

Step	<b>Algorithm:</b> $C := AB + C$ where $A$ is symmetric and stored in the lower triangular part
1a	$\{C = \widehat{C}$
4	$B \rightarrow \left( B_L \middle  B_R \right), C \rightarrow \left( C_L \middle  C_R \right)$ where $B_L$ has 0 columns, $C_L$ has 0 columns
2	$\left\{ \left( C_L \middle  C_R \right) = \left( AB_L + \widehat{C}_L \middle  \widehat{C}_R \right) \right\}$
3	<b>while</b> $n(B_L) < n(B)$ <b>do</b>
2,3	$\left\{ \left( C_L \middle  C_R \right) = \left( AB_L + \widehat{C}_L \middle  \widehat{C}_R \right) \wedge n(B_L) < n(B) \right\}$
5a	$\left( B_L \middle  B_R \right) \rightarrow \left( B_0 \middle  b_1 \ B_2 \right), \left( C_L \middle  C_R \right) \rightarrow \left( C_0 \middle  c_1 \ C_2 \right)$ where $b_1$ has 1 column, $c_1$ has 1 column
6	$\left\{ \left( C_0 \ c_1 \ C_2 \right) = \left( AB_0 + \widehat{C}_0 \quad \widehat{c}_1 \quad \widehat{C}_2 \right) \right\}$
8	$c_1 := Ab_1 + c_1$
7	$\left\{ \left( C_0 \ c_1 \ C_2 \right) = \left( AB_0 + \widehat{C}_0 \ Ab_1 + \widehat{c}_1 \quad \widehat{C}_2 \right) \right\}$
5b	$B \rightarrow \left( B_L \middle  B_R \right) \leftarrow \left( B_0 \ b_1 \middle  B_2 \right), C \rightarrow \left( C_L \middle  C_R \right) \leftarrow \left( C_0 \ c_1 \middle  C_2 \right)$
2	$\left\{ \left( C_L \middle  C_R \right) = \left( AB_L + \widehat{C}_L \middle  \widehat{C}_R \right) \right\}$
	<b>endwhile</b>
2,3	$\left\{ \left( C_L \middle  C_R \right) = \left( AB_L + \widehat{C}_L \middle  \widehat{C}_R \right) \wedge \neg(n(B_L) < n(B)) \right\}$
1b	$\{C = AB + \widehat{C}$

Step	Algorithm: $C := AB + C$ where $A$ is symmetric and stored in the lower triangular part
1a	{
4	where
2	{
3	while do
2,3	{ $\wedge$ }
5a	where
6	{
8	
7	{
5b	
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	endwhile
2,3	{ $\wedge \neg($ ) }
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3	while do
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	while $n(B_L) < n(B)$ do
	$\left( B_L \middle  B_R \right) \rightarrow \left( B_0 \middle  b_1 \ B_2 \right), \left( C_L \middle  C_R \right) \rightarrow \left( C_0 \middle  c_1 \ C_2 \right)$ where $b_1$ has 1 column, $c_1$ has 1 column
	$c_1 := Ab_1 + c_1$
	$B \rightarrow \left( B_L \middle  B_R \right) \leftarrow \left( B_0 \ b_1 \middle  B_2 \right), C \rightarrow \left( C_L \middle  C_R \right) \leftarrow \left( C_0 \ c_1 \middle  C_2 \right)$
	endwhile

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$B \rightarrow \left( B_L \middle  B_R \right), C \rightarrow \left( C_L \middle  C_R \right)$ where $B_L$ has 0 columns, $C_L$ has 0 columns while $n(B_L) < n(B)$ do $\left( B_L \middle  B_R \right) \rightarrow \left( B_0 \middle  b_1 \ B_2 \right), \left( C_L \middle  C_R \right) \rightarrow \left( C_0 \middle  c_1 \ C_2 \right)$ where $b_1$ has 1 column, $c_1$ has 1 column $c_1 := Ab_1 + c_1$ $B \rightarrow \left( B_L \middle  B_R \right) \leftarrow \left( B_0 \ b_1 \middle  B_2 \right), C \rightarrow \left( C_L \middle  C_R \right) \leftarrow \left( C_0 \ c_1 \middle  C_2 \right)$ endwhile