

① $A \in \mathbb{C}^{m \times n}$, $\text{rank}(A) = n$

$b \in \mathbb{C}^m$

$$\begin{bmatrix} I & A \\ A^* & 0 \end{bmatrix} \begin{bmatrix} r \\ x \end{bmatrix} = \begin{bmatrix} b \\ 0 \end{bmatrix};$$

$$I_{m \times m} \cdot r_{m \times 1} + A_{m \times n} \cdot x_{n \times 1} = b_{m \times 1}$$

$$A^*_{m \times m} \cdot r_{m \times 1} = 0$$

$$\therefore r_{m \times 1} = b_{m \times 1} - A_{m \times n} \cdot x_{n \times 1}$$

$$A^*_{m \times m} (b_{m \times 1} - A_{m \times n} \cdot x_{n \times 1}) = 0$$

$$A^*_{m \times m} A_{m \times n} x_{n \times 1} = A^*_{m \times m} b_{m \times 1}$$

② $A \in \mathbb{C}^{m \times m} \rightarrow$ non-singular;

Prove it has an LU Factorization iff $\forall k \in [1, m]$, the upper left $k \times k$ block is non-singular.

Proof: for a 2×2 matrix, we need, $a_{11} \neq 0$ which is the leading block matrix and we don't want it to be singular. Suppose

$$A \in \mathbb{C}^{m \times m} \quad \begin{bmatrix} a_{11} & a_{12} & \dots & a_{1n} & \dots & a_{mn} \\ | & | & \dots & | & \dots & | \\ a_{21} & a_{22} & \dots & a_{2n} & \dots & a_{mn} \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} & \dots & a_{mn} \end{bmatrix}$$

Suppose $a_{11} = 0$, then any row operations aren't valid to make column entries a_{21} to $a_{m1} = 0$. So the $\det(a_{11}) \neq 0$.

$$L_1 = \begin{bmatrix} 1 & \dots & 0 \\ -\frac{a_{21}}{a_{11}} & 1 & \dots \\ \vdots & \ddots & 1 \\ -\frac{a_{m1}}{a_{11}} & \dots & \dots & 1 \end{bmatrix}; \quad R_1 = \begin{bmatrix} a_{11} & \dots & a_{1m} \\ 0 & \frac{a_{21}}{a_{11}} \cdot a_{22} + a_{21} & \dots & \frac{a_{21}}{a_{11}} \cdot a_{2m} \\ 0 & \frac{a_{31}}{a_{11}} \cdot a_{32} + a_{31} & \dots & \frac{a_{31}}{a_{11}} \cdot a_{3m} \\ \vdots & \vdots & \ddots & \vdots \\ 0 & \frac{a_{m1}}{a_{11}} \cdot a_{m2} + a_{m1} & \dots & \frac{a_{m1}}{a_{11}} \cdot a_{mm} + a_{m1} \end{bmatrix} = L_1 A.$$

Now, for second case, we need $-\frac{a_{21}}{a_{11}} a_{22} + a_{22} \neq 0$.

@) in other words, $-a_{21} a_{11} + a_{11} a_{22} \neq 0$. (if $a_{11} \neq 0$), this is the determinant of the sub matrix $\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}$

$$\text{so, determinant of sub matrix } \underbrace{\begin{bmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{bmatrix}}_{\text{submatrix } A_{1,2,2,1}} = \begin{bmatrix} a_{11} & a_{12} \\ 0 & -\frac{a_{21}}{a_{11}} \cdot a_{12} + a_{22} \end{bmatrix} =$$

$$\text{determinant of } \begin{bmatrix} a_{11} & a_{12} \\ a_{m1} & a_{m2} \end{bmatrix}$$

Suppose we go till column k , then we know that determinant of $A^{(k \times k)} = \det(A_{1:k, 1:k})$; in which case, if the entry, $A_{kk}^{(k \times k)} = \frac{1}{\beta} \det(A_{1:k, 1:k})$, ($\beta \in \mathbb{C} \neq 0$), should not be 0 for us to proceed further.

Suppose LU decomposition isn't unique,

then for any $M \neq I$, $M \in \mathbb{C}^{m \times m}$, $LU = LDD^{-1}U$

where (L, D, U) : $L \cdot U^T$, $L \rightarrow$ lower tr., $U^T \rightarrow$ upper tr.

LD for any non-singular D will not be strictly lower tr unless D is a diagonal matrix, but if its diagonal, we can have L's diagonal to be 1 when diagonal entries of D are $\neq 1$. which means, there exists no $D \neq I$ that can guarantee this and hence, LU decomposition is unique.

③ Bandwidth = $2p+1$

$$p=1 \Rightarrow BW=3$$

$$a_{ij}=0 \text{ for } |i-j|>1 \\ \Rightarrow \text{for } A \in \mathbb{C}^{3 \times 3}, \quad \begin{bmatrix} a_{11} & a_{12} & 0 \\ a_{21} & a_{22} & a_{23} \\ 0 & a_{32} & a_{33} \end{bmatrix} \quad \leftarrow \quad \begin{bmatrix} 1 & 0 & 0 \\ -\frac{a_{21}}{a_{11}} & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \cdot A^{(3 \times 3)} = A^{(3)}$$

$$\text{for } A \in \mathbb{C}^{4 \times 4}, \quad \begin{bmatrix} a_{11} & a_{12} & 0 & 0 \\ a_{21} & a_{22} & a_{23} & 0 \\ 0 & a_{32} & a_{33} & a_{34} \\ 0 & 0 & a_{43} & a_{44} \end{bmatrix} \quad \leftarrow \quad \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & -\frac{a_{21}}{a_{11}} & 1 & 0 \\ 0 & 0 & a_{32} & 1 \end{bmatrix} \cdot A^{(4 \times 4)} = A^{(4)}$$

$$p=2 \Rightarrow BW=5 \quad |i-j|>2$$

$$A \in \mathbb{C}^{5 \times 5}, \quad \begin{bmatrix} a_{11} & a_{12} & 0 & 0 & 0 \\ a_{21} & a_{22} & a_{23} & a_{24} & 0 \\ a_{31} & a_{32} & a_{33} & a_{34} & a_{35} \\ 0 & a_{42} & a_{43} & a_{44} & a_{45} \\ 0 & 0 & a_{53} & a_{54} & a_{55} \end{bmatrix} \quad \leftarrow \quad \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & -\frac{a_{21}}{a_{11}} & 1 & 0 & 0 \\ 0 & 0 & a_{32} & 1 & 0 \\ 0 & 0 & 0 & a_{43} & 1 \end{bmatrix} \cdot A^{(5 \times 5)} = A^{(5)}$$

L, U will have same sparsity pattern of A which means, we can figure out the rows on which row reduction need not be performed and avoid excess flops.

④ a) if A is non-singular, such a factorization always exists.

Prof: We need the $\det(A_{ck, nk}) \neq 0$; so, even if that seems to be the case, this can be avoided through pivoting. We can permute those corresponding rows and we can make sure $AQ=LU$ is formed

ex: $A = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$, LU with pivoting possible.

b) if A is singular; $\begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$ (ex)

In this case, we don't have an LU Factorization at all even with pivoting. Hence, we can't always prove its existence

⑤ a) One way is to augment the matrix as following;

$$AA = [A \mid I] \quad \text{where } I \in \mathbb{C}^{m \times m}, A \in \mathbb{C}^{m \times m}$$

Now performing and reducing AA to its row echelon form until the reduced matrix becomes $AA^* = [I \mid B]$ will generate $A^{-1} = B$.

The number of divisions $\frac{(m+n)(m+n-1)}{2}$ don't change, but the number of multiplications and subtractions are on the rows of AA, which has 2m columns.

$$\text{so, } \textcircled{X} \rightarrow \frac{(2m)^3 + 3(2m)^2 - 5(2m)}{6} / 6 = \frac{8m^3}{3}$$

$$\textcircled{Y} \rightarrow \frac{(2m)^3 + 3(2m)^2 - 5(2m)}{6} / 6 = \frac{8m^3}{3}$$

\rightarrow Should reduce it to order $\sim 2m^3$

c) i) We can augment A to make it AA: $[A \mid B]$ and perform LU until it becomes $AA^* = [I \mid Y]$

Number of columns have become $m+n$

$$\text{so, } \textcircled{X} \rightarrow \frac{(m+n)^3 + 3(m+n)^2 - 5(m+n)}{6} / 6$$

$$\textcircled{Y} \rightarrow \frac{(m+n)^3 + 3(m+n)^2 - 5(m+n)}{6} / 6$$

$$\hookrightarrow \frac{2(m+n)^3}{6} \approx \frac{(m+n)^3}{3}$$

ii) Each entry of $A^{-1}b_i$ will have $2m$ flops.

(m, \textcircled{X}) and $m(\textcircled{Y})$. There are m entries as well

addressing $2m$, m times $\approx 2m^2$

Then we need to do this ' n ' times for this system of equations \Rightarrow Plops $\approx 2m^2n$

To find A^{-1} , we'll have $\approx \frac{8m^3}{3}$ flops

\therefore total operation count $\approx \frac{8m^3}{3} + 2m^2n$