

CSE 122 / 222C ; WES 269

WiFi – Return of the MAC

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WiFi MAC Goals

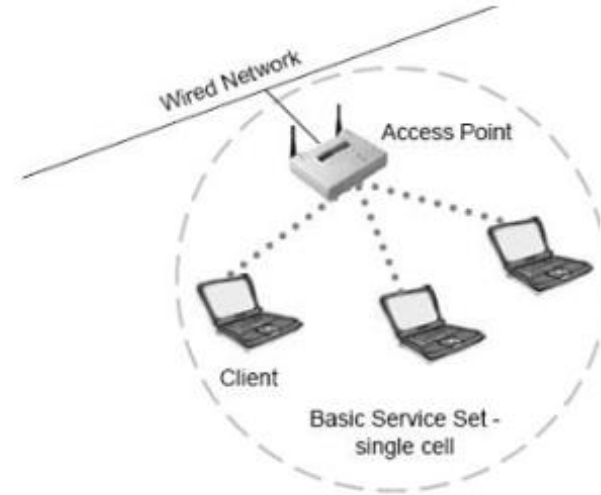
- Introduce MAC layer concepts in 802.11
- Understand what exists, what is used in practice, and why
- Microcontroller use of WiFi

Outline

- **802.11 Access Control**
- 802.11 Frame format
- 802.11e Improvements to MAC
- Microcontrollers and WiFi

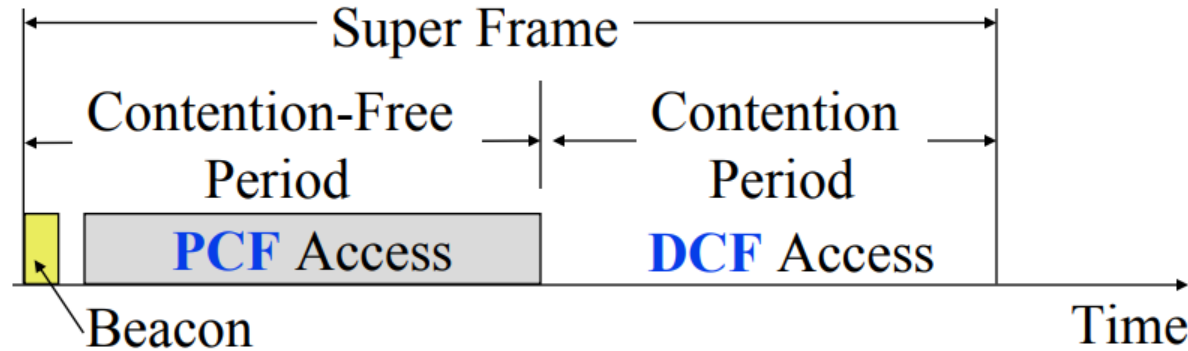
Basic WiFi network

- Star topology network
- Basic Service Set (BSS)
 - Access point(s)
 - Multiple connected clients
- Service Set ID (SSID)
 - Identifies network
 - Broadcast by access point in beacons



(original) WiFi superframe structure

- Beacon followed by contention-free period followed by contention
 - Repeats periodically (default ~100 ms)
 - 802.15.4 adopted a similar superframe



n.b. PCF is more hypothetical than real



bmcmurdo



Cisco Employee



09-04-2002 09:07 AM

PCF is not implemented on Cisco APs. PCF is not generally thought to be a very effective/useful QoS technique

The radio Access Point (and radio firmware) are "PCF aware" meaning if they

detect a PCF coordinator they will "play nice" in that area and back off a

little so it can do it's best effort.

IEEE 102.11e working group is currently working on a QoS standard but that has not yet been ratified.



0 Helpful

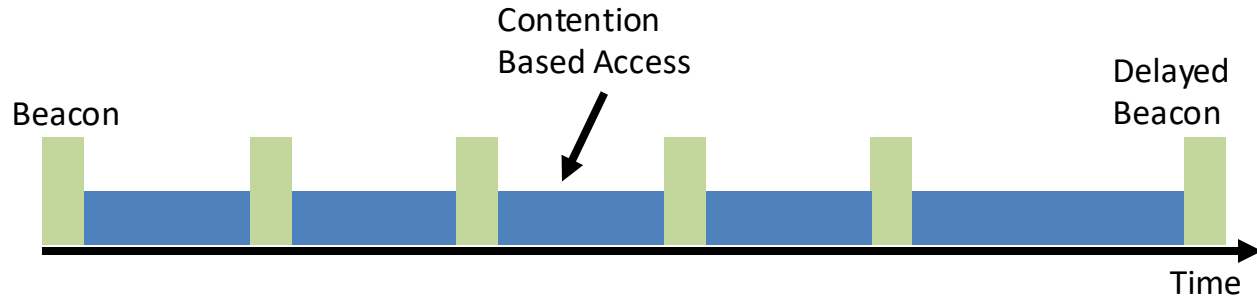
Reply

PCF superseded by Hybrid Coordination Function (HCF)

- 802.11e (2005 spec) defines HCCA
 - HCF Controlled Channel Access
- ... which is complicated, and only supported in specialized hardware and specialized use cases ...
- ... and HCCA is old; newer WiFi has more complex QoS protocols ...
- ... so that's all we're going to say about it

WiFi superframe in practice

- Continuous contention access period
 - DCF = “Distributed Coordination Function”
 - CSMA/CA + optional RTS/CTS
- Periodic beacons
 - Which also use CSMA and therefore may be delayed



802.11 beacons

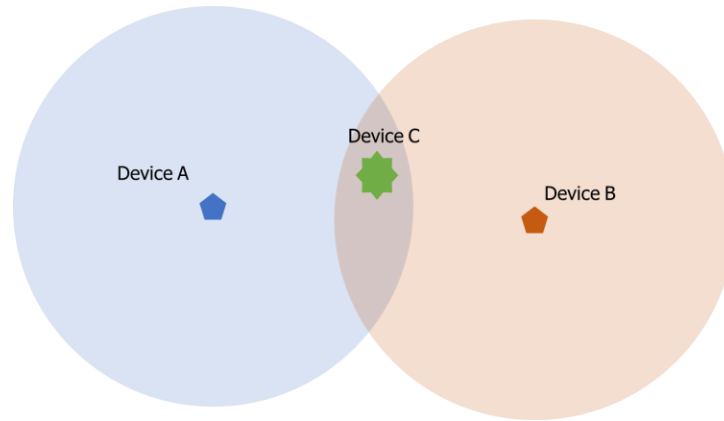
- Transmitted periodically (~100 ms by default)
 - Enable discovery of network
 - Contain capabilities and SSID for the network (802.11b/g/n/ac/ax...)
 - Assign contention-free slots if used
 - Notify devices of waiting packets
 - Traffic Indication Map (TIM) has a bitmap specifying which devices data is for
 - Enables devices to sleep, skipping a number of beacons
 - Handles broadcast/multicast messages
 - Every N beacons includes a notation of available broadcast messages
 - Messages are transmitted during next contention access period using normal CSMA
 - Defines maximum sleep period for devices (must listen to these beacons)

Contention-based access

- Known as Distributed Coordination Function (DCF)
 - Base communication method for WiFi
 - All packets are immediately ACK'd by receiving device
 - Uses CSMA/CA to determine when it can send
 - With random backoff
 - Problem: packets can be very long (up to 20 milliseconds)
 - Solution: Network Allocation Vector (NAV)
 - Packets include a notation of their duration
 - Sensing the beginning of a packet allows backoff to skip the whole packet duration before continuing

Reminder: hidden terminal problem

- Two devices communicating with Access Point may not be able to hear each other
 - CSMA fails and Access Point loses both messages



- A solution: RTS/CTS (Request/Clear To Send)

Drawbacks of RTS/CTS

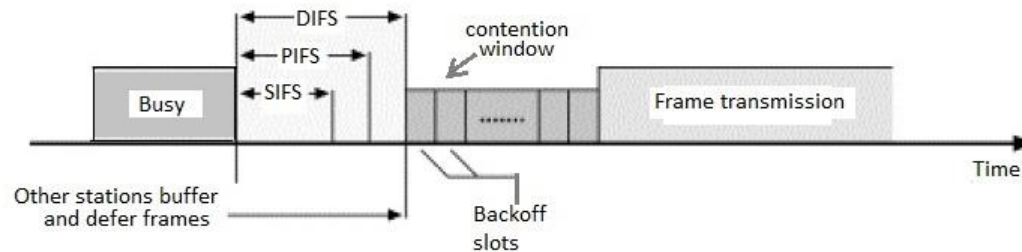
- Four packets per data (RTS, CTS, Data, Ack)
 - Could have just sent data instead of RTS
- Significant portion of traffic is application-layer Acks
 - Probably better to just have it fail and try again later
- RTS/CTS only used for very large packets in practice
 - *It's mentioned still in 802.11n and 802.11ac, so not entirely unused

Backoff in WiFi

- Listen for activity
 - If free
 - Wait for Inter Frame Spacing (IFS)
 - If still free, transmit
 - If busy
 - Randomly select a number of backoff **Slots**
 - Count down slots whenever medium is not busy
 - If busy when backoff completes:
 - Increase maximum backoff Slots
 - Repeat
- Slot time: basic time unit for protocol
 - Total time of: switch from Rx to Tx, plus processing time, plus propagation delay

Prioritizing packets with varying IFS

- Tiered Contention Multiple Access (TCMA)
 - Idea: assign different inter-frame spacing based on traffic class
 - Inherently prioritizes communication
- Acknowledgements sent with Short IFS (SIFS)
 - Will always transmit before new data clears CSMA check
- New data sent with DCF IFS (DIFS)

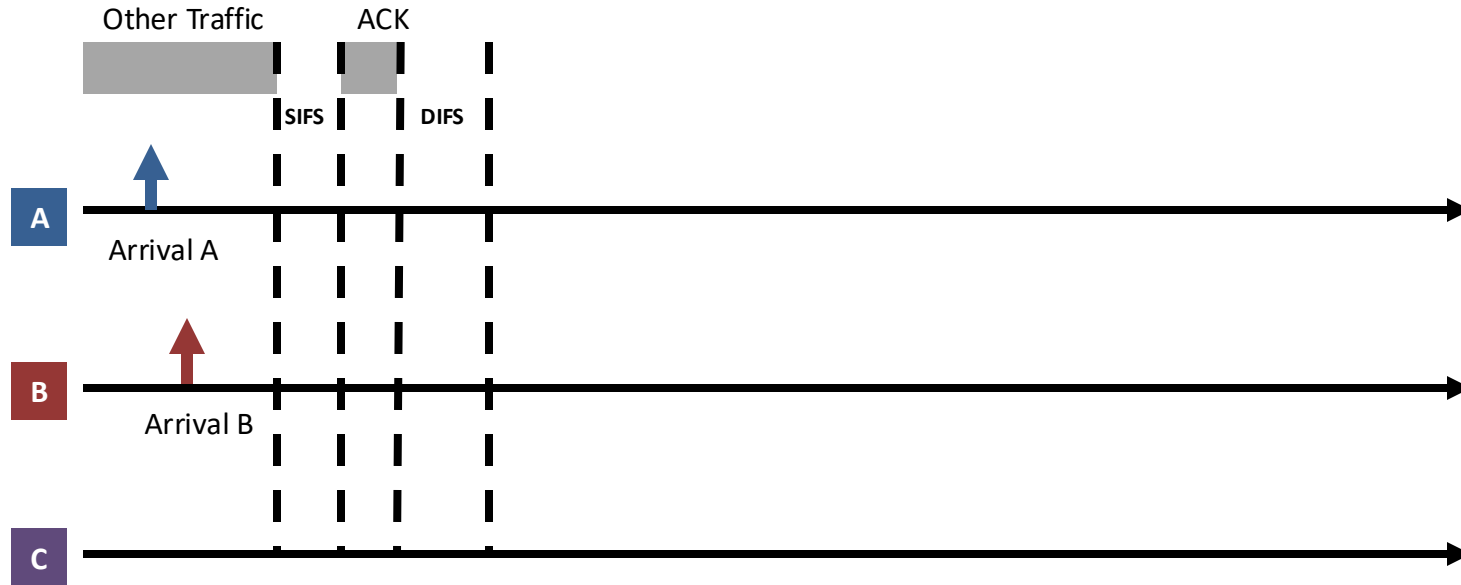


Putting backoff together

- Two variables
 - Contention Window (CW) – maximum backoff amount
 - Backoff Count (BO) – current remaining backoff
- When attempting to send, if busy Backoff selected in $[0, CW]$
 - Countdown Backoff slots whenever medium is not busy
 - At 0, attempt to transmit if not busy
 - If busy, double Window and select Backoff again
- 802.11g values:
 - Slot time= $20\ \mu\text{s}$, CWmin= 15 slots, CWmax= 1023 slots (20 ms)
 - SIFS= $10\ \mu\text{s}$, PIFS= $30\ \mu\text{s}$, DIFS= $50\ \mu\text{s}$

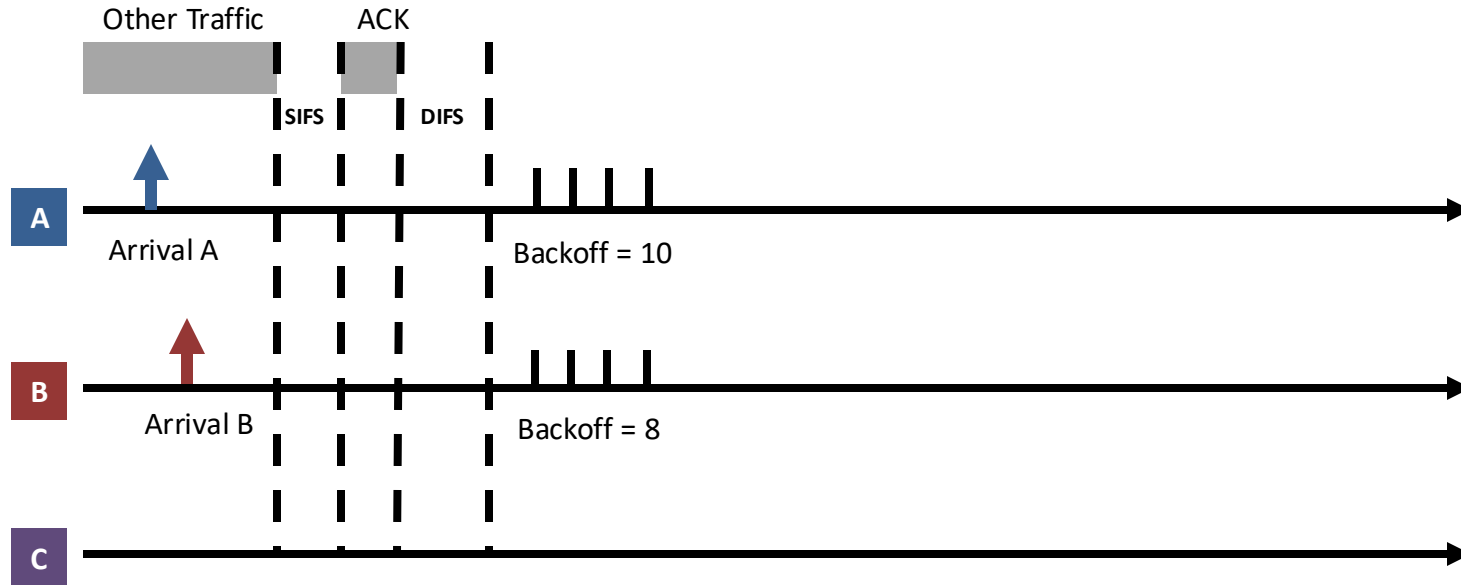
802.11 backoff example

- A and B want to send and see the medium is busy
 - Followed by an Acknowledgement after SIFS



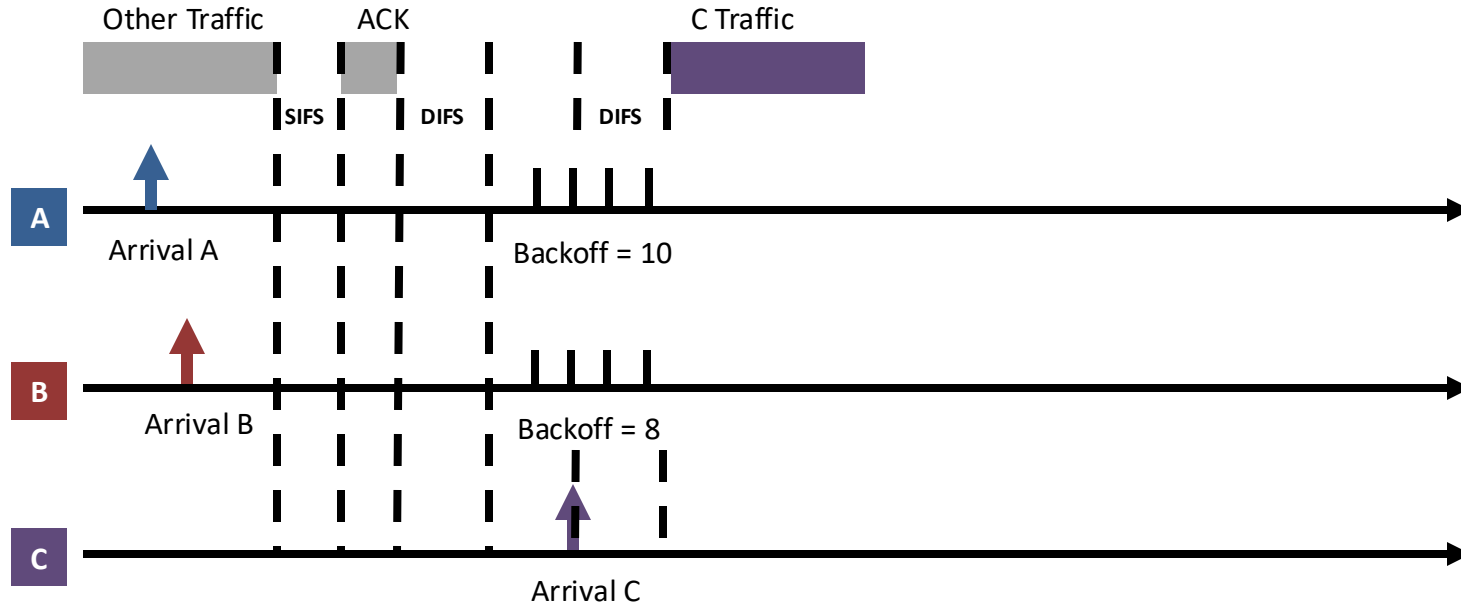
802.11 backoff example

- Each chooses a random backoff [0, CW] (we'll say CW is 32)
 - Start counting down backoff slots



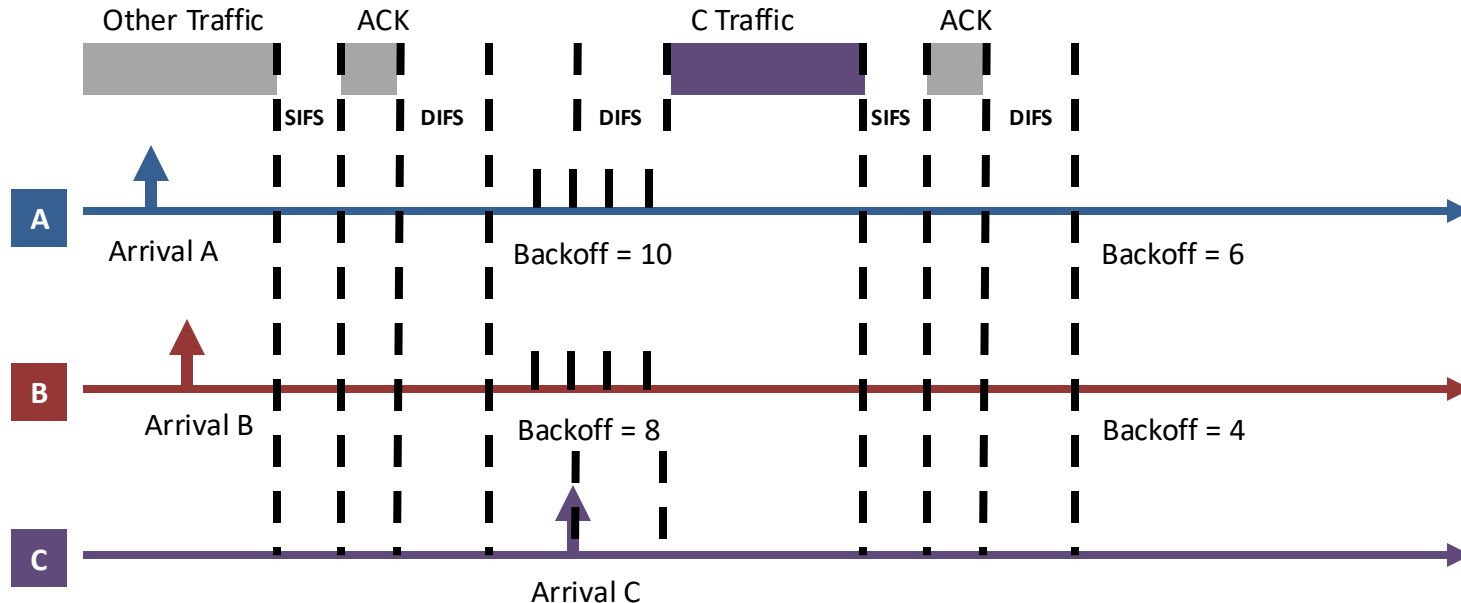
802.11 backoff example

- C wants to send, waits DIFS, and can send immediately
 - No other traffic is going on
 - A and B pause backoff for packet duration



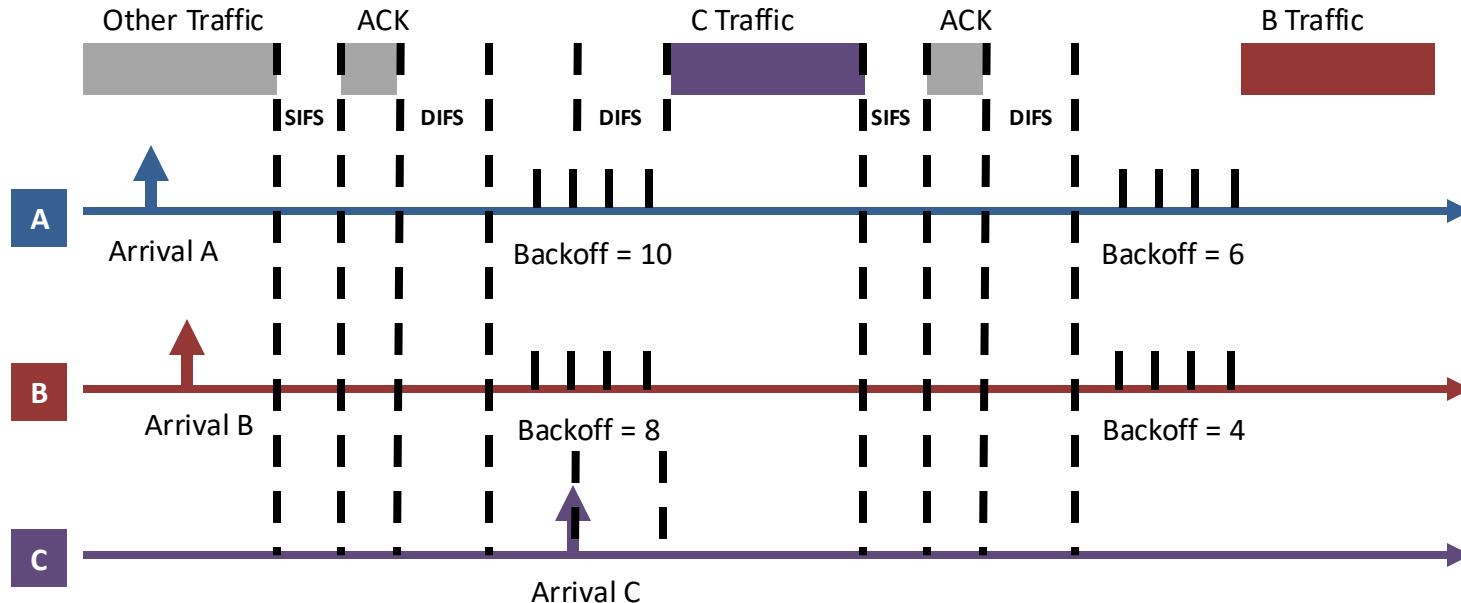
802.11 backoff example

- A and B used NAV to pause backoff for entire traffic plus ACK
 - After DIFS, resume backoff count from its previous value



802.11 backoff example

- B reaches zero backoff, finds channel empty, transmits
 - A pauses its backoff again for duration plus ACK



Break + Hacking

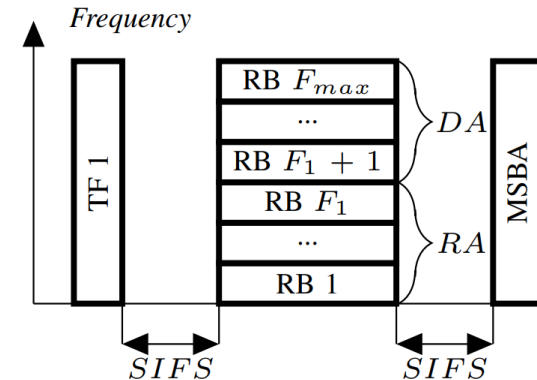
- If you wanted maximum data throughput on a WiFi radio, and you were willing to be non-standards-compliant, what would you do?

Break + Hacking

- If you wanted maximum data throughput on a WiFi radio, and you were willing to be non-standards-compliant, what would you do?
 - Never backoff at all. Just try during the next open period
 - Always be “device C” in our previous example
 - Use a shorter SIFS than other devices
 - If you start transmitting sooner, you get to keep transmitting!
 - Other devices will backoff on your transmission
 - Tragedy of the Commons: this utterly fails if many radios follow it

CSMA + OFDMA

- For WiFi 6 to be backwards compatible, the OFDMA stuff has to obey the normal CSMA/CA rules as well
- Downlink from AP
 - The AP wins contention and then transmits various data at various frequencies
- Uplink to AP
 - The AP wins contention, sends a “OFDMA uplink trigger frame”, then devices send their responses



Outline

- 802.11 Access Control
- **802.11 Frame format**
- 802.11e Improvements to MAC
- Microcontrollers and WiFi

802.11 frame

Field	Frame control	Duration, id	Address 1	Address 2	Address 3	Sequence control	Address 4	QoS control	HT control	Frame body	Frame check sequence
Length (Bytes)	2	2	6	6	6	0, or 2	6	0, or 2	0, or 4	Variable	4

- Frame control (various bits)
 - Type of packet (Control, Management, Data)
 - Subtype (Association, RTS, CTS, Ack, etc.)
 - Indication of to/from “distribution system” (Internet rather than intranet)
- Duration
 - Specifies on-air time of full packet in μs
 - Note: **no actual length field**

Surprising, but smart!

Recall MCS vary — but everyone needs to be able to parse header (for duration, for NAV)

Length can be very large (e.g. in ac: 5.5 ms max duration is 4.5 MB length!); sent at full data rate

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- Sequence control
 - 4-bit fragment number
 - 12-bit sequence number
- Quality of Service control
 - Identifies traffic category
- High Throughput Control
 - Configurations for selecting best data rate
- Frame body
 - Max size depends on PHY
 - ~2000 for lower rates
 - ~8000 for 802.11n
 - ~11000 for 802.11ac
- Frame check sequence
 - 32-bit CRC

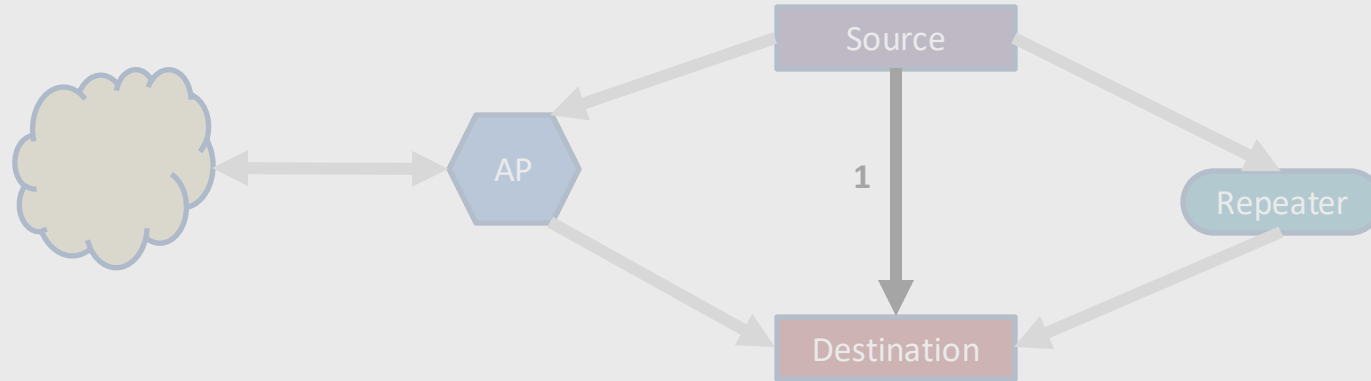
802.11 frames have four address fields... why?

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Wat?

Address field use cases



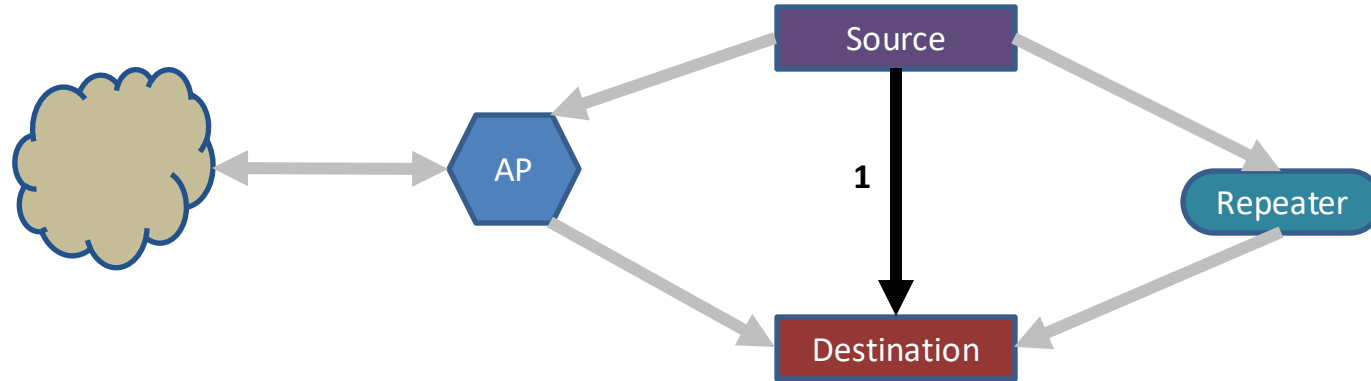
Devices filter on Address 1

	To DS	From DS	Address 1	Address 2	Address3	Address4	Use Case
1							Location
2							Internet
3							t
4							

DS = “Distribution System”

The infrastructure that connects APs. ‘Am I leaving this wireless area?’

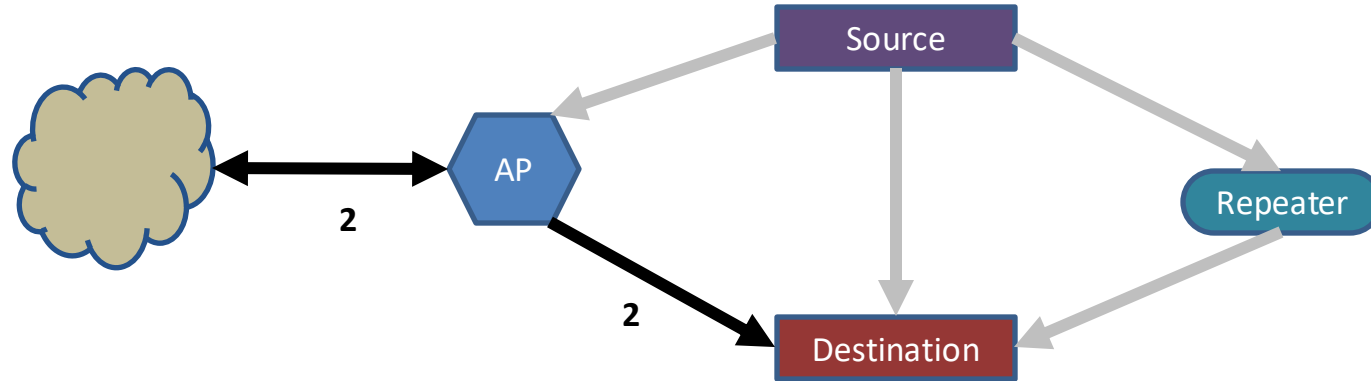
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1	0	0	Destination Addr	Source Addr	BSS ID	-	Direct communication
2	0	1	Destination Addr	BSS ID	Source Addr	-	Traffic from Internet
3	1	0	BSS ID	Source Addr	Destination Addr	-	Traffic to Internet
4	1	1	Receiver Addr	Transmitter Addr	Destination Addr	Source Addr	Repeater / Mesh

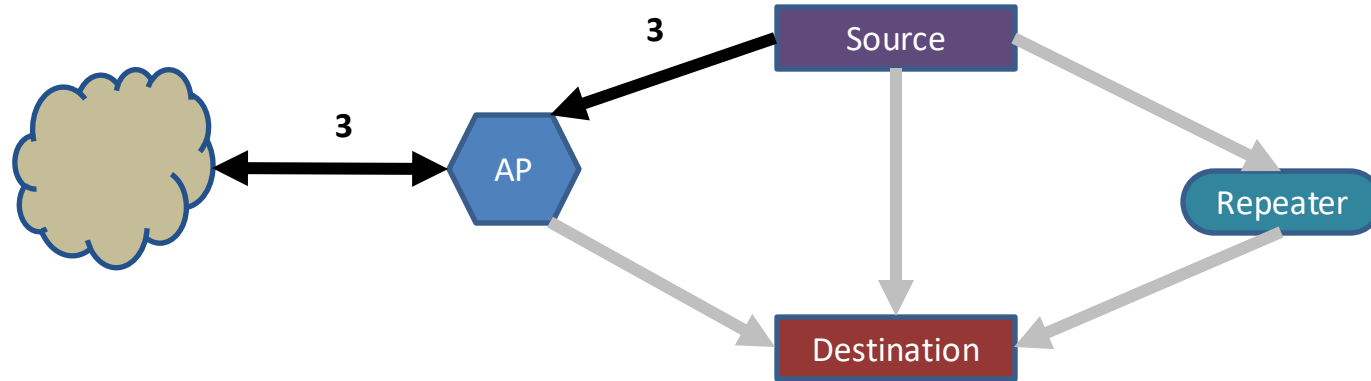
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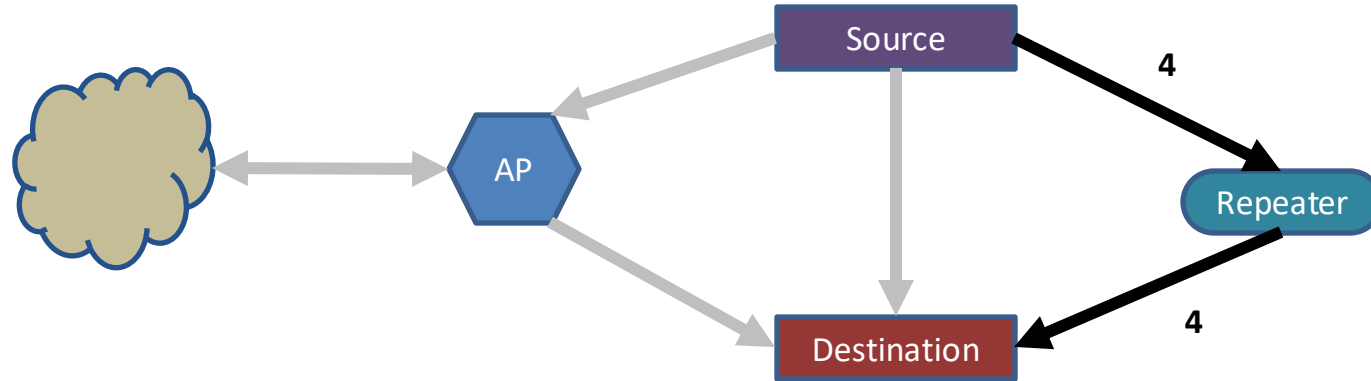
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Address field use cases

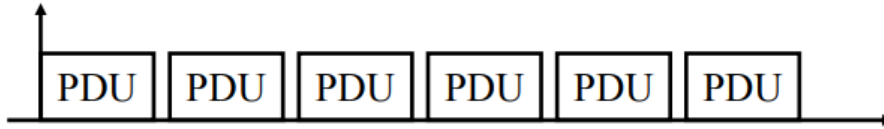


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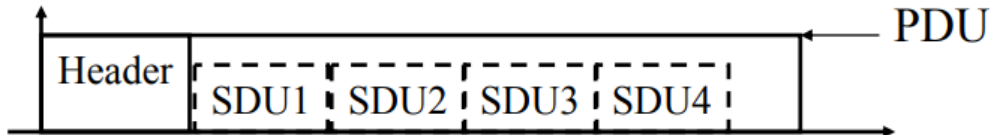
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Sending frames in WiFi

- Frame bursting
 - Transmit multiple frames in a row



- Frame fragmentation
 - Split service data over multiple frames
- Frame aggregation
 - Multiple service data in a single frame
 - Allows multiple packets to reach Access Point in a single transmission



Calculating packet durations

- Example duration for a 1500 byte 802.11g packet
 - 6 Mbps for header
 - 24 Mbps for payload
 - 566 μ s for total packet
 - Plus 10 μ s for SIFS
 - Plus 34 μ s for ACK
- https://sarwiki.informatik.hu-berlin.de/Package_transmission_time_in_802.11

Data transmission bitrate
(802.11g / a*):

	24	
	Mbps	

	Bitrate (Mbit/s)	Length (bits)	Time (μ s)
DIFS			28
D PHY header: PLCP preamble	-	-	16
A PHY header: PLCP header	6	24	4
MAC headers (28 bytes) + MAC			
T body	24	12246	512
A signal extension time			6
tx time data:			566
SIFS			10
A PHY header: PLCP preamble	-	-	16
C PHY header: PLCP header	6	24	4
K MAC headers + PHY pad	24	134	8
signal extension time			6
tx time ack:			44
tx time data + ack:			610

Implementation Drives Specification Sometimes

- SIFS nominally defined by processing time
 - Aside: Big challenge for SDRs
- Convolutional decoders need(ed) 16 μ s to finish processing
 - For highest-rate MCS (ERP-OFDM)
- Processing must finish before next packet starts
 - To be able to decode NAV in header

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(802.11g / a*):

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	Mbps	
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(Mbit/s)	(bits)	(μ s)

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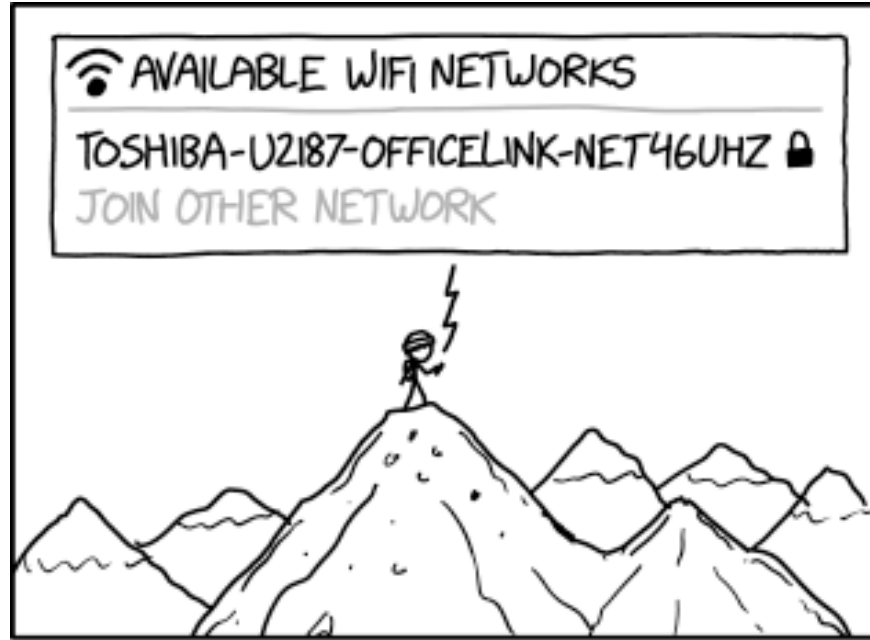
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tx time ack: 44

tx time data + ack: 610

Break + xkcd



TECH TRIVIA: NO ONE ACTUALLY KNOWS WHAT DEVICES PRODUCE THOSE CRYPTIC WIFI NETWORKS. THEY JUST APPEAR AT RANDOM ACROSS THE EARTH'S SURFACE.

Outline

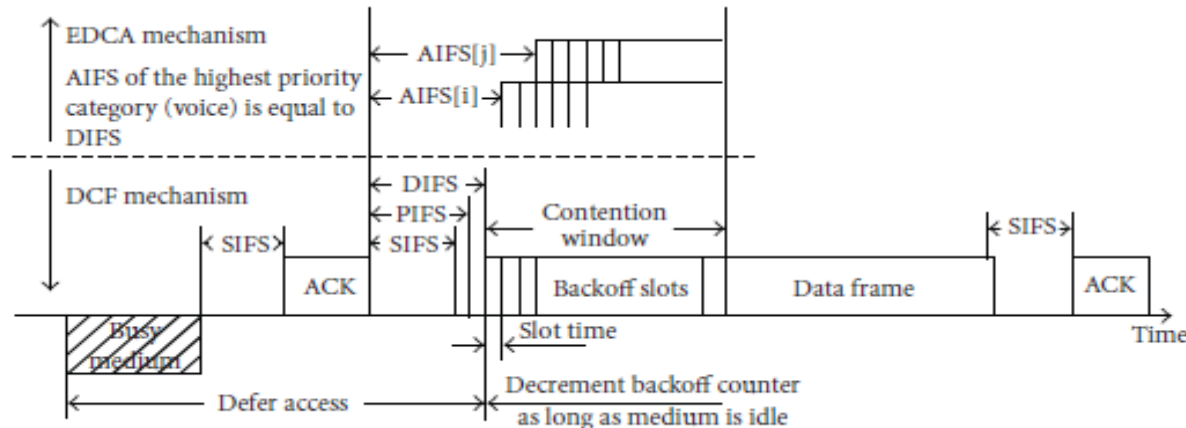
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- **802.11e Improvements to MAC**
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802.11e improves MAC layer

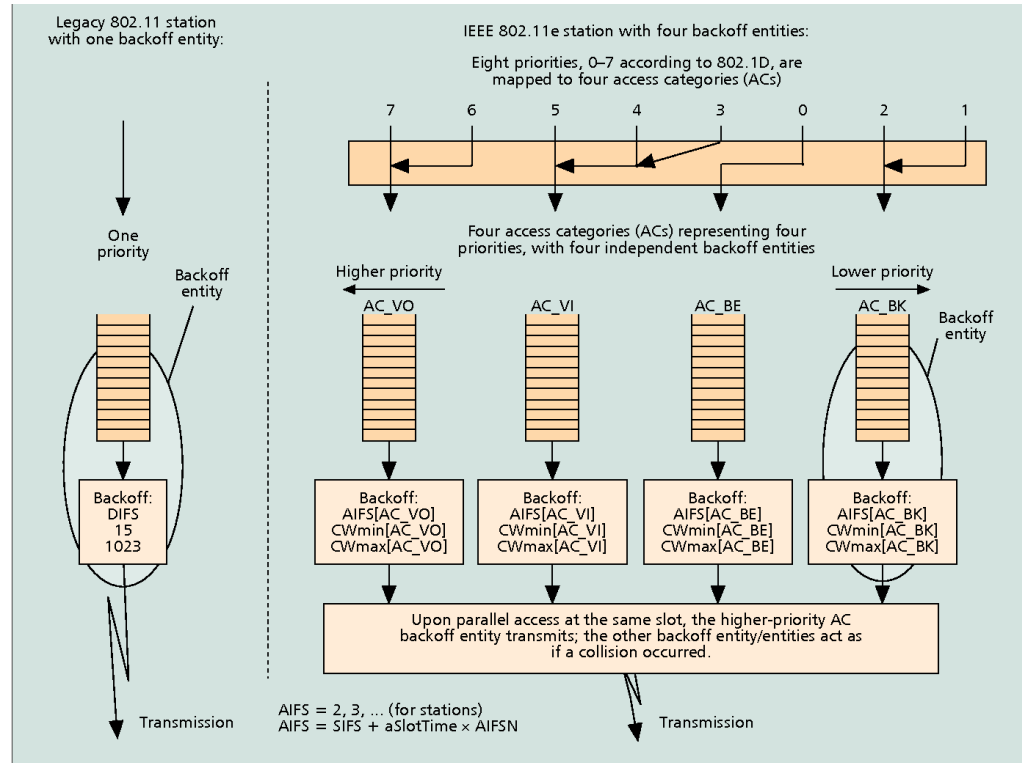
- Hybrid Coordination Function (HCF)
 - Modifies contention-free access (still no one uses it)
 - Modifies contention-based access: Enhanced Distributed Channel Access (EDCA)
- EDCA: Modifies Quality of Service based on application
 - Example of breaking layering for an optimization
 - Categories (lowest to highest priority):
 - Background
 - Best Effort
 - Video
 - Voice

Different priority for different application category

- Expand to more IFS lengths for different traffic categories
 - Smallest AIFS (equal to DIFS) goes to Voice, Largest to Background
 - Contention Window min and max also change for each category
 - Selects a *probability* that most important category goes first



Multiple queues within a single device



■ Figure 4. [3] Legacy 802.11 station and 802.11e station with four ACs within one station.

802.11e also adds maximum durations

- 802.11e also defines duration a device can transmit for
 - Based on PHY in use and Application category
 - Background/Best Effort: one frame per contention win
 - Example, up to 11 ms for Voice on 802.11ac
 - Could be one really big frame at a low data rate
 - Could be multiple frames in a row separated by SIFS

Outline

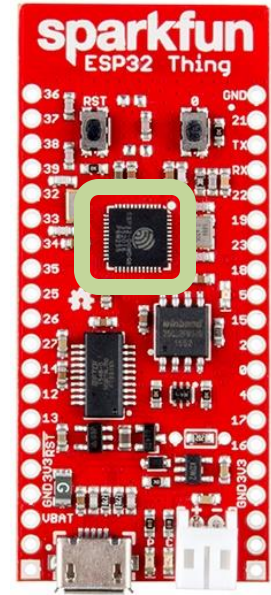
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- **Microcontrollers and WiFi**

Why, why not, talk WiFi in when thinking about wireless for resource-constrained computing (aka, “The IoT”)?

- Pros
 - Ubiquitous
 - High-performance
- Cons
 - Complex configuration
 - And security requirements
 - UCSD-DEVICE anyone?
 - Expensive in energy and money

WiFi capability in microcontrollers

- ESP32 (extremely popular; there are others)
 - Microcontroller plus WiFi radio in single chip
 - (Same idea as nRF52840)
- Capabilities
 - 802.11b/g/n 2.4 GHz only
 - 20 MHz or 40 MHz channels
 - Single antenna only (no MIMO)
 - MCS0-7
 - 7 Mbps – 150 Mbps
 - Tx power up to 20.5 dBm



Low power WiFi

- Question: should a microcontroller stay connected or reconnect?
 - Light sleep: stay connected always, only listening to beacons
 - Deep sleep: reconnect to network each time data is ready
- Answer for ESP32 depends on security and data interval
 - Resecuring during connection takes lots of energy
 - Crossover point is about 60 seconds
 - Insecure transmissions have a crossover of 5-15 seconds

<https://blog.voneicken.com/2018/lp-wifi-esp-comparison/#conclusions>

Next Up: Longer-Range Technologies