$(n-1) (1 \text{ remove + 1 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 1 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 2 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 2 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 2 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 2 add}) = 2 \times (n-1)$ $(n-1) (1 \text{ remove + 2 add}) = 2 \times (n-1)$

B1 $\begin{cases} n-1 > 1 \end{cases}$ B2 $\begin{cases} n-1 > 1 \end{cases}$ B2 $\begin{cases} n-1 > 1 \end{cases}$ B1 $\begin{cases} n-1 > 1 \end{cases}$ B2 $\begin{cases} n-1 > 1 \end{cases}$ B1 \begin{cases}

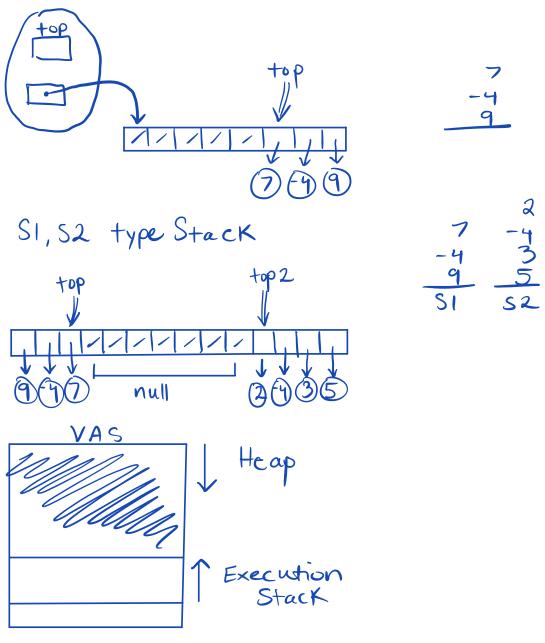
Stack ADT (LiFo Structure) Linear Collection w/ all approaches boolean is Empty() Object peck() throw... Void push() throw... Object pop () throw ... With list implementation... OR Peek = coll. get (coll. size() - 1); peek = coll. get (o); push = coll. add(coll. size()); or push = coll. add(0); or pop = coll. remove (coll.size()-1); pop = coll. remove (0); Stack ADT Implementations: Array Based O. Stack RA

```
Public stackRA()

( items = new Object[3];
    top = -1;
is Empty \rightarrow return top == -1;
peek -> if non empty return items[top]; if empty, throw exception
push (object item)
  if (top == items.length -1)

{
resize (); //resize from old lab
      items[++top] = item;
5
Object Pop ()
S if non empty
  result = items[top]
items[top--] = null;
  return result;
```

Array Based O. Stack SLS ---> return top == null; -> if non-empty, return top.getItem(); → top = new Node (item, top); push > if non-empty, result=top.getItem(); top = top. get Next(); pros: simple, no additional Stack RA Cons: copy upon resizing pre-allocation of mem Pros: on demand allocation, Still simple Cons: double amount of mem/item



Generic Classes ArrayList<String> list; Hashmap<Integer, String> hmp;

Stack RA, Stack SLS, Node

All occurrences of Object type replace
with T.

Except for array allocation

T[] items;

items = (T[]) new Object [3]; //don't replace
here

Use Node (T)