

CSCI 260 Notes

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<https://github.com/rvente/CSCI-260-Notes>



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Note: class topics are important because the textbook has incomplete information on this subject.

Functions in MIPS

(1) Example: Implement the following c code in MIPS

```
main() {
    d = foo(2*a, b);
}

int foo(int x, int y) {
    return (x+y);
}

#allocate a -> $s0, b -> $s1, d -> $s2

    add  $a0, $s0, $s0    #a0 <- 2a (argument 0 register)
    add  $a1, $s1, $zero  #a1 <- b (" ")
    jal  foo              # call foo (see note)
```

```

        add $s2, $v0, $zero # d <- Return Value (2a,b)
        ...
foo:    add $v0, $a0, $a1    # v0 <- x+y
        jr  $ra             # return

        # jal sticks result of next function into $ra
        # jr is transferring control to main

```

Note: we should not be writing foo at this point, this is a top down approach.

Issues

1. What if main and foo use same registers?

Have a caller and callee follow a convention to save registers somewhere special.

Table 1: MIPS Conventions

register	contents	Saved by
0	\$zero	
1	\$at assembler temp	
2, 3	return value	caller
4, 7	\$a0-3, arg regs	caller
8..15	\$t0-8, temp regs	caller
16.. 23	\$s0-7, s-registers	callee
24, 25	\$t3, \$t9 temp reg	caller
26, 27	\$k0, \$k1 reserved	
28	\$gp global ptr	callee
29	\$sp stack pointer	implicitly preserved
30	\$fp frame pointer	callee
31	\$ra return address	callee

2. What if foo calls bar (or is recursive)?

Last in First out Stacked function calls

Use a stack to store different values of \$ra.

```

# push s1 onto stack:
    addi $sp, $sp, -4 # because stack grows downward
    sw $s1, 0($sp)
# pop s5 off the stack:
    lw $s5, 0($sp)
    add $sp, $sp, 4

```

3. What if > 4 arguments?

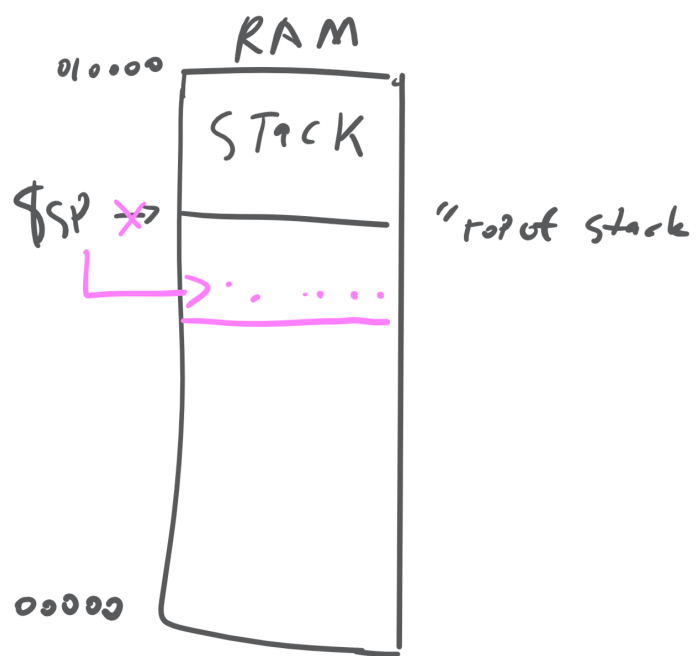


Figure 1: stack visualization

same as 4

4. Where do functions store local vars that don't fit in registers?

store extra args & local [non-static!!!] that don't fit in regs in an "activation record" (also called a "procedure frame"). place frame on the stack, with \$fp pointing to it.

access local vars as offsets to \$fp

Frame

caller -saved regs		
Extra args		
saved \$ra		<--- \$fp
saved \$fp		
other callee saved-regs		
Local Vars		
-----		< --- \$sp

For more information, observe the procedure in `procedure.pdf` supplied (and copyrighted) by Shankar.