## Algorithmics 3 Assessed Exercise Status and Implementation Reports

# Ryan Wells 1002253w

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### Status report

Based upon the results found in the empirical results at the end of this document I conclude that both my Dijkstra and Backtrack search algorithms are fully functional and operate in expeted runtimes.

### Implementation report

Below I show the implementation of the creation of the graph. Both variations of the algorithm use this creation method.

**Algorithm 1:** One constructor of a Graph

(a) I implemented the Dijkstra search algorithm based on the pseudo-code provided in the lecture notes, my exact implementation is below. I decided not to use the boolean Visited inside the Vertex class and opted for a faster method of using an array to store what Vertices have been visited.

```
input: A file name f containing data about the Graph and path, assumed to be correct
output: On standard output: Distance from float to sink, path and computation time
   fileScanner \leftarrow Scanner for f;
    nodeTotal \leftarrow next line of fileScanner;
   dGraph \leftarrow Graph(nodeTotal, f);
   float \leftarrow next line of fileScanner;
   sink \leftarrow next line of fileScanner;
   distance \leftarrow array containing the distance from root for each node;
   Visited ← array containing whether or not a node has been visited;
   for i \leftarrow 0 to nodeTotal do
        \mathsf{Visited}[i] \leftarrow 0;
        distance[i] \leftarrow distance from float to <math>vertex[i];
        if distance[i]! = -1 then Predecessor of vertex[i] \leftarrow float;
        else Predecessor of vertex[i] \leftarrow -1;
   end
   Visited[float] \leftarrow 1;
   active \leftarrow false;
   found \leftarrow false;
   if vertex[float] has outward connections in graph then active \leftarrow true;
    while active do
        vertex[i] \leftarrow node \ with \ shortest \ distance \ to \ float \ not \ yet \ visited;
        if no such node exists then
            active \leftarrow false;
            break:
        end
        \mathsf{Visited}[i] \leftarrow 1;
        if vertex[i] == vertex[sink] then
            \mathsf{active} \leftarrow false;
            found \leftarrow true;
        end
        for vertex[j] in dGraph not yet visited do
            if vertex[j] has not been reached yet and can be reached then
                Predecessor of vertex[j] \leftarrow vertex[i];
                distance[j] \leftarrow vertex[i] + distance from <math>vertex[i] to vertex[j];
            else if vertex[j] has been reached previously then
                if distance \ through \ vertex[i] \ to \ vertex[j] < distance[j] \ then
                    Predecessor of vertex[i] \leftarrow vertex[i];
                    distance[j] \leftarrow vertex[i] + distance from <math>vertex[i] to vertex[j];
                end
            end
        end
        active = false;
        if further unvisited nodes can be reached then
            active = true;
        end
   end
   \mathbf{if} \ \mathsf{found} \ \mathbf{then}
        Traverse dGraph from vertex[sink] to vertex[float] using the Predecessor, adding
        each vertex to a stack to reverse order;
   end
end
```

(b) I have implemented the Backtrack search using the pseudo-code given at the end of the assignment handout, however I do some preliminary checks before running the algorithm properly. To get the best path for backtrack, because it is such an exhaustive algorithm, you first call a function called bestPath(float, sink) which will only run the backtrack search if the float and sink nodes are different from the ones relating to the current path or if the current path is null. If this is not qualified then nothing will happen. A get method is needed to access the path, which is stored inside the Graph class. The bestPath is stored inside the Graph class.

I attempted to store the path as a doubly indexed integer array with the x value corresponding to the Vertex ID and the second as the accumulative distance so far so I could add in O(1) and remove in O(1) time. I encountered many problems which may have been solved but I decided that under the time constraints using a LinkedList of AdjListNodes was a more suitable option. Also, I decided to use each Vertex's Visited boolean since the function is recursive. The Backtrack algorithm here is assuming that this is the first time a path has been determined on this Graph, and on the first call of this method the workingPath only contains the Float vertex.

```
input : LinkedList<AdjListNode> workingPath, integer sink
begin
   current \leftarrow node referred to by the head of workingPath;
   for vertex in Adjacency List of current do
       if not yet visited then
            Visit and add to head of workingPath;
           if bestPath == null\ or\ workingPath\ is\ a\ shorter\ path\ than\ the\ bestPath\ then
               if head of workingPath is sink then
                   \mathsf{bestPath} \leftarrow \mathit{deep}\ \mathit{copy}\ \mathit{of}\ \mathsf{workingPath}\ ;
               else
                   Recursively call this algorithm on workingPath;
               end
           end
           Remove and un-visit the head of workingPath
       end
   end
end
To print path, print each Vertex starting from the end of the workingPath;
```

Algorithm 3: Backtrack Search Algorithm

## Empirical results

Below are the results from texting the algorithms on my personal  $3.4\mathrm{GHz}$  computer running Debain Squeeze.

Test Files	Dijkstra	Backtrack
data6.txt	Shortest distance from vertex 2 to vertex 5 is 11	Shortest distance from vertex 2 to vertex 5 is 11
	Shortest path: 2 1 0 5	Shortest path: 2 1 0 5
	Elapsed time: 15 milliseconds	Elapsed time: 15 milliseconds
data20.txt	Shortest distance from vertex 3 to vertex 4 is 1199	Shortest distance from vertex 3 to vertex 4 is 1199
	Shortest path: 3 0 4	Shortest path: 3 0 4
	Elapsed time: 25 milliseconds	Elapsed time: 29 milliseconds
data40.txt	Shortest distance from vertex 3 to vertex 4 is 1157	Shortest distance from vertex 3 to vertex 4 is 1157
	Shortest path: 3 36 4	Shortest path: 3 36 4
	Elapsed time: 54 milliseconds	Elapsed time: 129 milliseconds
data60.txt	Shortest distance from vertex 3 to vertex 4 is 1152	Shortest distance from vertex 3 to vertex 4 is 1152
	Shortest path: 3 49 4	Shortest path: 3 49 4
	Elapsed time: 80 milliseconds	Elapsed time: 470 milliseconds
data80.txt	Shortest distance from vertex 4 to vertex 3 is 1152	Shortest distance from vertex 4 to vertex 3 is 1152
	Shortest path: 4 49 3	Shortest path: 4 49 3
	Elapsed time: 100 milliseconds	Elapsed time: 11472 milliseconds
data1000.txt	Shortest distance from vertex 24 to vertex 152 is 17	Shortest distance from vertex 24 to vertex 152 is 17
	Shortest path: 24 582 964 837 152	Shortest path: 24 582 964 837 152
	Elapsed time: 901 milliseconds	Elapsed time: 35789 milliseconds