Type inference for PHP

The value of annotations in a dynamic language

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January 11, 2015, 37 pages

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Abstract

Dynamic language are generally hard to statically analyse because of run-time dependencies. Without running the program there are many things unknown. Because dynamic languages are PHP are widely used, the need for decent analysis tool grows. This research examines the value of adding annotations to PHP code to improve the analysability. In the results we see that annotations improve the analysability of software code (this is a guess). Here I should state something about the correctness of the annotations. And end with a general conclusion.

Preface

In this section I will thank everyone who has helped me. Maybe also introduce some anecdote on how this research came to be.

Introduction

1.1 PHP

PHP¹is a server-side scripting language created by Rasmus Lerdorf in 1995. The original name 'Personal Home Page' changed to 'PHP: Hypertext Preprocessor' in 1998. PHP source files are executed using the PHP Interpreter. The language is dynamically typed, which means that the types of variables are examined during run-time. In statically typed languages all variable types are known at compile time. PHP supports duck-typing, which allows variables to change types during execution.

Evolution The programming language PHP evolved after its creation in 1995. In the year 2000 Object-Oriented (OO) language structures were added to the langue with the release of PHP 4.0. The 5th version of PHP was released in 2004 and provided an improved OO structure. Namespace were added in PHP 5.3 in 2009, to be able to resolve class naming conflicts between library and create better readable class names. Namespaces are comparable to packages in JAVA. OPcache extension is added was added in PHP 5.5 and speeds up the performance of including files on run-time by storing precompiled script bytecode in shared memory. The most recent stable version is 5.6 and includes more internal performance optimisations and introduces a new debugger.

Popularity According to the Tiobe Index² of december 2014, PHP is the 6th most popular programming language. The language has been in the top 10 since its introduction in the Tiobe Index in 2001. More than 80 procent of the websites have a php backend³. The majority of these websites use PHP version 5, rather than version 4 or older versions. It is therefor wise to focus on PHP version from 5 and discard the older unused versions.

Analysability Although the popularity for more than a decade, there is still a lack of good PHP code analysis tools. Tools can help to reveal security vulnerabilities or find vulnerabilities or bugs in source code. The tools can also provide code completions or do automatic transformations which can be used to execute refactoring patterns. Source code analysis can be performed statically or dynamically or a combination of the two. More information on the analysability of php can be found in section 2.1.

1.2 Position

In this research we want to find out how we can improve the analysability of programs written in PHP. We will show that the use of annotated source code can help to improve the analysability. The correctness of the annotations can also be examined by checking the implementation of the code.

These annotations can help to improve the analysis. The results can be used to find security issues, and if they are highly reliable we can even make compiler optimisations.

¹http://php.net

²http://www.tiobe.com/index.php/content/paperinfo/tpci/index.html, December 2014

³http://w3techs.com/technologies/details/pl-php/all/all, December 2014

As far as we know, there is no constraint based type inference research like this one performed for PHP. That makes this research unique. There have been similar analysis for other dynamic languages, like smalltalk, ruby and javascript, but none like this.

1.3 Contribution

TODO Review this part when the result of the analysis are performed.

- Created an M3 for php.
- Constraint system.
- Show the value of annotations for analysis.

Some idea's are that this analysis can help IDE tools to perform transformations on the source code. (But the performance may not be sufficient.)

The creation of the M3 model can help to compare researchers compare PHP programs with other programming languages. For now only Java is implemented, but more can follow (unchecked statement).

1.4 Plan

The rest of this thesis is as follows: chapter 2 contains background and related work. Here we explain important language constructs and explain similar research. In the next chapter 3 the research method is explained, which contains the steps taken in this the research. Chapter 4 describes the actual research. The analysis is presented in chapter 5, with the results in chapter 6. This thesis ends with the conclusions in chapter 8.

Background and related work

This chapter explains language constructs that are important for this research. The next section explains Rascal and M^3 . The final section of this chapter describes the related work.

2.1 PHP Language Constructs

This section presents important language constructs for this research. Explanations of these constructs should help to understand the performed analysis.

Scoping In PHP, all classes and functions are globally accessible once they are declared. All classes and functions are implicitly public, inner classes are not allowed, and conditional functions (see paragraph about conditional classes and functions) will be available in the global scope. If a class or function is declared inside a namespace, their name will be prefixed with the namespace name. Closures (anonymous functions in PHP) have the same scoping rules as variables, but they can inherhit variables from outside their scope by providing them in the use statement. In this research we will not support closures because they are fairly new and not much used in practice.

For variables there are three scopes: global-, function-, and method-scope. There is an exception for some predefined global variables which are available everywhere. Examples are \$GLOBALS, \$_POST, and \$_GET. Variables inside a function or method can be aliased to a global variable by adding the keyword GLOBAL in front of the variable name. The variable will then be linked to the global variable in the symbol table¹.

Includes In PHP it is possible to include other PHP-files during execution of the program. These files will be loaded inline. This means that if you use an include in the middle of a file, the source code of this file will be inserted virtually at that place. In this research we will not perform an include analysis. Instead we will assume that all files in the project are included during execution.

According to the coding standard that is used in the php community², function- and class-name classes should not appear when using namespaces and autoloading. When a class which is not loaded in memory is instantiated, the autoloading will try to include a file and load the class. The structure of the autoloading is meant to include classes, interfaces, traits and functions and should not have inline code executions which would lead to side-effects.

Conditional classes and functions Once a file is included in the execution, all the classes and functions in the top scope are declared. All class and function declarations within condition statements or within a method or function scope are only declared when the code is executed.

An example of an conditional statement can be found in listing 2.1. If the class Foo or function bar to not exist before the statements is executed, then the class and function will not yet declared. When you try to use the class or function, the script will die with an fatal error (if the class or function was not defined before).

 $^{^{1}\}mbox{http://php.net/manual/en/language.variables.scope.php, July 2014}$

²http://www.php-fig.org/psr/psr-0/, July 2014

```
if (!class_exists("Foo"))
class Foo { /* ... */ }

if (!function_exists("bar"))
function bar() { /* ... */ }
```

Listing 2.1: Conditional class and function definitions

Listing 2.2 shows when functions and classes will be available. If the first call is g() as you can see in line 7, the script will result in a fatal error. When function f is executed, function g will be declared, but not yet class C. The class C will be declared once function g is executed. Once the functions and classes are declared, they are available in the top scope, possibly prefixed with the name of the namespace.

```
function f() {
  function g() {
    class C {}
}

g(); f(); // will fail because 'g(); ' is not declared yet
  f(); g(); // will work because 'g(); ' is declared when calling 'f(); '
  f(); new C(); // will fail because 'g(); ' needs to be called first
  f(); g(); new C(); // will work because 'g(); ' is called and has declared 'f(); '
```

Listing 2.2: Conditional function declaration

Dynamic features PHP includes some dynamic features like: include dynamic variables, dynamic class instantiations, dynamic function calls, dynamic function creation, reflection, and eval. A previous study by Mark Hills[HKV13] has shown that the dynamic features are not used too much. (please double check this!!) Our focus will not be on trying to analyse these features, because we would need constant propagation. The downside of these dynamic features is that it will probably lower our precision. (please check this, and move this to another section)

Late static binding Late static binding³ is implemented in PHP since version 5.3 by adding the keyword static to the language. It is similar to the keyword self, but it does not refer to the class it is declared in. The main difference is that self can be resolved statically, because it refers to the class it is declared in. static can only be resolved on runtime and represents the exact class that is instantiated.

Magic methods In PHP it is allowed to call methods or use properties that do not exists. Normally this would result in a fatal error, but not with the use of magic methods. One of the magic methods is het constructor method <code>__call</code>. This method is called when a non-accessible or non-existing method is called.

Dynamic class properties Although it is a good practice to define your class properties, it is not required. On runtime it is possible to add properties to classes, even without the implementation of magic methods. In listing 2.3 you can see a code sample of adding a property after instantiation of a class. The access of the non-existing property nonExistingProperty will result in a warning, but code execution will continue and will just return NULL. The code on line 5 is where the property is written. The object \$c will have the nonExistingProperty publicly available now. But in a new class instantiation, like you can see on line 6, will not have the property there.

```
class C {}

class C {}

$c = new C();

var_dump($c->nonExistingProperty); // NULL

$p = $c->nonExistingProperty = "property now exists";
```

³http://php.net/manual/en/language.oop5.late-static-bindings.php, July 2014

```
5  var_dump($p); // string(19) "property now exists"
6  $d = new C;
7  var_dump($d->nonExistingProperty); // NULL
```

Listing 2.3: Dynamic class property

2.2 Annotations

PHP has no native support for annotations. But PHP has a getDocComment ⁴method in the ReflectionClass. This getDocComment method returns the complete doc block of a certain element as a string. A doc block in php has the format /**...*/. Listing 2.4 shows an example of a doc block in PHP.

```
namespace Thesis;
2
3
4
   * Class Example
5
   * @package Thesis
6
   */
7
  class Example
8
  {
9
10
        * This is a description of the method getSomething
11
12
        * @param SomeTypeHint $someTypeHint
13
        * @return bool
        * @throws CustomException
14
15
       public function getSomething(SomeTypeHint $someTypeHint)
16
17
18
           if ($someTypeHint->getType() == 'something') {
19
                throw new CustomException(sprintf('%s is not supported', $someTypeHint
                    ->getType()));
20
21
22
           return true;
23
       }
24
  }
```

Listing 2.4: PHP DocBlock

In PHP Annotations are not part of the official language. They are however widely used. For instance in ZEND, Symfony and Doctrine you can write business logic rules in the form of annotations. These annotations will be parsed and used in real code.

Other annotations are placed on top of classes, methods, functions, and variables. These annotations will help the developers to better understand what the code does. For example you can see what kind of input and output is expected for a method. IDE's will also use this information to better analyse the source code.

Writing annotations is not yet in the PSR standards for PHP, but there is a proposal⁵. For this research we will only focus on the <code>@param</code>, <code>@return</code>, <code>@var</code>, and <code>@inheritDoc</code> annotations. The annotations <code>@return</code> and <code>@param</code> are only useful for functions and class methods. Type hints are described with <code>@var</code> and can be used on all structures, but mainly occur on variables and class fields.

There is no official standard for the use of annotations, but most projects follow the phpDocumentor syntax. For this research the following annotations are parsed:

 $\operatorname{Qreturn} = \Big\{ \operatorname{Qreturn} type, \quad \text{unconditionally read Qreturn } type.$ (2.1)

 $^{^4} http://php.net/manual/en/reflectionclass.getdoccomment.php \\$

 $^{^5 \}rm https://github.com/php-fig/fig-standards/pull/169/files, July 2014$

$$\operatorname{@var} = \begin{cases} \operatorname{@var} type \$var, & \text{if `@var} type \$var' \text{ occurs at least once.} \\ \operatorname{@var} \$var type, & \text{else if `@var } \$var type' \text{ occurs at least once.} \\ \operatorname{@var} type, & \text{otherwise try to match `@var } type'. \end{cases}$$

$$(2.3)$$

$$type = \begin{cases} type|type, & \text{if '|' in } type. \\ type, & \text{otherwise} \end{cases}$$
(2.4)

2.3 Rascal

Rascal is a meta programming language developed by the Centrum Wiskunde & Informatica (CWI)[KSV09]. Rascal is designed to analyse, transform and visualise source code. The language is build on top of JAVA and implements various constructs of existing programming languages. In this research, most of our code is implement in rascal.

2.4 M^3

The M^3 -model is a model which holds information of source code[Izm+13]. This model is created to analyse one single JAVA program or compare two or more JAVA systems with each other. The core of the M^3 -model contains containment, declarations, documentation, modifiers, names, types, uses, messages.

The declarations relation contains class, method, variable- information with their logical name and their real location. The type of the relation are locations and represent the logical name of the declaration and will be used in the rest of the M^3 . The containment relation has information on what declarations are contained in each other. For example a package can contain a class; a class can contain fields and methods or an inner class; a method can contain variables. The documentation relation contains all comments from the source code. The modifiers relation has information on the modifiers of declarations. Modifiers are abstract, final, public, protected or private. The names relation contains a simplified name of the full declarations. The types relation has information about the type of the source code elements. The uses relation describes what reference is using which object. For instance when a field of a class is used in some expression, the uses relation links the field in the expression to the declaration of the field in the class. And lastly, messages contains errors, warnings or info statements.

2.5 Type system

The type system defines how rules are applied to types. The system determines validates the type usage with type checking. The process of actually resolving types is called type inference. On one hand the system needs to determine the type of variables, on the other hand it will check the type of the variables.

Type checking Type checking of dynamic languages differs from static languages. In dynamic languages, most type checking is performed during run-time. The down side here is that you actually have to run the program to get feedback on incorrect usage of types. For example when you divide an integer by a boolean with value true. In PHP these operations are (most of the times) valid and will return you a value. In this case, the boolean will be internally concerted to the integer 1. This can result in unexpected behaviour and is discouraged to do.

On the other hand we have static type systems which allow type checking without executing the code.

Most analysis can be performed during compile time. These systems can catch programming errors related to typing issues and can avoid overhead of runtime type checks. In dynamic languages the type checking is done during run time, and this will decrease the performance.

Type inference In the previous section we explained the type checking. In order to check the types of variables and expression, we need to know the types of the variables and expressions first. The process of resolving the type of a variable or expressions is called type inference. In dynamic languages like PHP it can be difficult to resolve the type of a variable or expression without running the program. When you are able to resolve the type before running the program, you are able to optimise the execution code. This allows you to make performance improvements or avoid memory usage.

In this research we use constraint based type inference. Every tapeable language construct will have a set of possible types. At the start of the analysis everything will have all possible types, we call this Universe. Then we will read all the constraints and will try to solve them. Every time a constraint is solved, the number of possible types for one variable will decrease. This process will be repeated until we have solved all the constraints and can no longer limit the number of types for a variable.

Other forms of type inference are k-CFA (k-control flow analysis) and CPA (cartesian product algorithm). The analysis we perform is flow insensitive, which means that we do not analyse the flows of the program. We do take the context of the programs into account, making the analysis context sensitive. (todo, describe K-CFA and CPA in more details)

2.6 Related work

Describe these:

- 'The HipHop Compiler for PHP' [Zha+12] (not much information available on their type inference, only source code)
- 'Phantm: PHP analyzer for type mismatch' [KSK10a; KSK10b] (investigate this in more details, their focus is on finding type errors)
- PHPLint ⁶ (uses a different kind of annotations, not the java like phpdocs)
- 'Soft typing and analyses on PHP programs' [], code implementations: https://github.com/henkerik/typing and https://github.com/marcelosousa/soft-typing-PHP5 (created for php4, code for php5, should check this out, might be able to compare results with this)
- 'Design and Implementation of an Ahead-of-Time Compiler for PHP'[Big10] (to check in detail)

Also describe their differences with my research.

⁶http://www.icosaedro.it/phplint, july 2014

Research Method

In this research we will try to answer research questions. Section 3.1 describes our research questions. Section 3.2 explains in which context the research is executed.

3.1 Research question

The main research question we want to answer is:

Will the use of annotations¹ help to do better static source code analysis?

Subquestions:

- What is the accuracy (recall and precision) of our analysis?
- Are annotations used in programs reliable?

The subquestions are here to support the main question. In the first question we want to measure the precision and recall of this research. With precision we want to measure if the results we have are reliable. With recall we want to have the number of results we were able to reveal.

In the second subquestion we want to know if the annotations conform to the documentation. We will measure this by doing the analysis without using annotations, and see if the results fit the annotations.

3.2 Research context

In order to let our research take place, we need to make sure that some environment variables are constant.

Program Correctness In order to be able to execute this research we will need to assume that the program is correct and works as intended. We will assume that the system contains no bugs. This is needed to be able to say something about the programs we analyse.

File includes In this research we will assume that all file are included during runtime. When a PHP system is constructed of classes with namespaces, the files will be logically loaded using PHP's autoloader. Because most recent systems use namespaces, we will assume that all files are included. For legacy systems, this can influence the results of this research.

Register globals Register globals allows variables to be magically be created from GET and POST values. Since it is discouraged to use this setting, we will assume that all software products have this setting disabled.

 $^{^{1}\}mathrm{The}$ annotations are limited to: <code>Qparam</code>, <code>Qreturn</code>, <code>Qvar</code>, and <code>QinheritDoc</code>.

PHP Warnings For this research we will ignore all warnings. Warnings do not alter the behaviour of the program. In a normal production environment, the warnings are not shown.

Sensitivity Our analysis is flow-, control-, and context-insensitive. Flow-insensitive means that we do not look at messages between objects, and only look in the body of a method. Control-insensitive means that we ignore all control structures. Examples of control structures are if, else, and switch. Context-insensitive means that we do not look at the order of which code is executed.

Research

This research is executed in 4 main steps. First we created an M^3 model for PHP and is explained in section 4.1. The next section, section 4.2 explains how we define the type system. This will include the subtype relation between types. After we define the types, we gather facts from the source code in section 4.3. The next step is to include annotations in chapter 4.4. The final chapter 4.5 shows how we solve the constraints we got from the fact extracting.

4.1 M^3 for PHP

We have constructed a similar M^3 for PHP as was already done for JAVA. Because language constructs differ per language we had to check which elements were applicable for PHP and which were not. For the overlapping language constructs we created a similar structure. The model will be used to query the system for information about the system. The following steps are performed to create an M^3 for a system.

- Parse all PHP files, resulting in an AST.
- For each file, create an independent M^3 from the ast.
- Combine all the m3s into one M^3 for the project.
- For each file, add additional informationNow run more analysis when all facts are collected in the m3.

Our goals is to provide the results in an M3 model. Future research can use this to compare different programming languages.

The following core items are filled for M3:

- Containment
- Declarations (need to explain in more details)
- Modifiers
- Extends
- Uses

The containment has information about what elements logically contain other elements. For example, a property or method is contained in a class and a class is contained in a package. In declarations, you will find the relation between the logical name and the actual file location. For example the logical name of a class can be |php+class://SomeNameSpace/ClassX| while the actual location might be |file:///project/SomeNameSpace/ClassX.php|. The modifiers element contains information about the public, private, or protected modifiers of a certain method, class, function or property. Extends contains information about what classes and interface extend other classes of interfaces. Please note that

we do not hold information about which class implements which interface, because that information is contained in the implements relation. The uses relation holds information about the usages of certain elements. For example when you instantiate a class, in that case you 'use' that specific class.

The following php specific items are added:

- Implements (which class implements which interfaces)
- TraitUses (which class uses which trait)
- Parameters (methods and functions and their parameters)
- Constructors (which class uses which constructor, explain this more)
- Aliases (class aliases, for example the usage of class alias)
- Annotations (contraints annotations on classes, methods, fields and variables)

Implements holds information on what class implements which interfaces. TraitUses knows what traits a class uses. Parameters keeps track of the parameters of a method or function. Constructors lists all the constructors for classes. This is needed because it is not always clear what constructor is used, due to legacy PHP4 way of using class constructors. Aliases lists the class aliases, created by the class_alias function. Annotations has all the annotations for a certain element.

4.2 Types

PHP is dynamically typed and allow coercions and duck-typing. Coercion make it possible for variables to hold different types Duck-typing checks the object instead of the class.

4.2.1 PHP types

PHP has a similar class inheritance structure and interface implementation as Java. The main difference is that in PHP all class are public and that inner classes are not allowed in PHP.

The basis types in PHP are integers, floats (similar to doubles and reals), booleans, strings, arrays, resources and null. When variables are initialised without a values, they are null. The recourse type is a special one which is not important for this research.

4.2.2 Subtypes

Explain something about subtypes here. For now, only this figure 4.1

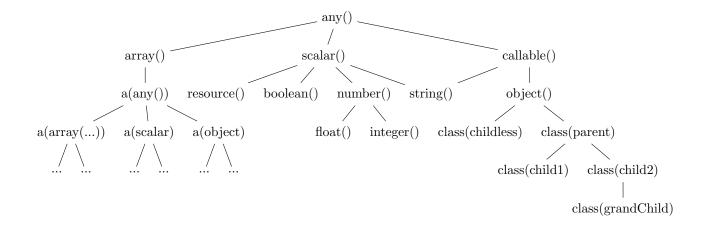


Figure 4.1: Subtype hierarchy

The subtype relation of class inheritance is a reflexive transitive closure relation. A class extension of class A on class C will define class A as a subtype of class C in our analysis, as you can see in figure 4.2. If a class does not extend another class, it will implicitly extend the stdClass class. You can see that this happens with class D in the example. The stdClass is represented as the type object() in our analysis.

The basic PHP types also contain a subtype relation. Integers are subtypes of floats.

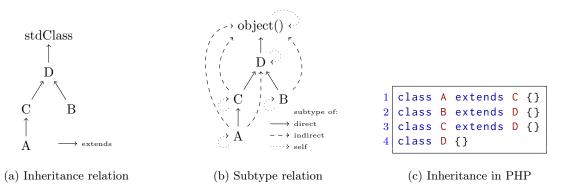


Figure 4.2: Relation of subtypes among classes

4.3 Fact extraction

We can extract fact about classes, class-constants/fields/methods, functions, parameters. For these facts, we can use a relation, so we have a many-to-many relation. On the left size we will have the class, function or method. On the right side we have their attribute.

Other facts that will be used:

```
anno rel[loc decl, PhpParams params] M3@parameters; // formal parameters of functions/methods anno rel[loc decl, loc to] M3@constructors; // 'decl' and its constructor 'to' anno rel[loc decl, Annotation annotation] M3@annotations; // result of parsed php docs
```

4.3.1 Type extraction

In order to define the subtype relations in class extensions, we will need to declare all existing class types. We can do this in rascal like is done in the example below:

```
visit (system) {
    case c:class(_, _, _, _, _): types += class(c@decl);
}
     Once all types are defined, we can add the subtype relation. We will need to have the subtype of
int() and float() and the class extensions. You can see that in the code below:
public rel[TypeSymbol, TypeSymbol] getSubTypes(M3 m3, System system)
{
    rel[TypeSymbol, TypeSymbol] subtypes
        // add int() as subtype of float()
        = { <\int(), float()> }
        // use the extends relation from M3
        + { <class(c), class(e)> | <c,e> <- m3@extends }
        // add subtype of object for all classes which do not extends a class
        + \{ \langle class(c@decl), object() \rangle \mid 1 \langle -system, /c:class(n,_,noName(),_,_) \langle -system[l] \}
};
    // compute reflexive transitive closure and return the result
    return subtypes*;
}
```

4.3.2 Constraint extraction

(based on these rules, we can add constraints to the source code)

Introduction is needed here... for now I will just list the types that I have found. Maybe this needs to be moved to a different chapter.

This is a list of items which are not supported (yet):

- References (in PHP they are symbol table aliases)
 - on expression assignments :: a = &b- on functions :: function f(a)- on parameters :: function f(a)
- Variable structures:
 - Variable variables :: \$\$a;
 - Variable class instantiation :: new \$a;
 - Variable method or function calls :: a();
- List assign :: list(\$a,\$b) = array("one","two"); (we can assume that the rhs is of type array, when the program is correct)
- Method or function parameters (including type hints)

- Class structures, method calls
- Class Constants
- The global statement (should be resolved by the usage relation from M3)
- Casts of expressions
- Parameters
- Predefined variables (\$this, self, parent, static)
- Eval (will not be supported)
- Closures (not used much in production code)
- Traits (not used much in production code)
- Callable (introduced in 5.4 as typehint, not used much in production code)
- For each (a as ... (=> ...)) => a is an array or an object;
- return; => return type is null (is added to the situation when there are no return statements)
- add predefined globals (and their type: \$[GLOBALS, _SERVER, _GET, _POST, _REQUEST, _COOKIE, _ENV, _SESSION, php_errormsg] (all in global scope))
- add magic constants: __[DIR, FILE, LINE, NAMESPACE, FUNCTION, CLASS, METHOD]__
- $\bullet \ \, \text{predefined constants: } \ \, \text{TRUE(b), } \ \, \text{FALSE(b), } \ \, \text{NAN(f), } \ \, \text{INF(f), } \ \, \text{NULL(n), } \ \, \text{STDIN(r), } \ \, \text{STDERR(r)}$
- define("name", value) mixed with constants (?out of scope?)
- keywords: self, parent, static in a class (is included in method and property calls)
- ADD CONSTANTS! RECORD THE TYPE OF THE DEFINED CONSTANTS AND TRY TO READ THEIR TYPE.

Legend

```
Equal to (type)
                                                      A class
<:
           Is subTypeOf
                                          \rightarrow c
                                                      A class constant
           An expression
E_k
                                          \rightarrow p
                                                      A class property
[E_k]
           Type of some expression
                                                 = A class method
                                          \rightarrow m
           A function
                                          [m]
                                                 = (Return) type of a method call
[f]
           (Return) type of a function
                                                     The n'th actual argument
                                          (A_n)
                                                 =
           Static property fetch
                                          (P_n)
                                                      The n'th formal parameter
:: c
           Static method call
                                                      Type hint
:: m
                                           th
           Static property fetch
                                                      Default value
:: p
                                           v
                                            Γ
Mfs
           Modifiers
                                                      Whole program
```

Table 4.1: Constraint legend

A list of predefined items can be found here:

- Constants
- Variables
- todo: functions
- todo: classes

Expressions

Normal assignment

```
E \equiv (E_1 = E_2)
[E_2] <: [E_1],
[E_1] <: [E]
```

Listing 4.1: Normal assignment

Ternary

$$E \equiv (E_1 ? E_2 : E_3)$$

$$[E] <: [E_2] \lor [E] <: [E_3]$$

$$E \equiv (E_1 ? : E_3)$$

$$[E] <: [E_1] \lor [E] <: [E_3]$$

```
$\frac{1}{2} \texpr ? \texpr ? \texpr : \texpr ? \texpr / \texpr (E) is subtypeOf(\text{$b}) or subtypeOf(\text{$c})
$\texpr ? : \text{$c}; // \text{typeOf(E) is subtypeOf(\text{$expr}) or subtypeOf(\text{$c})}$
```

Listing 4.2: Ternary

Assignments with operators (1) always resulting in ints

```
E_1 \& = E_2
E_1 \mid = E_2
E_1 \hat{\ } = E_2
E_1 <<= E_2
E_1 >>= E_2
E_1 \% = E_2
[E_1] = integer()
```

Listing 4.3: Assignments with operators (1)

Assignments with operators (2) string concat (. =)

$$\begin{split} \underline{E_1 := E_2} \\ \overline{[E_1] = string(),} \\ if([E_2] <: object()) => hasMethod([E_2], "__tostring") \end{split}$$

Listing 4.4: Assignments with operators (2)

Assignments with operators (3) resulting in int where rhs is no array

$$E_1 /= E_2$$

$$E_1 -= E_2$$

$$[E_1] = integer(),$$

$$[E_2] \neq array(any())$$

Listing 4.5: Assignments with operators (3)

Assignments with operators (4) resulting in int or float

$$E_1 *= E_2 \ E_1 += E_2 \ \hline [E_1] <: float()$$

Listing 4.6: Assignments with operators (4)

Unary operators

$$E \equiv (+E_1) \lor (-E_1)$$

$$[E] <: float(),$$

$$[E_1] \neq array(\backslash any())$$

$$\frac{E \equiv (!E_1)}{[E] = boolean()}$$

$$\frac{E \equiv (\sim E_1)}{[E_1] = float() \lor [E_1] = integer() \lor [E_1] = string(),}$$

$$[E] = integer() \lor [E] = string()$$

$$\frac{E \equiv (E_1 + +) \lor (E_1 - -)}{if([E_1] <: array(\backslash any()) => [E] <: array(\backslash any()),}$$

$$if([E_1] = boolean()) => [E] = boolean(),$$

$$if([E_1] = float()) => [E] = float(),$$

$$if([E_1] = integer()) => [E] = integer(),$$

$$if([E_1] = null()) => [E] = integer() \lor [E] = null,$$

$$if([E_1] <: object()) => [E] <: object(),$$

$$if([E_1] = resource()) => [E] = resource(),$$

$$if([E_1] = string()) => [E] = integer() \lor [E] = string()$$

```
E \equiv (+ + E_1) \lor (- - E_1)
if([E_1] \lessdot array( \land any()) \Rightarrow [E] \lessdot array( \land any()),
if([E_1] = boolean()) \Rightarrow [E] = boolean(),
if([E_1] = float()) \Rightarrow [E] = float(),
if([E_1] = integer()) \Rightarrow [E] = integer(),
if([E_1] = null()) \Rightarrow [E] = null,
if([E_1] \lessdot object()) \Rightarrow [E] \lessdot object(),
if([E_1] = resource()) \Rightarrow [E] = resource(),
if([E_1] = string()) \Rightarrow [E] = integer() \lor [E] = float() \lor [E] = string()
```

```
1 +$a // positive
2 -$a // negation
3 !$a // not
4 ~$a // bitwise not
5 $a++ // post increase
6 $a-- // post decrease
7 ++$a // pre increase
8 --$a // pre decrease
```

Listing 4.7: Unary operators

Binary operators

```
E \equiv (E_1 + E_2)
E \equiv (E_1 + E_2)
E \equiv (e_1 + e_2) \lor [E] \lt : array(), \forall [E] \vdash (e_1) \lor [E] \lor (e_2) \lor [E] \lor (e_1) \lor [E] \lor (e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (E_1 - E_2) \lor (E_1 \lor e_2) \lor (E_1 \lor e_2)
E \equiv (E_1 - E_2) \lor (E_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1 \lor e_2) \lor (e_1 \lor e_2)
E \equiv (e_1
```

```
1  $a + $b // addition
2  $a - $b // subtraction
3  $a * $b // mulitiplication
4  $a / $b // division
5  $a % $b // modulus
6  $a & $b // bitwise And
7  $a | $b // bitwise Or
8  $a ^ $b // bitwise Shift left
10  $a >> $b // bitwise shift right
```

Listing 4.8: Binary operators

Comparison operators

```
E \equiv (E_1 == E_2)
E \equiv (E_1 == E_2)
E \equiv (E_1 != E_2)
E \equiv (E_1 <= E_2)
E \equiv (E_1 <= E_2)
E \equiv (E_1 < E_2)
E \equiv (E_1 >= E_2)
```

Listing 4.9: Comparison operators

Logical operators

```
E \equiv (E_1 \text{ and } E_2)
E \equiv (E_1 \text{ or } E_2)
E \equiv (E_1 \text{ xor } E_2)
E \equiv (E_1 \&\& E_2)
E \equiv (E_1 \parallel E_2)
E \equiv boolean()
```

Listing 4.10: Logical operators

Array

Array value fetch

$$\begin{split} \underline{E &\equiv E_1[E_2]} \\ \hline [E_1] &\neq object(), \\ if([E_1] &= string()) => [E] = string(), \\ if([E_1] &= array(\{types\})) => [E] <: \{types\}, \\ if([E_1] &\neq string() \land [E_1] \neq array(_)) => [E] = null() \end{split}$$

Listing 4.11: Array value fetch

Array declaration

```
E', where E' is an array declaration [E'] <: array(any())
```

```
1 array(/*...*/); // typeOf() = array();
2 // Rascal: array(_) => array(\any())
```

Listing 4.12: Array declaration

Scalars

Scalars

$$\frac{E, E \text{ is a string}}{[E] = string()}$$

$$E, E$$
 is a float $[E] = float()$

$$E, E$$
 is a integer $[E] = integer()$

```
1  "Str" // string()
2  'abc' // string()
3  100 // integer()
4  1.4 // float()
```

Listing 4.13: Scalars

Encapsulated strings

$$\frac{E, E \text{ is an encapsed string*}}{[E] = string()}$$

* When a string contains expression(/variables), it is processed as encapsed.

```
1 "$var"
```

Listing 4.14: Encapsulated strings

Casts

Casts

Note: PHP Warnings are ignored

$$E \equiv (array)E_1$$

$$E \equiv (bool)E_1 \lor (boolean)E_1$$

$$E \equiv (bool)E_1 \lor (double)E_1 \lor (real)E_1$$

$$E \equiv (float)E_1 \lor (double)E_1 \lor (real)E_1$$

$$E \equiv (int)E_1 \lor (integer)E_1$$

$$E \equiv (int)E_1 \lor (integer)E_1$$

$$E \equiv (object)E_1$$

$$E \equiv (object)E_1$$

$$E \equiv (string)E_1$$

$$E \equiv (anset)E_1$$

$$E \equiv (unset)E_1$$

$$E \equiv (unset)E_1$$

$$E \equiv (unset)E_1$$

```
(array)$a // <: array(\any())
(bool)$a // boolean()
(float)$a // float()
(int)$a // integer()
(object)$a // object()
(string)$a // string(), when $a == object() the object needs to have __toString()
(unset)$a // null()</pre>
```

Listing 4.15: Casts

Clone

Clone

$$E \equiv clone(E_1)$$

$$E[E] <: object(), [E_1] <: object()$$

```
clone($a) // typeOf($a) = object, typeOf(clone($a)) = object
```

Listing 4.16: Clone

Class

Class instantiation (1) matching the class name

$$E \equiv new C_1()$$
$$[E] = class(C.decl)$$

```
l new C;
```

Listing 4.17: Class instantiation (1)

Class instantiation (2) of an expression

$$\frac{E \equiv new \ E_1}{[E] <: object(),}$$
$$[E_1] <: object() \lor [E_1] = string()$$

```
1  $c = "C";
new $c;
```

Listing 4.18: Class instantiation (2)

Special keywords self parent parent static

```
E \equiv \$this \in C
[E] <: object()
[E] = class(C) \lor [E] :> class(C)
\frac{E \equiv self \in C}{[E] <: object()}
[self] = class(C)
\frac{E \equiv parent \in C}{[E] <: object()}
[E] :> class(C)
\frac{E \equiv static \in C}{[E] <: object()}
([E] <: class(C) \lor [E] :> class(C))
```

```
// $this can only be used within a class

this // in class C -> class(C)

self // in class C -> class(C)

parent // in class C -> parentOf(class(C))

table transfer

static // in class C -> class(C) or parentOf(class(C))
```

Listing 4.19: Special keywords

Class property fetch

```
 \begin{array}{c} \$this \to E_1 \subseteq C_1 \\ \hline [E_1] = C_1. \\ \text{hasProperty}(E_1. \\ \text{name, static} \notin \text{Mfs}) \lor \\ [E_1] = C_1. \\ \text{parent.hasProperty}(E_1. \\ \text{name, public} \\ \text{protected} \in \text{Mfs} \land \\ \text{static} \notin \text{Mfs}) \lor \\ [E_1] = C_1[. \\ \text{parent].hasMethod}("\_\_get") \\ \end{array}
```

^{*} Possible add fact that the field E is declared in class C, when it is on the left side of an assignment.

$$self:: E_1 \subseteq C_1$$

$$\overline{[E_1] = C_1.\text{hasProperty}(E_1.\text{name, static} \in \text{Mfs})}$$

$$\underline{parent:: E_1 \subseteq m}$$

$$\overline{[E_1] = \text{C.parent.hasProperty}(E_1.\text{name, static} \in \text{Mfs})}$$

$$\underline{E_1 \rightarrow E_2 \subseteq C_1*}$$

$$\overline{[E_1] = C_1.\text{hasProperty}(E_2.\text{name, static} \notin \text{Mfs}) \vee}$$

$$[E_1] = C_1.\text{parent.hasProperty}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \vee}$$

$$[E_1] = C.\text{hasProperty}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \notin \text{Mfs})$$

*The same goes for static property fetches, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E_1 \to E_2 \not\subseteq C \subseteq \Gamma^*$$

$$[E_1] = \text{C.hasProperty}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \notin \text{Mfs})$$

*Property fetch outside a class scope, also for static properties.

```
$\fints -> \text{prop} // \text{ name} = \text{prop}, \text{vis} = \text{public|protected}, !static || \text{mm}$

self::\fints \text{prop} // static property in class

parent::\fints \text{prop} // static property in the parent(s)}

4 \fints -> \text{prop} // \text{non-static property fetch}

5 \fints ::\fints \text{pro} // static property fetch}
$=:\fints \text{pro} // static property fetch}
```

Listing 4.20: Class property fetch

Class property fetch variable

$$E \equiv E_1 \rightarrow E_2, E_2 \text{ is an expression}$$

 $[E_1] <: \text{ object()}$

```
1  $b = "b";
2  $a->$b
```

Listing 4.21: Class property fetch variable

Class method call

$$E \equiv \$this \rightarrow E_1 \subseteq C_1$$

$$[\$this] <: \text{object}(), [\$this] = \text{class}(C) \lor [E] :> \text{class}(C),$$

$$[E_1] \text{ isMethod}(), [E_1] \text{ hasName}(E_1.name \lor "__\text{call"}),$$

$$[E_1] = C_1.\text{hasMethod}(E_1.\text{name, static} \notin \text{Mfs}) \lor \\ [E] <: [E_1]$$

$$\underline{E \equiv self :: E_1 \subseteq C_1}$$

$$\overline{[E_1] = C_1.\text{hasMethod}(E_1.\text{name, static} \in \text{Mfs}) \lor}$$

$$[E_1] = C_1.\text{hasMethod}(E_1.\text{name, public|protected} \in \text{Mfs} \land \text{static} \in \text{Mfs}) \lor \\ [E_1] = C_1.\text{hasMethod}("__\text{callStatic"})$$

$$\underline{E \equiv parent :: E_1 \subseteq C_1}$$

$$\overline{[E_1] = C_1.\text{parent.hasMethod}(E_1.\text{name, public|protected} \in \text{Mfs}) \lor }$$

$$[E_1] = C_1.\text{parent.hasMethod}("__\text{callStatic"})$$

$$\underline{E \equiv E_1 \rightarrow E_2 \subseteq C_1^*}$$

$$\overline{[E_1] = C_1.\text{hasMethod}(E_2.\text{name, static} \notin \text{Mfs}) \lor }$$

$$[E_1] = C_1.\text{parent.hasMethod}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor }$$

$$[E_1] = C.\text{hasMethod}(E_2.\text{name, public|protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor }$$

*The same goes for static method calls, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E \equiv E_1 \to E_2 \not\subseteq C \subseteq \Gamma^*$$

$$[E_1] = \text{C.hasMethod}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \not\in \text{Mfs})$$

$$\begin{array}{c} \$this \rightarrow E_1 \subseteq C_1 \\ \hline [E_1] = C_1. \text{hasMethod}(E_1. \text{name, static} \notin \text{Mfs}) \lor \\ [E_1] = C_1. \text{parent.hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C_1[. \text{parent].hasMethod}("__\text{call"}) \\ \hline \underbrace{self :: E_1 \subseteq C_1} \\ \hline [E_1] = C_1. \text{hasMethod}(E_1. \text{name, static} \in \text{Mfs}) \lor \\ [E_1] = C_1. \text{hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \in \text{Mfs}) \lor \\ \hline [E_1] = C_1. \text{hasMethod}("__\text{callStatic"}) \\ \hline \underbrace{parent :: E_1 \subseteq C_1} \\ \hline [E_1] = C_1. \text{parent.hasMethod}(E_1. \text{name, public}| \text{protected} \in \text{Mfs}) \lor \\ \hline [E_1] = C_1. \text{parent.hasMethod}("__\text{callStatic"}) \\ \hline \underbrace{E_1 \rightarrow E_2 \subseteq C_1^*} \\ \hline [E_1] = C_1. \text{hasMethod}(E_2. \text{name, static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C_1. \text{parent.hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{static} \notin \text{Mfs}) \lor \\ \hline [E_1] = C. \text{hasMethod}(E_2. \text{name, public}| \text{protected} \in \text{Mfs} \land \text{max})$$

*The same goes for static method calls, except for the 'static ∉ Mfs' part: 'static ∈ Mfs'.

$$E_1 \to E_2 \not\subseteq \Gamma \subseteq \Gamma^*$$

$$[E_1] = \text{C.hasMethod}(E_2.\text{name, public} \in \text{Mfs} \land \text{static} \not\in \text{Mfs})$$

*method call outside a class scope, also for static methods.

```
1 | $this->methodCall();
 self::methodCall();
 parent::methodCall();
 $a->methodCall();
 $a::methodCall();
```

Listing 4.22: Class method call

Class method call variable

$$E \equiv E_1 \to E_2(), E_2 \text{ is an expression}$$
$$[E_1] <: \text{object()}$$

\$a->\$methodCall()

Listing 4.23: Class method call variable

Class constants (needs to be reviewed)

$$\frac{\text{self}::c_1 \subseteq \Gamma}{[\text{self}::c_1] = C_1.\text{hasConstant}(E_2.\text{name}) \vee}$$

$$[\text{self}::c_1] = C_1.\text{parent.hasConstant}(E_2.\text{name}, \text{public}|\text{protected} \in \text{Mfs})$$

$$\frac{\text{parent}::c_1 \subseteq \Gamma}{[\text{self}::c_1] = C_1.\text{parent.hasConstant}(E_2.\text{name}, \text{public}|\text{protected} \in \text{Mfs})}$$

$$\frac{E_1::c_1 \subseteq \Gamma}{[E_1] = \text{object}()}$$

```
self::CONST
parent::CONST
SOMECLASS:: CONST
```

Listing 4.24: Class constants (needs to be reviewed)

Parameters

Parameters in class instantiation

*These parameters are just examples for what happens if they have typeHints (th), default values(v) or none *The constructor can be found in the m3 model (@constructors(loc classDecl, loc constructorMethodDecl))

```
new \ C_1 \ (A_1, A_2, \dots, A_k) \subseteq \Gamma
\$a \to m() \ (A_1, A_2, \dots, A_k) \subseteq \Gamma
function_1 \ (A_1, A_2, \dots, A_k) \subseteq \Gamma
function_1 \ (A_1, A_2, \dots, A_k) \subseteq \Gamma
function_1 \ (A_1, A_2, \dots, A_k) \subseteq \Gamma
function \ (th_1 \ P_1, P_2 = v, \dots, P_k) \subseteq \Gamma
function \ (th_1 \ P_1, P_2 = v, \dots, P_k) \subseteq \Gamma
[P_1] <: [A_1], [A_1] <: [th_1], [P_1] <: [th_1], \text{ hasRequiredParam}(P_1), \text{ hasRequiredParam}(P_k)
```

```
new C($foo);
```

Listing 4.25: Parameters in class instantiation

Scope

Type of a certain variable within some scope this applies to global- class- function- and method- scope

$$\frac{E, E', E'', E''' \dots \ etc \subseteq f \qquad E \text{ is a variable}}{[E] = [E] \vee [E'] \vee [E''] \vee [E'''] \dots \ etc}$$

```
function f() {
    $a = 1;
    $a = "true";
}
// typeOf($a) is typeOf($a1, $a2, ..., $an);
```

Listing 4.26: Type of a certain variable within some scope

Return type of function or method (1) having no return statements or return;

$$return \not\subseteq f \lor return; \subseteq f$$
$$[f] = null()$$

```
function f() {} // no return = null()
function f() { return; } // return; = null()
```

Listing 4.27: Return type of function or method (1)

Return type of function or method (2) every exit path ends with a return statement

```
(\text{return } E_1) \lor (\text{return } E_2) \lor \cdots \lor (\text{return } E_k) \subseteq f[f] <: [E_1] \lor [E_2] \lor \cdots \lor [E_k]
```

```
function f() {
   if (rand(0,1))
   return $a;
else
   return $b;
}
// returns typeOf($a) or typeOf($b)
```

Listing 4.28: Return type of function or method (2)

Return type of function or method (3) possible no return value

$$(\text{return } E_1) \vee (\text{return } E_2) \vee \cdots \vee (\text{return } E_k) \vee (\neg \text{ return}) \subseteq f$$
$$[f] <: [E_1] \vee [E_2] \vee \cdots \vee [E_k] \vee null()$$

```
function f() {
   if (rand(0,1))
    return $a;
   else if (rand(0,1))
    return $b;
}
// returns typeOf($a) or typeOf($b) or null()
```

Listing 4.29: Return type of function or method (3)

Function calls

Function call

```
\frac{f() \subseteq \Gamma}{[f()] <: \text{return of } [f]}
```

```
1 function f() {}
2 f();
```

Listing 4.30: Function call

Function call variable

```
\begin{split} \underline{E \equiv E_1() \subseteq \Gamma} \\ \underline{[E] = any(),} \\ [E_1] <: object() \vee [E_1] = string(), \\ \text{i} f([E_1] <: object()) => hasMethod("\__invoke") \end{split}
```

```
function f() {}

f
```

Listing 4.31: Function call variable

How to resolve expressions:

- Find all expressions which are defined above and annotate them with @type.
- Annotate the rest of the expressions with @type = any(); (should only be for relevant expressions)

4.4 Annotations

After we gathers all the facts from the source code, we will add additional information which we read from the annotations. For this we use regex to match @return type and @param type var for methods and functions. We read @var type for variables and class attributes. In our first analysis we do not include the facts we gathered from the annotations. In the second analysis we do include the facts. This way we can compare the end results.

In order to gain some knowledge about the reliability of the annotations, we compare the result our or initial analysis with the provided provided annotation information. Here the implementation should comply to the used annotations.

4.5 Constraint solving

When we have all the constraints from the source code as facts, we will solve the constraints until we can no longer solve any constraints. The result will be a list of possible types for each class, method, fields, functions, variable and expression. The first step is to initialise all type-able objects we want to solve. In this initial phase the status of all objects can be of any type, because we do not know anything about them yet. When we solve the constraints step by step, we will be able to limit the number of possible types for a certain object. We do this by taking the intersection of the constraint result and the possible types we have. This way there should be less and less possible types for each variable.

When we take the intersection of none overlapping types, we have a type error. There is no possible type for this object. Because PHP allows ducktyping, allowing objects to be of more than one type, we will use widening in these cases. Widening means that we pick the union instead of the intersection, so we can continue our analysis.

Some of the constraints only give us information that a certain object should be a subtype of a certain type. When we solve these constraints, we will take the least common ancestor.

Analysis

In order to validate the performed research we have tested them on the most popular packages of Packagist¹, which are listed in table 5.1. The statistics are generated using phploc². All packages have between 2 and 6 million downloads. Packagist is a repository for Composer³ projects. Composer is a dependency manager for PHP projects. All external plugins for PHP projects can be managed via a composer json file. You only need know the name and a version of the external package in the repository of your project.

Product				iles	C	bjects		Lines of code				
Vendor	Project*	Version	$\overline{\mathrm{D}^1}$	F^2	C^3	I^4	T^5	Total ↑	Logical		Global ⁶	
doctrine	lexer	v1.0	2	7	3	0	0	733	128	(17.46%)	13	(10.16%)
phpunit	php-timer	1.0.5	5	11	5	0	0	740	117	(15.81%)	17	(14.53%)
phpunit	php-text-template	1.2.2	5	11	5	0	0	768	125	(16.28%)	15	(12.00%)
doctrine	inflector	v1.0	2	7	3	0	0	853	130	(15.24%)	13	(10.00%)
psr-fig	log	1.0.0	3	15	8	2	2	1039	155	(14.92%)	22	(14.19%)
phpunit	php-file-iterator	1.3.4	5	13	7	0	0	1071	176	(16.43%)	15	(8.52%)
symfony	filesystem	v2.5.3	3	11	5	2	0	1 090	193	(17.71%)	19	(9.84%)
symfony	yaml	v2.5.3	3	16	11	1	0	2270	509	(22.42%)	28	(5.50%)
phpunit	php-token-stream	1.2.2	6	13	169	0	0	2360	377	(15.97%)	15	(3.98%)
doctrine	collections	v1.2	3	18	11	3	0	2504	394	(15.73%)	33	(8.38%)
symfony	process	v2.5.3	3	19	14	1	0	3198	604	(18.89%)	37	(6.13%)
symfony	finder	v2.5.3	8	43	36	3	0	4976	909	(18.27%)	80	(8.80%)
symfony	dom-crawler	v2.5.3	12	63	53	6	0	7825	1296	(16.56%)	157	(12.11%)
symfony	translation	v2.5.3	21	121	97	20	0	12345	2299	(18.62%)	257	(11.18%)
symfony	console	v2.5.3	17	84	66	13	2	13546	2556	(18.87%)	246	(9.62%)
symfony	http-foundation	v2.5.3	16	90	76	10	0	14179	2262	(15.95%)	154	(6.81%)
twig	twig	v1.16.0	18	172	148	19	0	14689	2630	(17.90%)	15	(0.57%)
symfony	event-dispatcher	v2.5.3	27	170	133	31	3	20230	3629	(17.94%)	418	(11.52%)
swiftmailer	swiftmailer	v5.2.1	37	238	170	52	0	28965	4645	(16.04%)	144	(3.10%)
phpunit	php-code-coverage	2.0.1	62	259	381	24	0	50371	6579	(13.06%)	87	(1.32%)
phpunit	phpunit	4.2.2	65	270	388	26	0	51516	6764	(13.13%)	129	(1.91%)
phpunit	phpunit-mock-objects	2.2.0	66	271	393	27	0	51735	6 801	(13.15%)	132	(1.94%)
doctrine	annotations	v1.2.0	69	306	423	28	0	57325	7718	(13.46%)	188	(2.44%)
doctrine	common	v2.4.2	76	337	440	45	0	62406	8 3 2 6	(13.34%)	298	(3.58%)
symfony	http-kernel	v2.5.3	96	565	471	90	3	79294	14169	(17.87%)	1449	(10.23%)
doctrine	cache	1.3.0	152	687	729	102	2	103024	16667	(16.18%)	1355	(8.13%
doctrine	dbal	v2.4.2	121	557	628	63	0	104 630	15234	(14.56%)	1 033	(6.78%)
guzzle	guzzle	v3.9.2	150	832	828	141	7	117699	19772	(16.80%)	1787	(9.04%)
doctrine	orm	v2.4.4	175	1007	875	119	2	158530	27932	(17.62%)	2866	(10.26%
monolog	monolog	1.10.0	350	1911	1904	135	2	288507	31415	(10.89%)	4221	(13.44%)
werkspot	old-Website	07-2014	928	6225	4907	224	0	1054686	167978	(15.93%)	22693	(13.51%)

^{*}This is a list of the 30 most popular packages of packagist ordered by total lines of code, in July 2014. 1 = Directories, 2 = Files, 3 = Classes, 4 = Interfaces, 5 = Traits, 6 = Not in class or function

Table 5.1: List of analysed projects.

To collect the source code for each project, we have executed the following steps:

1. git clone the github repo.

¹https://packagist.org/explore/popular, July 2014

²https://github.com/sebastianbergmann/phploc, July 2014

 $^{^3}$ https://getcomposer.org/, July 2014

- 2. Run composer install, and the source code including dependencies will be downloaded in the /vendor folder.
- 3. Remove the autoload.php and composer folder, as we don't need them.
- 4. Remove the test folders by removing all folders matching Tests or tests.

To measure the coverage:

- 1. git clone the github repo.
- 2. Run composer install, and the source code including dependencies will be downloaded in the /vendor folder.
- 3. Run the unittests and use xdebug to resolve the types.
- 4. Compare the results.

Results

For the results we picked X software products to see how it performs. For each product we performed the type inference with and without reading annotations from the doc blocks.

6.1 Results

Show the results...

6.2 Validation of the results (or something)

Say something about:

- Soundness (what we measured, is it correct?)
- Completeness (how much did we measure?)
- Accuracy (how precise are the results?)

6.3 Annotations

The results of the analysis when adding the annotations to the analysis. Compare the results with the results of the analysis without the annotation information.

Case Study

Explain how the case study is performed.

This chapter will show the case study, but I just need to place this information somewhere. A list of the 40 most popular packages from packages.

- ullet create composer file
- composer install
- mkdir phploc
- \bullet list all packages: find ./vendor/* -maxdepth 1 -mindepth 1 -type d -exec ls -d ""
- \bullet prefix the list with "phploc" and postfix with " > phploc/<file>.phploc"

Conclusion

Summary of the whole work, with conclusions. T.B.A.

8.1 Conclusion

8.2 Future work

These items will not be covered by the analysis (maybe add this to threats/future work)

- Analysis is flow insensitive
- \bullet Closure
- References
- Variable constructs (variable -variable, -method/function calls, -class instantiation, eval) :: todo: explain WHY not.
- Yields

Explain something about combining this analysis to other analysis (like dead code elimination, constant folding/propagation resolve, alias analysis, array analysis) to gain more precise results.

Something about performance optimisations... Explain what is already done to boost the performance and what still can be done.

Use a bigger corpus to gains better results of the analysis by doing analysis on more programs.

Glossary

AST

An abstract representation of the structure of the source code. .

Rascal

Rascal is a meta-programming language developed by SWAT (Software analyse and transformation) team at CWI in the Netherlands. See http://www.rascal-mpl.org/ for more information.

reflexive transitive closure

A relation is transitive if $\langle a, b \rangle \in R$ then $\langle b, a \rangle \in R$.

A relation is reflexive if $\langle a,b\rangle \in R$ and $\langle b,c\rangle \in R$ then $\langle a,c\rangle \in R$.

A reflexive transitive closure can be established by creating direct paths for all indirect paths and adding self references, until a fixed point is reached.

$\mathbf{stdClass}$

A predefined class in the PHP library. The class is the root of the class hierarchy. It is comparable to the Object class in Java.

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