



CSc 28

Discrete Structures

Chapter 9

Recurrence Relations

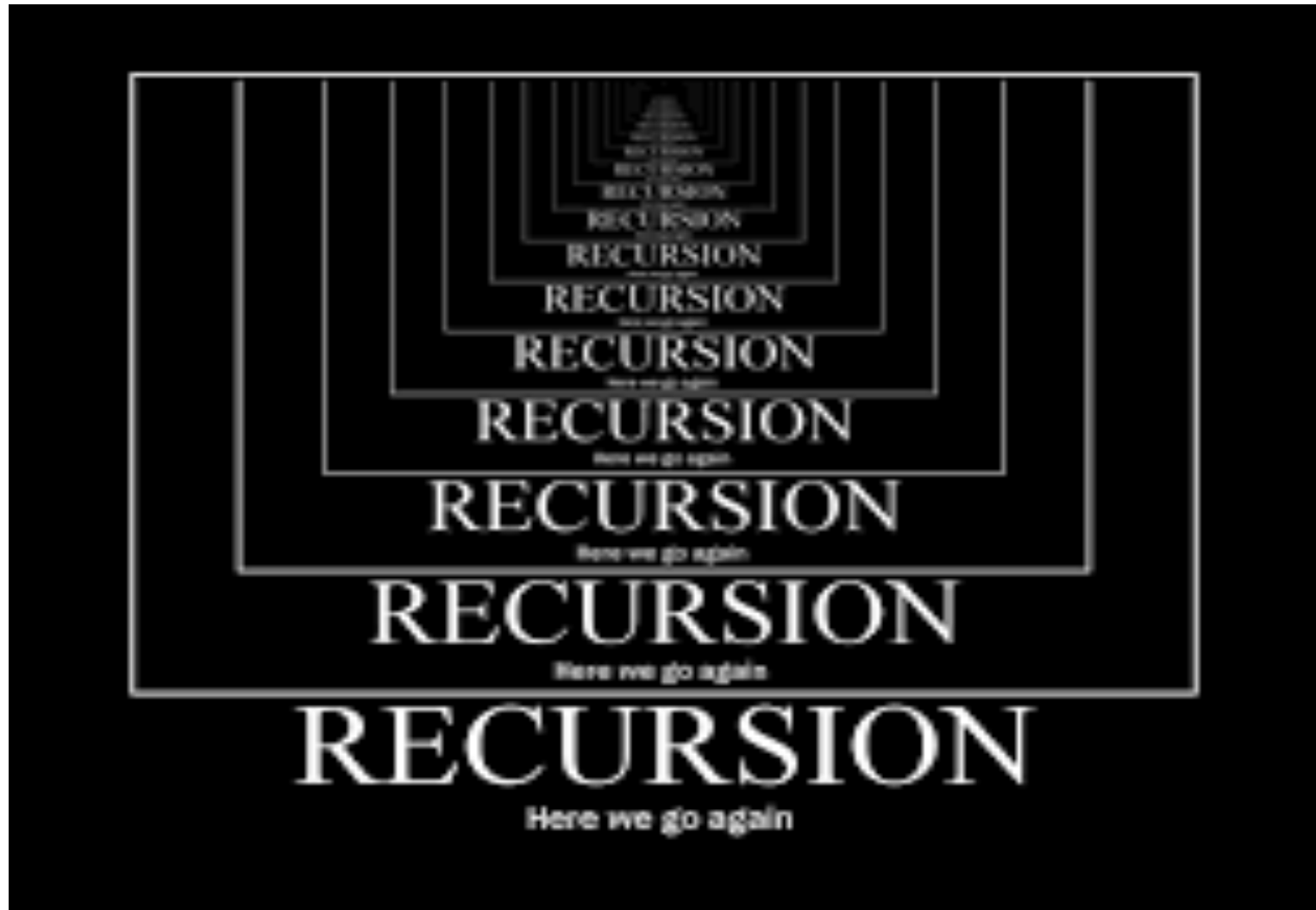
Herbert G. Mayer, CSU CSC
Status 1/1/2021

Syllabus

- **Recursion**
- **Difference**
- **Recurrence Relations (RR)**
- **Modeling with RR**
- **Solving Recurrence Relations**
- **References**

Recurrence Relations & Recursion

Recursion



Difference

- **Recurrence relations** look similar to **recursive** formulae
- **Difference 1:** A Recurrence Relation **do not** provide initial values; Recursive Definition **must** do so!
- **Difference 2:** Recurrence Relation defines mathematical formula, yet Recursive Definitions may define complex algorithms, way beyond complexity of math formulae
- **Difference 3:** Recurrence always **directly** expressed in some formula; recursion may be **indirect**. I.e. some recursive **function f1()** may not have any direct, recursive reference to **f1()** in its program body
- In that case we say the recursion is **indirect**, via some intermediate function **f2()**, or even further functions

Difference

- Example from [1]: Consider the **recurrence relation**
$$a_n = 5 a_{n-1} - 6 a_{n-2}$$
- It is a recurrence relation, not a recursive definition
- Why? Because no initial condition is provided! Hence: Is a **recurrence relation**!
- Once we provide initial conditions, e.g. $a_0 = 1$, $a_1 = 2$, then we get a directly **recursive formula**:
 $a_0 = 1, a_1 = 2, a_n = 5 a_{n-1} - 6 a_{n-2} \text{ for } n \geq 2$
- $a_0 = 1, a_1 = 2, a_2 = 5 a_1 - 6 a_0 = 5*2 - 6*1$ i.e. $a_2 = 4$

And a_3 and a_4 ?

Difference

- Example from [1]: Consider the **recurrence relation**

$$a_n = 5 a_{n-1} - 6 a_{n-2}$$

- It is a recurrence relation, not a recursive definition
- Why? Because no initial condition is provided! Hence: Is a **recurrence relation**!
- Once we provide initial conditions, e.g. $a_0 = 1$, $a_1 = 2$, then we get a directly **recursive formula**:

$$a_0 = 1, a_1 = 2, a_n = 5 a_{n-1} - 6 a_{n-2} \text{ for } n \geq 2$$

- $a_0 = 1$, $a_1 = 2$, $a_2 = 5 a_1 - 6 a_0 = 5*2 - 6*1$ i.e. $a_2 = 4$
- $a_1 = 2$, $a_2 = 4$, $a_3 = 5 a_2 - 6 a_1 = 5*4 - 6*2$ i.e. $a_3 = 8$
- $a_2 = 4$, $a_3 = 8$, $a_4 = 5 a_3 - 6 a_2 = 5*8 - 6*4$ i.e. $a_4 = 16$
- Yields sequence: 1, 2, 4, 8, 16, . . . **Powers of 2. You saw?**

Difference

- Continue analyzing the **recurrence relation**

$$a_n = 5 a_{n-1} - 6 a_{n-2}$$

- If we provide a **different initial condition**, for example values **$a_0 = 1$** , **$a_1 = 3$** , we also get a recursive formula, yet a different one:

$$a_0 = 1, a_1 = 3, a_n = 5 a_{n-1} - 6 a_{n-2} \text{ for } n \geq 2$$

- $a_0 = 1, a_1 = 3, a_2 = 5 a_1 - 6 a_0 = 5*3 - 6*1$ i.e. **$a_2 = 9$**
- And **a_3** and **a_4** ?

Difference

- Continue analyzing the **recurrence relation**

$$a_n = 5 a_{n-1} - 6 a_{n-2}$$

- If we provide a **different initial condition**, for example values **$a_0 = 1$, $a_1 = 3$** , we also get a recursive formula, yet a different one:

$$a_0 = 1, a_1 = 3, a_n = 5 a_{n-1} - 6 a_{n-2} \text{ for } n \geq 2$$

- $a_0 = 1, a_1 = 3, a_2 = 5 a_1 - 6 a_0 = 5 \cdot 3 - 6 \cdot 1$ i.e. **$a_2 = 9$**
- $a_1 = 3, a_2 = 9, a_3 = 5 a_2 - 6 a_1 = 5 \cdot 9 - 6 \cdot 3$ i.e. **$a_3 = 27$**
- $a_2 = 9, a_3 = 27, a_4 = 5 a_3 - 6 a_2 = 5 \cdot 27 - 6 \cdot 9$ i.e. **$a_4 = 81$**
- Yielding sequence: 1, 3, 9, 27, 81, . . . **Powers of 3**
- One recurrence relation used here generates 2 distinct recursive definitions; with **many more** feasible!

Difference

- When recursive formulae focus on mathematical **formulae**, the notions of *recurrence relation* and *recursion* appear very similar; almost identical!
- Key message: **Initial conditions** are not provided for **recurrence relations**!
- Once recursion is applied to say, the parsing of a complex context-free grammar, there is little (but some) resemblance left between recursion and recurrence relations
- For our simple purpose here, the 2 math concepts look quite similar

Recurrence

Recurrence Relations



Recursion? Recurrence? Infinite Regress?

Recurrence Relations

- From [2] Math Insight:
- **Definition of Recurrence Relation**: An equation that **defines a sequence** that can compute the next term as a function of a number of previous terms; **previous** implying at least **one** term
- In its simplest form, a **recurrence relation** depends only on the immediately previous term, i.e. **one single term**
- If we denote the **n^{th} term** in the sequence by x_n , then such a simple recurrence relation is of form:

$$x_{n+1} = f(x_n)$$

for some function **$f()$**

Recurrence Relations

- Here another simple Recurrence Relation: $x_{n+1} = 2 x_n$
- However, recurrence relations can be of higher order, in which case term x_{n+1} not only depends on previous term x_n but also on x_{n-1} , x_{n-2} etc., so could be of form:

$$x_{n+1} = f(x_n, x_{n-1})$$

- or some function $f(n)$ with multiple inputs!
- **Here a practical example:** the recurrence relation

$$x_{n+1} = x_n + x_{n-1}$$

- Simply generates the well known **Fibonacci** sequence; assuming correct initial conditions for x_0 and x_1

Recurrence Relations

- Similar definition:
- A recurrence relation for **sequence $\{a_n\}$** is an equation that expresses **a_n** in terms of *one or more* previous terms of the sequence, for example a_0, a_1, \dots, a_{n-1} , for all **n** , with index **n** being a nonnegative integer
- A sequence is called a **solution** of a recurrence relation if its **terms satisfy** the recurrence relation

Recurrence Relations

- A **recurrence relation** is reminiscent of a recursively defined sequence, but does not specify initial values; AKA: does not specify **initial conditions**
- Therefore, the same recurrence relation can have (and usually has) multiple incarnations (i.e. solutions)
- If both the **initial conditions** and the **recurrence relation** are specified, then the sequence is uniquely determined to evolve into a **recursive formula**

Recurrence Relations

Example:

- Consider **recurrence relation**: $a_n = 2a_{n-1} - a_{n-2}$ for values $n = 2, 3, 4, \dots$
- Is the sequence $\{a_n\}$ with $a_n = 3n$ some **solution** of this recurrence relation?

For $n \geq 2$: $a_n = 2a_{n-1} - a_{n-2} = 2(3(n - 1)) - 3(n - 2)$

$$a_n = 2a_{n-1} - a_{n-2} = 6n - 6 - 3n + 6 = 3n$$

- Therefore, $\{a_n\}$ with $a_n = 3n$ is a solution of the above recurrence relation

Modeling with Recurrence Relations



For example: Compute Compound Interest

Modeling with Recurrence Relations

Example:

- Someone deposits **\$100,000** in a savings account at a bank yielding **5% per year** interest, compounded annually
- How much money will be in the account after 30 years?

Solution:

- Let P_n denote the amount in the account after n years
- How can we determine P_n as a function of P_{n-1} i.e. compute P_n as a function of the previous year?

Modeling with Recurrence Relations

Solution Cont'd: Derive the **interest recurrence relation**:

- $P_n = P_{n-1} + 0.05 P_{n-1} = 1.05 P_{n-1}$
- The initial condition is $P_0 = 100,000$

Then we have:

- $P_1 = 1.05 P_0$
- $P_2 = 1.05 P_1 = (1.05)^2 P_0$
- $P_3 = 1.05 P_2 = (1.05)^3 P_0$
- ...
- $P_n = 1.05 P_{n-1} = (1.05)^n P_0$
- We now have a formula to calculate P_n for any natural number **n** and can avoid iteration altogether!

Modeling with Recurrence Relations

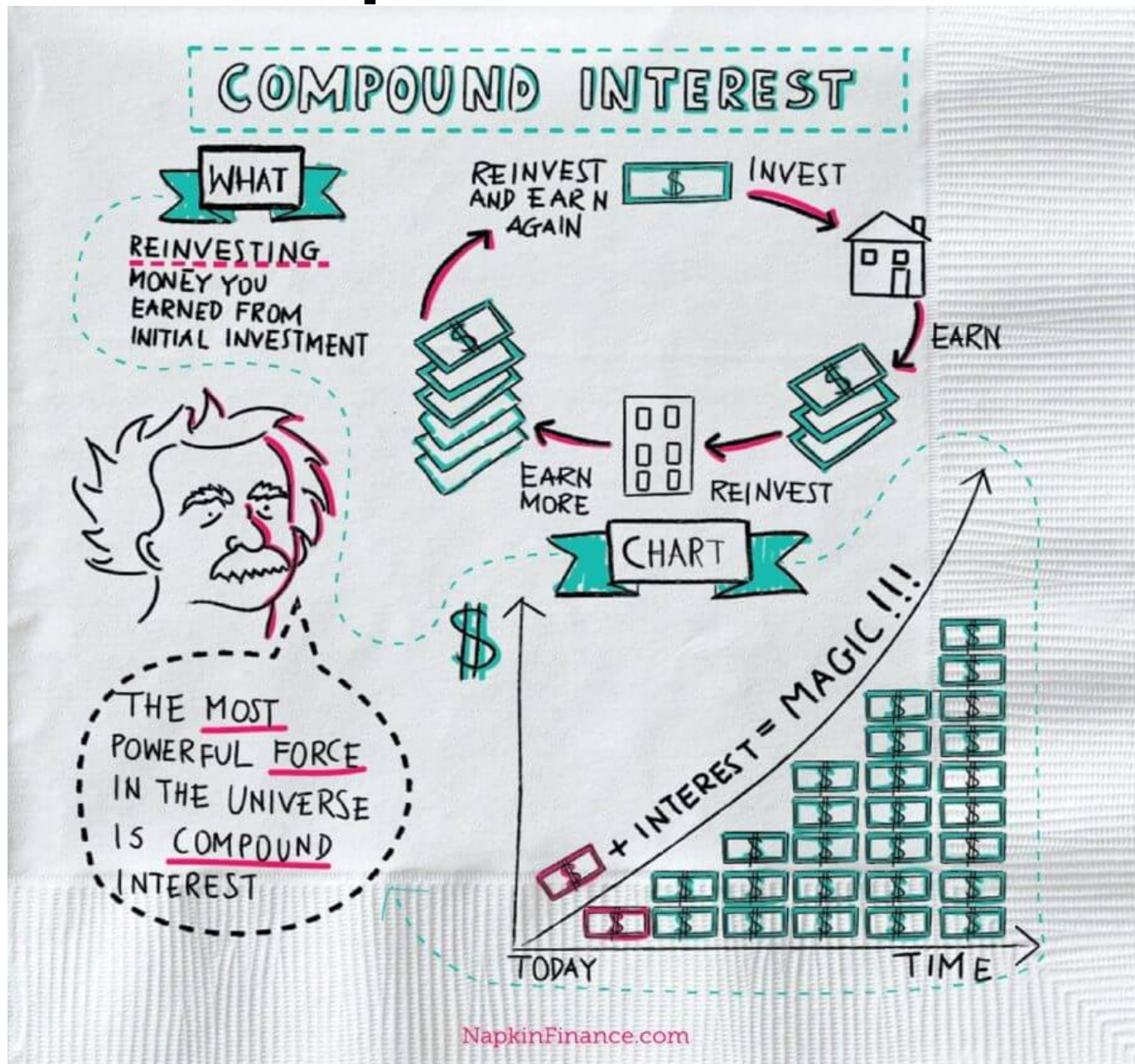
- Use the formula just found to identify P_{30} under initial condition:

$$P_0 = 100,000$$

- $P_{30} = (1.05)^{30} * 100,000 = 432,194.20$
- After 30 years, the account contains \$432,194.20

But who wants to wait that long? ☺

Compound Interest



Modeling with Recurrence Relations

Another example:

- Let a_n denote the number of bit strings of length n that do **not have two consecutive 0s** (“valid strings”)
- String **10001** also excluded: has ≥ 2 consecutive 0s!
- Find a recurrence relation and give initial conditions for the sequence $\{a_n\}$

Solution Idea:

- The number of valid strings is the number of strings ending with a 0 (i.e. the bit before is not also a 0) . . .
- Plus the number of strings ending with a 1 (o.k. if the bit before that ending 1 is a 0)

Modeling with Recurrence Relations

- We don't know the number a_n
- Let us assume that we do know the solution a_{n-1} of valid strings of length $n-1$
- Then how large is the subset of valid strings of length n , if these strings end with digit 1 ?
- There are a_{n-1} such strings, namely the set of all valid strings of length $n-1$ with 1 appended to them
- Whenever we append a 1 to a valid string, that string remains valid, even if there are **two 1s** in a row or a **01** at the end, which is ok

Modeling with Recurrence Relations

- We still need to know: How many valid strings of length n are there that end with digit 0?
- Valid strings of length n ending with 0 must have a 1 as their $(n-1)^{\text{st}}$ bit, otherwise they would end with 00, and that would be invalid
- So what is the total number of valid strings of length $n-1$ that end with digit 1?
- We already know that there are a_{n-1} strings of length n that end with digit 1
- Therefore, there are a_{n-2} strings of length $n-1$ that end with digit 1

Modeling with Recurrence Relations

- So there are a_{n-2} valid strings of length n that end with a 0: i.e. all valid strings of length $n-2$ with 10 appended to them
- Reiterate: The total number of valid strings is the number of valid strings ending with a 0, plus the number of strings ending with a 1
- The solution is the recurrence relation:

$$a_n = a_{n-1} + a_{n-2}$$

Modeling with Recurrence Relations

- What are the initial conditions?

$a_1 = 2$ 0 and 1, neither has 2 consecutive 0s ☺

$a_2 = 3$ 01, 10, and 11

$a_3 = a_2 + a_1 = 3 + 2 = 5$

$a_4 = a_3 + a_2 = 5 + 3 = 8$

$a_5 = a_4 + a_3 = 8 + 5 = 13 \dots$

- This sequence satisfies the same recurrence relation as the **Fibonacci sequence**
- Since $a_1 = f_3$ and $a_2 = f_4$, we have $a_n = f_{n+2}$
- Recall f_i sequence: $f_0 = 0, f_1 = 1, f_2 = 1, f_3 = 2, f_4 = 3, \dots$

Solving Recurrence Relations

- In general, we would prefer an **explicit formula** to compute the value of a_n rather than conducting n iteration steps
- For some class of recurrence relations, we can obtain such formulae in a systematic way
- Those are the recurrence relations that express the terms of a sequence **as a linear combination** of previous terms

Solving Recurrence Relations

- **Definition:** A **linear homogeneous recurrence relation** of degree k with constant coefficients is a recurrence relation of the form:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

- Where c_1, c_2, \dots, c_k are real numbers, and $c_k \neq 0$
- A sequence satisfying such a recurrence relation is uniquely determined by that relation and **k initial conditions**

Solving Recurrence Relations

Examples:

- The recurrence relation $P_n = (1.05)P_{n-1}$ is a linear homogeneous recurrence relation of degree **one**
- Recurrence relation $f_n = f_{n-1} + f_{n-2}$ is a linear homogeneous recurrence relation of degree **two**
 - But due to index **n-2**, not due to the **two** summands ☺
 - Requiring 2 initial conditions
- The recurrence relation $a_n = a_{n-5}$ is a linear homogeneous recurrence relation of degree **five**
 - As it goes back **5** generations of prior solutions
 - Requiring 5 initial conditions

Solving Recurrence Relations

- When solving such recurrence relations, we aim to find solutions of the form $a_n = r^n$, where r is constant
- $a_n = r^n$ is a recurrence relation solution $a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$ if and only if:

$$r^n = c_1 r^{n-1} + c_2 r^{n-2} + \dots + c_k r^{n-k}$$

- Divide this equation by r^{n-k} and subtract the right-hand side from the left:
- $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_{k-1} r - c_k = 0$
- This is called the **characteristic equation** of the recurrence relation

Solving Recurrence Relations

- Solutions of this equation are called the **characteristic roots** of the recurrence relation
- Let us consider linear homogeneous recurrence relations of degree two
- Theorem: Let c_1 and c_2 be real numbers. Suppose that $r^2 - c_1r - c_2 = 0$ has two distinct roots r_1 and r_2
- Then the sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1a_{n-1} + c_2a_{n-2}$ if and only if $a_n = \alpha_1r_1^n + \alpha_2r_2^n$ for $n = 0, 1, 2, \dots$, where α_1 and α_2 are constants

Solving Recurrence Relations

- **Example:** What is the solution of recurrence relation $a_n = a_{n-1} + 2a_{n-2}$ with $a_0 = 2$ and $a_1 = 7$?
- **Solution:** The characteristic equation of the recurrence relation is $r^2 - r - 2 = 0$
- Its roots are $r = 2$ and $r = -1$
- Hence, the sequence $\{a_n\}$ is a solution to the recurrence relation if and only if:
- $a_n = \alpha_1 2^n + \alpha_2 (-1)^n$ for some constants α_1 and α_2

Solving Recurrence Relations

- Given the equation $a_n = \alpha_1 2^n + \alpha_2 (-1)^n$ and the initial conditions $a_0 = 2$ and $a_1 = 7$, it follows that

$$a_0 = 2 = \alpha_1 + \alpha_2$$

$$a_1 = 7 = \alpha_1 * 2 + \alpha_2 * (-1)$$

- Solving these two equations gives $\alpha_1 = 3$ and $\alpha_2 = -1$
- Therefore, the solution to the recurrence relation and initial conditions is the sequence $\{a_n\}$ with

$$a_n = 3 * 2^n - (-1)^n$$

Solving Recurrence Relations

- $a_n = r^n$ is a solution of the **linear homogeneous recurrence relation**

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

- if and only if
- $r^n = c_1 r^{n-1} + c_2 r^{n-2} + \dots + c_k r^{n-k}.$
- Divide this equation by r^{n-k} and subtract the right-hand side from the left:
- $r^k - c_1 r^{k-1} - c_2 r^{k-2} - \dots - c_{k-1} r - c_k = 0$
- This is called the **characteristic equation** of the recurrence relation

Solving Recurrence Relations

- Solutions of this equation are called the **characteristic roots** of the recurrence relation
- Consider linear homogeneous recurrence relations of degree two
- Theorem: Let c_1 and c_2 be real numbers. Suppose that $r^2 - c_1r - c_2 = 0$ has two distinct roots r_1 and r_2
- Then the sequence $\{ a_n \}$ is a solution of the recurrence relation $a_n = c_1a_{n-1} + c_2a_{n-2}$ if and only if $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$ for $n = 0, 1, 2, \dots$, where α_1 and α_2 are constants

Solving Recurrence Relations

- **Example:** Give an explicit formula for **Fibonacci Numbers**

- **Solution:** Fibonacci Numbers satisfy recurrence relation

$$f_n = f_{n-1} + f_{n-2} \text{ with initial conditions } f_0 = 0 \text{ and } f_1 = 1$$

- **The characteristic equation is:** $r^2 - r - 1 = 0$

- **Its roots are** $r_1 = \frac{1 + \sqrt{5}}{2}, \quad r_2 = \frac{1 - \sqrt{5}}{2}$

Solving Recurrence Relations

Therefore, the Fibonacci numbers are given by

$$f_n = \alpha_1 \left(\frac{1 + \sqrt{5}}{2} \right)^n + \alpha_2 \left(\frac{1 - \sqrt{5}}{2} \right)^n$$

for some constants α_1 and α_2 . We can determine values for these constants so that the sequence meets

$$f_0 = 0 \text{ and } f_1 = 1:$$

$$f_0 = \alpha_1 + \alpha_2 = 0 \quad f_1 = \alpha_1 \left(\frac{1 + \sqrt{5}}{2} \right) + \alpha_2 \left(\frac{1 - \sqrt{5}}{2} \right) = 1$$

Solving Recurrence Relations

The unique solution to this system of two equations and two variables is:

$$\alpha_1 = \frac{1}{\sqrt{5}}, \quad \alpha_2 = -\frac{1}{\sqrt{5}}$$

Explicit formula for the Fibonacci numbers:

$$f_n = \frac{1}{\sqrt{5}} \left(\frac{1 + \sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left(\frac{1 - \sqrt{5}}{2} \right)^n$$

Solving Recurrence Relations

- What happens if the **characteristic equation** has just one root?
- How can we match our equation with initial conditions a_0 and a_1 ?
- **Theorem:** Let c_1 and c_2 be real numbers with $c_2 \neq 0$. Suppose that $r^2 - c_1r - c_2 = 0$ has only one root r_0
- A sequence $\{a_n\}$ is a solution of the recurrence relation $a_n = c_1a_{n-1} + c_2a_{n-2}$ if and only if
- $a_n = \alpha_1 r_0^n + \alpha_2 n r_0^n$, for $n = 0, 1, 2, \dots$, where α_1 and α_2 are constants

Solving Recurrence Relations

- **Example:** What is the solution of the recurrence relation $a_n = 6a_{n-1} - 9a_{n-2}$ with $a_0 = 1$ and $a_1 = 6$?
- **Solution:** The only root of $r^2 - 6r + 9 = 0$ is $r_0 = 3$. Hence, the recurrence relation solution is
- $a_n = \alpha_1 3^n + \alpha_2 n 3^n$ for some constants α_1 and α_2
- To match the initial condition, we need

$$a_0 = 1 = \alpha_1$$

$$a_1 = 6 = \alpha_1 * 3 + \alpha_2 * 3$$

- Solving these equations yields $\alpha_1 = 1$ and $\alpha_2 = 1$
- Consequently, the overall solution is given by

$$a_n = 3^n + n * 3^n$$

Summary

- Recursion and recurrence relations share strong similarities
- Yet **recurrence relations** don't provide initial values
- Thus are open to “customization” for any particular formula
- **Recursion** in contrast must define initial conditions
- Moreover, **recursion is not constrained to mathematical** formulae
- Recursion can be applied to complex algorithms; not constraint to formal, mathematical expressions
- Typical applications of recursive definitions in addition to mathematical formulations: **Language definitions**, spec. programming languages!

References

- 1. Recurrence Relation vs. recursion: http://discrete.openmathbooks.org/dmoi3/sec_recurrence.html**
- 2. Recurrence: https://mathinsight.org/definition/recurrence_relation**
- 3. Wiki on recurrence: https://en.wikipedia.org/wiki/Recurrence_relation#Definition**