Database Lab 2 Report

Course: Database Management

Lab Number: Lab 2
Date: 2025-01-29
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1. Objective

• More practice getting around in the PostgreSQL and pgAdmin environments.

- Familiarize yourself with the CAP database data.
- · Get more easy lab points.

2. Lab Setup

Create the CAP database in PostgreSQL.

3. Procedure

Part 1: Querying CAP

Execute the following queries (one at a time), take a screenshot, and verify its correctness.

select *
from People;

```
[19:20] zsh
~
) sudo -u postgres psql
[sudo] password for ryan:
psql (17.2)
Type "help" for help.
postgres=# \c cap
You are now connected to database "cap" as user "postgres".
cap=# select *
cap-# from People;
pid | prefix | firstname | lastname | suffix | homecity |
                                                                     dob
                                          | Piano Man | Oyster Bay | 1949-05-09
  1 | Mr.
              | Billy
                         Joel
  2 | Ms.
              Renee
                           Rosnes
                                                     | Regina | 1962-03-24
              | Elton
  3 | Sir
                           John
                                          | Esq.
                                                     | Pinner
                                                                 1947-03-25
                                                     | Pinner
  4 | Mr.
              | Reginald | Dwight
                                                                | 1947-03-25
              | Michael | McDonald
  5 | Mr.
                                                    | St. Louis | 1952-02-12
              Ray
  6 | Mr.
                           Charles
                                          I MD
                                                     | in Georgia | 1930-09-23
              | Stevie | Wonder
  7 | Dr.
                                          | Ph.D.
                                                     | Saginaw | 1950-01-12
 8 | Ms. | Yuja | Wang (∭∭) |
10 | Dr. (Hon) | Diana | Krall |
                                                     | Beijing
                                                                | 1987-02-10
                                                     Nanaimo
                                                                 | 1960-11-16
(9 rows)
```

```
select *
from Customers;
```

```
select *
from Agents;
```

```
select *
from Products;
```

```
cap=# select *
cap-# from Products;
prodid |
                              city
                                       | qtyonhand | priceusd
                name
      | Kurzweil PC2R
                              | Dallas
                                               47 I
                                                       67.76
p01
       | Yamaha CP-80
                                                       51.50
p02
                              Newark
                                             2399
p03
      | Apple //+
                             Duluth
                                             1979
                                                       65.02
      | LCARS module
                             Duluth
p04
                                                3 I
                                                       17.01
      | Roland 808
                             | Dallas |
                                           8675309 I
                                                       16.61
p05
p06
      | PDP-11 operator panel | Beijing |
                                               88 I
                                                       88.00
      | Flux Capacitor
                             Newark
                                             1007 I
                                                        1.00
p07
       | HAL 9000 memory chip | Newark |
                                                        1.25
p08
                                              200 |
       | Oberheim OB-Xa
p09
                              | Regina |
                                                1 | 37900.42
(9 rows)
```

```
select *
from Orders
```

cap=# select *					
<pre>cap-# from Orders; ordernum dateordered</pre>	custid	agentid	prodid	quantityordered	totalusd
	-+	-+	-+	+	-+
1011 2024-01-22	1 1	2	p01	1100	58794.00
1012 2023-01-23	4] 3	l p03	1200	76096.81
1015 2022-01-23	1 5	3	p05	1000	15771.20
1016 2021-01-23	1 7	3	p01	1000	66404.80
1017 2023-02-14	1	3	p03	500	25643.98
1018 2023-02-14	1	3	p04	600	8050.49
1019 2023-02-14	1	2	p02	400	16249.28
1020 2023-02-14	4	6	p07	600	585.18
1021 2023-02-14	4	6	p01	1000	66086.33
1022 2023-03-15	1	3	p06	450	31236.48
1023 2023-03-15	1	2	p05	500	6550.98
1024 2023-03-15	5	2	p01	880	56671.55
1025 2022-04-01	1 7] 3	p07	888	870.24
1026 2022-05-04	1 7	6	p03	808	47277.29
(14 rows)					

Part 2: Keys

Explain the distinctions among the terms primary key, candidate key, and superkey.

A superkey is any field or set of fields in a table that can uniquely identify every row in said table. It often has unnecessary attributes.

A candidate key is a minimal superkey. Like the superkey, it can uniquely identify each row in the table. "Minimal superkey" means that it is the superkey with a small set of columns/attributes required to ensure the uniqueness.

The primary key is the candidate key selected to provide uniqueness. The primary key is the most minimal candidate and can be used for table relations. The primary keys should not be null nor should they be defined by the user. We should use artificial keys like CWID wherever possible.

Summary: superkey is any set of fields that ensures uniqueness, candidate key is minimal superkey, and primary key is the best candidate key.

Part 3: Data Types

Write a short essay on data types. Select a topic for which you might create a table. Name the table and list its fields (columns). For each field, give its data type and whether or not it is nullable

Data types good. Essay done. Despite the obvious value of data types, we need to appreciate that they restrict the kind of data that can be stored in each column of a table. The data in each column fits specific format and is comparable. Common data types in databases include integers, floating-point numbers, strings, dates, and booleans.

Topic: Job applicants (its always on my mind...)

Applicants: (This isn't a DB table its a reverse table describing the columns)

Field Name Data Type Nullable Description

Field Name	Data Type	Nullable	Description
applicant_id	INT	No	Unique applicant identifier (Primary Key)
first_name	STRING	No	First name
last_name	STRING	No	Last name
email	STRING	No	Email
phone_number	STRING	Yes	Phone - not required
position_applied	STRING	No	Position applied for
application_date	DATE	No	Date application was submitted
cover_letter	TEXT	Yes	Cover letter - optional

Part 4: Relational Rules

1. First Normal Form

1. This rule dictates that there are to be no multi-valued attributes or values with internal structure at any intersection of a row and a column in a table (original description... definitely not the slides talking). All values at the intersection must be atomic (cannot be divided further). An example of this would be a column for email that has someone's home and work emails both listed separated by a comma. This field is not atomic, and should be divided into "home_email" and "work email" attributes.

2. Access Rows by Content Only

1. This rule states that data should be queried based on *what it is* (content) rather than *where is is* (row/col). For example, querying "What is the first name of pid 007?" is valid because this is content based and will work no matter how the table is ordered. However, asking, "What is the name in the first row?" references an arbitrary and temporary table location. As new data is added, the table is ordered, or defragmentation occurs, the table rows will shift positions. Your location based query will no longer work (or even worse - return a different record!).

3. All Rows must be Unique

1. This rule declares (without trumpets) that all rows need to be distinguishable from every other row. The uniqueness of rows is essential to data integrity and querying. For example, if two rows were the same, it would be unclear which row a query actually pertained to. This rule is not as heavily enforced, as we saw in class. However, restrictions such as primary keys will often prevent duplicate rows from being created.