



Set Data Structures

Part 13

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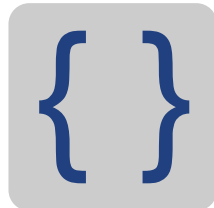
Importance of Sets

Organizing Information

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Importance of Sets

- A *set* is an unordered collection of “objects”
- Sets are used in computer science in a wide variety of ways



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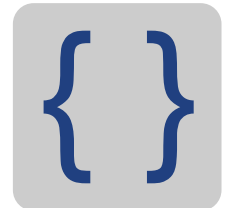
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Importance of Sets

- Depending on the attributes of the set, and how we use it, there are different approaches
- Programmers must choose the best model given how it will be used



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Set Review: Membership

- Set notation uses a special symbols to denote if an object is a member of a set
- Below, the set V contains vegetables

$\text{potato} \in V$
 $\text{bacon} \notin V$

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Set Review: Union

- A union of two sets combines all members of each set into a new one
- So, the result is two merged sets



$$A \cup B = \{ x \mid x \in A \text{ or } x \in B \}$$

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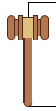
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Set Review: Intersection

- The intersection of two sets contains only those elements that are found in **both** sets
- So, the result is where the two sets overlap



$$A \cap B = \{ x \mid x \in A \text{ and } x \in B \}$$

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Set Review: Difference

- **Difference** (aka exclusion) removes all items found in set from another
- Typically, it is written either $A - B$ or $A \setminus B$



$$A - B = \{ x \mid x \in A \text{ and } x \notin B \}$$

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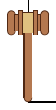
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Set Review: Complement

- The **complement** of a set A , is all elements in the Universe, not in A
- Typically written as A' or A^c
- Alternatively written using an overbar



$$A' = \{ x \mid x \notin A \}$$

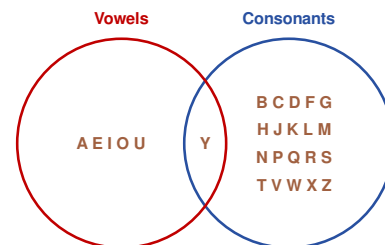
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Example: Letters in English



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Multiset



- A **multiset** is closely related to a set, but permits duplicate elements (as does a tuple)
- The Bag ADT, from the beginning of the semester, supports a basic multiset

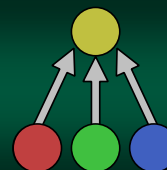
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Union-Find

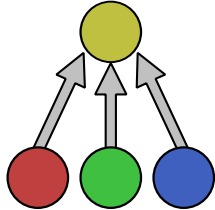


When it absolutely, positively has to be unioned overnight ... er... $O(1)$

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Union-Find

- Often there are situations where we have a number of known objects
- ... i.e. a finite countable set
- ... and we want to quickly union them together



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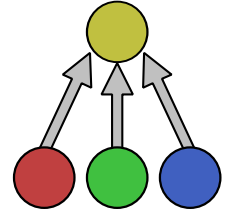
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Union-Find

- Union-Find data structure maintains a list of nodes partitioned into disjoint subsets
- e.g. $\{ \{a\} \{b\} \{c\} \{d\} \{e\} \}$ is a full partition of $\{a,b,c,d,e\}$



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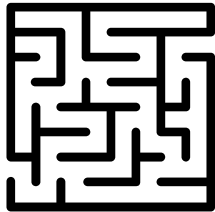
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Example Applications

- Kruskal's Tree-spanning Algorithm – uses a set of known edges
- Maze creation algorithms use it to track sets of connected paths – so the maze is always solvable



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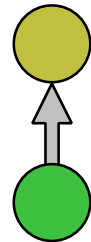
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The Approach

- Sets are stored as a variation of the classic tree
- However, in this approach children link to their parents
- So, the branches point "backwards" from standard trees – i.e. upwards



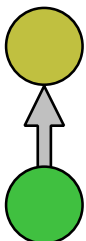
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Arrangement of Nodes



- A parent node can have multiple children - this is not a binary tree!
- Every node - in the tree - is part of the same set

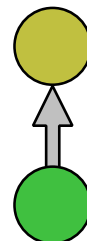
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Arrangement of Nodes



- The root is called the representative of the set
- This node is not special nor is its value special
- ... it just happens to be the node the represents the set

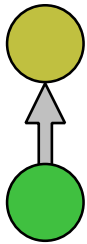
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Arrangement of Nodes



- Any node can "find" its representative – by following the upwards branches
- If any two nodes have the same representative, then they are in the same set

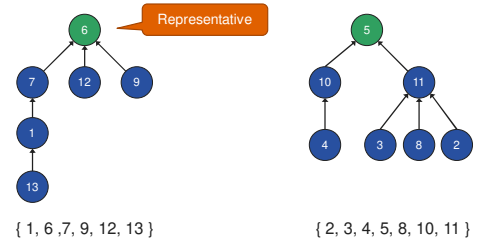
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Examples



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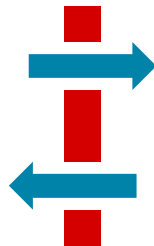
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Union-Find Operations

- The Union-Find data structure contains 3 operations that handle sets
- All 3 can modify the structure of the tree – which automatically optimizes itself into $O(1)$



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Union-Find API

```
public class UnionFind
{
    void makeSet(object value)
    void union(object a, object b)
    object find(object a)
}
```

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Makeset



- Creates a set (within the Union-Find instance) that contains only a single element
- Also known as a singleton
- Essentially, this is the same as an "add" for the ADT

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Find



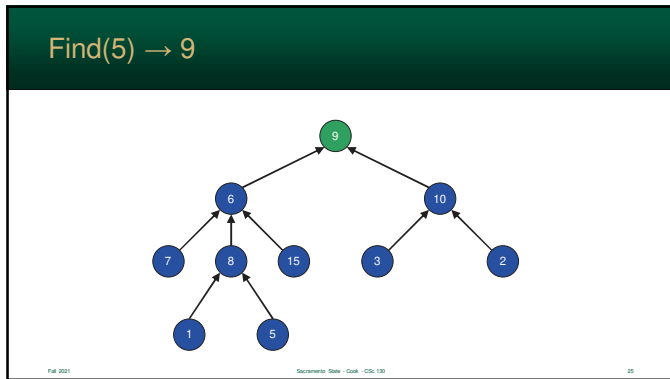
- The find operation starts with a node and locates its representative
- Travels up the tree until it finds a node without a link (which is the representative)

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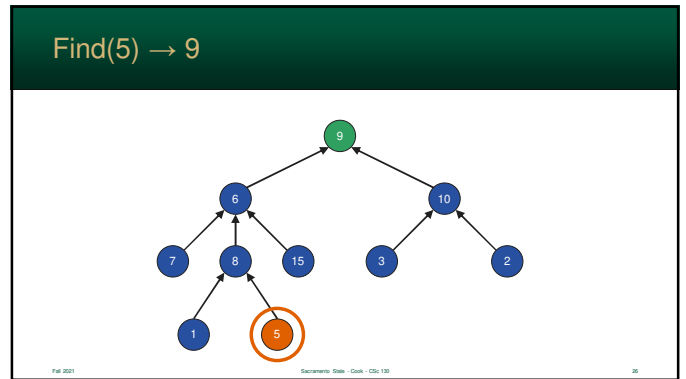
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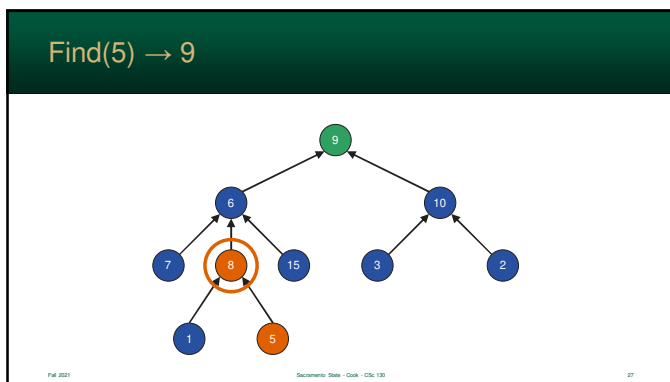
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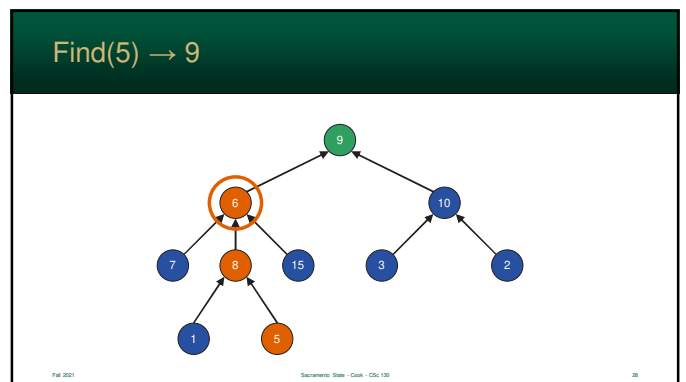
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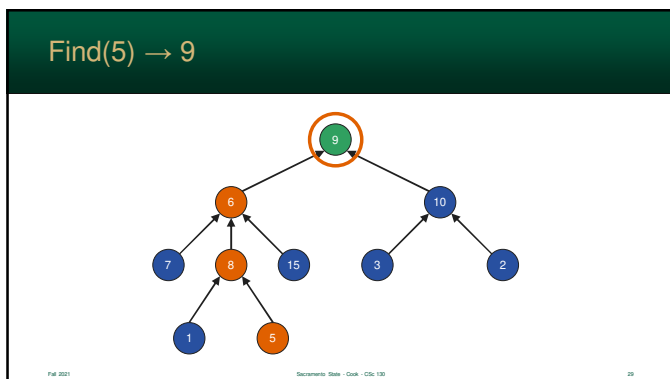
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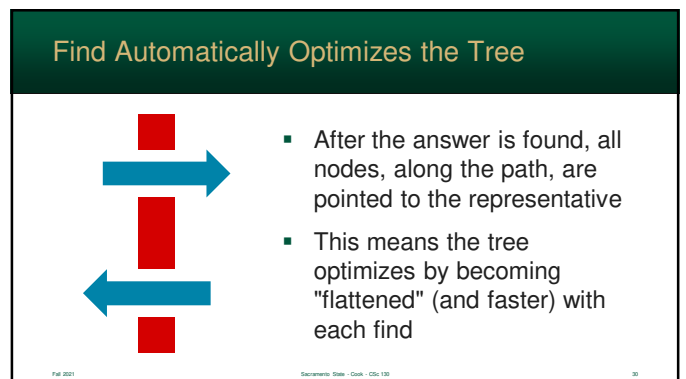
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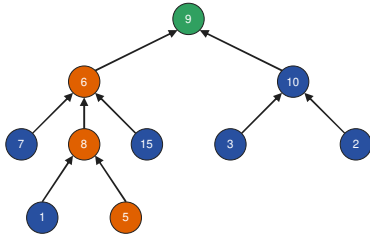


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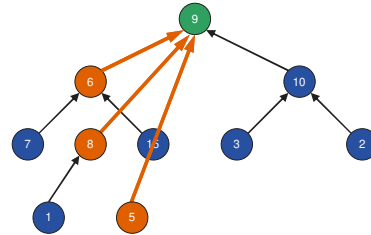
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Connect Each Node to Representative



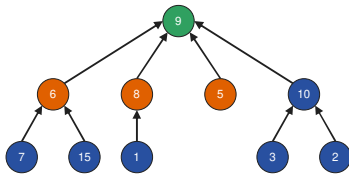
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Connect Each Node to Representative



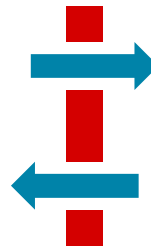
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Updated, Flattened, Tree



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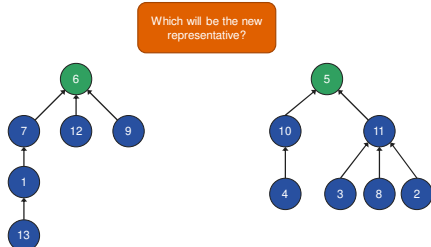
Union



- Combines two sets into a single set
- This is accomplished by updating a single parent link
- So, one representative will be pointed at another representative

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So, which one do we choose?



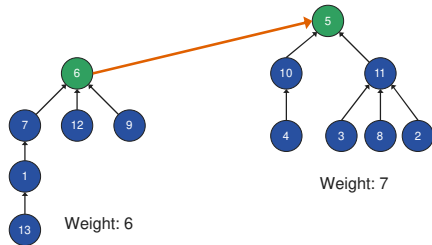
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Approach #1: Union by Weight

- Using *Union by Weight*, the tree with the fewer nodes (**not edge weight**) is made a subtree of the larger tree
- This helps create a more balanced union in terms of weight
- Again, the Find operation will optimize the tree

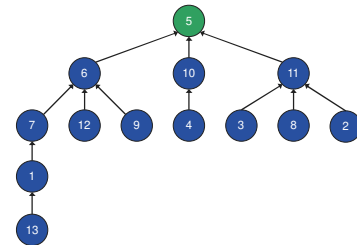
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So, which one do we choose?



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Approach #1: Result



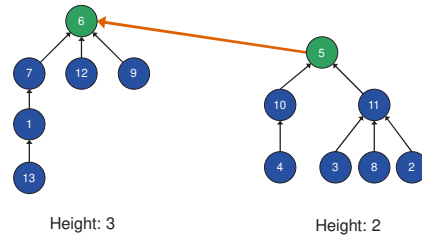
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Approach #2: Union By Height

- Using *Union by Height*, the tree with the smaller height is made a subtree of the taller tree
- This helps create a more balanced union in terms of height
- But, Find, will fix this automatically

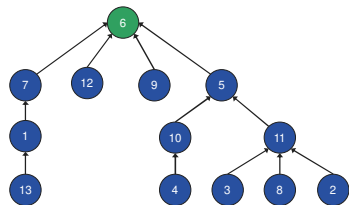
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Approach #2: Union By Height

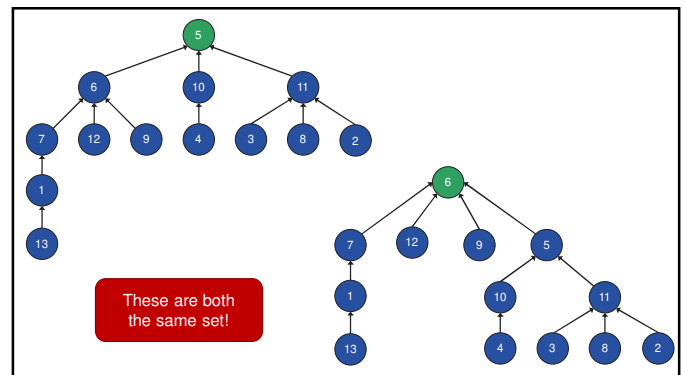


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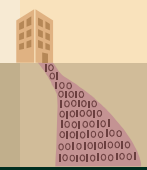
Approach #2: Result



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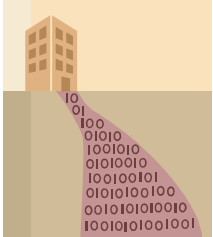
Bit Vectors

Sets and Bits

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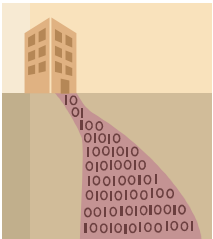
Bit Vectors

- A *bit vector* can store finite, countable sets using bits
- Also known as a *bit array*, *bit set*, and *bit map*
- Compact format that can perform a set operations with a single operation (*fast!*)



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Bit Vectors



- Each object in the universe is represented as a single bit in the string of bits
- If $x \in A$, then the bit is **1**, otherwise **0**

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Example 1

- $U = \{ \text{fry, leela, zoidberg, bender, hermes} \}$
- $A = \{ \text{fry, leela, bender} \}$

U = 11111
A = 11010

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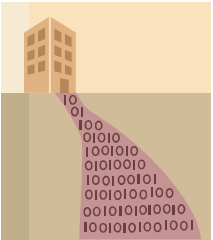
Example 2

- $U = \{ 2, 3, 5, 7, 11, 13, 17, 19 \}$
- $A = \{ 3, 5, 11, 19 \}$

U = 11111111
A = 01101001

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Why this is useful

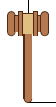


- Computers can easily perform **and** & **or** operations on bytes (or multiple bytes)
- This means set operations can be performed amazingly fast

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Let's look at the definitions again...

- The definitions of union and intersection are nearly identical
- The relationship between the elements is defined using an **AND** or **OR**



$$A \cup B = \{x \mid x \in A \text{ or } x \in B\}$$

$$A \cap B = \{x \mid x \in A \text{ and } x \in B\}$$

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Let's look at the definitions again...

- We can apply **bit-wise-and** & **bit-wise-or**
- It will apply the operation to each of the bits in matching columns



$$A \cup B = \{x \mid x \in A \text{ or } x \in B\}$$

$$A \cap B = \{x \mid x \in A \text{ and } x \in B\}$$

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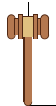
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Let's look at the definitions again...

- So, each bit in **A** will be compared to its matching bit in **B**
- Bit match can do sets!



$$A \cup B = \{x \mid x \in A \text{ or } x \in B\}$$

$$A \cap B = \{x \mid x \in A \text{ and } x \in B\}$$

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Example: Union (using or)

$U = \{a, b, c, d, e, f, g\}$
 $A = \{b, c, d\} = 0111000$
 $B = \{d, e, f\} = 0001110$

0111000
 or 0001110
 $0111110 = \{b, c, d, e, f\}$

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Example: Intersection (using and)

$U = \{a, b, c, d, e, f, g\}$
 $A = \{b, c, d\} = 0111000$
 $B = \{d, e, f\} = 0001110$

0111000
 and 0001110
 $0001000 = \{d\}$

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Complement

- Then, how do we do a complement of a set **A**?
- We must flip all the bits from 1 to 0, and 0 to 1
- We can use a **binary-not** or the **XOR** operation



$$A' = \{x \mid x \notin A\}$$

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Example: Complement (using not)

$U = \{a, b, c, d, e, f, g\}$

$A = \{b, c, d\} = 0111000$

not 0111000
1000111 = $\{a, e, f, g\}$

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Exclusion

- Finally, how do we do set difference?
- The "subtract" operator will not work
- Let's look at the definition a bit more closely

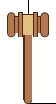


$$A - B = \{x \mid x \in A \text{ and } x \notin B\}$$

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Exclusion

- It's essentially the definition of *intersection*
- Except, the second operand is the definition of *complement*.



$$A - B = \{x \mid x \in A \text{ and } x \notin B\}$$

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Example: Intersection (using and)

$U = \{a, b, c, d, e, f, g\}$

$A = \{b, c, d, f\} = 0111010$

$B = \{d, e, f\} = 0001110$

$B' = \text{not } 0001110 = 1110001$

0111010
and 1110001
0110000 = $\{b, c\}$

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Java/C Code

Intersection : $a \ \& \ b$
Union : $a \ | \ b$
Complement : $\sim a$
Exclusion : $a \ \& \ \sim b$

$\&\&$ and $||$ are
Boolean operators.
These are bit-wise.

The tilde \sim is a
bitwise not.

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