CryptoHomework5

February 25, 2020

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     CS 764 Blockchains and Cryptocurrencies
     Module 5, Homework 5
     February 18, 2020
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# Q1. Blind Digital Signatures #
     # (i) Given p=5, q=11 determine public key (e,n) and private key (d,n) Where e=7
     p = 5
     q = 11
     e = 7# Public Key (e,n)
     n = p*q# base (55)
     Phi_n = (p-1)*(q-1)
     print('n = %d' %n)
     print('Phi(n) = %d' % Phi_n)
     # Note: d == e^{-1} (mod(Phi_n)), or e*d = 1 mod Phi_n, 7*d = 1 mod 40
     # Define a helper function for obtaining modular inverse
     # Refference: egcd() and modinv() were obtained from following site
     # https://stackoverflow.com/questions/4798654/
      \rightarrow modular-multiplicative-inverse-function-in-python
     def egcd(a, b):
         if a == 0:
             return (b, 0, 1)
         else:
             g, y, x = egcd(b \% a, a)
             return (g, x - (b // a) * y, y)
     def modinv(a, m):
         g, x, y = egcd(a, m)
         if g != 1:
             raise Exception('modular inverse does not exist')
         else:
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return x % m
     # Determine the Value of the Private Key
     d = modinv(e, Phi_n)
     print('d = %d' % d)
     n = 55
     Phi(n) = 40
     d = 23
[40]: | # (ii) Given: b=6 determine blinding factor using derived public key
     b = 6 # randomly chosen number used to blind the message
     BF = pow(b, e) \% n
     print('Blinding Factor = %d' % BF)
     Blinding Factor = 41
[41]: # (iii) Given: m=11, determine blinded message
     m = 11# Plain Text Message
     BM = m*BF % n# Blinded Message
     print('Blinded Message = %d' % BM)
     Blinded Message = 11
[42]: | # (iv) Determine signature on blinded message using private key
     BM_signed = pow(BM, d) % n# Bob's signature on blinded message
     print('Signed Blinded Message = %d' % BM_signed)
     Signed Blinded Message = 11
[43]: # (v) Unblind the signature on the blinded message to obtain signature on
     SignedMessage = BM_signed*modinv(b, n) % n# Original Message signed by Bob
     print('Signed Message = %d' % SignedMessage)
     Signed Message = 11
[44]: \# (vi) Verify blind signature matches signature on m using (d,n)
     VerifiedSignature = pow(m, d) % n
     print('Verified Signature = %d' % VerifiedSignature)
     Verified Signature = 11
# Q2, Bit-commitment protocol #
     ####################################
     import hashlib
      # Secret Nonces for each person (Alice, Bob, Carol, David)
     NA = 2
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ND = 5
     # Secret Predicted Winner John:0, Jane:1
     PWA = 0
     PWB = 0
     PWC = 1
     PWD = 1
     # Public Hashes
     HashA = hashlib.md5()
     HashA.update((str(NA) + str(PWA)).encode('utf-8'))
     HashB = hashlib.md5()
     HashB.update((str(NB) + str(PWB)).encode('utf-8'))
     HashC = hashlib.md5()
     HashC.update((str(NC) + str(PWC)).encode('utf-8'))
     HashD = hashlib.md5()
     HashD.update((str(ND) + str(PWD)).encode('utf-8'))
     # Review Hashes
     print('HashA: ', HashA.digest())
     print('HashB: ', HashB.digest())
     print('HashC: ', HashC.digest())
     print('HashD: ', HashD.digest())
     HashA: b'\x98\xf17\x08!\x01\x94\xc4uh{\xe6\x10j;\x84'}
     HashB: b'|\xb\xc4\t\xec\x99\x0f\x19\xc7\x8cu\xbd\x1e\x06\xf2\x15'
     HashC: b'\xc1jS \xfaGU0\xd9X<4\xfd5n\xf5'
     HashD: b''(8\x02:w\x8d\xfa\xec\xdc!'\x08\xf7!\xb7\x88"
 []:
# Q3. Zero-Knoledge proofs #
     # Remote client C proves to server S it knows password
     # wihtout actually sending it across internet
     # (i) C Sends username to S
     username = 'Alice'
     \# (ii) S Verifies username and sends nonce N to C
     if(username == 'Alice'):
         Nonce = 2357
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NB = 7NC = 3

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# (iii) C Computes hash H=MD5(N||P||N) and sends H to S
HashZero = hashlib.md5()
HashZero.update((str(Nonce) + "SecretPassword123" + str(Nonce)).encode())
HashZero = HashZero.digest()
print('HashZero: ', HashZero)
```

HashZero: $b'\x15$; $C\x10~imA\xd1S\x1a\xfd$; $\x7f\x95'$

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[34]: # (iv) S independently computes G = MD5(N//P//N) and compares
HashG = hashlib.md5()
HashG.update((str(Nonce) + "SecretPassword123" + str(Nonce)).encode())
HashG = HashG.digest()
print('HashG: ', HashG)
```

HashG: $b'\x15$; $C\x10\sim imA\xd1S\x1a\xfd$; $\x7f\x95'$

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[35]: if HashZero == HashG:
    print("Verified!")
else:
    print("Invalid!")
```

Verified!

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# Q4. Discrete Zero Knoledge Proofs #
    # Given: Alice and Bob know public values p=17 and A=5
    # Alice also knows x=10
    # Show: How Alice can prrove she knows x w/o revealing it to Bob
    # Public Variables
    p = 17
    A = 5
    # Alice's Private Variables
    x = 10
    # 1.) A, B, and p are public
    # 2.) Alice Computes A^x mod p
    # Note: A \hat{x} = B \mod p
    B = pow(A, x) \% p
    print('B = \%d' \% B)
    # Helper Function
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from random import randint
# 3.) Alice chooses random number r < p and sends q = A^r \mod p to Bob
r = randint(0, p-1)
print('r = %d' %r)
q = pow(A, r) \% p
print('q = %d' % q)
# 4.) Bob sends random bit i
i = randint(0, 1)
print('i = %d' % i)
# 5.) Alice Computes s=(r+i*x) \mod p and sends s to Bob
s = (r + i*x) \% (p - 1)
print('s = %d' % s)
# 6.) Bob computes C=A \hat{}s mod p, which chould be equal to D=q*B \hat{}i
# If it doesnt then he knows Alice doesn't know x
C = pow(A, s) \% p
print('C = %d' % C)
D = q*pow(B, i) \% p
print('D = %d' % D)
if(C == D):
    print('Alice knows the value of x')
else:
    print('Alice does NOT know the value of x')
B = 9
r = 15
q = 7
i = 1
s = 9
C = 12
D = 12
Alice knows the value of x
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