

Exercise 1

Rick Veens Studentno: 0912292 Huib Donkers Studentno: 0769015
r.veens@student.tue.nl h.t.donkers@student.tue.nl

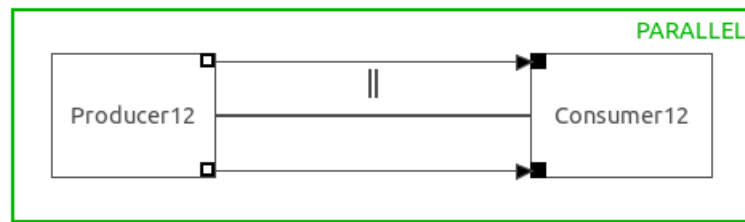
May 28, 2015

1

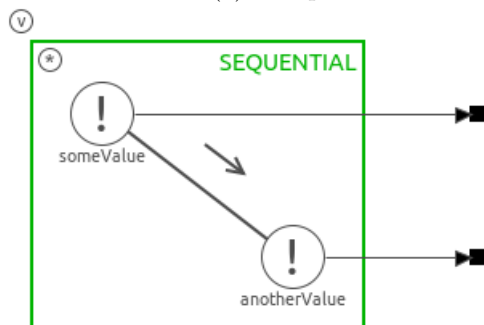
1.1

1.1.1

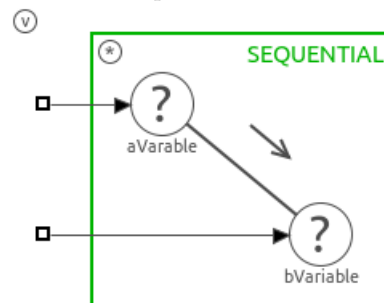
We modelled the deadlock free system as show in figure 1, and the deadlocked system as shown in figure 3. As expected, FDR accepts the deadlock free system (firtgure 2), but not the deadlocked system (figure 4).



(a) Composition of the Producer and Consumer processes.



(b) The producer process.



(c) The consumer process.

Figure 1: The producer-consumer (DF) model.

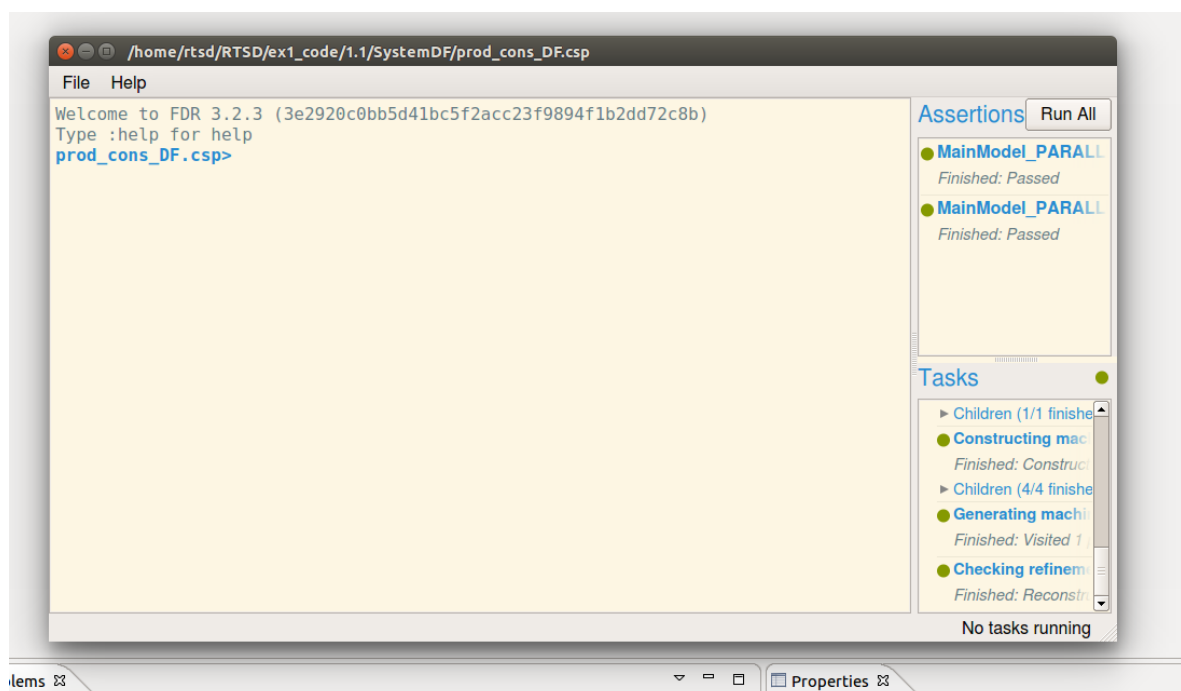
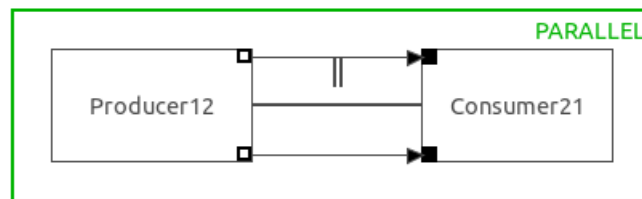
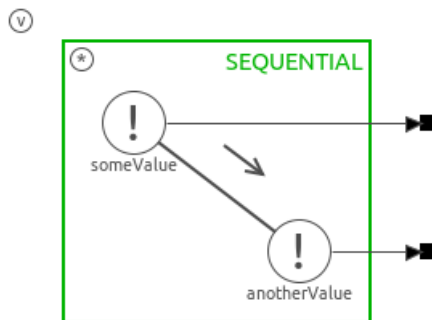


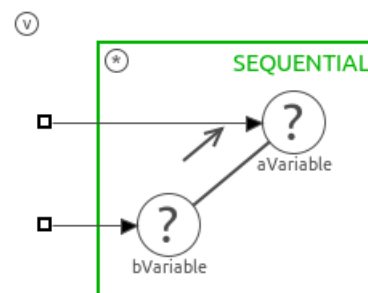
Figure 2: Successfully passes tests in FDR.



(a) Composition of Producer and Consumer processes.



(b) The producer process.



(c) The consumer process.

Figure 3: The producer-consumer (DC) model.

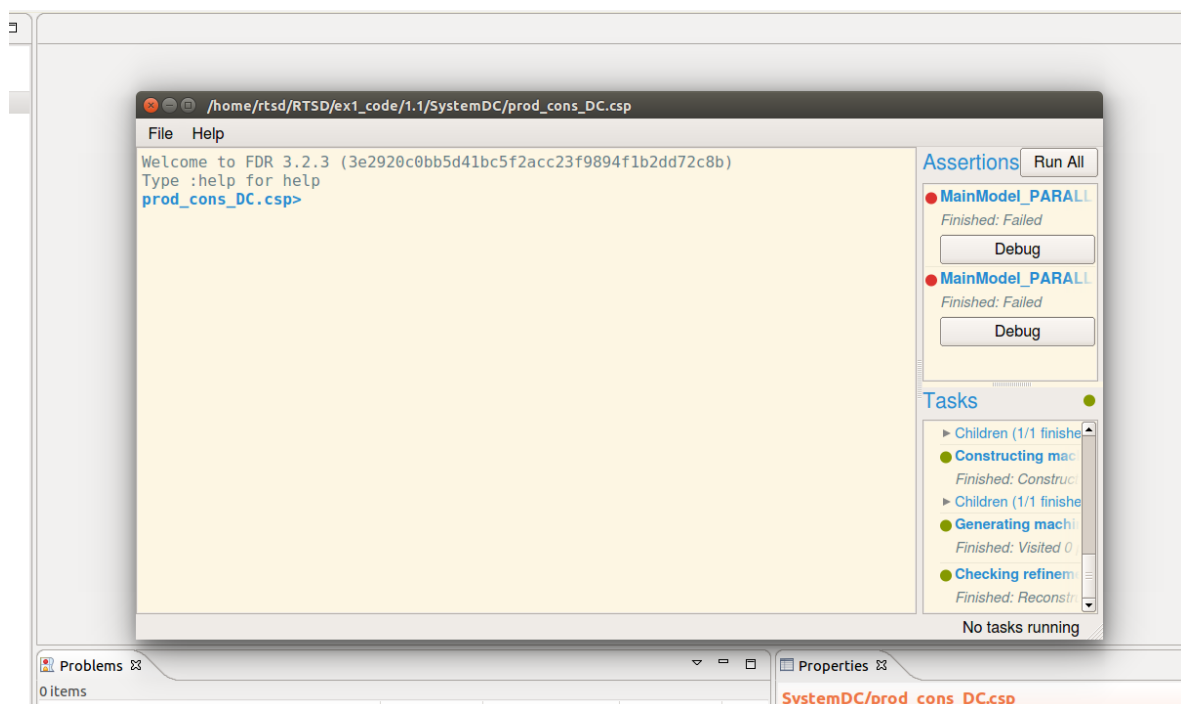


Figure 4: Fails the tests in FDR.

1.1.2

We extended the producer consumer (DF) model from figure 1 with code blocks, and added C++ implementation. The extended model is shown in figure 5 and the code is included in blocks 1 and 2. The initial few lines of the output is shown in figure 6.

Codeblock 1: Consumer12/CCode.cpp

```

1  /**
2  * Source file for the CCode model
3  * Generated by the TERRA CSPm2LUNA generator version 1.1.1
4  *
5  * protected region document description on begin
6  *
7  * protected region document description end
8  */
9
10 #include "Consumer12/CCode.h"
11 // protected region additional headers on begin
12 // Each additional header should get a corresponding dependency in the Makefile
13 // protected region additional headers end
14
15 namespace MainModel { namespace Consumer12 { namespace CCode {
16
17 CCode::CCode(int &CCode_aVariable, int &CCode_bVariable) :
18     CodeBlock(), CCode_aVariable(CCode_aVariable), CCode_bVariable(CCode_bVariable){
19     SETNAME(this, "CCode");
20
21     // protected region constructor on begin
22     // protected region constructor end
23 }
24
25 CCode::~~CCode()
26 {
27     // protected region destructor on begin
28     // protected region destructor end
29 }
30
31 void CCode::execute()
32 {
33     // protected region execute code on begin
34     if (this->CCode_aVariable == -1 || this->CCode_bVariable == -1)
35         exit();
36     else {
37         printf("Receiving: CCode_aVariable: \t'%c'\n", this->CCode_aVariable);
38         printf("Receiving: CCode_bVariable: \t'%c'\n", this->CCode_bVariable);
39
40         printf("\n");
41     }
42     // protected region execute code end
43 }
44
45 // protected region additional functions on begin
46 // protected region additional functions end
47
48 // Close namespace(s)
49 } } }

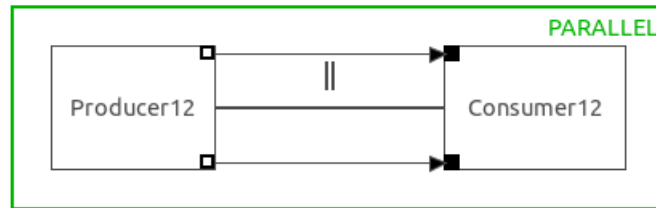
```

Codeblock 2: Producer12/PCode.cpp

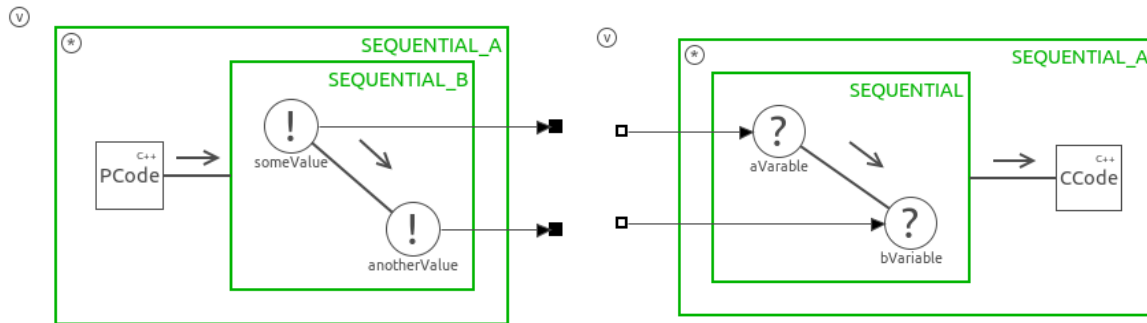
```

1  /**
2  * Source file for the PCode model

```



(a) Overview diagram of the producer consumer system.



(b) The producer process.

(c) The consumer process.

Figure 5: The producer-consumer (DF) model extended with code.

```

3  * Generated by the TERRA CSPm2LUNA generator version 1.1.1
4  *
5  * protected region document description on begin
6  *
7  * protected region document description end
8  */
9
10 #include "Producer12/PCode.h"
11 #include "string.h"
12 // protected region additional headers on begin
13 // Each additional header should get a corresponding dependency in the Makefile
14 // protected region additional headers end
15
16 namespace MainModel { namespace Producer12 { namespace PCode {
17
18 PCode::PCode(int &PCode_anotherValue, int &PCode_someValue) :
19     CodeBlock(), PCode_anotherValue(PCode_anotherValue), PCode_someValue(&PCode_anotherValue){
20     SETNAME(this, "PCode");
21
22     // protected region constructor on begin
23     // protected region constructor end
24 }
25
26 PCode::~PCode()
27 {
28     // protected region destructor on begin
29     // protected region destructor end
30 }
31
32 void PCode::execute()
33 {
34     // protected region execute code on begin

```

```
35
36     static int index = 0;
37
38     char *stuff = "Appelflap";
39
40     if (index == strlen(stuff)) {
41         this->PCode_someValue = -1;
42         this->PCode_anotherValue = -1;
43     } else {
44         this->PCode_someValue = stuff[index];
45         this->PCode_anotherValue = stuff[index++];
46
47         printf("Sending: PCode_someValue: \t'%c'\n", this->PCode_someValue);
48         printf("Sending: PCode_anotherValue: \t'%c'\n", this->PCode_anotherValue);
49     }
50
51
52     // protected region execute code end
53 }
54
55 // protected region additional functions on begin
56 // protected region additional functions end
57
58 // Close namespace(s)
59 } } }
```

```

rtsd@rtsd-ubuntu-64bit: ~/RTSD/ex1_code/1.2/prod_cons_DF
g++ -L/opt/LUNA/target-x86_64-linux_x86_64_CSP/lib64 obj/main.o obj/MainModel.o
obj/Producer12.o obj/Producer12/PCode.o obj/Consumer12.o obj/Consumer12/CCode.o
-llUNA -lpthread -o prod_cons_DF
rtsd@rtsd-ubuntu-64bit:~/RTSD/ex1_code/1.2/prod_cons_DF$ ./prod_cons_DF
Sending: PCode_someValue: 'A'
Sending: PCode_anotherValue: 'A'
Sending: PCode_someValue: 'p'
Sending: PCode_anotherValue: 'p'
Receiving: CCode_aVariable: 'A'
Receiving: CCode_bVariable: 'A'

Sending: PCode_someValue: 'p'
Sending: PCode_anotherValue: 'p'
Receiving: CCode_aVariable: 'p'
Receiving: CCode_bVariable: 'p'

Sending: PCode_someValue: 'e'
Sending: PCode_anotherValue: 'e'
Receiving: CCode_aVariable: 'p'
Receiving: CCode_bVariable: 'p'

Sending: PCode_someValue: 'l'
Sending: PCode_anotherValue: 'l'
Receiving: CCode_aVariable: 'e'
Receiving: CCode_bVariable: 'e'

Sending: PCode_someValue: 'f'
Sending: PCode_anotherValue: 'f'
Receiving: CCode_aVariable: 'l'
Receiving: CCode_bVariable: 'l'

Sending: PCode_someValue: 'l'
Sending: PCode_anotherValue: 'l'
Receiving: CCode_aVariable: 'f'
Receiving: CCode_bVariable: 'f'

Sending: PCode_someValue: 'a'
Sending: PCode_anotherValue: 'a'
Receiving: CCode_aVariable: 'l'
Receiving: CCode_bVariable: 'l'

Sending: PCode_someValue: 'p'
Sending: PCode_anotherValue: 'p'
Receiving: CCode_aVariable: 'a'
Receiving: CCode_bVariable: 'a'

Receiving: CCode_aVariable: 'p'
Receiving: CCode_bVariable: 'p'

```

Figure 6: The output that our system produces.

1.1.3

The consumer process cannot start before something is written on the channels. The producer process can start with the execution of its C++ code. The parallel composition of the two processes can therefore only start with the execution of the C++ code of the producer, hence the first two lines of the output originate from the producer.

Now, for the first channel, both the reader and the writer are available, and no other actions can be done by either process. The data is sent over the channel. After that, the same goes for the second channel.

Next, the producer can start over, starting with the execution of the C++ code, but now the consumer can also execute its C++ code. We expect either two lines of output produced by the producer's C++ code, or two lines produced by the consumer's C++ code. And indeed we see the output of the producer's C++ code.

Next, the writers in the producer cannot write to the channels, because the readers of the consumer are not ready. First the consumer's C++ code needs to be executed before the process can recur and execute its readers. Hence the next two lines are produced by the consumer.

Now, the readers and writers do their thing and, again, either the producer's C++ code can be executed, or the consumer's. We see for the complete initial part of the run that the producer's C++ code is executed before the consumer's C++ code.

1.1.4

ProBE is an interpreter for CSPm scripts, the user can manually inspect the possible execution paths. When a process is deadlock free, its execution paths are infinite. ProBE will never see the end of these execution paths, nor will it know whether it has an end. FDR recognises this infinite behaviour, or rather, it recognises finite behaviour. When it finds the process is able to terminate, it shows that the process fails the deadlock free-test. Similarly it can check for other properties without needing to simulate the complete execution paths like ProBE does. ProBE can be used to inspect how a process deadlocks.

1.1.5

If we adapt the C++ code of the producer to start at the beginning of the word again when all characters have been sent, we can observe the behaviour of the process when it runs for a while. We saw previously that the consumer only continued when the producer had to wait. This does not seem fair parallel. But when we watch what happens when the code runs for a while, we occasionally see the consumer receiving four values in a row, meaning that its C++ code executed before the producer's C++ code could prepare the next two values for the writers. So when both processes can execute, the parallel composition does not consistently execute one of them before the other. This seems fair.

Give proof that it is fair using the tools

1.1.6

The composition $A|||B$ does not deadlock. Whichever process gets to execute its C++ code first, it will need to wait for the other process directly after. When both have executed their

C++ code, execution can continue by writing and reading on the channels. At no point do the processes wait for the other.

1.2

1.2.1

See Figure 7 for some screenshots of the model. We chose to implement three channels.

1.2.2

See Codeblock 3 for the code of C++ code block PCode. The other code blocks in the program are of no interest, because they merely print their input variable.

Figure 9 shows some of the output of the program after extending it with C++ code.

The decision was made to have the program send each of the three values once. After sending the prompt appears again.

The guards on the producer side (see Figure 7b) are configured in the model in such a way that a value '1' gets send to channel 1, value '2' to channel 2 and value '3' to channel 3. Any other value that is entered will result in the "Prullenbak" C++ code being executed.

Note that some of the messages are out of order, the same stuff is happening as discussed in exercise 1.1.3.

In short, order we see occurs because the top producer and consumer process are linked in parallel. The communication of the consumer/producer processes are synchronized, but not the code being executed in their internal C++ code blocks.

1.2.3

Note that in Figure 7b a C++ codeblock is added called "Prullenbak". This "Prullenbak" C++ code block is executed only when none of the other guard expressions in the ALTERNATIVE block apply. The function of "Prullenbak" is to catch and print input data that is not to be send to the consumer, it is a garbage can.

1.2.4

It is wise to start with the process-structure before implementing C++ code.

Generally, if you have to choose, you want most of the functionality in the model instead of in C++ code blocks. The model also serves as documentation. Making sure the model is of decent quality first helps in containing most of the functionality in the model instead of the C++ code.

1.2.5

The method we chose to implement is to add another guard that is triggered when none of the other guards are triggered. It is a garbage bin exit.

Another possible method is to copy the guard logic of the alternative blocks in the producer to the "PCode" C++ code block, and do a check on the input data to make sure it checks out with your guard logic. This idea is less desired, because now you have the guard logic in *two* places: in the model and in the C++ code.

You could also remove the alternative structure, have only one instead of three channels, and put all your guard logic in the first C++ code block (PCode). This idea makes the program harder to understand, so not recommended.

Codeblock 3: Producer/PCode.cpp

```

1  /**
2  * Source file for the PCode model
3  * Generated by the TERRA CSPm2LUNA generator version 1.1.1
4  *
5  * protected region document description on begin
6  *
7  * protected region document description end
8  */
9
10 #include "Producer/PCode.h"
11 // protected region additional headers on begin
12 #include <stdio.h>
13 #include <iostream>
14 #include <algorithm>
15 // Each additional header should get a corresponding dependency in the Makefile
16 // protected region additional headers end
17
18 namespace MainModel { namespace Producer { namespace PCode {
19
20 PCode::PCode(int &peer) :
21     CodeBlock(), peer(peer){
22     SETNAME(this, "PCode");
23
24     // protected region constructor on begin
25     // protected region constructor end
26 }
27
28 PCode::~PCode()
29 {
30     // protected region destructor on begin
31     // protected region destructor end
32 }
33
34 void PCode::execute()
35 {
36     // protected region execute code on begin
37     static int index = 3;
38
39     if (index > 2) {
40         index = 0;
41         while (!initialize()) {
42             printf("Numbers not distinct.\n");
43         }
44     }
45     printf("P1->Sending %d\n", vars[index]);
46     this->peer = vars[index++];
47     // protected region execute code end
48 }
49
50 // protected region additional functions on begin
51 bool PCode::initialize(void)
52 {
53     printf("Enter 3 numbers\n");
54     while (scanf("%i", &vars[0]) != 1) {}
55     while (scanf("%i", &vars[1]) != 1) {}
56     while (scanf("%i", &vars[2]) != 1) {}
57
58     std::sort(vars, vars+3);

```

```
59
60     printf("I got: %i, %i, %i\n", vars[0], vars[1], vars[2]);
61     if (vars[0] == vars[1] || vars[0] == vars[2] || vars[1] == vars[2])
62         return false;
63     else
64         return true;
65 }
66 // protected region additional functions end
67
68 // Close namespace(s)
69 } } }
```

1.3

1.4

1.4.1

Figure 10 presents the JIWIY controller model as presented in the slides, with the addition of extra readers and writers that: simulate the input of the joystick, the output to the robot and the feedback from the robot.

This is added because otherwise it was not possible to generate CSPm: one cannot leave input and output ports unconnected in the top model.

Figure 11 informs us that the model results in a deadlock in FDR.

1.4.2

Figure 12a shows the small fix to remove the deadlock and this is proven with a screenshot of FDR as seen in Figure 12a.

In the deadlock situation, observe the following:

- The Check process must read from the Vertical process first, before it can read from the Horizontal process.
- The Vertical process must read from the Horizontal process before it can write to the Check process.
- The Horizontal process must first write to the Check process before it can write to the Vertical process.

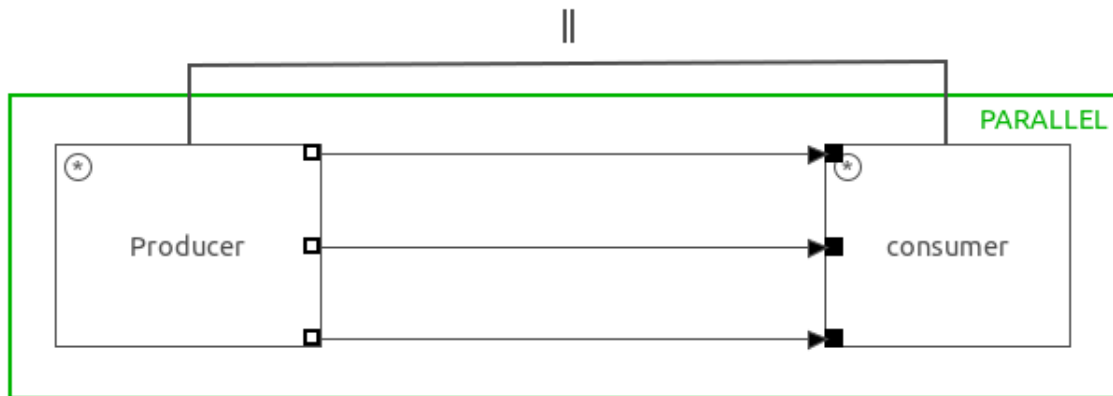
This results in the deadlock, we have a loop.

Making sure that the Check process has to read first from the Horizontal process instead of the Vertical process will fix the problem.

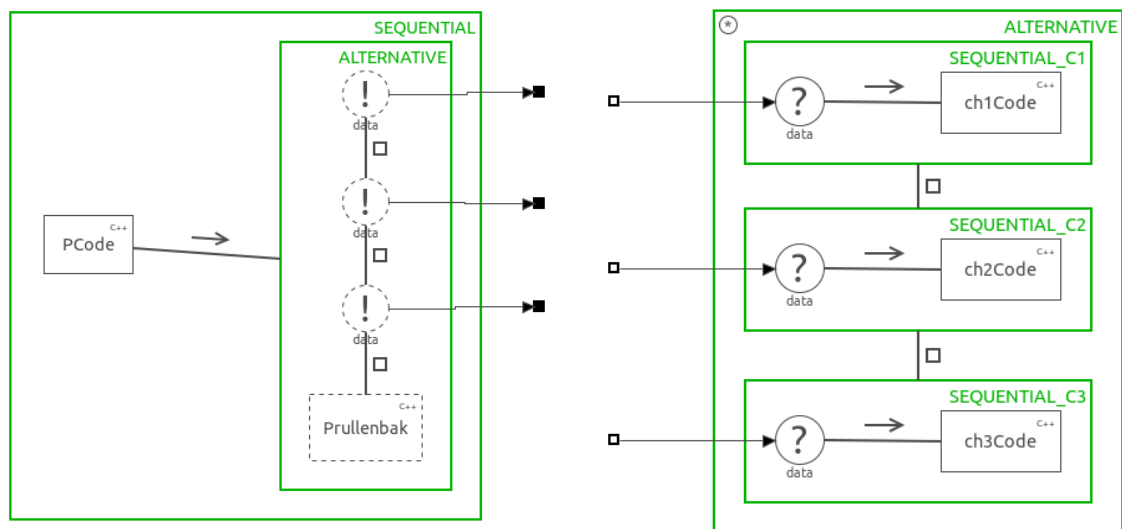
If the Check process first reads from the Horizontal process and then to the Vertical process; the Horizontal Process is first able to write to the Check process and then to the Vertical process (which is waiting on Horizontal), and this in turn allows the Vertical process to write to the Check process.

The Check process is able to receive the write from the Vertical process because it first reads from Horizontal and then from Vertical.

So this small change fixes the deadlock, see Figure 12a. Note the arrow is actually in the wrong direction, this is a visual bug in TERRA.



(a) Overview diagram of the multiple channel producer consumer system.



(b) The producer process.

(c) The consumer process.

Figure 7: The CSP model of exercise 1.2.

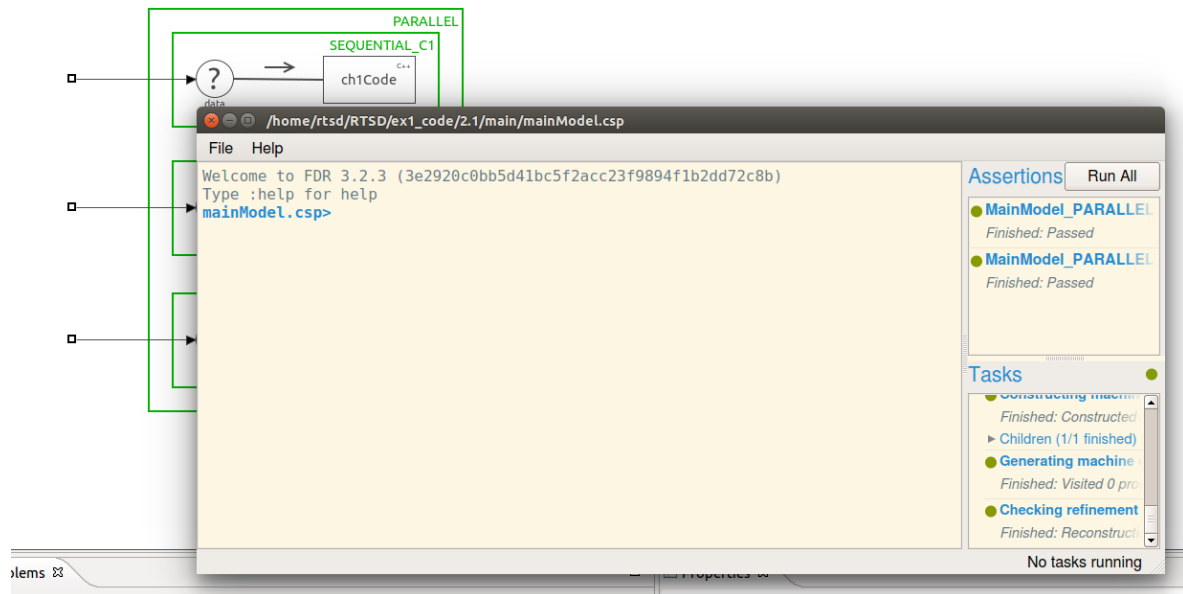
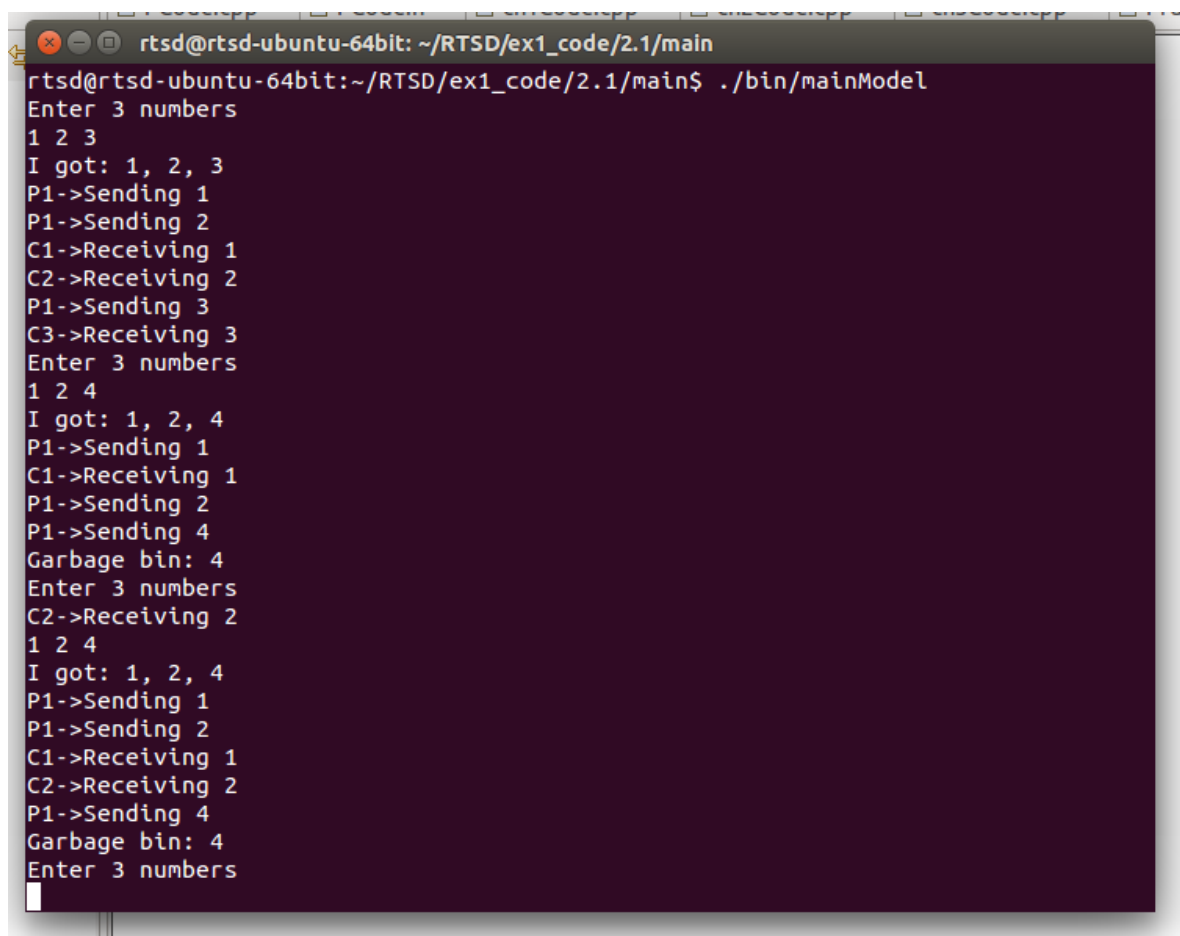


Figure 8: Deadlock free.

1.4.3

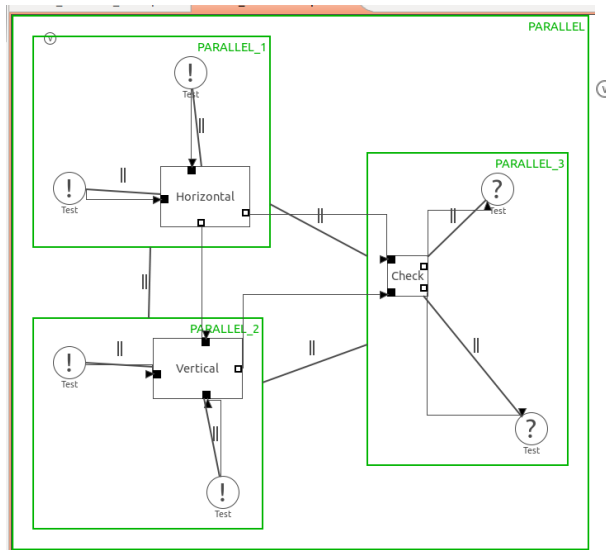
Figure 13 shows the IO SEQ pattern used on the JIWI model and Figure 14 displays the output of FDR, which does not deadlock.

1.4.4

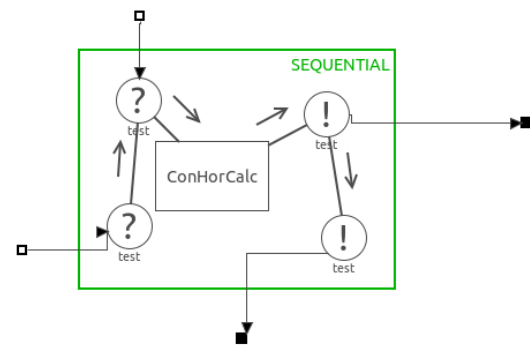


```
rtsd@rtsd-ubuntu-64bit: ~/RTSD/ex1_code/2.1/main
rtsd@rtsd-ubuntu-64bit:~/RTSD/ex1_code/2.1/main$ ./bin/mainModel
Enter 3 numbers
1 2 3
I got: 1, 2, 3
P1->Sending 1
P1->Sending 2
C1->Receiving 1
C2->Receiving 2
P1->Sending 3
C3->Receiving 3
Enter 3 numbers
1 2 4
I got: 1, 2, 4
P1->Sending 1
C1->Receiving 1
P1->Sending 2
P1->Sending 4
Garbage bin: 4
Enter 3 numbers
C2->Receiving 2
1 2 4
I got: 1, 2, 4
P1->Sending 1
P1->Sending 2
C1->Receiving 1
C2->Receiving 2
P1->Sending 4
Garbage bin: 4
Enter 3 numbers
```

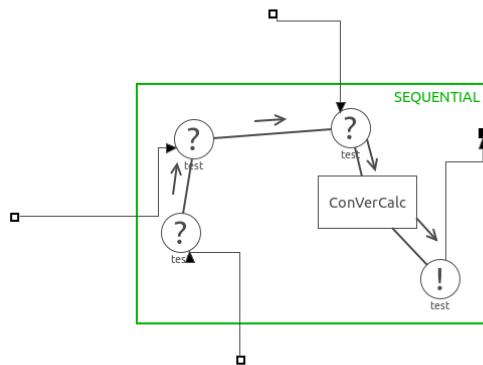
Figure 9: Output of the program exercise 1.2.



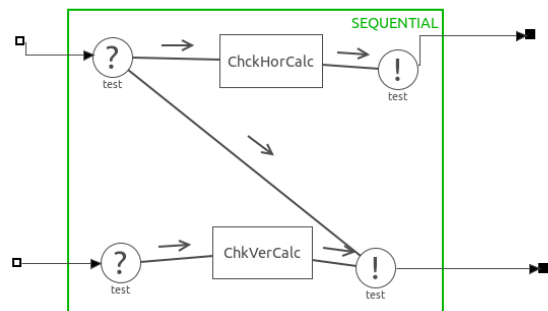
(a) Overview diagram of JIWIY model.



(b) Inside the Horizontal process.



(c) Inside the Vertical process.

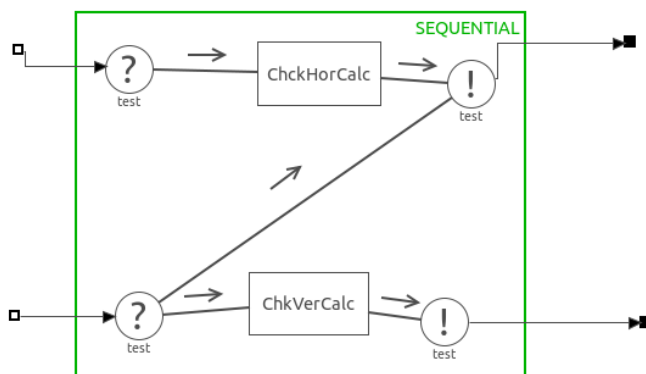


(d) Inside the Check process.

Figure 10: JIWIY controller model (deadlock).



Figure 11: The model deadlocks.

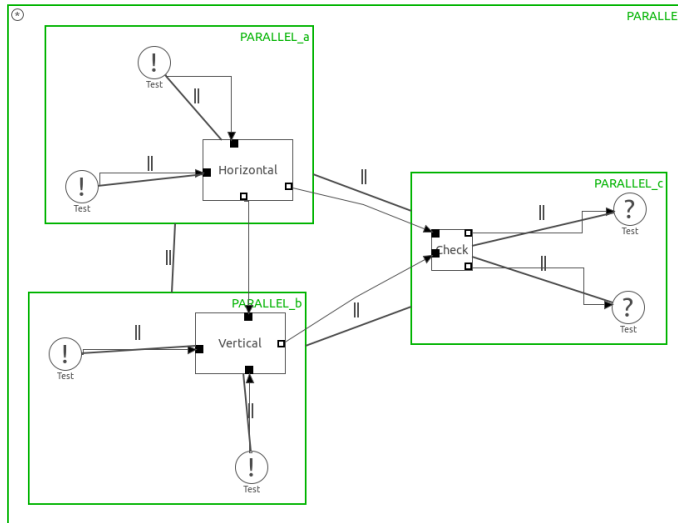


(a) Minimal deadlock fix.

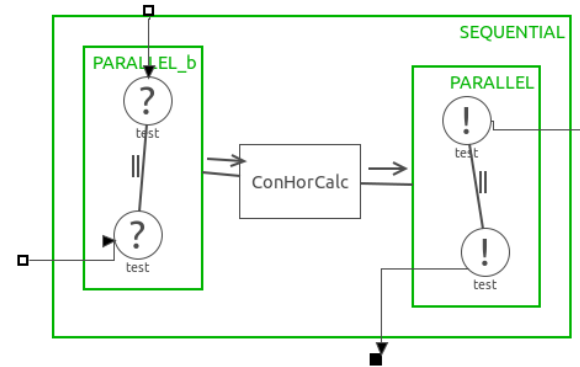


(b) No deadlock according to FDR.

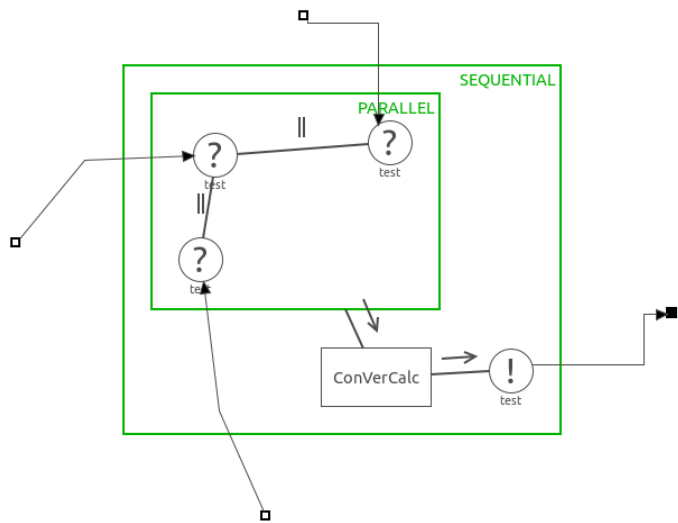
Figure 12: Fixing the deadlock.



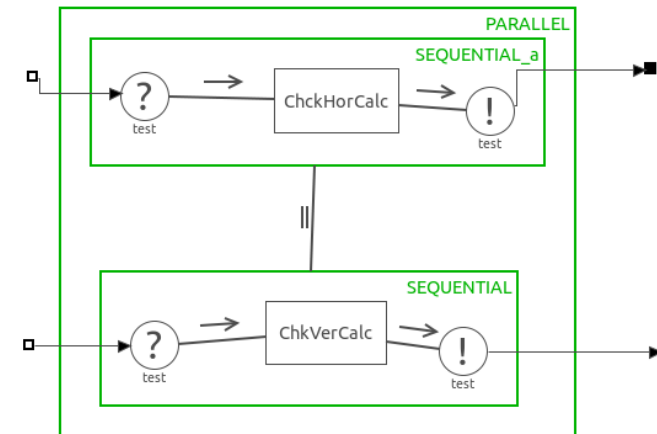
(a) Overview diagram of JIWIY model.



(b) Inside the Horizontal process.



(c) Inside the Vertical process.



(d) Inside the Check process.

Figure 13: JIWIY controller model IO SEQ pattern.



Figure 14: No deadlock with the IO SEQ pattern.