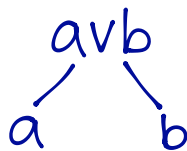


## 5. Truth Trees

- A **truth tree** is an alternative structure for examining all the ways that a compound proposition, say  $X$ , can be true.
- **Unlike a truth table**, the size of a **truth tree** does not grow exponentially as a function of the number of propositional variables in  $X$ ; instead, the size of a truth tree varies depending on the number of logical connectives in  $X$ , and the order in which we **grow the tree**.
- We place  $X$  at the **root** of a truth tree: the **root** is at the top of the tree and the rest of the tree **"grows"** down from there using **branching rules**.

### BRANCHING RULES FOR TRUTH TREES (A.K.A. SEMANTIC TABLEAUX)

#### Splitting Rule



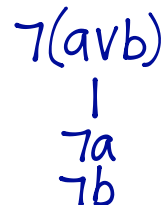
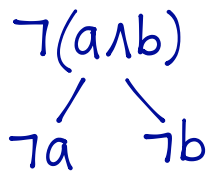
#### Non-Splitting Rule



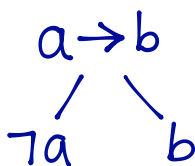
#### Splitting Rule

#### (De Morgan's Laws)

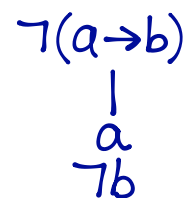
#### Non-Splitting Rule



#### Splitting Rule



#### Non-Splitting Rule



#### (Implication Law)

## One More Non-Splitting Rule

$$\neg\neg a$$

$$\downarrow$$

$$a$$

(Double Negation Law)

## Two More Splitting Rules

+ one unofficial Splitting Rule

$$a \leftrightarrow b$$

$$\swarrow \quad \searrow$$

$$a \quad \neg a$$

$$b \quad \neg b$$

$$\neg(a \leftrightarrow b)$$

$$\swarrow \quad \searrow$$

$$a \quad \neg a$$

$$\neg b \quad b$$

$$a \oplus b$$

$$\swarrow \quad \searrow$$

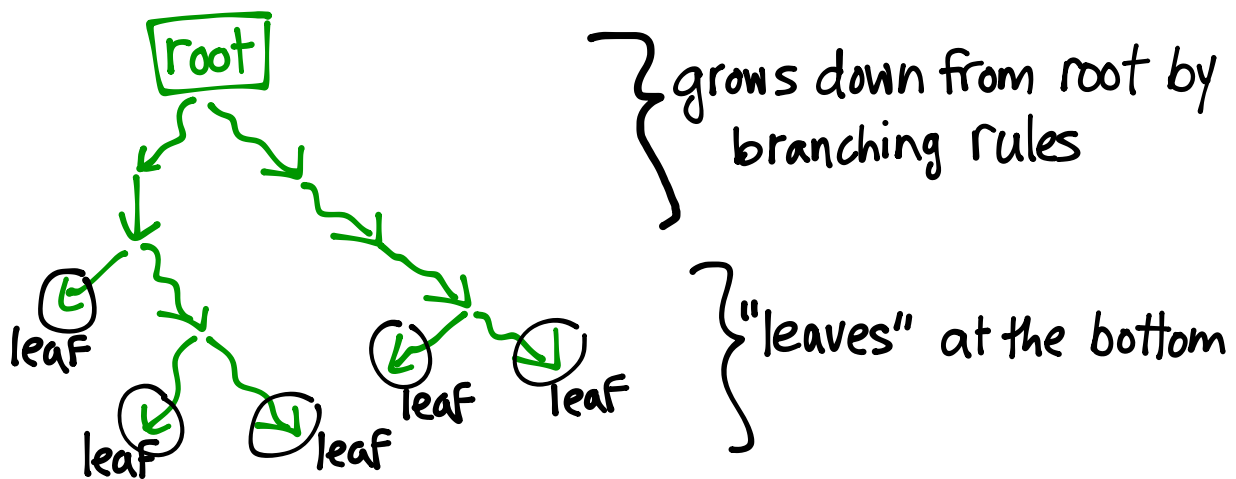
$$a \quad \neg a$$

$$\neg b \quad b$$

(Biconditional Law  $a \leftrightarrow b \equiv (a \wedge b) \vee (\neg a \wedge \neg b)$ )

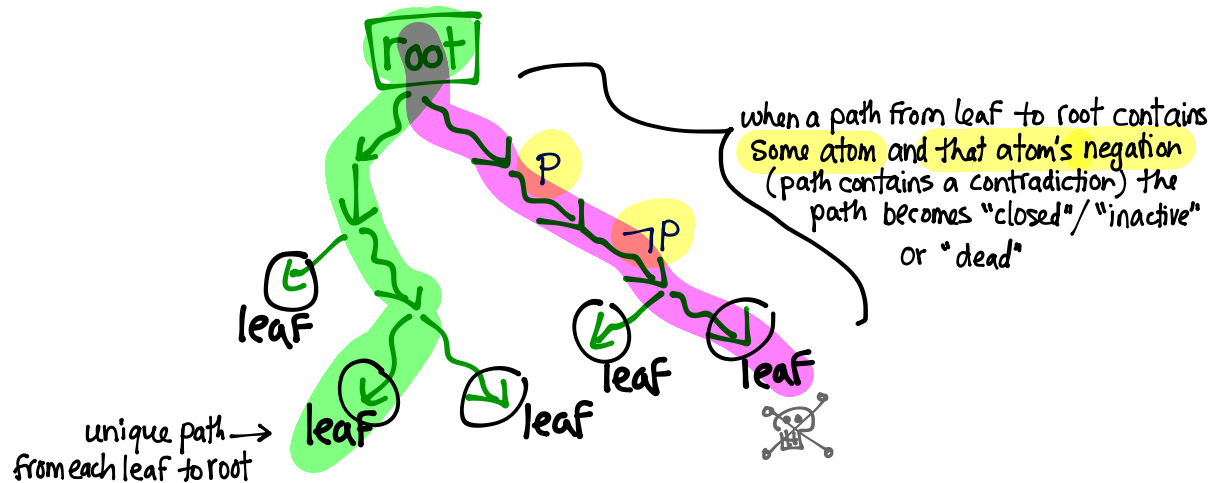
## HOW TO GROW A TRUTH TREE

- Truth trees grow **from the root** (at the top), **by branching rules**, **down to the leaves** (at the bottom).



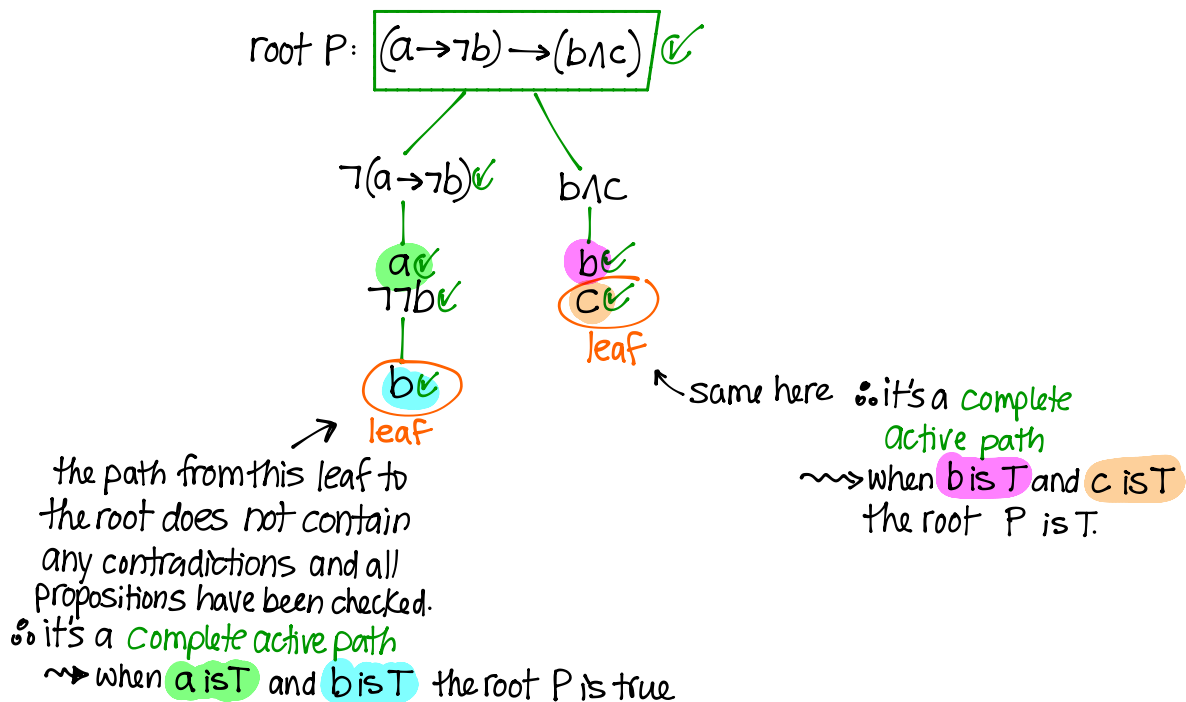
- Each proposition in a truth tree is called **unchecked** until
  - it is a literal (just an atom or just the negation of an atom), or
  - its branching rule has been applied to all paths stemming down from the proposition's location in the tree.
- Starting from the top, one unchecked proposition at a time, apply the branching rule of each unchecked proposition to all paths that stem down from the unchecked proposition.
- Once the branching rule has been applied, the proposition becomes **checked** ✓

- ▷ **Fact.** From each **leaf** (at the bottom of the tree so far), there is a unique path going back up to the root.



- ▷ A path from a leaf back to the root is called **inactive** or **closed** if it contains an atom as well as that atom's negation; otherwise the path is called **active** or **open**.
- ▷ A path from a leaf back to the root is called **complete** if there are no unchecked propositions on that path.
- ▷ Each **complete open path** tells us one way that will make the root true.
- ▷ The tree is **complete** (i.e. **done growing**) when each path from leaf to root is **closed** or **complete**.

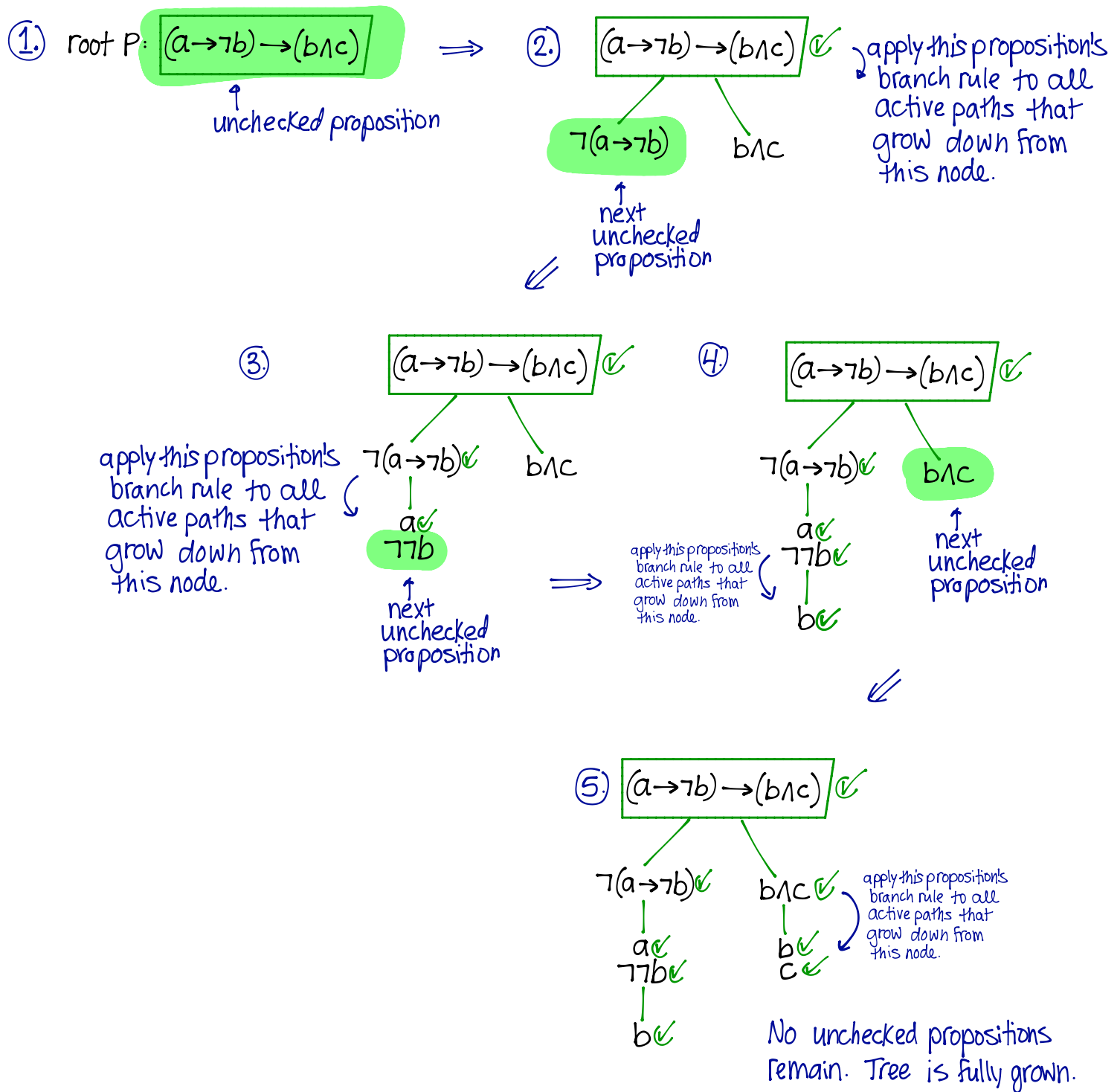
**Example 5.1.** Grow a complete truth tree for the compound proposition  $P : (a \rightarrow \neg b) \rightarrow (b \wedge c)$ .



In summary, the complete active paths tell us that  $P$  is T when

$a \wedge b$  is T or when  $b \wedge c$  is T i.e.  $P \equiv (a \wedge b) \vee (b \wedge c)$  ← what is this?!  
 a DNF for  $P$ !

Here is Example 5.1 again, showing one step at a time:



## FINDING A DNF FOR THE ROOT OF A TRUTH TREE

To summarize:

- To find a DNF for a compound proposition  $X$ , we grow a complete truth tree with  $X$  at its root.
- Each complete open path gives one **conjunctive clause** consisting of the conjunction of the literals found on that path going from the leaf back up to the root.
- The disjunction of all such conjunctive clauses gives a DNF for  $P$ .

## DETERMINING WHETHER THE ROOT OF A TRUTH TREE IS A CONTRADICTION

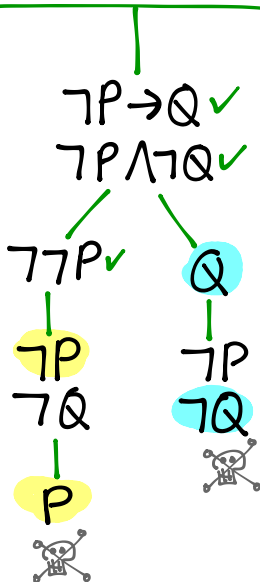
**Question to Ponder.** Suppose we grow a complete truth tree with  $X$  at its root. What does it mean if **all paths are closed/inactive** ("dead") ?

It means the root can never be true.

i.e. the root is a contradiction.

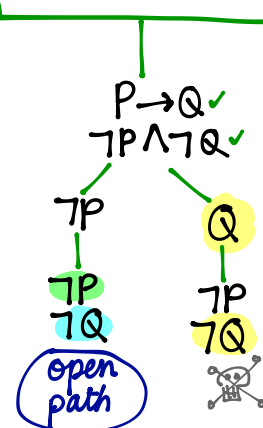
**Example 5.2.** Use a truth tree to determine whether each of the following compound propositions is a contradiction. If it is not a contradiction, give all counterexamples (that is, give all truth assignments that make the proposition true, thereby certifying that the proposition is not a contradiction).

i.  $(\neg P \rightarrow Q) \wedge (\neg P \wedge \neg Q)$  root



Since all paths are "dead" the root is never True. ∴ the root is a contradiction

ii.  $(P \rightarrow Q) \wedge (\neg P \wedge \neg Q)$  root



Since there is an open path, the root can be True. ∴ the root is not a contradiction.

Counterexample (to prove root is not a contradiction)  
• When  $P=F$  and  $Q=F$ , the root is True

## DETERMINING WHETHER A SET OF COMPOUND PROPOSITIONS IS CONSISTENT

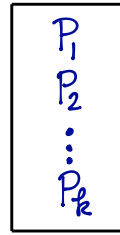
- A set  $\{P_1, P_2, \dots, P_n\}$  of compound propositions is called **consistent** if there exists at least one truth assignment that makes all the propositions  $P_1, \dots, P_n$  true at the same time.
- If  $\{P_1, P_2, \dots, P_n\}$  is not consistent, then, for each possible truth assignment, at least one of the propositions  $P_i$  is false; in this case, the set  $\{P_1, P_2, \dots, P_n\}$  is called **inconsistent**. In this case, the conjunction  $P_1 \wedge \dots \wedge P_n$  is a contradiction.

Equivalently, a set  $\{P_1, P_2, \dots, P_n\}$  is...

- ▷ **consistent** if the conjunction  $P_1 \wedge \dots \wedge P_n$  is *not* a contradiction.
- ▷ **inconsistent** if the conjunction  $P_1 \wedge \dots \wedge P_n$  is a contradiction.

♠ To test whether the set  $\{P_1, P_2, \dots, P_k\}$  is consistent

- grow a tree with **root**

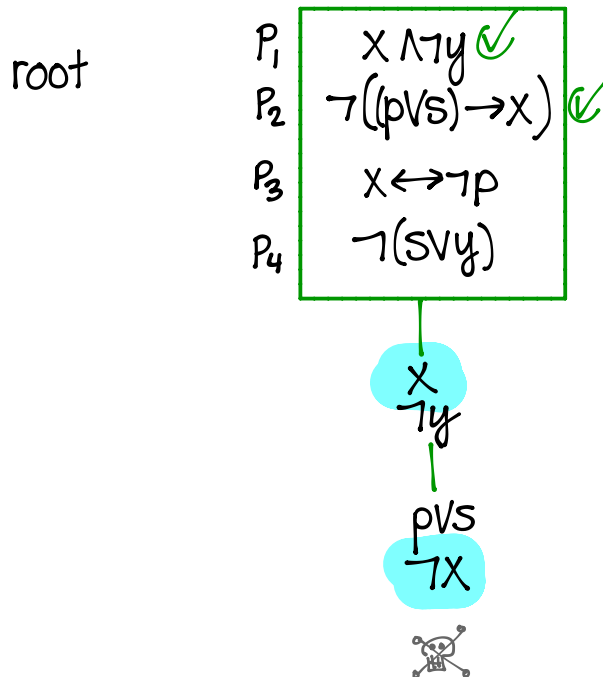


← this root corresponds to non-splitting rule for the conjunction  $P_1 \wedge \dots \wedge P_k$

- if all paths are dead, then  $(P_1 \wedge P_2 \wedge \dots \wedge P_k)$  is a contradiction, which means  $(P_1 \wedge P_2 \wedge \dots \wedge P_k)$  can never be true, which means  $\{P_1, P_2, \dots, P_k\}$  is inconsistent.
- if one path (or more) is alive (complete open path), then these paths give the truth assignments that prove that  $(P_1 \wedge P_2 \wedge \dots \wedge P_k)$  can be true, which translates to showing that  $(P_1 \wedge P_2 \wedge \dots \wedge P_k)$  is consistent.

**Example 5.3.** Use a truth tree to determine whether the following set of four compound

propositions is consistent:  $\left\{ \underbrace{x \wedge \neg y}_{P_1}, \underbrace{\neg((p \vee s) \rightarrow x)}_{P_2}, \underbrace{x \leftrightarrow \neg p}_{P_3}, \underbrace{\neg(s \vee y)}_{P_4} \right\}$

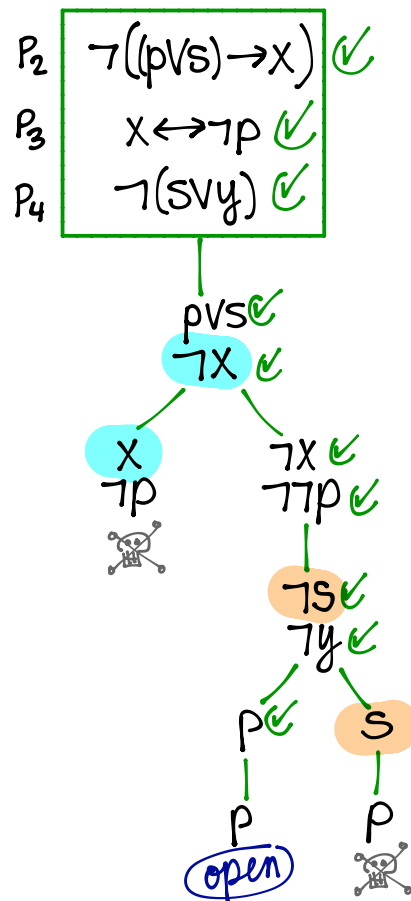


The tree's only path so far is already dead, so the root (which corresponds to the conjunction  $P_1 \wedge P_2 \wedge P_3 \wedge P_4$ ) is a contradiction

∴ the set  $\{P_1, P_2, P_3, P_4\}$  is inconsistent.

**Exercise 5.4.** Consider the same propositions  $P_1, P_2, P_3, P_4$  from Example 5.3. Determine whether the set  $\{P_2, P_3, P_4\}$  is consistent, first, using a truth tree, then using a truth table.

root



Since there exists at least one complete active path the root can be true

∴ the set  $\{P_2, P_3, P_4\}$  is consistent.

The open path tells us a truth assignment for which  $P_2 \wedge P_3 \wedge P_4$  is true, namely, when  $P=T, y=F, s=F, x=F$

**Exercise** Verify the above answer using a truth table.

## STUDY GUIDE

### Important terms and concepts:

- ◇ truth trees (semantic tableaux) branching rules open vs. closed paths
- ◇ using a truth tree to find a DNF for a given proposition
- ◇ using a truth tree to check whether a proposition is a contradiction
- ◇ using a truth tree to determine whether a set of propositions is consistent/inconsistent

Exercises	Sup.Ex. §1 # 4b, 7b
	Sup.Ex. §1 # 1 Is $\{P_1, P_2, P_3, P_4\}$ consistent? Is $\{P_1, P_2, P_3\}$ consistent? Is $\{P_2, P_3\}$ consistent?
	Rosen §1.2 # 9, 11