| Step | Algorithm: $C := AB + C$ |
|------|--|
| 1a | $\{C = \widehat{C}\}$ |
| 4 | $A \to \left(A_L \middle A_R \right), B \to \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows |
| 2 | $\left \left\{ C = A_L B_T + \widehat{C} \right. \right.$ |
| 3 | while $n(A_L) < n(A)$ do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ |
| 5a | Determine block size b $ \left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right), \left(\begin{array}{c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right) $ |
| 6 | where A_1 has b columns, B_1 has b rows $ \left\{ C = A_0 B_0 + \widehat{C} \right\} $ |
| | |
| 8 | $C := A_1 B_1 + C$ |
| 7 | $\left\{ C = A_0 B_0 + A_1 B_1 + \widehat{C} \right\}$ |
| 5b | $A \to \left(A_L \middle A_R \right) \leftarrow \left(A_0 \middle A_1 \middle A_2 \right), \left(\frac{B_T}{B_B} \right) \leftarrow \left(\frac{B_0}{B_1} \right)$ |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg (n(A_L) < n(A)) \right\}$ |
| 1b | $\left\{ C = AB + \widehat{C} \right\}$ |

| Step | Algorithm: $C := AB + C$ |
|------|--------------------------|
| 1a | { |
| 4 | where |
| 2 | { |
| 3 | while do |
| 2,3 | \ \ \ |
| | Determine block size b |
| 5a | wh one |
| 6 | where { |
| 8 | |
| 7 | { |
| 5b | |
| 2 | { |
| | endwhile |
| 2,3 | { |
| 1b | \ |

| Step | Algorithm: $C := AB + C$ |
|------|---|
| 1a | ${C = \widehat{C}}$ |
| 4 | where |
| 2 | { |
| 3 | while do |
| 2,3 | \ \ \ |
| | Determine block size b |
| 5a | |
| 6 | where { |
| 8 | |
| 7 | { |
| 5b | |
| 2 | { |
| | endwhile |
| 2,3 | { |
| 1b | $\left\{ C = AB + \widehat{C} \right\}$ |

| Step | Algorithm: $C := AB + C$ |
|------|---|
| 1a | ${C = \widehat{C}}$ |
| 4 | where |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| 3 | while do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \right.$ |
| | Determine block size b |
| 5a | |
| 6 | where { |
| 8 | |
| 7 | <i>\</i> |
| • | |
| 5b | |
| 2 | $\left\{ \qquad C = A_L B_T + \widehat{C} $ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg () \right\}$ |
| 1b | $ \begin{cases} C = A_L B_T + \widehat{C} \land \neg() \\ C = AB + \widehat{C} \end{cases} $ |

| Step | Algorithm: $C := AB + C$ |
|------|--|
| 1a | ${C = \widehat{C}}$ |
| | |
| 4 | |
| | where |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| 3 | while $n(A_L) < n(A)$ do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ |
| | Determine block size b |
| | |
| 5a | |
| | 1 |
| 6 | where { |
| 8 | |
| 7 | { } |
| 1 | |
| F1 | |
| 5b | |
| | |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg (n(A_L) < n(A)) \right\}$ |
| 1b | $\left\{ C = AB + \widehat{C} \right\}$ |

| Step | Algorithm: $C := AB + C$ |
|------|---|
| 1a | $\{C = \widehat{C}$ |
| 4 | $A \to \left(A_L \middle A_R \right), B \to \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| 3 | while $n(A_L) < n(A)$ do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ |
| | Determine block size b |
| 5a | |
| | where |
| 6 | { |
| 8 | |
| 7 | { |
| | |
| 5b | |
| | |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg (n(A_L) < n(A)) \right\}$ |
| 1b | $\left \left\{ C = AB + \widehat{C} \right) \right $ |

| Step | Algorithm: $C := AB + C$ |
|------|--|
| 1a | $\{C = \widehat{C}\}$ |
| 4 | $A \to \left(A_L \middle A_R \right), B \to \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| 3 | while $n(A_L) < n(A)$ do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ |
| 5a | Determine block size b $ \left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right), \left(\begin{array}{c c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right) $ |
| _ | where A_1 has b columns, B_1 has b rows |
| 6 | { |
| 8 | |
| 7 | \ |
| 5b | $A \to \left(A_L \middle A_R \right) \leftarrow \left(A_0 \middle A_1 \middle A_2 \right), \left(\frac{B_T}{B_B} \right) \leftarrow \left(\frac{B_0}{B_1} \right)$ |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg (n(A_L) < n(A)) \right\}$ |
| 1b | $\left\{ C = AB + \widehat{C} \right\}$ |

| Step | Algorithm: $C := AB + C$ |
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| 1a | $\{C = \widehat{C}\}$ |
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| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| 3 | while $n(A_L) < n(A)$ do |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ |
| 5a | Determine block size b $ \left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right), \left(\begin{array}{c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right) $ |
| C | where A_1 has b columns, B_1 has b rows |
| 6 | $\left\{ C = A_0 B_0 + \widehat{C} \right\}$ |
| 8 | |
| 7 | \ |
| 5b | $A \to \left(A_L \middle A_R \right) \leftarrow \left(A_0 \middle A_1 \middle A_2 \right), \left(\frac{B_T}{B_B} \right) \leftarrow \left(\frac{B_0}{B_1} \right)$ |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right\}$ |
| | endwhile |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge \neg (n(A_L) < n(A)) \right\}$ |
| 1b | $\left\{ C = AB + \widehat{C} \right\}$ |

| Step | Algorithm: $C := AB + C$ | |
|------|--|---|
| 1a | $\{C=\widehat{C}$ | } |
| 4 | $A \to \left(A_L \middle A_R \right), B \to \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows | |
| 2 | $\left\{C = A_L B_T + \widehat{C}\right\}$ | } |
| 3 | while $n(A_L) < n(A)$ do | |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ | |
| 5a | Determine block size b $ \left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right), \left(\begin{array}{c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right) $ | |
| | where A_1 has b columns, B_1 has b rows | |
| 6 | $\left\{ C = A_0 B_0 + \widehat{C} \right.$ | } |
| 8 | | |
| 7 | $\left\{ C = A_0 B_0 + A_1 B_1 + \widehat{C} \right.$ | } |
| 5b | $A \to (A_L \mid A_R) \leftarrow (A_0 \mid A_1 \mid A_2), (\frac{B_T}{B_B}) \leftarrow (\frac{B_0}{B_1})$ | |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right.$ | } |
| | endwhile | |
| 2,3 | $\{C = A_L B_T + \widehat{C} \land \neg (n(A_L) < n(A))\}$ | } |
| 1b | $\left\{C = AB + \widehat{C}\right\}$ | } |

| Step | Algorithm: $C := AB + C$ | |
|------|--|---------------|
| 1a | $\{C = \widehat{C}$ | } |
| 4 | $A \to \left(A_L \middle A_R \right), B \to \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows | |
| 2 | $\left\{C = A_L B_T + \widehat{C}\right\}$ | } |
| 3 | while $n(A_L) < n(A)$ do | |
| 2,3 | $\left\{ C = A_L B_T + \widehat{C} \wedge n(A_L) < n(A) \right\}$ | |
| 5a | Determine block size b $ \left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right), \left(\begin{array}{c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right) $ | |
| | where A_1 has b columns, B_1 has b rows | $\frac{1}{2}$ |
| 6 | $\left\{ C = A_0 B_0 + \widehat{C} \right.$ | } |
| 8 | $C := A_1 B_1 + C$ | |
| 7 | $\left\{ C = A_0 B_0 + A_1 B_1 + \widehat{C} \right\}$ | } |
| 5b | $A \to (A_L \mid A_R) \leftarrow (A_0 \mid A_1 \mid A_2), (\frac{B_T}{B_B}) \leftarrow (\frac{B_0}{B_1})$ | |
| 2 | $\left\{ C = A_L B_T + \widehat{C} \right.$ | } |
| | endwhile | |
| 2,3 | $\{C = A_L B_T + \widehat{C} \land \neg (n(A_L) < n(A))\}$ | } |
| 1b | $\left\{C = AB + \widehat{C}\right\}$ | } |

| Algorithm: $C := AB + C$ |
|--|
| |
| $A 	o \left(A_L \middle A_R \right), B 	o \left(\frac{B_T}{B_B} \right)$ where A_L has 0 columns, B_T has 0 rows |
| whole 11 has a columns, 21 has a law |
| while $n(A_L) < n(A)$ do |
| |
| Determine block size b |
| $\left(\begin{array}{c c} A_L & A_R \end{array}\right) \to \left(\begin{array}{c c} A_0 & A_1 & A_2 \end{array}\right) , \left(\begin{array}{c} B_T \\ \hline B_B \end{array}\right) \to \left(\begin{array}{c} B_0 \\ \hline B_1 \\ B_2 \end{array}\right)$ |
| where A_1 has b columns, B_1 has b rows |
| |
| $C := A_1 B_1 + C$ |
| |
| $A \to \left(A_L \middle A_R \right) \leftarrow \left(A_0 \middle A_1 \middle A_2 \right), \left(\frac{B_T}{B_B} \right) \leftarrow \left(\frac{B_0}{B_2} \right)$ |
| |
| endwhile |
| |
| |

Algorithm: C := AB + C

$$A \to \left(A_L \mid A_R \right), B \to \left(\frac{B_T}{B_B} \right)$$

where A_L has 0 columns, B_T has 0 rows

while $n(A_L) < n(A)$ do

Determine block size b

$$\left(A_L \middle| A_R \right) \to \left(A_0 \middle| A_1 \middle| A_2 \right), \left(\frac{B_T}{B_B} \right) \to \left(\frac{B_0}{B_1} \middle| B_2 \right)$$

where A_1 has b columns, B_1 has b rows

$$C := A_1 B_1 + C$$

$$A \to \left(A_L \middle| A_R \right) \leftarrow \left(A_0 \middle| A_1 \middle| A_2 \right), \left(\frac{B_T}{B_B} \right) \leftarrow \left(\frac{B_0}{B_1} \right)$$

endwhile