

{ w | w contains twice as many 0's as 1's }

Ex: 011000

011000 Find a "1"
011000 change 1 to X
0X1000 move L
0X1000 look for left edge (HASH MARK)
0X1000 scan right for a zero
0X1000 change 0 to a Y
YX1000 scan right for second zero
YX1000
YX1000 change 0 to a Y
YX1000 scan left to X
YX1Y00 find a "1"
YX1Y00 chg to X
YXXY00 scan for a Y
YXXY00 scan right for a zero
YXXY00 → FIND 'L', FIRST → REJECT
YXXY00 chg to X
YXXY00 scan right for second 0
YXXY00 chg to Y
YXXYY scan left for X
YXXYY scan right for 1
YXXYYY NO 1 FOUND
YXXYYY look for remaining zeros
YXXYYY scan for # AND SEE ONLY X'S & Y'S
YXXYYY See a # → reject
YXXYYY See the # → accept

$$(q_0, \#) \rightarrow (q_1, \#, R)$$

$$(q_1, \circ) \rightarrow (q_1, \emptyset, R) \quad (q_1, \sqcup) \rightarrow (q_1, \sqcup, \sqcup)$$

$$(q_1, \top) \rightarrow (q_2, \times, \sqcup)$$

$$(q_2, \#) \rightarrow (q_3, \#, R)$$

$$(q_2, Y) \rightarrow (q_3, Y, R)$$

$$(q_2, \emptyset) \rightarrow (q_2, \emptyset, \sqcup)$$

$$(q_2, \top) \rightarrow (q_2, \top, \sqcup)$$

$$(q_3, \times) \quad \} \rightarrow (q_3, \frac{\times}{c}, R) \quad \text{HALT}$$

$$Y \quad \}$$

$$(q_3, \emptyset) \rightarrow (q_4, Y, R)$$

$$(q_3, \sqcup) \rightarrow \text{REJECT} \rightarrow (q_r, \sqcup, \sqcup)$$

$$(q_4, \times) \rightarrow (q_4, \times, R) \quad (q_4, \sqcup) \rightarrow \text{REJECT}$$

$$(q_4, \emptyset) \rightarrow (q_5, Y, \sqcup)$$

$$(q_5, \emptyset) \rightarrow (q_5, \emptyset, \sqcup)$$

$$(q_5, \times) \rightarrow (q_1, \times, R)$$

$$M_{\text{2x}} \{ Q = \{ q_0, \dots, q_n \} \}$$

$$\Sigma = \{ \sigma_1, \dots, \sigma_k \}$$

$$\Gamma = \Sigma \cup \{ x, y \} \cup L$$

S = See above

$$q_0 = q_0$$

$q_{\text{accept}} = q \in Q$ is the accept state

$q_{\text{reject}} = q \in Q$ is the reject state

$$q_{\text{accept}} \neq q_{\text{reject}}$$

$$1 \rightarrow x_j L$$



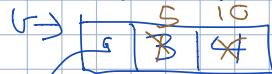
Rust - Collections, Iterators, mutability, and reference counting

Collections

- VECTORS

- SIZE EXPAND GROW ON RIGHT
- Composed of ANY TYPE or object (same type)

LET $v = \text{VEC}([0, 1, 2])$



$[0 | 1 | 2 |]$

$v.push(12);$

$v.push(13);$

$[0 | 1 | 2 | 13 |]$

LET $v2 = \text{VEC}(0..10);$

- deque

↳ FIFO QUEUE

↳ Add to end, pull from other end

- HASHMAP

- key / value pairs

- LET $h = \text{HashMap::new}(<\text{String}, \text{u32}>());$

- $h.insert("key struc", 42);$

ITERATOR :

- WALKING THROUGH A COLLECTION

LET $\sigma = \text{vec}([0, 1, 2])$ LET $h = \text{Hashmap}.\text{New}(\dots)$

FOR i IN σ {

FOR (k, v) IN h {

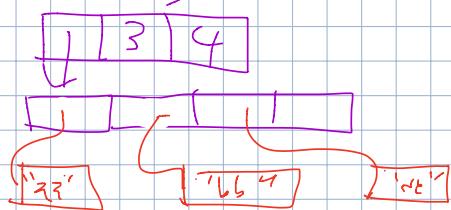
}

}

MUTABILITY

LET $\sigma = \text{vec}(['a', 'b', 'c'])$

LET MUT $v = \&v$;



INTERIOR MUTABILITY

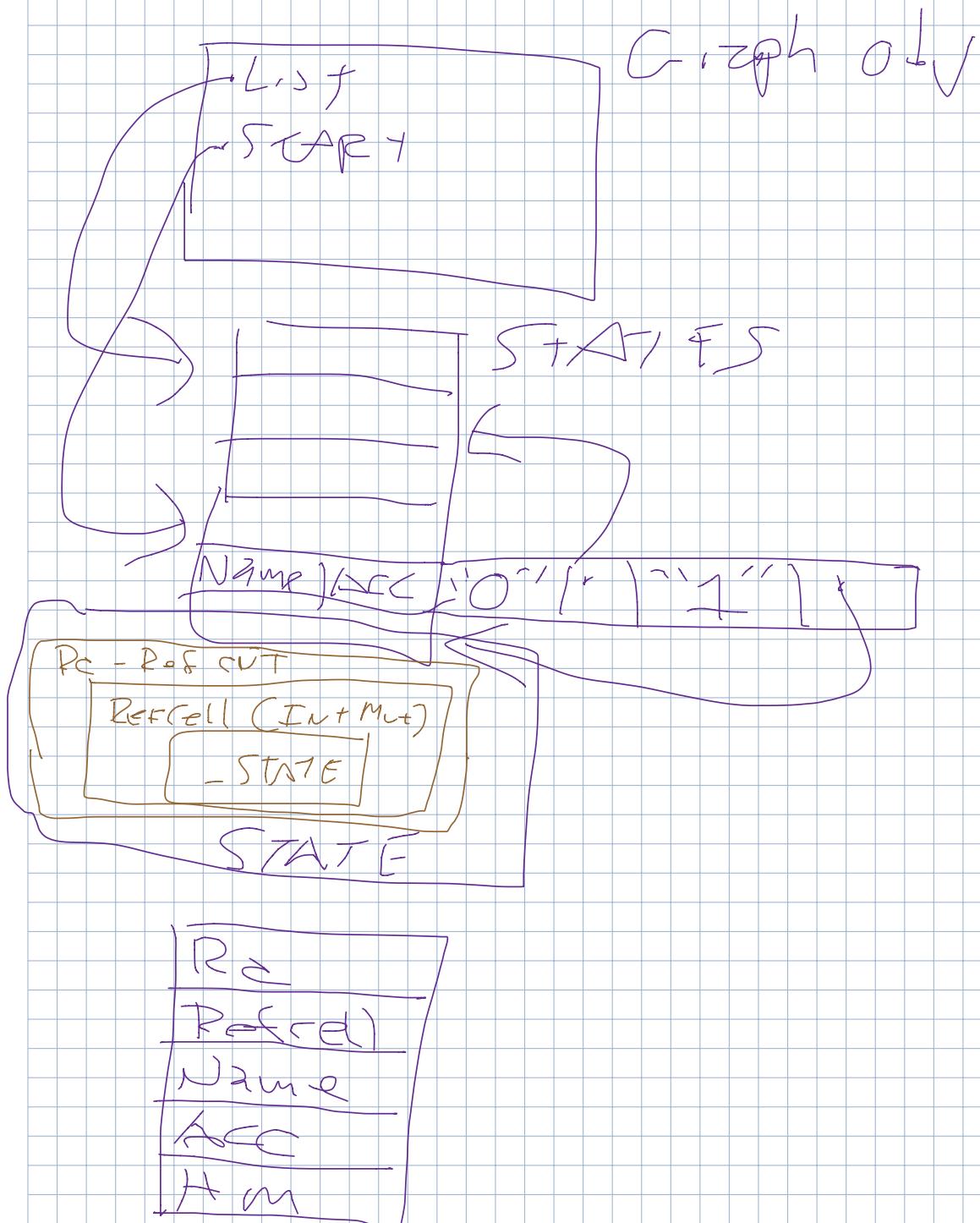
- MUTABLE REFERENCE TO SOMETHING INSIDE AN IMMUTABLE OBJECT.

REFERENCE COUNTED REFERENCES

- MULTIPLE MUT/IMMUT REFERENCES TO SOMETHING

LIST OF STATES

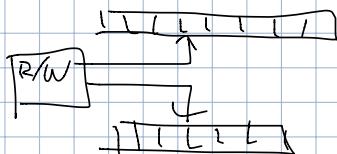
- START STATE \rightarrow ONE STATE
- MARK STATES AS ACCEPT STATES
- TRANSITION STATED AS PREFERENCES



VARIANTS OF TM'S

- MULTIPLE TAPES

$$\delta(q_h, \Gamma_i) \rightarrow (q_{h+1}, \Gamma_i, Y_k)$$



$$\delta(q_0, \Gamma_i, \Gamma_j) \rightarrow$$

$$(q_{n+1}, \Gamma_i^{\tau_1}, \Gamma_j^{\tau_2}, \dots)$$

$$(q_{n+1}, \Gamma_{i+1}^{\tau_1}, \Gamma_{j+1}^{\tau_2}, \dots)$$

- Provable that any multi-tape TM can be expressed as a single tape TM.

- ENUMERATOR

- 2 TAPES (CR mode)

- ONE READ/ WRITE

- WRITE ONLY TAPE

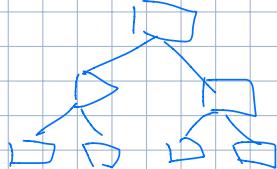
$$L = \{a^n b^n c^n | n \geq 0\}$$

a b c
a a b b c c
a a a b b b c c c

Non-Deterministic T.M.

$$\delta(q_n, \Sigma) \rightarrow (q_{n+1})$$

NFA: $\delta(q_n, \Sigma) \rightarrow \{(q_{n+1}), (q_{n+2})\}$



$$\delta(q_n, \Gamma_n) \rightarrow \{(q_{n+1}, \Gamma_{n+1}^{1/p/s}), (q_{n+2}, \Gamma_{n+1}^{2/p/s})\}$$

- Tree

- At least one branch accepts \Rightarrow accept

- True Decidable

- All possible branches halt

$$L = \{\partial^m \mid m \text{ is a compound number}\}$$

Compound number has at least 2 factors not involving 1
- m is not prime

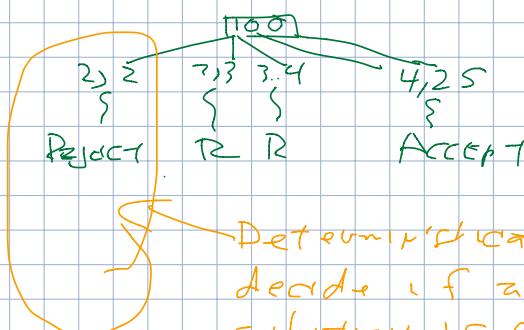
NTM - is m a compound number

- Select 2 El's p & q

- Deterministic T.M to see if p & q are factors of

- If so ACCEPT

- Otherwise REJECT



Deterministically decide if a single solution is correct