://github.com/php-fig/fig-standards/pull/169/files>, July 2014

There is no official standard for the use of annotations, but most projects follow the phpDocumentor syntax. For this research the following annotations are parsed:

$$\text{@return} = \left\{ \begin{array}{l} \text{@return } type, \quad \text{unconditionally read @return } type. \end{array} \right. \quad (2.1)$$

$$\text{@param} = \left\{ \begin{array}{ll} \text{@param } type \$var, & \text{if '@param } type \$var' \text{ occurs at least once.} \\ \text{@param } \$var type, & \text{else if '@param } \$var type' \text{ occurs at least once.} \\ \text{@param } type, & \text{otherwise try to match '@param } type'. \end{array} \right. \quad (2.2)$$

$$\text{@var} = \left\{ \begin{array}{ll} \text{@var } type \$var, & \text{if '@var } type \$var' \text{ occurs at least once.} \\ \text{@var } \$var type, & \text{else if '@var } \$var type' \text{ occurs at least once.} \\ \text{@var } type, & \text{otherwise try to match '@var } type'. \end{array} \right. \quad (2.3)$$

$$type = \left\{ \begin{array}{ll} type|type, & \text{if '|' in } type. \\ type, & \text{otherwise} \end{array} \right. \quad (2.4)$$

2.3 Rascal

[Rascal](#) is a meta programming language developed by the Centrum Wiskunde & Informatica (CWI)[KSV09]. Rascal is designed to analyse, transform and visualise source code. The language is build on top of Java and implements various constructs of existing programming languages. In this research, most of our code is implement in rascal.

2.4 M^3

The M^3 -model is a model which holds information of source code[Izm+13]. This model is created to analyse one single Java program or compare two or more Java systems with each other. The core of the M^3 -model contains containment, declarations, documentation, modifiers, names, types, uses, messages.

The **declarations** relation contains class, method, variable- information with their logical name and their real location. The type of the relation are **locations** and represent the logical name of the declaration and will be used in the rest of the M^3 . The **containment** relation has information on what declarations are contained in each other. For example a package can contain a class; a class can contain fields and methods or an inner class; a method can contain variables. The **documentation** relation contains all comments from the source code. The **modifiers** relation has information on the modifiers of declarations. Modifiers are abstract, final, public, protected or private. The **names** relation contains a simplified name of the full declarations. The **types** relation has information about the type of the source code elements. The **uses** relation describes what reference is using which object. For instance when a field of a class is used in some expression, the **uses** relation links the field in the expression to the declaration of the field in the class. And lastly, **messages** contains errors, warnings or info statements.

2.5 Type system

The type system defines how rules are applied to types. The system determines validates the type usage with type checking. The process of actually resolving types is called type inference. On one hand the system needs to determine the type of variables, on the other hand it will check the type of the variables.

Type checking Type checking of dynamic languages differs from static languages. In dynamic languages, most type checking is performed during run-time. The down side here is that you actually have to run the program to get feedback on incorrect usage of types. For example when you divide an integer by a boolean with value true. In PHP these operations are (most of the times) valid and will return you a value. In this case, the boolean will be internally converted to the integer 1. This can result in unexpected behaviour and is discouraged to do.

On the other hand we have static type systems which allow type checking without executing the code. Most analysis can be performed during compile time. These systems can catch programming errors related to typing issues and can avoid overhead of runtime type checks. In dynamic languages the type checking is done during run time, and this will decrease the performance.

Type inference In the previous section we explained the type checking. In order to check the types of variables and expression, we need to know the types of the variables and expressions first. The process of resolving the type of a variable or expressions is called type inference. In dynamic languages like PHP it can be difficult to resolve the type of a variable or expression without running the program. When you are able to resolve the type before running the program, you are able to optimise the execution code. This allows you to make performance improvements or avoid memory usage.

In this research we use constraint based type inference. Every tapeable language construct will have a set of possible types. At the start of the analysis everything will have all possible types, we call this Universe. Then we will read all the constraints and will try to solve them. Every time a constraint is solved, the number of possible types for one variable will decrease. This process will be repeated until we have solved all the constraints and can no longer limit the number of types for a variable.

Other forms of type inference are k -CFA (k -control flow analysis) and CPA (cartesian product algorithm). The analysis we perform is flow insensitive, which means that we do not analyse the flows of the program. We do take the context of the programs into account, making the analysis context sensitive.

(todo, describe K-CFA and CPA in more details)

2.6 Related work

Describe these:

- ‘The HipHop Compiler for PHP’[Zha+12] (not much information available on their type inference, only source code)
- ‘Phantom: PHP analyzer for type mismatch’[KSK10a; KSK10b] (investigate this in more details, their focus is on finding type errors)
- PHPLint ⁶(uses a different kind of annotations, not the Java like PhpDocs)
- ‘Soft typing and analyses on PHP programs’[], code implementations: <https://github.com/henkerik/typing> and <https://github.com/marcelosousa/soft-typing-PHP5> (created for php4, code for php5, should check this out, might be able to compare results with this)
- ‘Design and Implementation of an Ahead-of-Time Compiler for PHP’[Big10] (to check in detail)

Also describe their differences with my research.

⁶<http://www.icosaedro.it/phplint>, July 2014

Chapter 3

Research Method

In this research we will try to answer research questions. Section 3.1 describes the research questions. In section 3.2 we will explain more about types, subtypes and their relationship. Section 3.3 explains in which context the research is executed.

3.1 Research question

The main research question we want to answer is:

Will the use of annotations¹ help to do better static source code analysis?

Subquestions:

- What is the accuracy (recall and precision) of our analysis?
- Are annotations used in programs reliable?

The subquestions are here to support the main question. In the first question we want to measure the precision and recall of this research. With precision we want to measure if the results we have are reliable. With recall we want to have the number of results we were able to reveal.

In the second subquestion we want to know if the annotations conform to the documentation. We will measure this by doing the analysis without using annotations, and see if the results fit the annotations.

3.2 Types

PHP is dynamically typed and allow coercions and duck-typing. Coercion make it possible for variables to hold different types Duck-typing checks the object instead of the class.

```
module lang::php::m3::TypeSymbol

data TypeSymbol
  = \any()      | array(TypeSymbol arrayType)
  | boolean()   | class(loc decl)
  | float()     | integer()
  | null()      | object()
  | resource()  | string()
;
```

¹The annotations are limited to: @param, @return, @var, and @inheritDoc.

3.2.1 PHP types

PHP has a similar class inheritance structure and interface implementation as Java. The main difference is that in PHP all class are **public** and that inner classes are not allowed in PHP.

The basis types in PHP are integers, floats (similar to doubles and reals), booleans, strings, arrays, resources and null. When variables are initialised without a values, they are null. The recourse type is a special one which is not important for this research.

3.2.2 Subtypes

Explain something about subtypes here. For now, only this figure 3.1

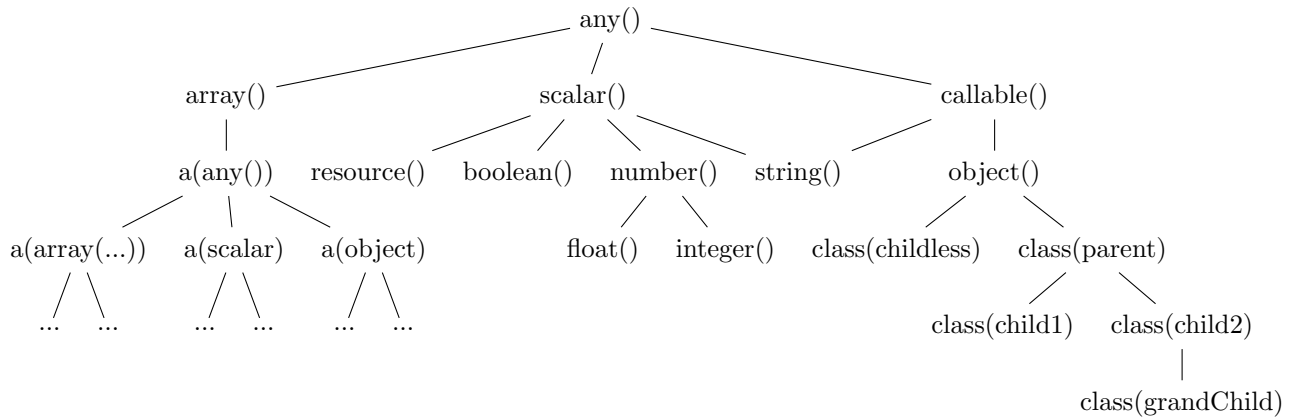


Figure 3.1: Subtype hierarchy

The subtype relation of class inheritance is a **reflexive transitive closure** relation. A class extension of class A on class C will define class A as a subtype of class C in our analysis, as you can see in figure 3.2. If a class does not extend another class, it will implicitly extend the **stdClass** class. You can see that this happens with class D in the example. The **stdClass** is represented as the type **object()** in our analysis.

The basic PHP types also contain a subtype relation. Integers are subtypes of floats.

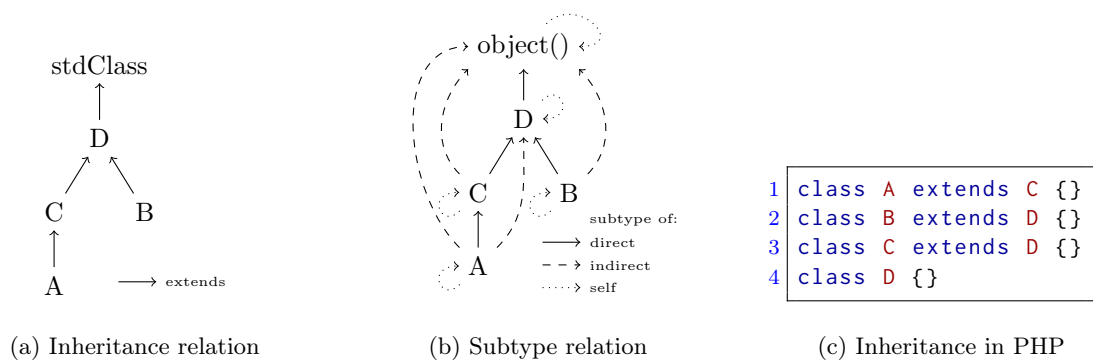


Figure 3.2: Relation of subtypes among classes

3.3 Research context

In order to let our research take place, we need to make sure that some environment variables are constant.

Program Correctness In order to be able to execute this research we will need to assume that the program is correct and works as intended. We will assume that the system contains no bugs. This is needed to be able to say something about the programs we analyse.

File includes In this research we will assume that all file are included during runtime. When a PHP system is constructed of classes with namespaces, the files will be logically loaded using PHP's autoloader. Because most recent systems use namespaces, we will assume that all files are included. For legacy systems, this can influence the results of this research.

Register globals Register globals allows variables to be magically be created from GET and POST values. Since it is discouraged to use this setting, we will assume that all software products have this setting disabled.

PHP Warnings For this research we will ignore all warnings. Warnings do not alter the behaviour of the program. In a normal production environment, the warnings are not shown.

Sensitivity Our analysis is flow-, control-, and context-insensitive. Flow-insensitive means that we do not look at messages between objects, and only look in the body of a method. Control-insensitive means that we ignore all control structures. Examples of control structures are `if`, `else`, and `switch`. Context-insensitive means that we do not look at the order of which code is executed.

Chapter 4

Research

This research is executed in 3 main steps. The first step is to create an M^3 model for a PHP program, containing various facts about the program, which is explained in section 4.1 in more details. Once the M^3 model is constructed, constraints are extracted from the M^3 model which is explained in section 4.2. In the third step, in section 4.3, the constraints are solved and resolve the types of variables used in the program. The final chapter 4.4 explains how annotations can be included in the process to gain more precise results.

4.1 M^3 for PHP

Constructing a M^3 for PHP is differs from construction a M^3 for Java. For Java to constructing a We have constructed a similar M^3 for PHP as was already done for Java. Because language constructs differ per language we had to check which elements were applicable for PHP and which were not. For the overlapping language constructs we created a similar structure. The model will be used to query the system for information about the system. The following steps are performed to create an M^3 for a system.

- Parse all PHP files, resulting in an [AST](#).
- For each file, create an independent M^3 from the ast.
- Combine all the m3s into one M^3 for the project.
- For each file, add additional informationNow run more analysis when all facts are collected in the m3.

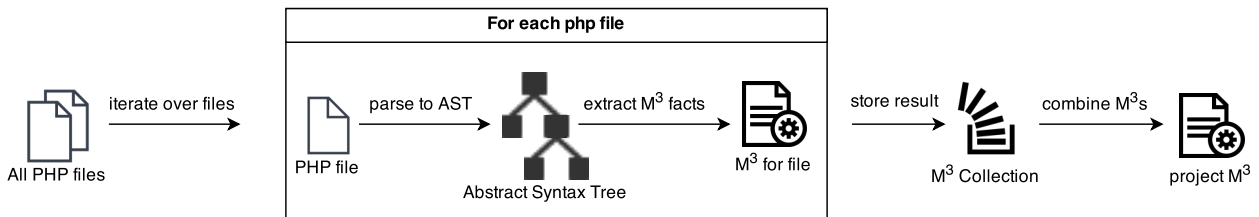


Figure 4.1: M^3 Creation

Our goals is to provide the results in an M3 model. Future research can use this to compare different programming languages.

The following core items are filled for M3:

- Containment

- Declarations (need to explain in more details)
- Modifiers
- Extends
- Uses

The **containment** has information about what elements logically contain other elements. For example, a property or method is contained in a class and a class is contained in a package. In **declarations**, you will find the relation between the logical name and the actual file location. For example the logical name of a class can be `|php+class://SomeNameSpace/ClassX|` while the actual location might be `|file:///project/SomeNameSpace/ClassX.php|`. The **modifiers** element contains information about the public, private, or protected modifiers of a certain method, class, function or property. **Extends** contains information about what classes and interface extend other classes or interfaces. Please note that we do not hold information about which class implements which interface, because that information is contained in the **implements** relation. The **uses** relation holds information about the usages of certain elements. For example when you instantiate a class, in that case you 'use' that specific class.

The following php specific items are added:

- Implements (which class implements which interfaces)
- TraitUses (which class uses which trait)
- Parameters (methods and functions and their parameters)
- Constructors (which class uses which constructor, explain this more)
- Aliases (class aliases, for example the usage of `class_alias`)
- Annotations (constraints annotations on classes, methods, fields and variables)

Implements holds information on what class implements which interfaces. **TraitUses** knows what traits a class uses. **Parameters** keeps track of the parameters of a method or function. **Constructors** lists all the constructors for classes. This is needed because it is not always clear what constructor is used, due to legacy PHP4 way of using class constructors. **Aliases** lists the class aliases, created by the `class_alias` function. **Annotations** has all the annotations for a certain element.

Algorithm 1: PHP program to M^3

Input: PHP files of a program

Output: M^3 for the PHP program

```

1 m3Collection = [];
2 forall the file ∈ program do
3     ast = parseUsingPhpParserAndReturnRascalAST(file);
4     ast = addPublicModifierWhenNotProvided(ast);
5     m3 = createEmptyM3(ast);
6     m3 = addDeclarations(m3, ast);
7     ast = addScopeInformation(m3, ast);
8     m3 = addContainment(m3, ast);
9     m3 = addExtendsAndImplements(m3, ast);
10    m3 = addModifiers(m3, ast);
11    m3 = addRawDocBlocksAndAnnotations(m3, ast);
12    m3 = calculateAliasesFlowInsensitive(m3, ast);
13    m3 = calculateUsesFlowInsensitive(m3, ast);
14    m3Collection += m3;
15 end
16 projectM3 = composePhpM3(m3Collection);
17 return projectM3;
```

4.2 Fact extraction

We can extract fact about classes, class-constants/fields/methods, functions, parameters. For these facts, we can use a relation, so we have a many-to-many relation. On the left side we will have the class, function or method. On the right side we have their attribute.

Other facts that will be used:

```
alias PhpParams = lrel[loc decl, set[loc] typeHints, bool isRequired, bool byRef];
data Annotation = returnType(set[TypeSymbol]) | parameterType(loc var, set[TypeSymbol])
  | varType(loc var, set[TypeSymbol]);

anno rel[loc from, loc to] M3@containment;           // 'from' directly contains 'to'
anno rel[loc from, loc to] M3@extends;              // 'from' extends 'to'
anno rel[loc from, loc to] M3@implements;           // 'from' implements 'to'
anno rel[loc decl, PhpParams params] M3@parameters; // formal parameters of functions/methods
anno rel[loc decl, loc to] M3@constructors;         // 'decl' and its constructor 'to'
anno rel[loc decl, Annotation annotation] M3@annotations; // result of parsed php docs
```

4.2.1 Type extraction

In order to define the subtype relations in class extensions, we will need to declare all existing class types. We can do this in rascal like is done in the example below:

```
visit (system) {
  case c:class(_, _, _, _, _): types += class(c@decl);
}
```

Once all types are defined, we can add the subtype relation. We will need to have the subtype of `int()` and `float()` and the class extensions. You can see that in the code below:

```
public rel[TypeSymbol, TypeSymbol] getSubTypes(M3 m3, System system)
{
  rel[TypeSymbol, TypeSymbol] subtypes
    // add int() as subtype of float()
    = { <\int(), float(> }
    // use the extends relation from M3
    + { <class(c), class(e)> | <c,e> <- m3@extends }
    // add subtype of object for all classes which do not extends a class
    + { <class(c@decl), object(> | l <- system, /c:class(n,_,noName(),_,_) <- system[l]
};

  // compute reflexive transitive closure and return the result
  return subtypes*;
}
```

4.2.2 Constraint extraction

(based on these rules, we can add constraints to the source code)

Introduction is needed here... for now I will just list the types that I have found. Maybe this needs to be moved to a different chapter.

This is a list of items which are not supported (yet):

- References (in PHP they are symbol table aliases)

- on expression assignments :: `$a = &$b`
- on functions :: `function &f() {...}`
- on parameters :: `function f(&$a) {...}`
- Variable structures:
 - ~~Variable variables~~ :: `$$a;`
 - ~~Variable class instantiation~~ :: `new $a;`
 - ~~Variable method or function calls~~ :: `$a();`
- List assign :: `list($a,$b) = array("one","two");` (we can assume that the rhs is of type array, when the program is correct)
- ~~Method or function parameters (including type hints)~~
- ~~Class structures, method calls~~
- ~~Class Constants~~
- The global statement (should be resolved by the usage relation from M3)
- ~~Casts of expressions~~
- Parameters
- Predefined variables (`$this`, `self`, `parent`, `static`)
- ~~Eval~~ (will not be supported)
- ~~Closures~~ (not used much in production code)
- ~~Traits~~ (not used much in production code)
- ~~Callable~~ (introduced in 5.4 as typehint, not used much in production code)
- `Foreach($a as ... (=> ...)) =>` `$a` is an array or an object;
- `return; =>` `return` type is null (is added to the situation when there are no return statements)
- add predefined globals (and their type: `$_GLOBALS`, `$_SERVER`, `$_GET`, `$_POST`, `$_REQUEST`, `$_COOKIE`, `$_ENV`, `$_SESSION`, `php_errormsg`) (all in global scope))
- add magic constants: `__`[`DIR`, `FILE`, `LINE`, `NAMESPACE`, `FUNCTION`, `CLASS`, `METHOD`]`__`
- ~~predefined constants: `TRUE(b)`, `FALSE(b)`, `NAN(f)`, `INF(f)`, `NULL(n)`, `STDIN(r)`, `STDOUT(r)`, `STDERR(r)`~~
- `define("name", value)` mixed with constants (?out of scope?)
- ~~keywords: `self`, `parent`, `static` in a class~~ (is included in method and property calls)
- ADD CONSTANTS! RECORD THE TYPE OF THE DEFINED CONSTANTS AND TRY TO READ THEIR TYPE.

Legend

$=$	=	Equal to (type)	C	=	A class
$<:$	=	Is subTypeOf	$\rightarrow c$	=	A class constant
E_k	=	An expression	$\rightarrow p$	=	A class property
$[E_k]$	=	Type of some expression	$\rightarrow m$	=	A class method
f	=	A function	$[m]$	=	(Return) type of a method call
$[f]$	=	(Return) type of a function	(A_n)	=	The n'th actual argument
$:: c$	=	Static property fetch	(P_n)	=	The n'th formal parameter
$:: m$	=	Static method call	th	=	Type hint
$:: p$	=	Static property fetch	v	=	Default value
Mfs	=	Modifiers	Γ	=	Whole program

Table 4.1: Constraint legend

A list of predefined items can be found here:

- [Constants](#)
- [Variables](#)
- todo: functions
- todo: classes

Expressions

Normal assignment

$$\frac{E \equiv (E_1 = E_2)}{\begin{array}{l} [E_2] <: [E_1], \\ [E_1] <: [E] \end{array}}$$

```
1 $a = $b; // [$b] <: [$a], [$a] <: [$a=$b]
2 $c = $d = $e; // [$e] <: [$d], [$d] <: [$c],
3           // [$d] <: [$d=$e], [$c] <: [$c=$d=$e]
```

Listing 4.1: Normal assignment

Ternary

$$\frac{E \equiv (E_1 ? E_2 : E_3)}{[E] <: [E_2] \vee [E] <: [E_3]}$$

$$\frac{E \equiv (E_1 ? : E_3)}{[E] <: [E_1] \vee [E] <: [E_3]}$$

```
1 $expr ? $b : $c; // typeOf(E) is subtypeOf($b) or subtypeOf($c)
2 $expr ? : $c; // typeOf(E) is subtypeOf($expr) or subtypeOf($c)
```

Listing 4.2: Ternary

Assignments with operators (1) always resulting in ints

$$\begin{array}{l} E_1 \& = E_2 \\ E_1 | = E_2 \\ E_1 \wedge = E_2 \\ E_1 << = E_2 \\ E_1 >> = E_2 \\ E_1 \% = E_2 \end{array}$$

$$\frac{}{[E_1] = integer()}$$

```

1 $a &= $b; /* $a = integer() */
2 $a |= $b; /* $a = integer() */
3 $a ^= $b; /* $a = integer() */
4 $a <=<= $b; /* $a = integer() */
5 $a >>= $b; /* $a = integer() */
6 $a %= $b; /* $a = integer() */

```

Listing 4.3: Assignments with operators (1)

Assignments with operators (2) string concat (.=)

$$\frac{E_1 .= E_2}{[E_1] = \text{string()},}$$

$$\text{if}([E_2] <: \text{object}()) \Rightarrow \text{hasMethod}([E_2], \text{"__toString"})$$

```

1 $a .= $b; /* $a = string() */
2 // An error occurs when $b is of type object() and
3 // __toString is not defined or does not return a string

```

Listing 4.4: Assignments with operators (2)

Assignments with operators (3) resulting in int where rhs is no array

$$\frac{E_1 /= E_2}{[E_1] = \text{integer()},}$$

$$\frac{E_1 -= E_2}{[E_1] \neq \text{array}(\text{any}())}$$

```

1 $a /= $b; /* $a = integer() */
2 $a -= $b; /* $a = integer() */
3 // An error occurs when $b is of type array() for /= and -=
4 // Fatal error: Unsupported operand types

```

Listing 4.5: Assignments with operators (3)

Assignments with operators (4) resulting in int or float

$$\frac{E_1 *= E_2}{[E_1] <: \text{float}(),}$$

$$\frac{E_1 += E_2}{[E_1] <: \text{float}(),}$$

```

1 $a *= $b; /* when $b == (boolean()|integer()|null()) */ /* $a = integer() */
2 $a *= $b; /* when $b != (boolean()|integer()|null()) */ /* $a = float() */
3 $a += $b; /* when $b == (boolean()|integer()|null()) */ /* $a = integer() */
4 $a += $b; /* when $b != (boolean()|integer()|null()) */ /* $a = float() */

```

Listing 4.6: Assignments with operators (4)

Unary operators

$$\frac{E \equiv (+E_1) \vee (-E_1)}{[E] <: \text{float}(),}$$

$$[E_1] \neq \text{array}(\backslash \text{any}())$$

$$\frac{E \equiv (!E_1)}{[E] = \text{boolean}()}$$

$$\begin{array}{c}
\frac{E \equiv (\sim E_1)}{[E_1] = \text{float}() \vee [E_1] = \text{integer}() \vee [E_1] = \text{string}(), \\
[E] = \text{integer}() \vee [E] = \text{string}()} \\
\\
\frac{E \equiv (E_1 + +) \vee (E_1 - -)}{\begin{array}{l} \text{if}([E_1] <: \text{array}(\backslash \text{any}()) \Rightarrow [E] <: \text{array}(\backslash \text{any}()), \\ \text{if}([E_1] = \text{boolean}()) \Rightarrow [E] = \text{boolean}(), \\ \text{if}([E_1] = \text{float}()) \Rightarrow [E] = \text{float}(), \\ \text{if}([E_1] = \text{integer}()) \Rightarrow [E] = \text{integer}(), \\ \text{if}([E_1] = \text{null}()) \Rightarrow [E] = \text{integer}() \vee [E] = \text{null}, \\ \text{if}([E_1] <: \text{object}()) \Rightarrow [E] <: \text{object}(), \\ \text{if}([E_1] = \text{resource}()) \Rightarrow [E] = \text{resource}(), \\ \text{if}([E_1] = \text{string}()) \Rightarrow [E] = \text{integer}() \vee [E] = \text{float}() \vee [E] = \text{string}() \end{array}} \\
\\
\frac{E \equiv (+ + E_1) \vee (- - E_1)}{\begin{array}{l} \text{if}([E_1] <: \text{array}(\backslash \text{any}()) \Rightarrow [E] <: \text{array}(\backslash \text{any}()), \\ \text{if}([E_1] = \text{boolean}()) \Rightarrow [E] = \text{boolean}(), \\ \text{if}([E_1] = \text{float}()) \Rightarrow [E] = \text{float}(), \\ \text{if}([E_1] = \text{integer}()) \Rightarrow [E] = \text{integer}(), \\ \text{if}([E_1] = \text{null}()) \Rightarrow [E] = \text{null}, \\ \text{if}([E_1] <: \text{object}()) \Rightarrow [E] <: \text{object}(), \\ \text{if}([E_1] = \text{resource}()) \Rightarrow [E] = \text{resource}(), \\ \text{if}([E_1] = \text{string}()) \Rightarrow [E] = \text{integer}() \vee [E] = \text{float}() \vee [E] = \text{string}() \end{array}}
\end{array}$$

```

1 +$a // positive
2 -$a // negation
3 !$a // not
4 ~$a // bitwise not
5 $a++ // post increase
6 $a-- // post decrease
7 ++$a // pre increase
8 --$a // pre decrease

```

Listing 4.7: Unary operators

Binary operators

$$\begin{array}{c}
\frac{E \equiv (E_1 + E_2)}{[E] <: \text{array}(_) \vee [E] <: \text{float}(), \\
\text{if}([E_1] <: \text{array}(_) \wedge [E_2] <: \text{array}(_)) \Rightarrow [E] <: \text{array}(_), \\
\text{if}([E_1]! <: \text{array}(_) \vee [E_2]! <: \text{array}(_)) \Rightarrow [E] <: \text{float}()} \\
\\
\frac{E \equiv (E_1 - E_2) \vee (E_1 * E_2) \vee (E_1 / E_2)}{[E] <: \text{float}(), \\
[E_1] \neq \text{array}(_), \\
[E_2] \neq \text{array}(_)} \\
\\
\frac{E \equiv (E_1 \% E_2) \vee (E_1 << E_2) \vee (E_1 >> E_2)}{[E] = \text{integer}()}
\end{array}$$

$$\frac{E \equiv (E_1 \& E_2) \vee (E_1 | E_2) \vee (E_1 \wedge E_2)}{[E] = \text{string}() \vee [E] = \text{integer}(),}$$

$$\text{if}([E_1] = \text{string}() \wedge [E_2] = \text{string}()) \Rightarrow [E] = \text{string}(),$$

$$\text{if}([E_1] \neq \text{string}() \vee [E_2] \neq \text{string}()) \Rightarrow [E] = \text{integer}()$$

```

1 $a + $b // addition
2 $a - $b // subtraction
3 $a * $b // multiplication
4 $a / $b // division
5 $a % $b // modulus
6 $a & $b // bitwise And
7 $a | $b // bitwise Or
8 $a ^ $b // bitwise Xor
9 $a << $b // bitwise shift left
10 $a >> $b // bitwise shift right

```

Listing 4.8: Binary operators

Comparison operators

$$E \equiv (E_1 == E_2)$$

$$E \equiv (E_1 === E_2)$$

$$E \equiv (E_1 != E_2)$$

$$E \equiv (E_1 <> E_2)$$

$$E \equiv (E_1 !== E_2)$$

$$E \equiv (E_1 < E_2)$$

$$E \equiv (E_1 > E_2)$$

$$E \equiv (E_1 <= E_2)$$

$$E \equiv (E_1 >= E_2)$$

$$[E] = \text{boolean}()$$

```

1 $a == $b /* boolean() */
2 $a === $b /* boolean() */
3 $a != $b /* boolean() */
4 $a <> $b /* boolean() */
5 $a !== $b /* boolean() */
6 $a < $b /* boolean() */
7 $a > $b /* boolean() */
8 $a <= $b /* boolean() */
9 $a >= $b /* boolean() */

```

Listing 4.9: Comparison operators

Logical operators

$$E \equiv (E_1 \text{ and } E_2)$$

$$E \equiv (E_1 \text{ or } E_2)$$

$$E \equiv (E_1 \text{ xor } E_2)$$

$$E \equiv (E_1 \&\& E_2)$$

$$E \equiv (E_1 || E_2)$$

$$[E] = \text{boolean}()$$

```

1 $a and $b /* boolean() */
2 $a or $b  /* boolean() */
3 $a xor $b /* boolean() */
4 $a && $b  /* boolean() */
5 $a || $b  /* boolean() */

```

Listing 4.10: Logical operators

Array

Array value fetch

$$\begin{array}{c}
\frac{E \equiv E_1[E_2]}{[E_1] \neq \text{object()},} \\
\text{if}([E_1] = \text{string}()) \Rightarrow [E] = \text{string}(), \\
\text{if}([E_1] = \text{array}(\{\text{types}\})) \Rightarrow [E] <: \{\text{types}\}, \\
\text{if}([E_1] \neq \text{string}() \wedge [E_1] \neq \text{array}(_)) \Rightarrow [E] = \text{null}()
\end{array}$$

```

1 $a[0];
2 // typeof($a) != object()
3 // when typeof($a) == string() => typeof($a[/...*/]) is string()
4 // when typeof($a) == array() => typeof($a[/...*/]) is mixed()
5 // when typeof($a) != string|array => typeof($a[/...*/]) is null()

```

Listing 4.11: Array value fetch

Array declaration

$$\frac{E', \text{ where } E' \text{ is an array declaration}}{[E'] <: \text{array}(\text{any}())}$$

```

1 array(/...*/); // typeof() = array();
2 // Rascal: array(_) => array(\any())

```

Listing 4.12: Array declaration

Scalars

Scalars

$$\frac{E, E \text{ is a string}}{[E] = \text{string}()}$$

$$\frac{E, E \text{ is a float}}{[E] = \text{float}()}$$

$$\frac{E, E \text{ is an integer}}{[E] = \text{integer}()}$$


```

1 "Str" // string()
2 'abc' // string()
3 100 // integer()
4 1.4 // float()

```

Listing 4.13: Scalars

Encapsulated strings

$$\frac{E, E \text{ is an encapsled string}^*}{[E] = \text{string}()}$$

* When a string contains expression(/variables), it is processed as encapsled.

```

1 "$var"

```

Listing 4.14: Encapsulated strings

Casts

Casts

Note: PHP Warnings are ignored

$$\frac{E \equiv (\text{array})E_1}{[E] <: \text{array}(\text{any}())}$$

$$\frac{E \equiv (\text{bool})E_1 \vee (\text{boolean})E_1}{[E] = \text{boolean}()}$$

$$\frac{E \equiv (\text{float})E_1 \vee (\text{double})E_1 \vee (\text{real})E_1}{[E] = \text{float}()}$$

$$\frac{E \equiv (\text{int})E_1 \vee (\text{integer})E_1}{[E] = \text{integer}()}$$

$$\frac{E \equiv (\text{object})E_1}{[E] <: \text{object}()}$$

$$\frac{E \equiv (\text{string})E_1}{[E] = \text{string}(),}$$

$$if([E_1] <: \text{object}()) \Rightarrow \text{hasMethod}([E_1], \text{'__toString'})$$

$$\frac{E \equiv (\text{unset})E_1}{[E] = \text{null}()}$$

```

1 (array)$a // <: array(\any())
2 (bool)$a  // boolean()
3 (float)$a // float()
4 (int)$a   // integer()
5 (object)$a // object()
6 (string)$a // string(), when $a == object() the object needs to have __toString()
7 (unset)$a // null()

```

Listing 4.15: Casts

Clone

Clone

$$\frac{E \equiv \text{clone}(E_1)}{[E] <: \text{object}(), [E_1] <: \text{object}()}$$

```
1 clone($a) // typeof($a) = object, typeof(clone($a)) = object
```

Listing 4.16: Clone

Class

Class instantiation (1) matching the class name

$$\frac{E \equiv \text{new } C_1()}{[E] = \text{class}(C.\text{decl})}$$

```
1 new C;
```

Listing 4.17: Class instantiation (1)

Class instantiation (2) of an expression

$$\frac{E \equiv \text{new } E_1}{[E] <: \text{object}(), [E_1] <: \text{object}() \vee [E_1] = \text{string}()}$$

```
1 $c = "C";  
2 new $c;
```

Listing 4.18: Class instantiation (2)

Special keywords self parent parent static

$$\frac{E \equiv \$this \in C}{[E] <: \text{object}()}$$
$$[E] = \text{class}(C) \vee [E] :> \text{class}(C)$$

$$\frac{E \equiv \text{self} \in C}{[E] <: \text{object}()}$$
$$[\text{self}] = \text{class}(C)$$

$$\frac{E \equiv \text{parent} \in C}{[E] <: \text{object}()}$$
$$[E] :> \text{class}(C)$$

$$\frac{E \equiv \text{static} \in C}{[E] <: \text{object}()}$$
$$([E] <: \text{class}(C) \vee [E] :> \text{class}(C))$$

```

1 // $this can only be used within a class
2 $this // in class C -> class(C)
3 self // in class C -> class(C)
4 parent // in class C -> parentOf(class(C))
5 static // in class C -> class(C) or parentOf(class(C))

```

Listing 4.19: Special keywords

Class property fetch

* Possible add fact that the field E is declared in class C, when it is on the left side of an assignment.

$$\frac{\$this \rightarrow E_1 \subseteq C_1}{[E_1] = C_1.\text{hasProperty}(E_1.\text{name}, \text{static} \notin \text{Mfs}) \vee [E_1] = C_1.\text{parent}.\text{hasProperty}(E_1.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs}) \vee [E_1] = C_1[\text{.parent}].\text{hasMethod}(\text{"__get"})}$$

$$\frac{self :: E_1 \subseteq C_1}{[E_1] = C_1.\text{hasProperty}(E_1.\text{name}, \text{static} \in \text{Mfs})}$$

$$\frac{parent :: E_1 \subseteq m}{[E_1] = C.\text{parent}.\text{hasProperty}(E_1.\text{name}, \text{static} \in \text{Mfs})}$$

$$\frac{E_1 \rightarrow E_2 \subseteq C_1^*}{[E_1] = C_1.\text{hasProperty}(E_2.\text{name}, \text{static} \notin \text{Mfs}) \vee [E_1] = C_1.\text{parent}.\text{hasProperty}(E_2.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs}) \vee [E_1] = C.\text{hasProperty}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})}$$

*The same goes for static property fetches, except for the ‘static \notin Mfs’ part: ‘static \in Mfs’.

$$\frac{E_1 \rightarrow E_2 \not\subseteq C \subseteq \Gamma^*}{[E_1] = C.\text{hasProperty}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})}$$

*Property fetch outside a class scope, also for static properties.

```

1 $this->prop // name = prop, vis = public|protected, !static || mm
2 self::$prop // static property in class
3 parent::$prop // static property in the parent(s)
4 $a->prop // non-static property fetch
5 $a::$pro // static property fetch

```

Listing 4.20: Class property fetch

Class property fetch variable

$$\frac{E \equiv E_1 \rightarrow E_2, E_2 \text{ is an expression}}{[E_1] <: \text{object}()}$$

```

1 $b = "b";
2 $a->$b

```

Listing 4.21: Class property fetch variable

Class method call

$$\frac{E \equiv \$this \rightarrow E_1 \subseteq C_1}{[\$this] <: \text{object}(), [\$this] = \text{class}(C) \vee [E] >: \text{class}(C), [E_1] \text{ isMethod}(), [E_1] \text{ hasName}(E_1.\text{name} \vee \text{"__call"}), [E_1] = C_1.\text{hasMethod}(E_1.\text{name}, \text{static} \notin \text{Mfs}) \vee [E] <: [E_1]}$$

$$\begin{array}{c}
\frac{E \equiv \text{self} :: E_1 \subseteq C_1}{[E_1] = C_1.\text{hasMethod}(E_1.\text{name}, \text{static} \in \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(E_1.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \in \text{Mfs}) \vee \\
[E_1] = C_1.\text{hasMethod}(__\text{callStatic}) \\
\\
\frac{E \equiv \text{parent} :: E_1 \subseteq C_1}{[E_1] = C_1.\text{parent}.\text{hasMethod}(E_1.\text{name}, \text{public|protected} \in \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(__\text{callStatic}) \\
\\
\frac{E \equiv E_1 \rightarrow E_2 \subseteq C_1^*}{[E_1] = C_1.\text{hasMethod}(E_2.\text{name}, \text{static} \notin \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(E_2.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs}) \vee \\
[E_1] = C_1.\text{hasMethod}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})
\end{array}$$

*The same goes for static method calls, except for the ‘static \notin Mfs’ part: ‘static \in Mfs’.

$$\frac{E \equiv E_1 \rightarrow E_2 \not\subseteq C \subseteq \Gamma^*}{[E_1] = C.\text{hasMethod}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})}$$

*method call outside a class scope, also for static methods.
This stuff is old!!!!!!

$$\begin{array}{c}
\frac{\$this \rightarrow E_1 \subseteq C_1}{[E_1] = C_1.\text{hasMethod}(E_1.\text{name}, \text{static} \notin \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(E_1.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs}) \vee \\
[E_1] = C_1[\text{parent}].\text{hasMethod}(__\text{call}) \\
\\
\frac{\text{self} :: E_1 \subseteq C_1}{[E_1] = C_1.\text{hasMethod}(E_1.\text{name}, \text{static} \in \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(E_1.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \in \text{Mfs}) \vee \\
[E_1] = C_1.\text{hasMethod}(__\text{callStatic}) \\
\\
\frac{\text{parent} :: E_1 \subseteq C_1}{[E_1] = C_1.\text{parent}.\text{hasMethod}(E_1.\text{name}, \text{public|protected} \in \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(__\text{callStatic}) \\
\\
\frac{E_1 \rightarrow E_2 \subseteq C_1^*}{[E_1] = C_1.\text{hasMethod}(E_2.\text{name}, \text{static} \notin \text{Mfs}) \vee} \\
[E_1] = C_1.\text{parent}.\text{hasMethod}(E_2.\text{name}, \text{public|protected} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs}) \vee \\
[E_1] = C_1.\text{hasMethod}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})
\end{array}$$

*The same goes for static method calls, except for the ‘static \notin Mfs’ part: ‘static \in Mfs’.

$$\frac{E_1 \rightarrow E_2 \not\subseteq C \subseteq \Gamma^*}{[E_1] = C.\text{hasMethod}(E_2.\text{name}, \text{public} \in \text{Mfs} \wedge \text{static} \notin \text{Mfs})}$$

*method call outside a class scope, also for static methods.

```

1 $this->methodCall();
2 self::methodCall();
3 parent::methodCall();
4 $a->methodCall();
5 $a::methodCall();

```

Listing 4.22: Class method call

Class method call variable

$$\frac{E \equiv E_1 \rightarrow E_2(), E_2 \text{ is an expression}}{[E_1] <: \text{object}()}$$

```
1 $a->$methodCall()
```

Listing 4.23: Class method call variable

Class constants (needs to be reviewed)

$$\frac{\text{self}::c_1 \subseteq \Gamma}{[\text{self}::c_1] = C_1.\text{hasConstant}(E_2.\text{name}) \vee [\text{self}::c_1] = C_1.\text{parent}.\text{hasConstant}(E_2.\text{name}, \text{public|protected} \in \text{Mfs})}$$

$$\frac{\text{parent}::c_1 \subseteq \Gamma}{[\text{self}::c_1] = C_1.\text{parent}.\text{hasConstant}(E_2.\text{name}, \text{public|protected} \in \text{Mfs})}$$

$$\frac{E_1::c_1 \subseteq \Gamma}{[E_1] = \text{object}()}$$

```
1 self::CONST
2 parent::CONST
3 SOMECLASS::CONST
```

Listing 4.24: Class constants (needs to be reviewed)

Parameters

Parameters in class instantiation

*These parameters are just examples for what happens if they have typeHints (*th*), default values(*v*) or none *The constructor can be found in the m3 model (@constructors(loc classDecl, loc constructorMethodDecl))

$$\frac{\text{new } C_1 (A_1, A_2, \dots, A_k) \subseteq \Gamma \quad \text{class } C (th_1 P_1, P_2 = v, \dots, P_k) \subseteq \Gamma}{\text{\$a} \rightarrow m() (A_1, A_2, \dots, A_k) \subseteq \Gamma \quad \text{public function } m() (th_1 P_1, P_2 = v, \dots, P_k) \subseteq \Gamma}$$

$$\frac{\text{function}_1 (A_1, A_2, \dots, A_k) \subseteq \Gamma \quad \text{function } (th_1 P_1, P_2 = v, \dots, P_k) \subseteq \Gamma}{[P_1] <: [A_1], [A_1] <: [th_1], [P_1] <: [th_1], \text{hasRequiredParam}(P_1), \text{hasRequiredParam}(P_k)}$$

```
1 new C($foo);
```

Listing 4.25: Parameters in class instantiation

Scope

Type of a certain variable within some scope

this applies to global- class- function- and method- scope

$$\frac{E, E', E'', E''' \dots \text{etc} \subseteq f \quad E \text{ is a variable}}{[E] = [E] \vee [E'] \vee [E''] \vee [E'''] \dots \text{etc}}$$

```
1 function f() {
2     $a = 1;
3     $a = "true";
4 }
5 // typeof($a) is typeof($a1, $a2, ..., $an);
```

Listing 4.26: Type of a certain variable within some scope

Return type of function or method (1) having no return statements or return;

$$\frac{\text{return} \not\subseteq f \vee \text{return}; \subseteq f}{[f] = \text{null}()}$$

```

1 function f() {} // no return = null()
2 function f() { return; } // return; = null()

```

Listing 4.27: Return type of function or method (1)

Return type of function or method (2) every exit path ends with a return statement

$$\frac{(\text{return } E_1) \vee (\text{return } E_2) \vee \dots \vee (\text{return } E_k) \subseteq f}{[f] <: [E_1] \vee [E_2] \vee \dots \vee [E_k]}$$

```

1 function f() {
2   if (rand(0,1))
3     return $a;
4   else
5     return $b;
6 }
7 // returns typeof($a) or typeof($b)

```

Listing 4.28: Return type of function or method (2)

Return type of function or method (3) possible no return value

$$\frac{(\text{return } E_1) \vee (\text{return } E_2) \vee \dots \vee (\text{return } E_k) \vee (\neg \text{return}) \subseteq f}{[f] <: [E_1] \vee [E_2] \vee \dots \vee [E_k] \vee \text{null}()}$$

```

1 function f() {
2   if (rand(0,1))
3     return $a;
4   else if (rand(0,1))
5     return $b;
6 }
7 // returns typeof($a) or typeof($b) or null()

```

Listing 4.29: Return type of function or method (3)

Function calls

Function call

$$\frac{f() \subseteq \Gamma}{[f()] <: \text{return of } [f]}$$

```

1 function f() {}
2 f();

```

Listing 4.30: Function call

Function call variable

$$\frac{E \equiv E_1() \subseteq \Gamma}{[E] = \text{any}(),}$$

$$[E_1] <: \text{object}() \vee [E_1] = \text{string}(),$$

$$\text{if}([E_1] <: \text{object}()) \Rightarrow \text{hasMethod}(\text{"__invoke"})$$

```

1 function f() {}
2 $f = "f";
3 $f(); // unknown what function will be called
4 // [$f] <: object(with __invoke method) | [$f] = string()

```

Listing 4.31: Function call variable

How to resolve expressions:

- Find all expressions which are defined above and annotate them with @type.
- Annotate the rest of the expressions with @type = any(); (should only be for relevant expressions)

4.3 Constraint solving

When we have all the constraints from the source code as facts, we will solve the constraints until we can no longer solve any constraints. The result will be a list of possible types for each class, method, fields, functions, variable and expression. The first step is to initialise all type-able objects we want to solve. In this initial phase the status of all objects can be of any type, because we do not know anything about them yet. When we solve the constraints step by step, we will be able to limit the number of possible types for a certain object. We do this by taking the intersection of the constraint result and the possible types we have. This way there should be less and less possible types for each variable.

When we take the intersection of none overlapping types, we have a type error. There is no possible type for this object. Because PHP allows ducktyping, allowing objects to be of more than one type, we will use widening in these cases. Widening means that we pick the union instead of the intersection, so we can continue our analysis.

Some of the constraints only give us information that a certain object should be a subtype of a certain type. When we solve these constraints, we will take the least common ancestor.

4.4 Annotations

After we gathers all the facts from the source code, we will add additional information which we read from the annotations. For this we use regex to match @return type and @param type var for methods and functions. We read @var type for variables and class attributes. In our first analysis we do not include the facts we gathered from the annotations. In the second analysis we do include the facts. This way we can compare the end results.

In order to gain some knowledge about the reliability of the annotations, we compare the result our or initial analysis with the provided provided annotation information. Here the implementation should comply to the used annotations.

Chapter 5

Analysis

In order to validate the performed research we have tested them on the most popular packages of Packagist¹, which are listed in table 5.1. The statistics are generated using phploc². All packages have between 2 and 6 million downloads. Packagist is a repository for Composer³ projects. Composer is a dependency manager for PHP projects. All external plugins for PHP projects can be managed via a composer.json file. You only need know the name and a version of the external package in the repository of your project.

Product			Files		Objects			Lines of code			
Vendor	Project*	Version	D ¹	F ²	C ³	I ⁴	T ⁵	Total ↑	Logical	Global ⁶	
doctrine	lexer	v1.0	2	7	3	0	0	733	128 (17.46%)	13 (10.16%)	
phpunit	php-timer	1.0.5	5	11	5	0	0	740	117 (15.81%)	17 (14.53%)	
phpunit	php-text-template	1.2.2	5	11	5	0	0	768	125 (16.28%)	15 (12.00%)	
doctrine	inflector	v1.0	2	7	3	0	0	853	130 (15.24%)	13 (10.00%)	
psr-fig	log	1.0.0	3	15	8	2	2	1 039	155 (14.92%)	22 (14.19%)	
phpunit	php-file-iterator	1.3.4	5	13	7	0	0	1 071	176 (16.43%)	15 (8.52%)	
symfony	filesystem	v2.5.3	3	11	5	2	0	1 090	193 (17.71%)	19 (9.84%)	
symfony	yaml	v2.5.3	3	16	11	1	0	2 270	509 (22.42%)	28 (5.50%)	
phpunit	php-token-stream	1.2.2	6	13	169	0	0	2 360	377 (15.97%)	15 (3.98%)	
doctrine	collections	v1.2	3	18	11	3	0	2 504	394 (15.73%)	33 (8.38%)	
symfony	process	v2.5.3	3	19	14	1	0	3 198	604 (18.89%)	37 (6.13%)	
symfony	finder	v2.5.3	8	43	36	3	0	4 976	909 (18.27%)	80 (8.80%)	
symfony	dom-crawler	v2.5.3	12	63	53	6	0	7 825	1 296 (16.56%)	157 (12.11%)	
symfony	translation	v2.5.3	21	121	97	20	0	12 345	2 299 (18.62%)	257 (11.18%)	
symfony	console	v2.5.3	17	84	66	13	2	13 546	2 556 (18.87%)	246 (9.62%)	
symfony	http-foundation	v2.5.3	16	90	76	10	0	14 179	2 262 (15.95%)	154 (6.81%)	
twig	twig	v1.16.0	18	172	148	19	0	14 689	2 630 (17.90%)	15 (0.57%)	
symfony	event-dispatcher	v2.5.3	27	170	133	31	3	20 230	3 629 (17.94%)	418 (11.52%)	
swiftmailer	swiftmailer	v5.2.1	37	238	170	52	0	28 965	4 645 (16.04%)	144 (3.10%)	
phpunit	php-code-coverage	2.0.1	62	259	381	24	0	50 371	6 579 (13.06%)	87 (1.32%)	
phpunit	phpunit	4.2.2	65	270	388	26	0	51 516	6 764 (13.13%)	129 (1.91%)	
phpunit	phpunit-mock-objects	2.2.0	66	271	393	27	0	51 735	6 801 (13.15%)	132 (1.94%)	
doctrine	annotations	v1.2.0	69	306	423	28	0	57 325	7 718 (13.46%)	188 (2.44%)	
doctrine	common	v2.4.2	76	337	440	45	0	62 406	8 326 (13.34%)	298 (3.58%)	
symfony	http-kernel	v2.5.3	96	565	471	90	3	79 294	14 169 (17.87%)	1 449 (10.23%)	
doctrine	cache	1.3.0	152	687	729	102	2	103 024	16 667 (16.18%)	1 355 (8.13%)	
doctrine	dbal	v2.4.2	121	557	628	63	0	104 630	15 234 (14.56%)	1 033 (6.78%)	
guzzle	guzzle	v3.9.2	150	832	828	141	7	117 699	19 772 (16.80%)	1 787 (9.04%)	
doctrine	orm	v2.4.4	175	1007	875	119	2	158 530	27 932 (17.62%)	2 866 (10.26%)	
monolog	monolog	1.10.0	350	1911	1 904	135	2	288 507	31 415 (10.89%)	4 221 (13.44%)	
werkspot	old-Website	07-2014	928	6225	4 907	224	0	1 054 686	167 978 (15.93%)	22 693 (13.51%)	

*This is a list of the 30 most popular packages of packagist ordered by total lines of code, in July 2014.

¹ = Directories, ² = Files, ³ = Classes, ⁴ = Interfaces, ⁵ = Traits, ⁶ = Not in class or function

Table 5.1: List of analysed projects.

To collect the source code for each project, we have executed the following steps:

1. `git clone` the github repo.

¹<https://packagist.org/explore/popular>, July 2014

²<https://github.com/sebastianbergmann/phploc>, July 2014

³<https://getcomposer.org/>, July 2014

2. Run `composer install`, and the source code including dependencies will be downloaded in the `/vendor` folder.
3. Remove the `autoload.php` and `composer` folder, as we don't need them.
4. Remove the test folders by removing all folders matching `Tests` or `tests`.

To measure the coverage:

1. `git clone` the github repo.
2. Run `composer install`, and the source code including dependencies will be downloaded in the `/vendor` folder.
3. Run the unittests and use `xdebug` to resolve the types.
4. Compare the results.

Chapter 6

Results

For the results we picked X software products to see how it performs. For each product we performed the type inference with and without reading annotations from the doc blocks.

6.1 Results

Show the results...

6.2 Validation of the results (or something)

Say something about:

- Soundness (what we measured, is it correct?)
- Completeness (how much did we measure?)
- Accuracy (how precise are the results?)

6.3 Annotations

The results of the analysis when adding the annotations to the analysis. Compare the results with the results of the analysis without the annotation information.

Chapter 7

Case Study

Explain how the case study is performed.

This chapter will show the case study, but I just need to place this information somewhere.

A list of the 40 most popular packages from packages.

- create composer file
- composer install
- mkdir phploc
- list all packages: `find ./vendor/* -maxdepth 1 -mindepth 1 -type d -exec ls -d ""`
- prefix the list with "phploc" and postfix with " > phploc/<file>.phploc"

Chapter 8

Conclusion

Summary of the whole work, with conclusions. T.B.A.

8.1 Conclusion

8.2 Future work

These items will not be covered by the analysis (maybe add this to threats/future work)

- Analysis is flow insensitive
- Closure
- References
- Variable constructs (variable -variable, -method/function calls, -class instantiation, eval) :: todo: explain WHY not.
- Yields

Explain something about combining this analysis to other analysis (like dead code elimination, constant folding/propagation resolve, alias analysis, array analysis) to gain more precise results.

Something about performance optimisations... Explain what is already done to boost the performance and what still can be done.

Use a bigger corpus to gains better results of the analysis by doing analysis on more programs.

8.3 Threats to validity

When relying on annotations for analysis, the annotations need to be correct and sound. The hard part is that the annotations are not examined on runtime, and are therefor not easily validated.

Glossary

AST

An abstract representation of the structure of the source code. .

Rascal

Rascal is a meta-programming language developed by SWAT (Software analyse and transformation) team at CWI in the Netherlands. See <http://www.rascal-mpl.org/> for more information.

reflexive transitive closure

A relation is transitive if $\langle a, b \rangle \in R$ then $\langle b, a \rangle \in R$.

A relation is reflexive if $\langle a, b \rangle \in R$ and $\langle b, c \rangle \in R$ then $\langle a, c \rangle \in R$.

A reflexive transitive closure can be established by creating direct paths for all indirect paths and adding self references, until a fixed point is reached.

stdClass

A predefined class in the PHP library. The class is the root of the class hierarchy. It is comparable to the Object class in Java.

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