# C335 Computer Structures

## Number Representation Part #2

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Adapted from Morgan Kaufmann and others

## **How to Represent Negative Numbers?**

- □ So far, <u>un</u>signed numbers
- Obvious solution: define leftmost bit to be sign!
  - 0 ⇒ +, 1 ⇒ -
  - Rest of bits can be numerical value of number
- Representation called <u>sign and magnitude</u>
- □ MIPS uses 32-bit integers. +1<sub>ten</sub> would be:
  - <u>0</u>000 0000 0000 0000 0000 0000 0001
- □ And −1<sub>ten</sub> in sign and magnitude would be:
  - 1000 0000 0000 0000 0000 0000 0000 0001

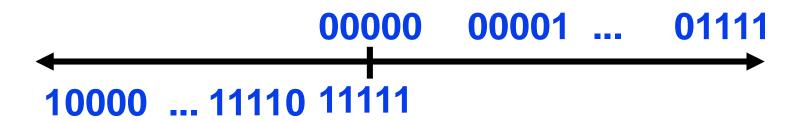
## **Shortcomings of sign and magnitude?**



- Arithmetic circuit complicated
  - Special steps depending whether signs are the same or not
  - Extra step to set the sign
- □ Also, two zeros
  - $0x00000000 = +0_{ten}$
  - $0x80000000 = -0_{ten}$
  - What would two 0s mean for programming?
- Therefore sign and magnitude abandoned

#### Another try: complement the bits

- □ Example:  $7_{10} = 00111_2$   $-7_{10} = 11000_2$
- Called One's Complement
- Note: positive numbers have leading 0s, negative numbers have leadings 1s.



- What is -00000 ?
  - Answer: 11111
- How many positive numbers in N bits?
- How many negative numbers?

#### **Shortcomings of One's complement?**



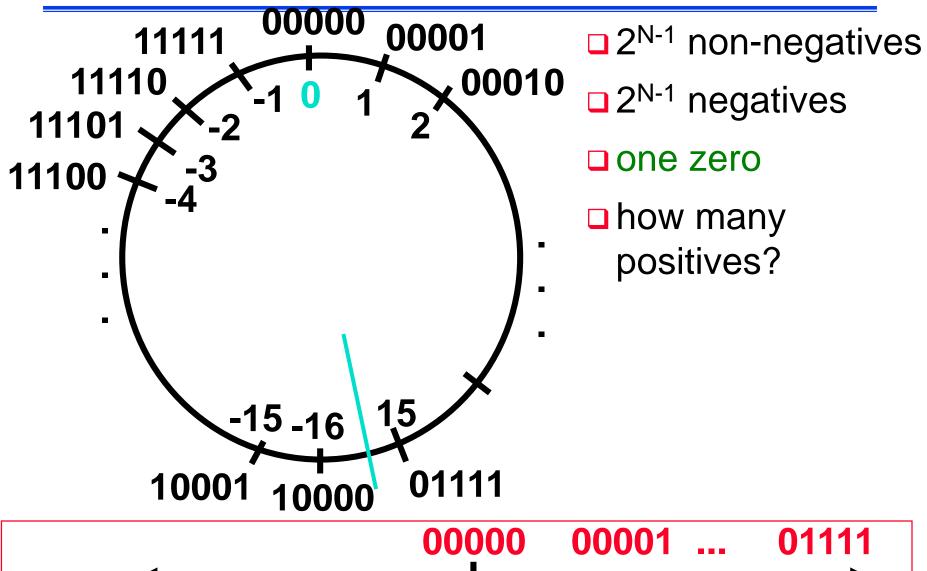
- Arithmetic still somewhat complicated.
- □ Still two zeros
  - $0 \times 0 0 0 0 0 0 0 0 = +0_{ten}$
- Although used for awhile on some computer products, one's complement was eventually abandoned because another solution was better.

#### Standard Negative Number Representation



- ■What is result for unsigned numbers if trie to subtract large number from a small one?
  - Would try to borrow from string of leading 0s, so result would have a string of leading 1s
    - $3-5 \Rightarrow 00...0011 00...0101 = 11...1110$
  - With no obvious better alternative, pick representation that made the hardware simple
  - As with 1's complement, leading 0s ⇒ positive, leading 1s ⇒ negative
    - 000000...xxx is  $\ge 0$ , 1111111...xxx is < 0
    - except 1...1111 is -1, not -0 (as in 1's complement.)
- ■This representation is <u>Two's Complement</u>

#### 2's Complement Number "line": N = 5



10000 11110 11111 Liqiang Zhang Indiana University South Bend

#### Two's Complement for N=32

0000 0000 0000 0000	$0000_{two} =$	$O_ten$
0000 0000 0000 0000	$0001_{two} =$	1 <sub>ten</sub>
0000 0000 0000 0000	$0010_{two} =$	2 <sub>ten</sub>
0111 1111 1111 1111	$1101_{two} =$	2,147,483,645 <sub>ten</sub>
0111 1111 1111 1111	$1110_{two} =$	2,147,483,646 <sub>ten</sub>
0111 1111 1111 1111	1111 <sub>two</sub> =	2,147,483,647 <sub>ten</sub>
1000 0000 0000 0000	$0000_{two} =$	-2,147,483,648 <sub>ten</sub>
1000 0000 0000 0000	$0001_{two} =$	-2,147,483,647 <sub>ten</sub>
1000 0000 0000 0000	$0010_{two} =$	-2,147,483,646 <sub>ten</sub>
1111 1111 1111 1111	$1101_{two} =$	-3 <sub>ten</sub>
1111 1111 1111 1111	$1110_{two} =$	-2 <sub>ten</sub>
1111 1111 1111 1111	1111 <sub>two</sub> =	-1 <sub>ton</sub>

- □One zero; 1st bit called sign bit
- □1 "extra" negative:no positive 2,147,483,648<sub>ten</sub>

### **Two's Complement Formula**

□ Can represent positive and negative numbers in terms of the bit value times a power of 2:

$$d_{31} \times (2^{31}) + d_{30} \times 2^{30} + ... + d_2 \times 2^2 + d_1 \times 2^1 + d_0 \times 2^0$$

□ Example: 1101<sub>two</sub>

$$= 1x-(2^3) + 1x2^2 + 0x2^1 + 1x2^0$$

$$= -2^3 + 2^2 + 0 + 2^0$$

$$= -8 + 4 + 0 + 1$$

$$= -8 + 5$$

$$=-3_{ten}$$

#### **Two's Complement shortcut: Negation**



- Change every 0 to 1 and 1 to 0 (invert or complement), then add 1 to the result
- Proof\*: Sum of number and its (one's) complement must be 111...111<sub>two</sub>

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However, 111...111_{two} = -1_{ten}
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Let  $x' \Rightarrow$  one's complement representation of x

Then 
$$x + x' = -1 \Rightarrow x + x' + 1 = 0 \Rightarrow -x = x' + 1$$

■ Example: -3 to +3 to -3



□ Convert 0b11100111 to decimal (assume 2's complement is used)

## Two's comp. shortcut: Sign extension



- Convert 2's complement number rep. using n bits to more than n bits
- Simply replicate the most significant bit (sign bit) of smaller to fill new bits
  - 2's comp. positive number has infinite 0s
  - 2's comp. negative number has infinite 1s
  - Binary representation hides leading bits;
     sign extension restores some of them
  - 16-bit -4<sub>ten</sub> to 32-bit:

1111 1111 1111 1100<sub>two</sub>

1111 1111 1111 1111 1111 1111 1111 1100<sub>two</sub>

### What if too big?



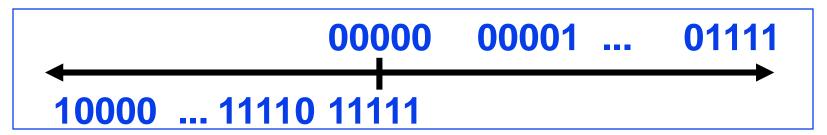
- Binary bit patterns above are simply representatives of numbers. Strictly speaking they are called "numerals".
- Numbers really have an ∞ number of digits
  - with almost all being same (00...0 or 11...1) except for a few of the rightmost digits
  - Just don't normally show leading digits
- □ If result of add (or -, \*, /) cannot be represented by these rightmost bits, <u>overflow</u> is said to have occurred.

00000 00001 00010 11110 11111 unsigned

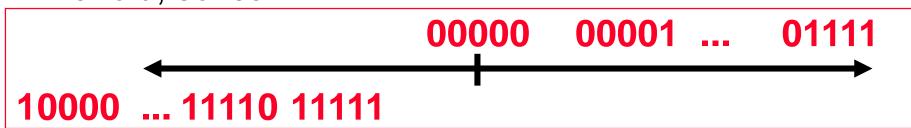
#### Number summary...



- We represent "things" in computers as particular bit patterns: N bits  $\Rightarrow$  2<sup>N</sup>
- Decimal for human calculations, binary for computers, hex to write binary more easily
- 1's complement mostly abandoned



2's complement universal in computing: cannot avoid, so learn



Overflow: numbers ∞; computers finite,errors!



 $Y = 0011 1011 1001 1010 1000 1010 0000 0000_{two}$ 

- A. X > Y (if signed)
- B. X > Y (if unsigned)
- c. To identify 30 registers, you need at least 5 bits

ABC

1: FFF

2: **FF**T

3: FTF

4: FTT

5: TFF

6: TFT

7: TTF

8: TTT

□ Convert -98₁₀ to binary (8 bits), then to hexadecimal (assume 1's complement is used)

□ Convert -98₁₀ to binary (8 bits), then to hexadecimal (assume 2's complement is used)

Convert 0xFFB5 to binary then to Decimal (assume 1's complement is used)

Convert 0xFFB5 to binary then to Decimal (assume 2's complement is used)