## Exercise 1 (Morphological Operations).

In the lecture, we introduced the closing operation  $\bullet$  and the opening operation  $\circ$ . We want to prove some properties of these two operations in this exercise.

- a) Let  $A, B \subset \mathbb{R}^n$ . Show that  $A \circ B \subset A$  and  $A \subset A \bullet B$ .
- b) Prove the following statement: If  $A_1 \subset A_2 \subset \mathbb{R}^n$ , we have

$$A_1 \oplus B \subset A_2 \oplus B$$
 and  $A_1 \ominus B \subset A_2 \ominus B$ .

Use this property and part a) to verify the equations

$$(A \bullet B) \bullet B = A \bullet B$$
 and  $(A \circ B) \circ B = A \circ B$ .

(A) (i) by definition: AOB:= (AOB) BB

Let X & AOB.

Then X + be A for be B, thus X & A

Now let X & AOB = (AOB) BB

Then X = y + b for y & AOB and be B

Because X & A from A > AOB,

X & A from X & AOB

Thus AOB C A

(ii) by difficition: A.B:= (A @B) @B

Ut x = A, thin X = A @ B by difficition

(A@B) @B:= 9x: x + b = A @B for b = B ?

thus x = A @B => x = (A @B) @B

hence, x = A => x = (A @B) @B

and A = A & B

(b)  $A_1 \otimes b := \{a_1 + b : a_1 \in A_2, b \in B\}$  since  $A_1 \in A_2$ =  $A_2 \otimes B$ 

mus A, DB C A2 B

 $A, \Theta B := \{x: x + b \in A, 3 \in \{x: x + b \in A_2\}\} = A_2 \Theta B$  since  $A, \in A$ thus  $A, \Theta B \subseteq A_2 \Theta B$ 

(i) mow (A.B) . B = A.B.

from pur (a), AOBCA

also from part (a), ACAOB

TWIS, ABB < (ABB) .B = (A 0 B) & B

8 408 = (A 0 B) O B

dilating both sides:

((A-B)9B) 0 B = (NOB) (PB

=> (A0B)0B = A0B

量

(ii) mow (A. 6) . B = H. B

 $(A \cdot B) \cdot b = ((A \cdot B)^{c} \cdot B^{c})^{c}$   $= ((A^{c} \cdot B^{c}) \cdot B^{c})^{c}$   $= (A^{c} \cdot B^{c})^{c}$ 

> A.B

## two (A.B) . B = A . B

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Exercise 3 (Discrete Convolution of Finite Sequences).

In this exercise, we want to take a look at the convolution of finite sequences. For infinite sequences from the space

$$\ell^{1}(\mathbb{Z}) = \left\{ x = (x_{n})_{n \in \mathbb{Z}} \, \Big| \|x\|_{1} = \sum_{n \in \mathbb{Z}} |x_{n}| < \infty \right\},$$

the convolution  $*: \ell^1(\mathbb{Z}) \times \ell^1(\mathbb{Z}) \to \ell^1(\mathbb{Z})$  is defined by the formula

$$(x * y)_n := \sum_{k \in \mathbb{Z}} x_k \cdot y_{n-k}$$
 for  $n \in \mathbb{Z}$ 

for all  $x, y \in \ell^1(\mathbb{Z})$ . However, we will mostly deal with finite data in practical cases, so that we want to take a look at suitable convolution methods for finite sequences.

a) Given a finite sequence  $x=(x_1,...x_l)$  with  $L\in\mathbb{N}$ , we can define the zero-padded version  $x^{(0)}\in\ell^1(\mathbb{Z})$  of x as

$$x_n^{(0)} = \begin{cases} x_n & \text{if } n \in \{1, ..., L\} \\ 0 & \text{if } n \in \mathbb{Z} \setminus \{1, ..., L\}. \end{cases}$$

For finite sequences  $x = (x_1, ... x_L)$  and  $y = (y_1, ... y_M)$ , we can perform a convolution by computing the convolution  $x^{(0)} * y^{(0)}$ . Show that we have

$$(x^{(0)} * y^{(0)})_n = 0$$
 for all  $n \in \mathbb{Z} \setminus \{2, ..., L + M\}$ .

$$(x'^{(0)} * y'^{(0)})_{n} = \sum_{v \in \mathbb{Z}} x_{v}^{(0)} \cdot y_{n-v}^{(0)}$$

$$= \sum_{v \in \mathbb{Z}} x_{v}^{(0)} \cdot y_{n-v}^{(0)} + \sum_{v \in \mathbb{Z}} x_{v}^{(0)} \cdot y_{n-v}^{(0)}$$

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but 
$$x_{\nu}^{(o)} = 0$$
 for  $\nu \in \{12, ..., \nu \}$   
 $\tau_{nus}, (x_{\nu}^{(o)} * y_{\nu}^{(o)})_{n} = 0$  for  $n \in \{21, ..., \nu \}$ 

b) In the previous part, we saw that the support of the convolution  $x^{(0)} * y^{(0)}$  is contained in the set  $\{2,...,L+M\} \subset \mathbb{Z}$ . Since we want the support of the convolution to be in  $\{1,...,L+M-1\}$ , we consider the **linear convolution**  $\star$  which is defined as

$$(x \star y)_n := \sum_{k=1}^{L} x_k \cdot y_{n-k+1}^{(0)}$$
 for all  $n \in \{1, ..., L+M-1\}$ .

Prove the following equation:

$$(x \star y)_n = (x^{(0)} * y^{(0)})_{n+1}$$
 for all  $n \in \{1, ..., L + M - 1\}$ .

c) Now assume that we have two finite sequences  $\tilde{x}=(\tilde{x}_1,...\tilde{x}_N)$  and  $y=(\tilde{y}_1,...\tilde{y}_N)$  with the same length. The **circular convolution**  $\tilde{x}\otimes\tilde{y}$  of  $\tilde{x}$  and  $\tilde{y}$  is then defined as

$$(\tilde{x} \otimes \tilde{y})_n = \sum_{k=1}^N \tilde{x}_k \cdot \tilde{y}_{n-k+1}^{(p)},$$

where  $\tilde{y}^{(p)}$  is the periodic continuation of  $\tilde{y}$ , i.e.

$$\tilde{y}_n^{(p)} = \tilde{y}_{(n-1)\%N+1} \quad \text{with} \quad (n-1)\%N = (n-1) \bmod N \qquad \text{for all } n \in \mathbb{Z}.$$

For the sequences  $x=(x_1,...x_L)$  and  $y=(y_1,...y_M)$ , where we might have  $L\neq M$ , we consider the zero-padded versions  $\tilde{x}, \tilde{y}$  of length L+M-1 given by

$$\tilde{x}_n = \begin{cases} x_n & \text{if } n \in \{1, ..., L\} \\ 0 & \text{if } n \in \{L+1, ... L+M-1\} \end{cases} \quad \text{and} \quad \tilde{y}_n = \begin{cases} y_n & \text{if } n \in \{1, ..., M\} \\ 0 & \text{if } n \in \{M+1, ... L+M-1\}. \end{cases}$$

Show that  $x\star y=\tilde{x}\otimes \tilde{y}$  holds in this case, which means that we can compute the linear convolution of x and y by a circular convolution.

consider two general sequences,  $x_n$  and  $y_n$ .

Ut  $x_n = \{x_1, x_2, x_3, x_4\}$  with t = 4  $y_n = \{y_1, y_2, y_3\}$  with m = 3

Lomputing & @ 3:

$$\begin{bmatrix} x_1 & 0 & 0 & x_4 & x_3 & x_2 \\ x_2 & x_1 & 0 & 0 & x_4 & x_3 \\ x_3 & x_1 & x_1 & 0 & 0 & x_4 \\ x_4 & x_3 & x_2 & x_1 & 0 & 0 \\ 0 & x_4 & x_3 & x_2 & x_1 & 0 \\ 0 & 0 & x_4 & x_3 & x_2 & x_1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix} =$$

Thus  $x + y = \hat{x} \otimes \hat{y}$  for our example but can be governfited for  $x_n$  and  $y_n$  of any length L and M.